



Stopping Safely Hierarchical Distributed Components: Application to GCM

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Agenda

- **Introduction and objectives**
 - Assumptions
 - Characteristics of the algorithm
- **Description of algorithm**
 - Example
- **Conclusion**
 - Extension and futures works



Introduction

- When is a system reconfigurable?
- How to stop a hierarchical components assembly?
 - An algorithm for reaching a *safe state*.
- First phase of the reconfiguration of components.



Objectives

Stopping a distributed component system.

- Without incoherencies in the system.
- Each inner components is in a *safe state* after the process.
- Stopping a system without livelock and deadlock.
- Don't block the system if it can't be stopped safely.



Assumptions

- Components do not share memory.
- Components communicate by asynchronous requests.
- Communications are performed over some bindings.
- It is possible to mark messages.
- It is possible to identify the *safe state* of each component.
- Requirements
 - Absence of deadlocks.
 - Absence of livelocks.
 - Fairness.



Characteristics of the algorithm

- Not specialized for a given configuration.
- Supports replies (futures).
- If the system can't be stopped safely, the algorithm never finishes.



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Algorithm description

- The *master component* receives initially the stop request.
- First phase: preparation
 - Only the *master component*.
 - The *master* marks all requests it sends.
- Second phase: stopping system.
 - The *master* blocks the requests it receives.

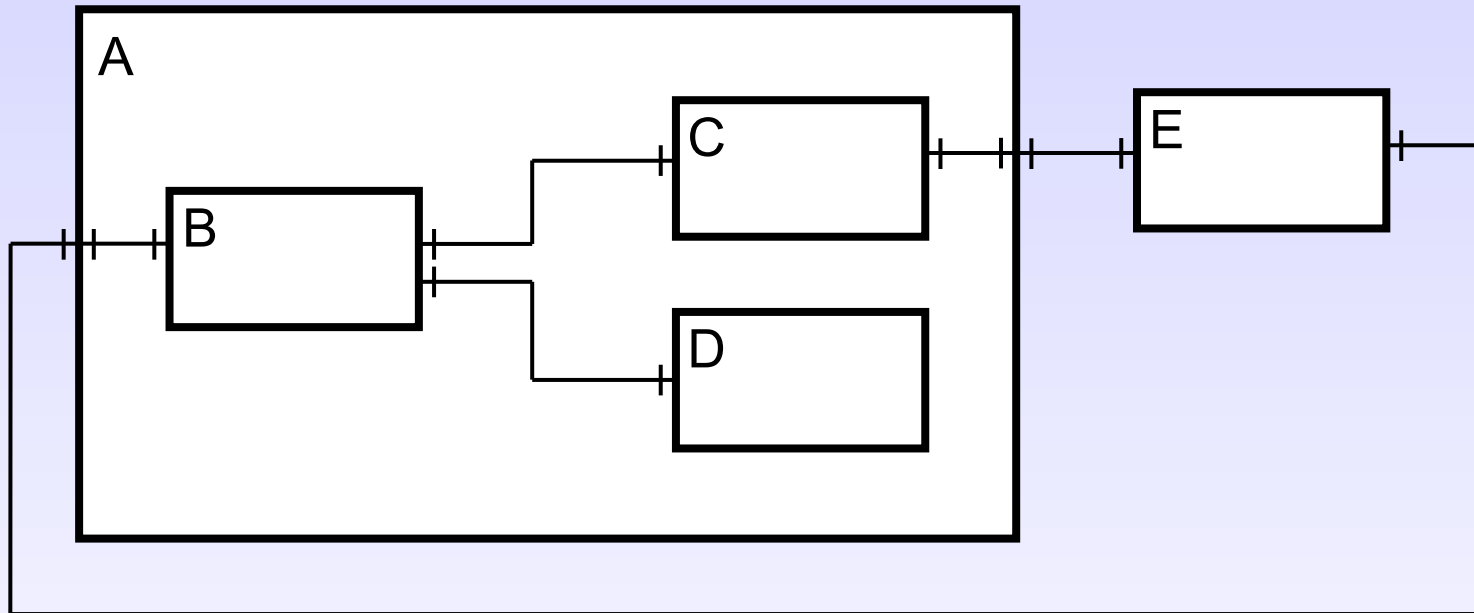


Algorithm description

- States of the components:
 - Started.
 - Wait for stopping (only *master component*).
 - Stopping.
 - Ready to stop. (R2S)
 - Stopped.



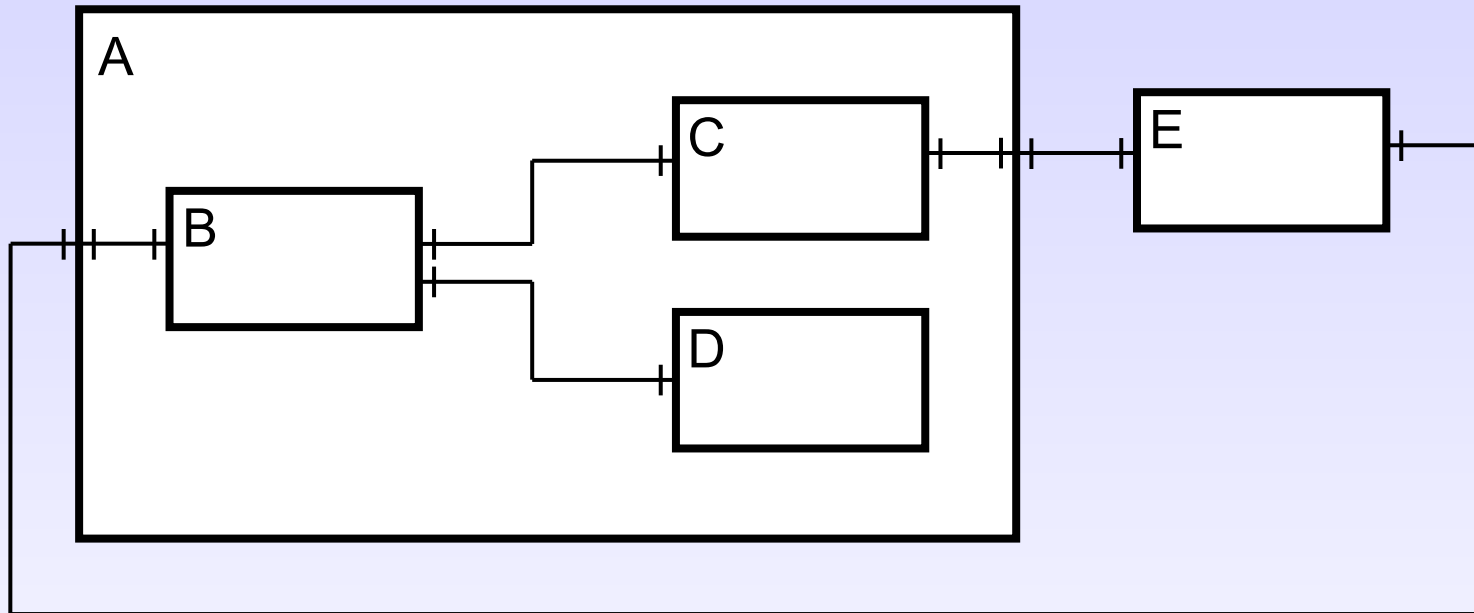
Example





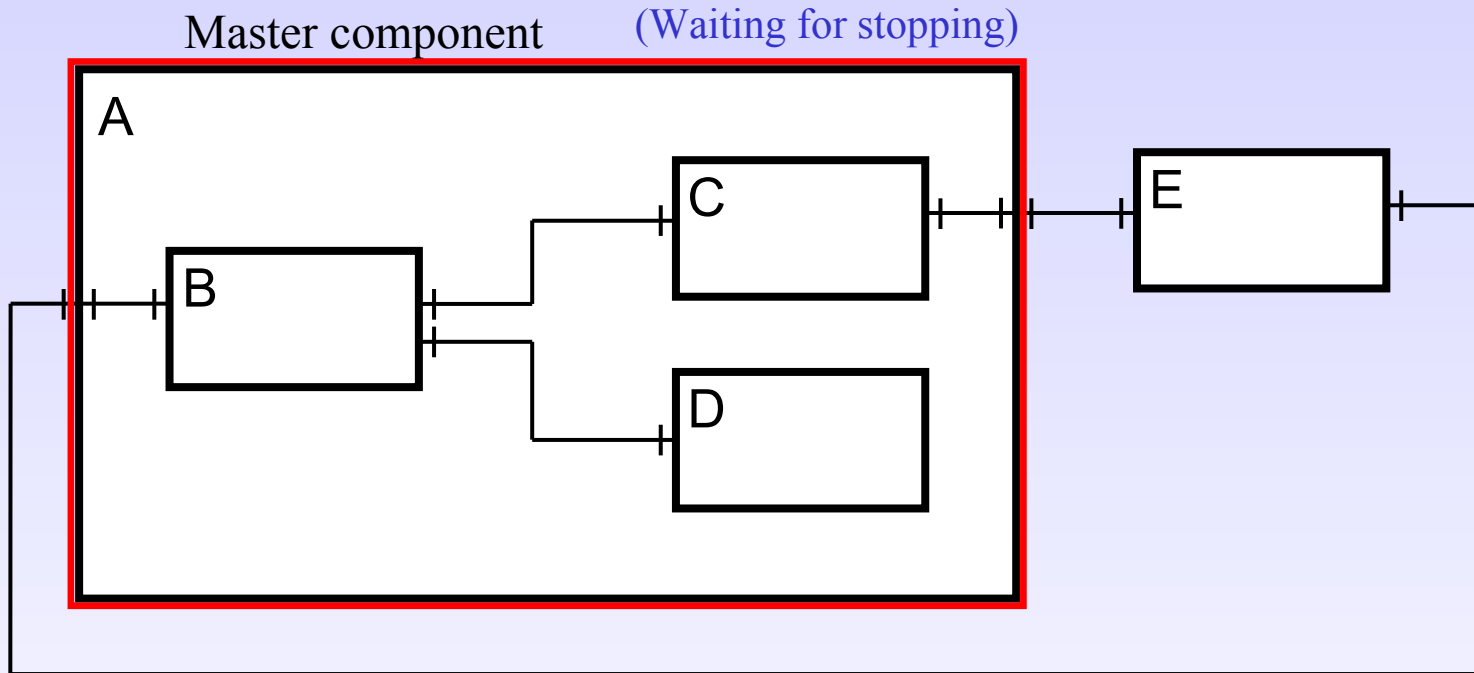
Example

Stop signal
↓



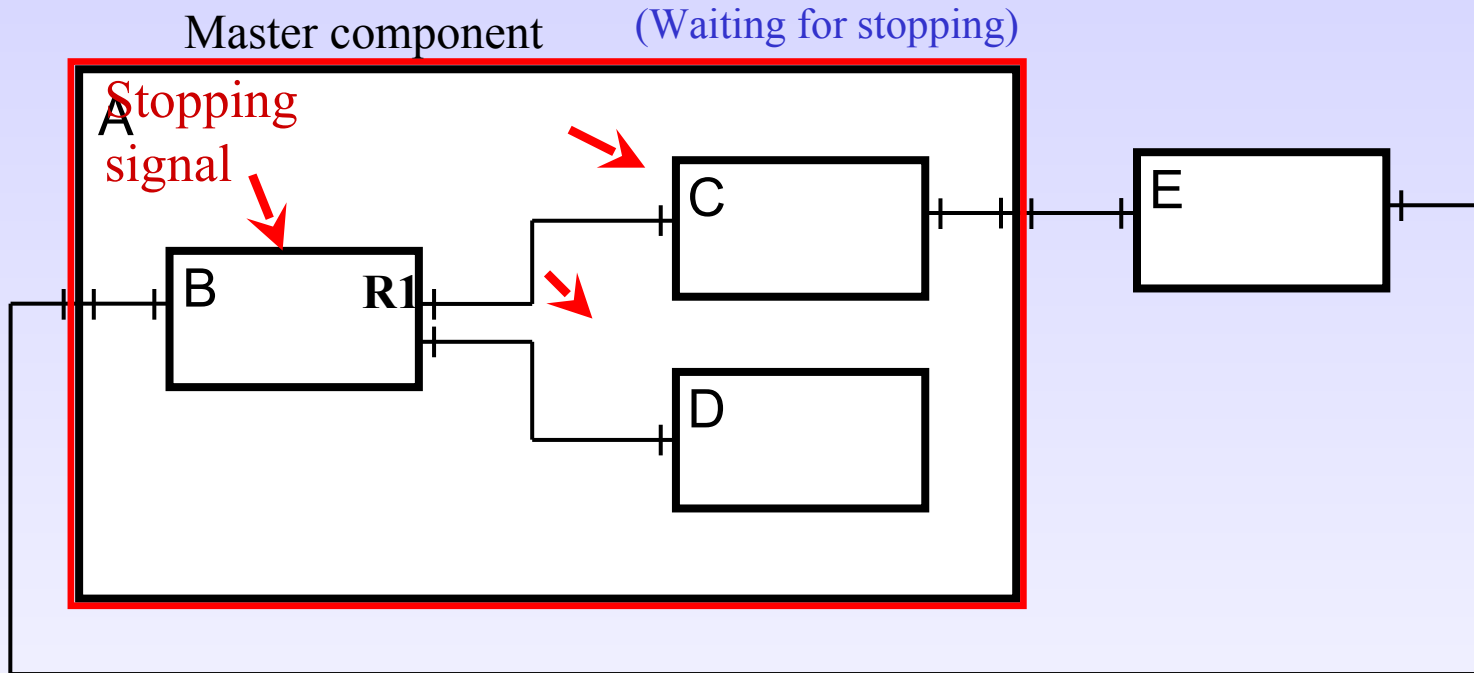


Example



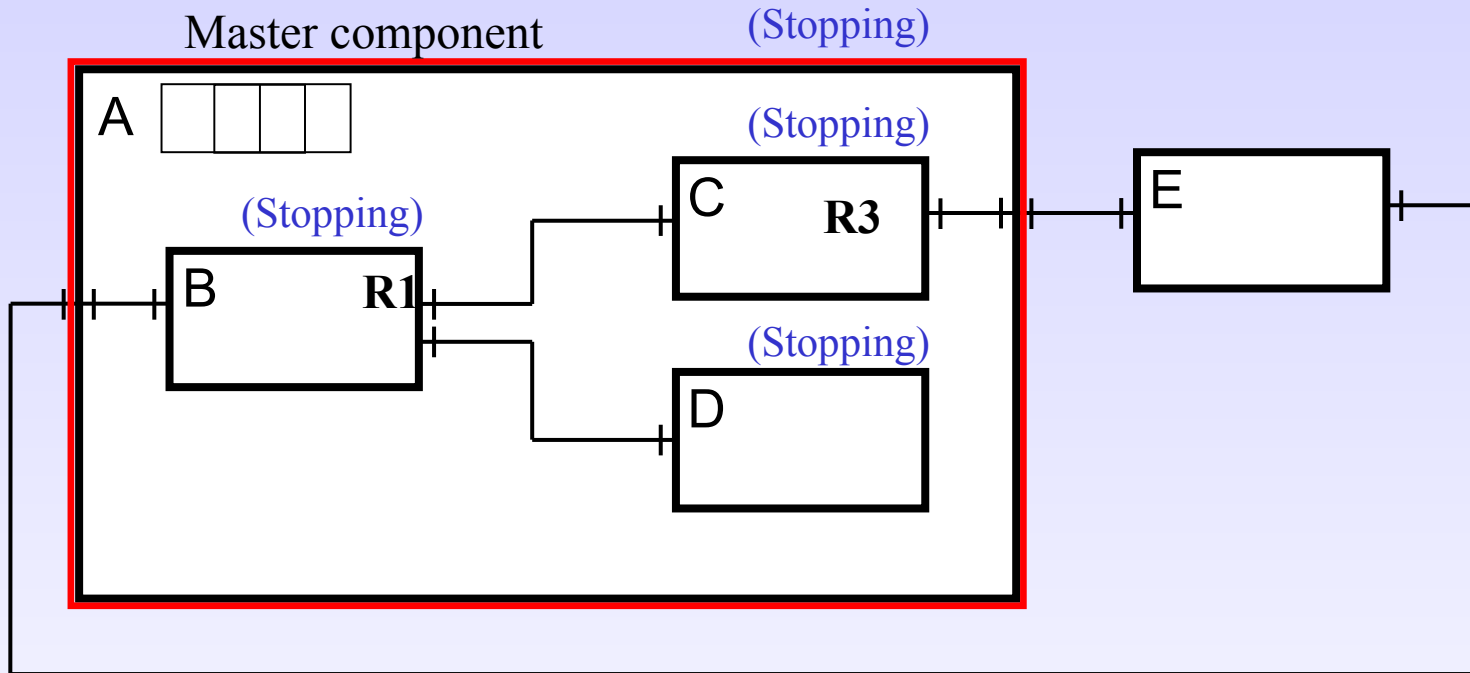


Example



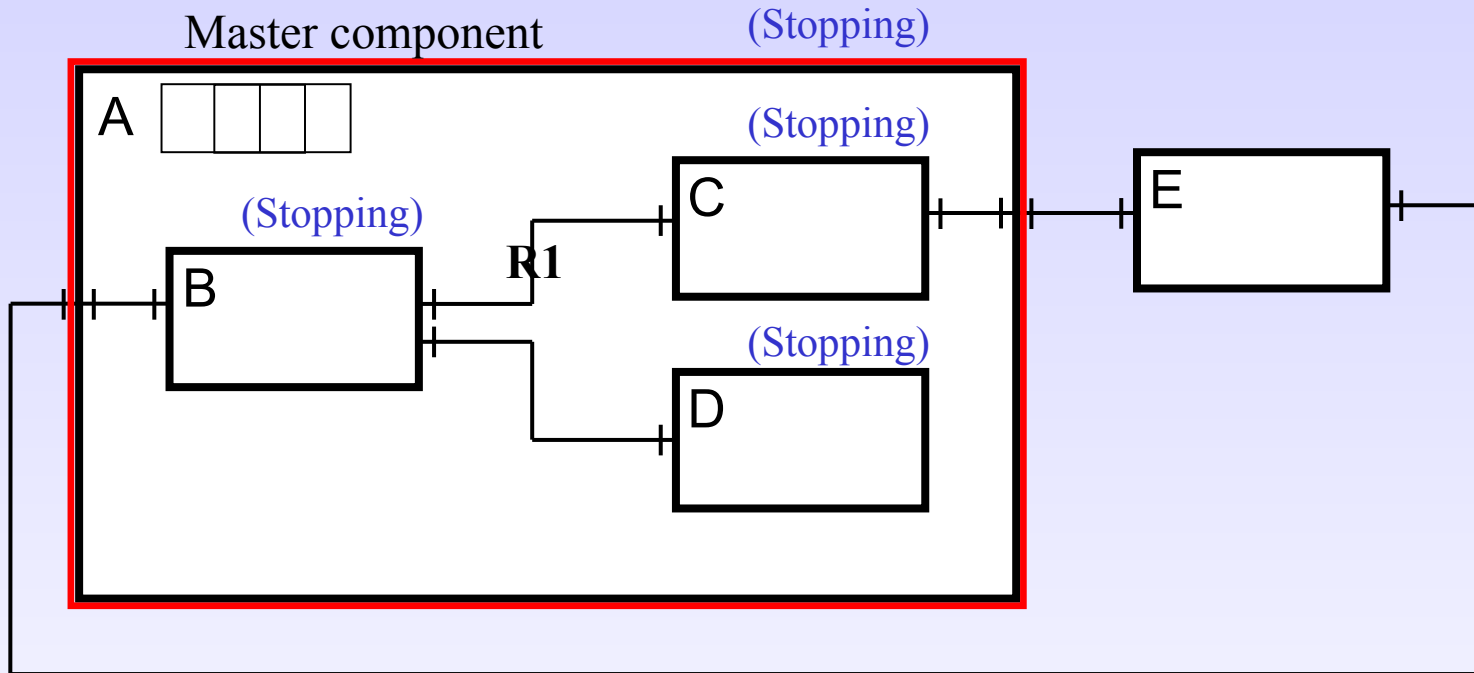


Example



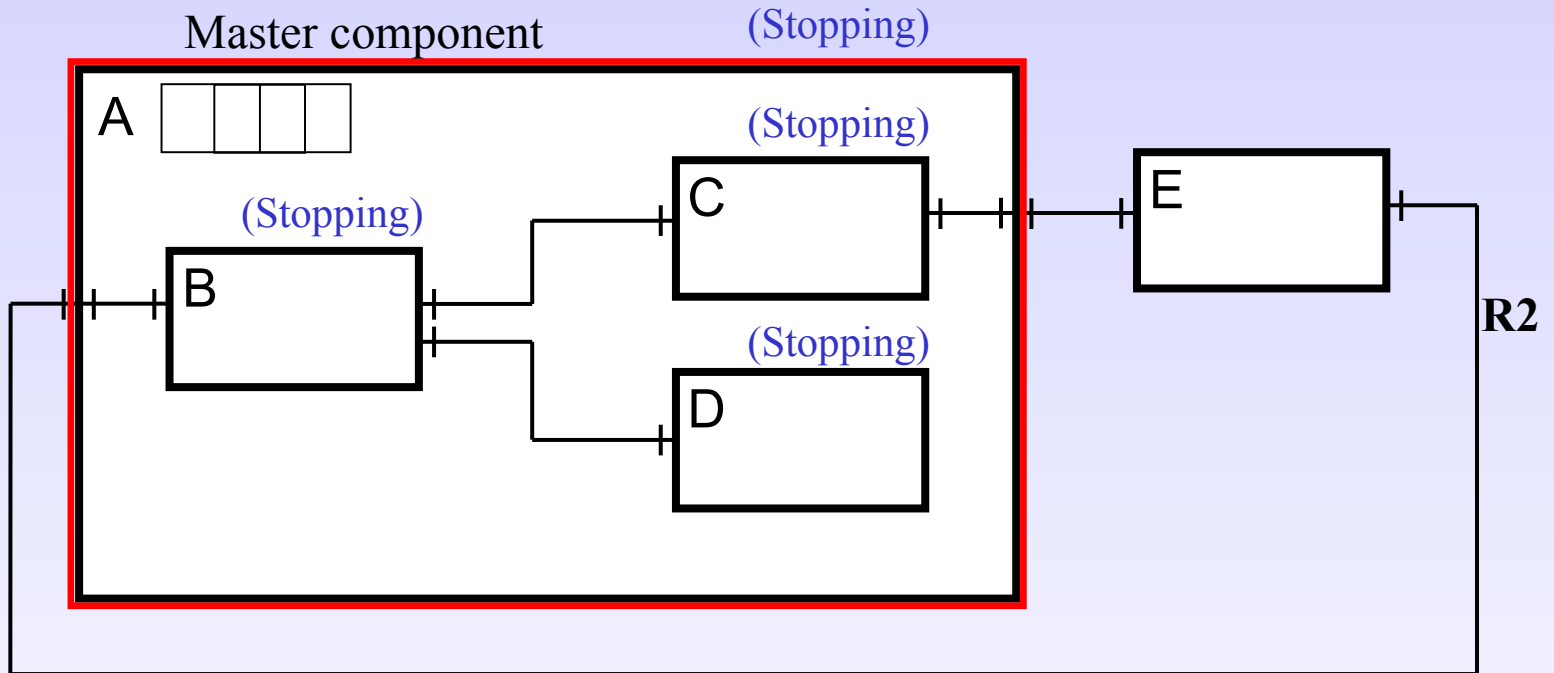


Example



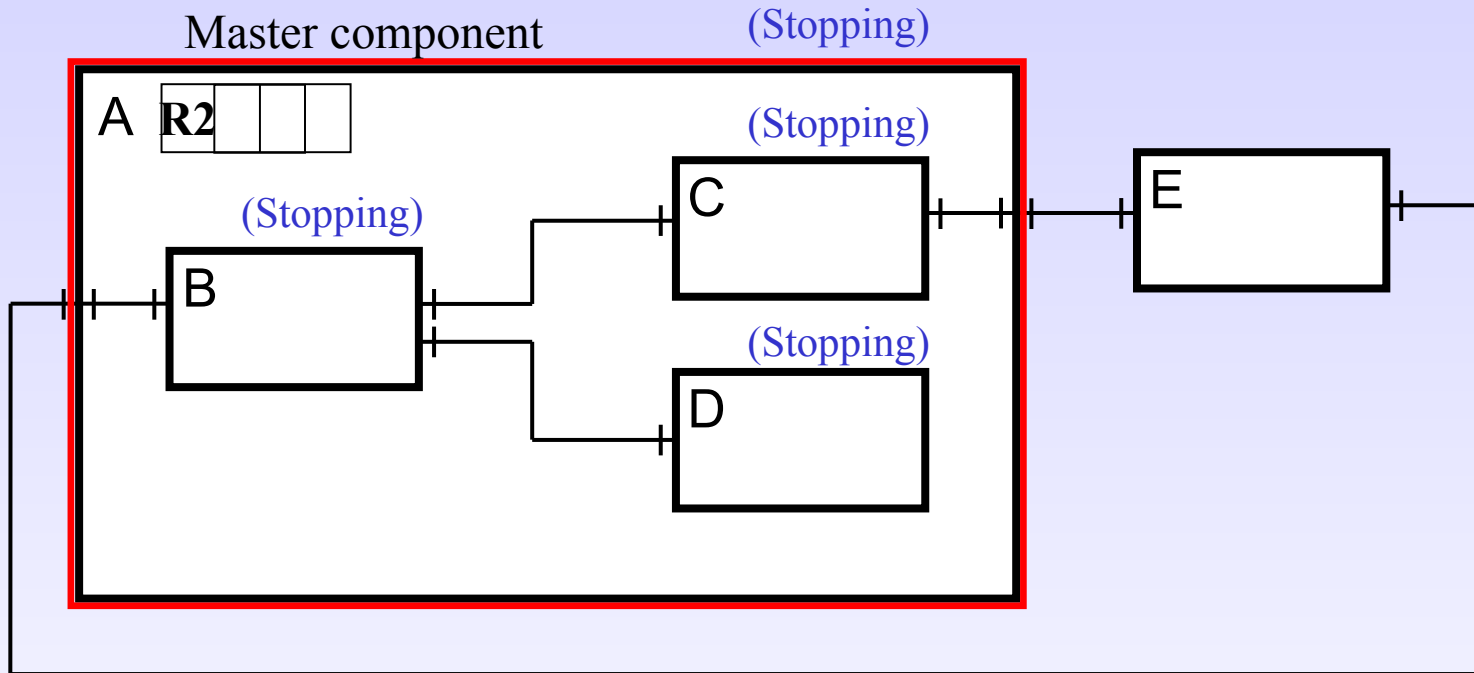


Example



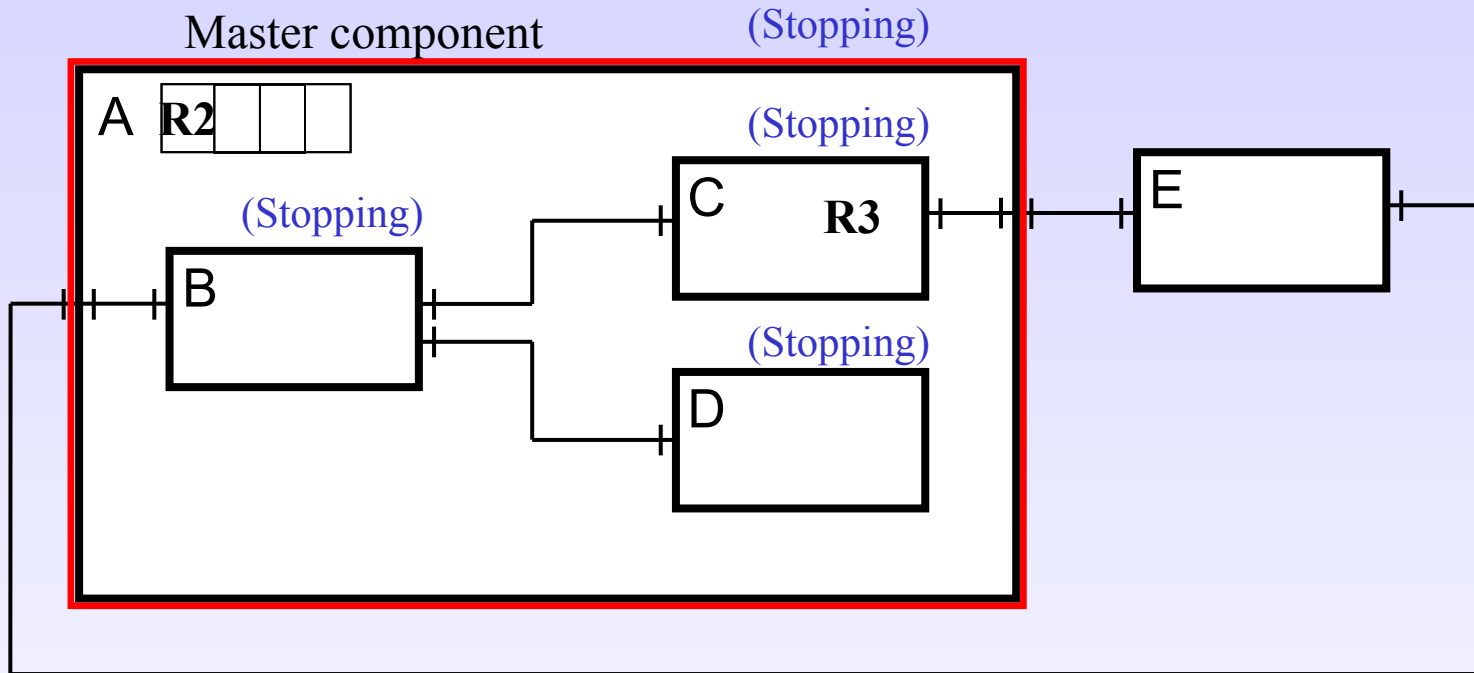


Example



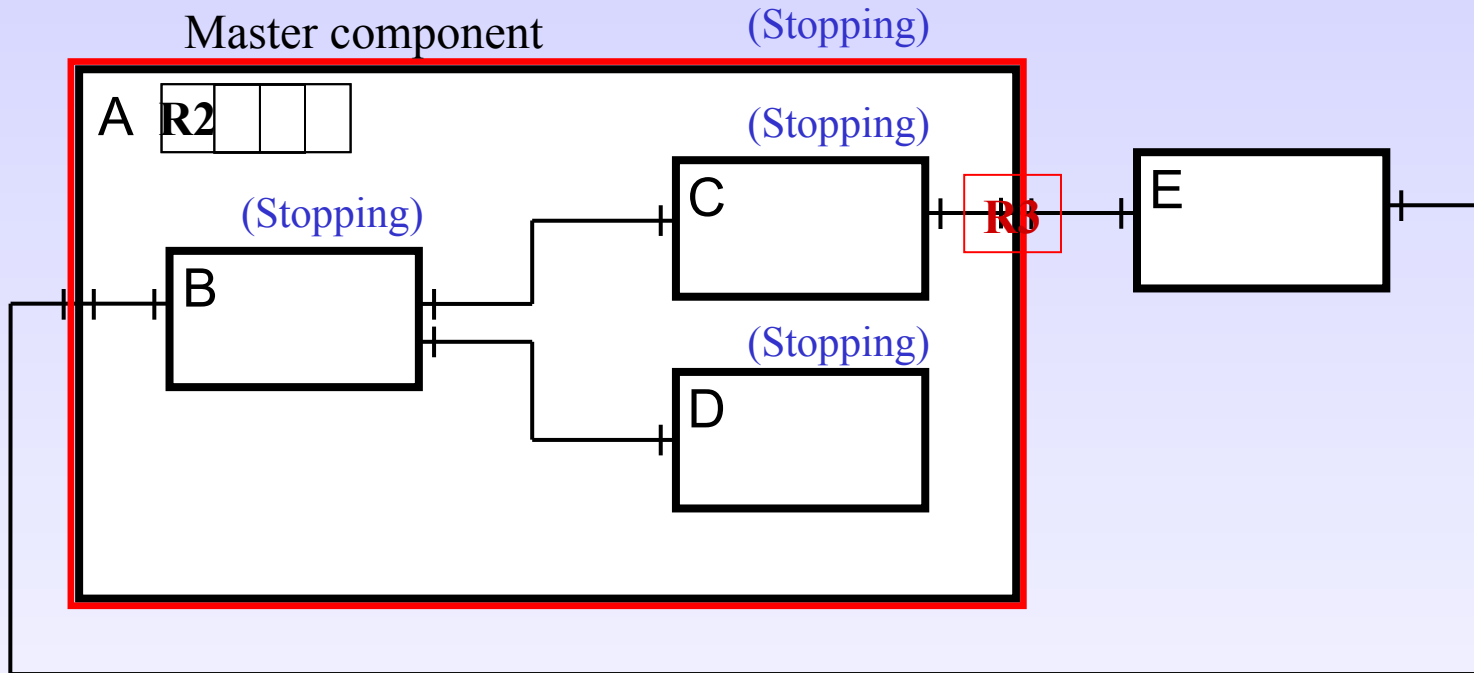


Example



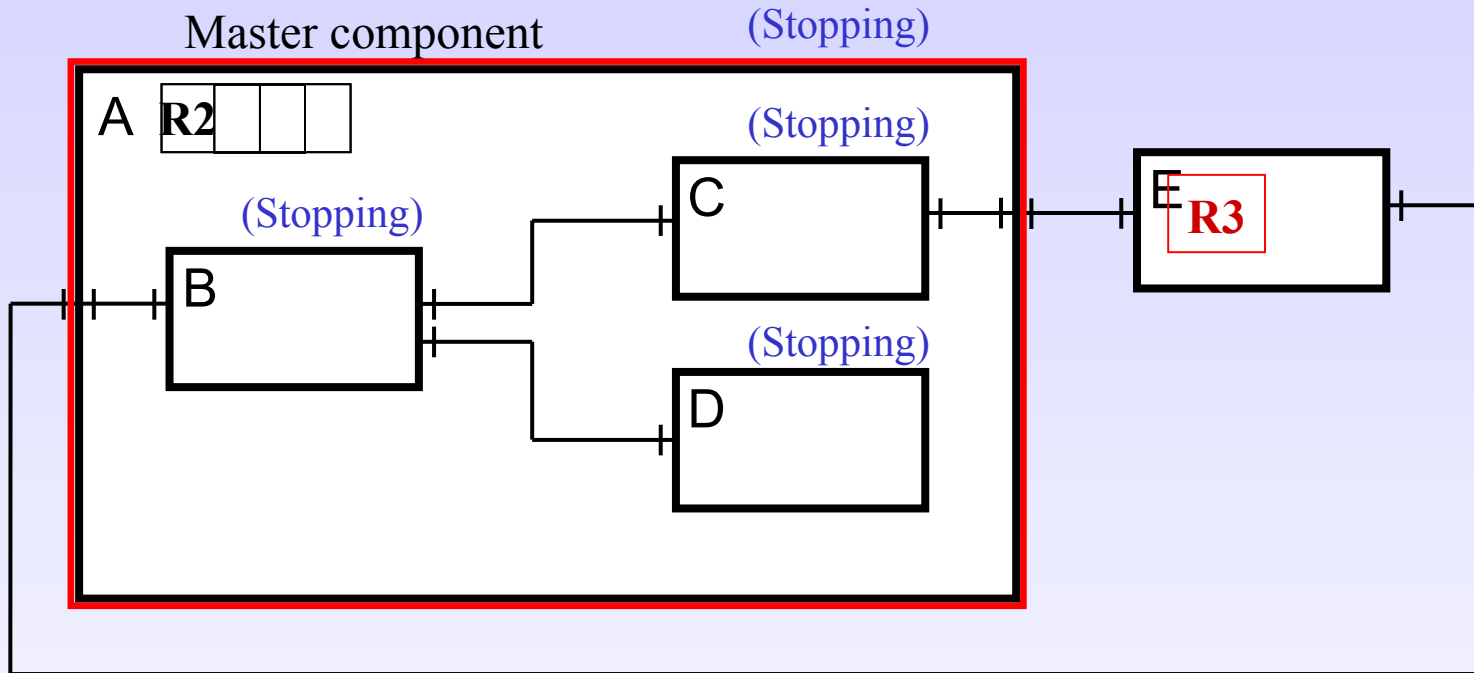


Example



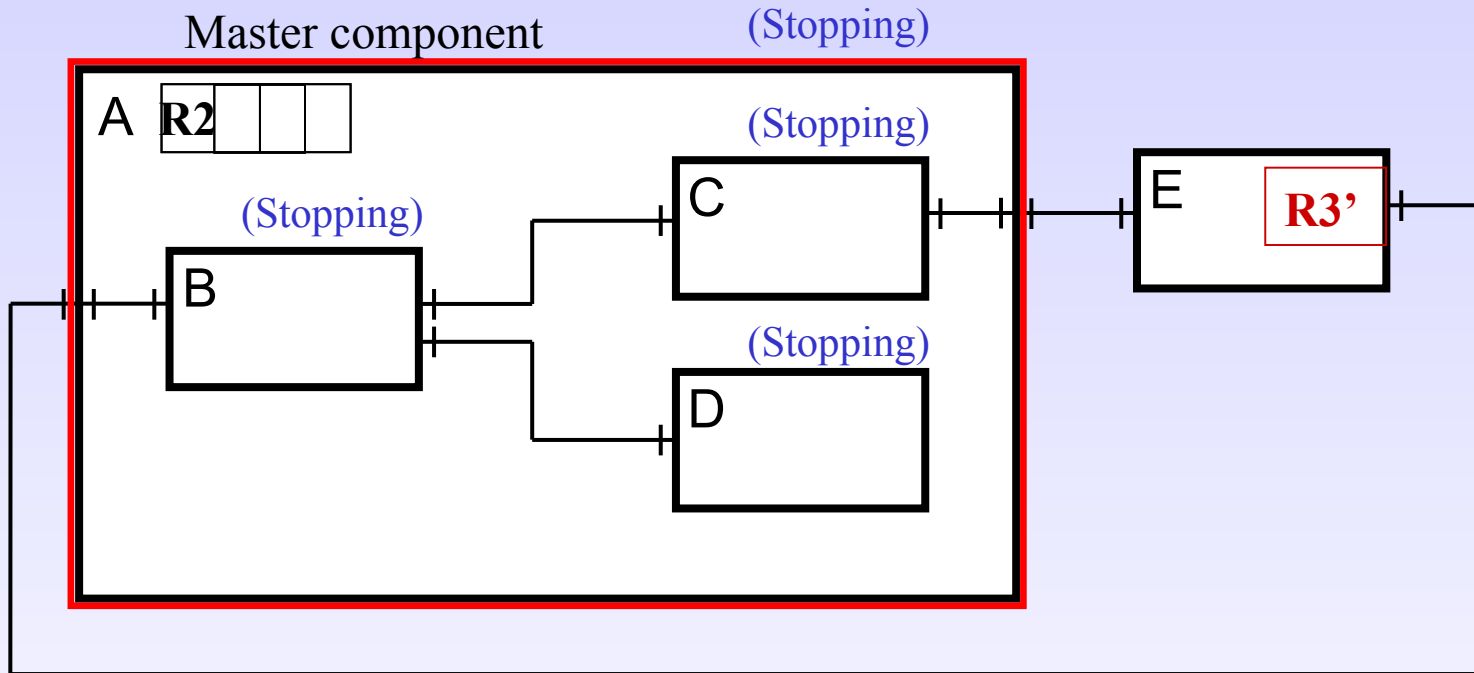


Example



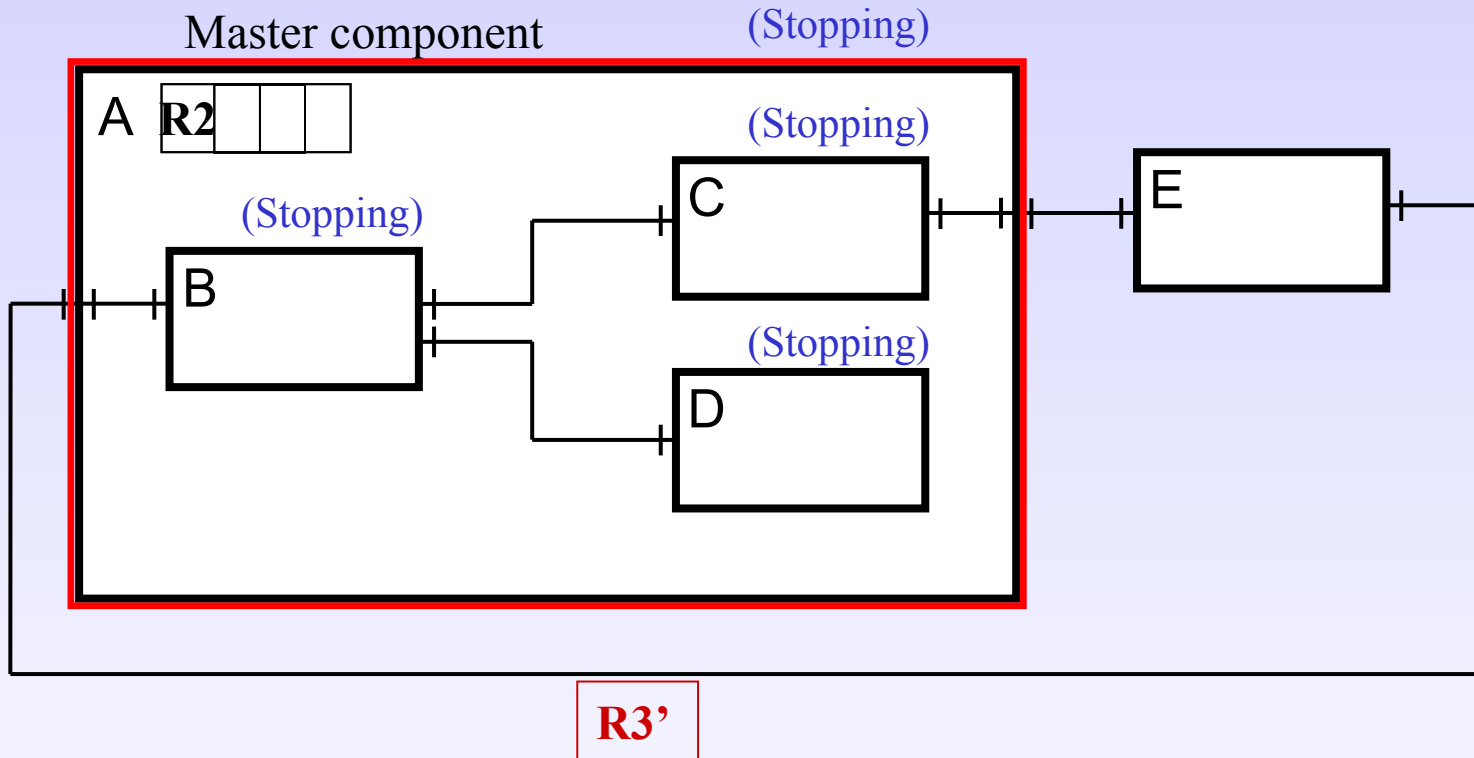


Example



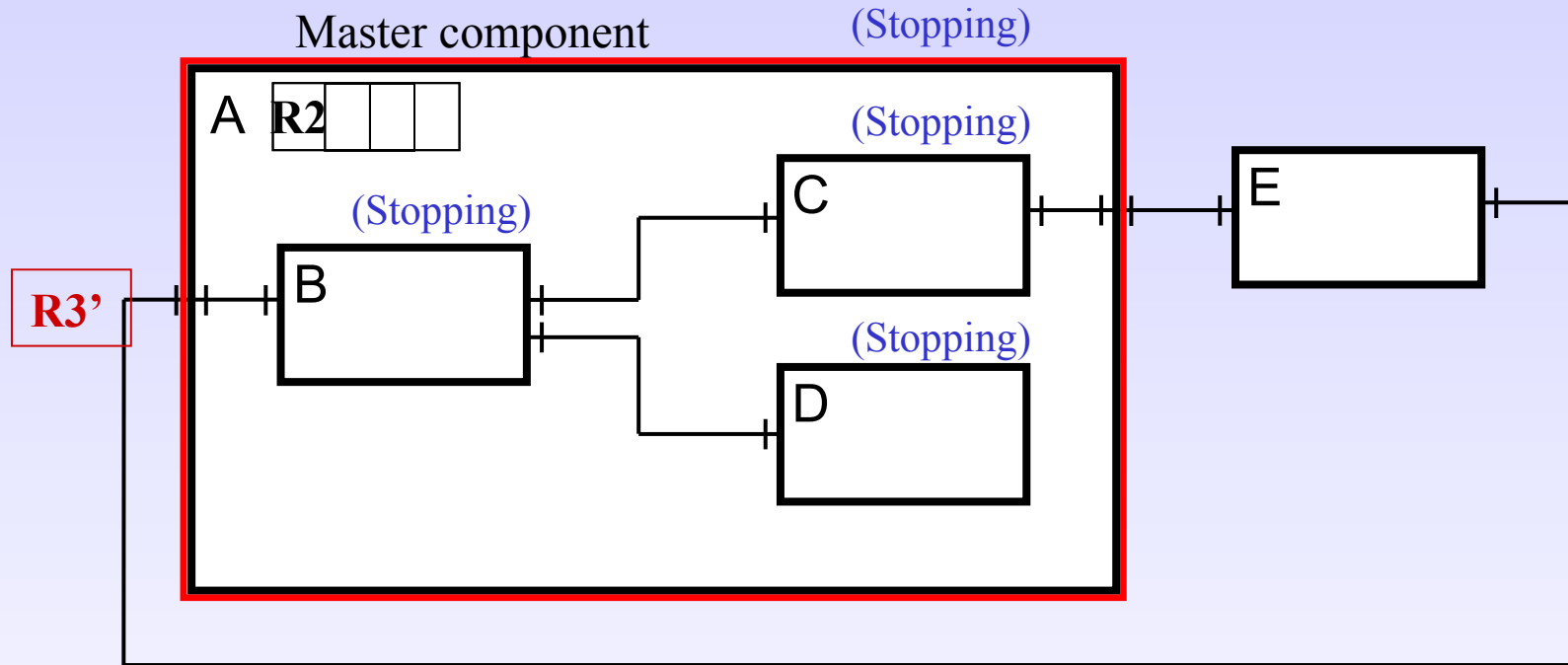


Example



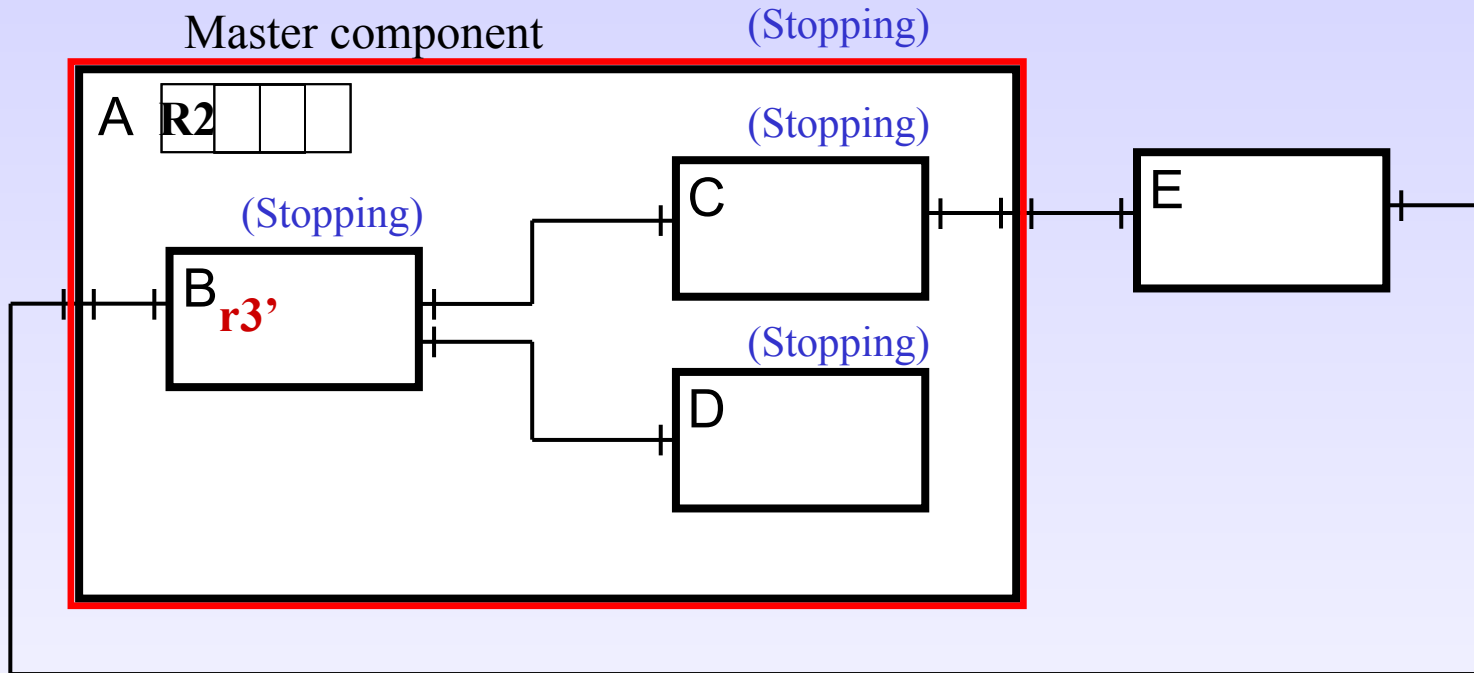


Example



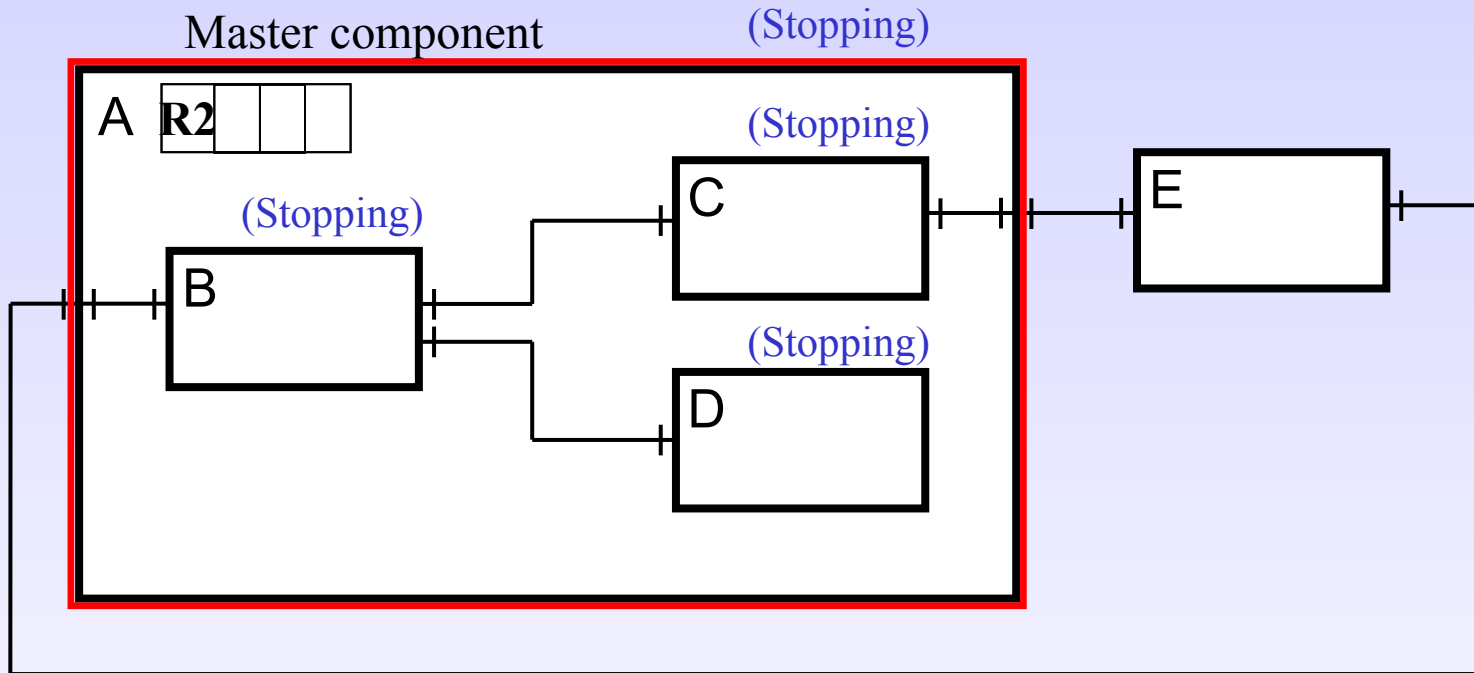


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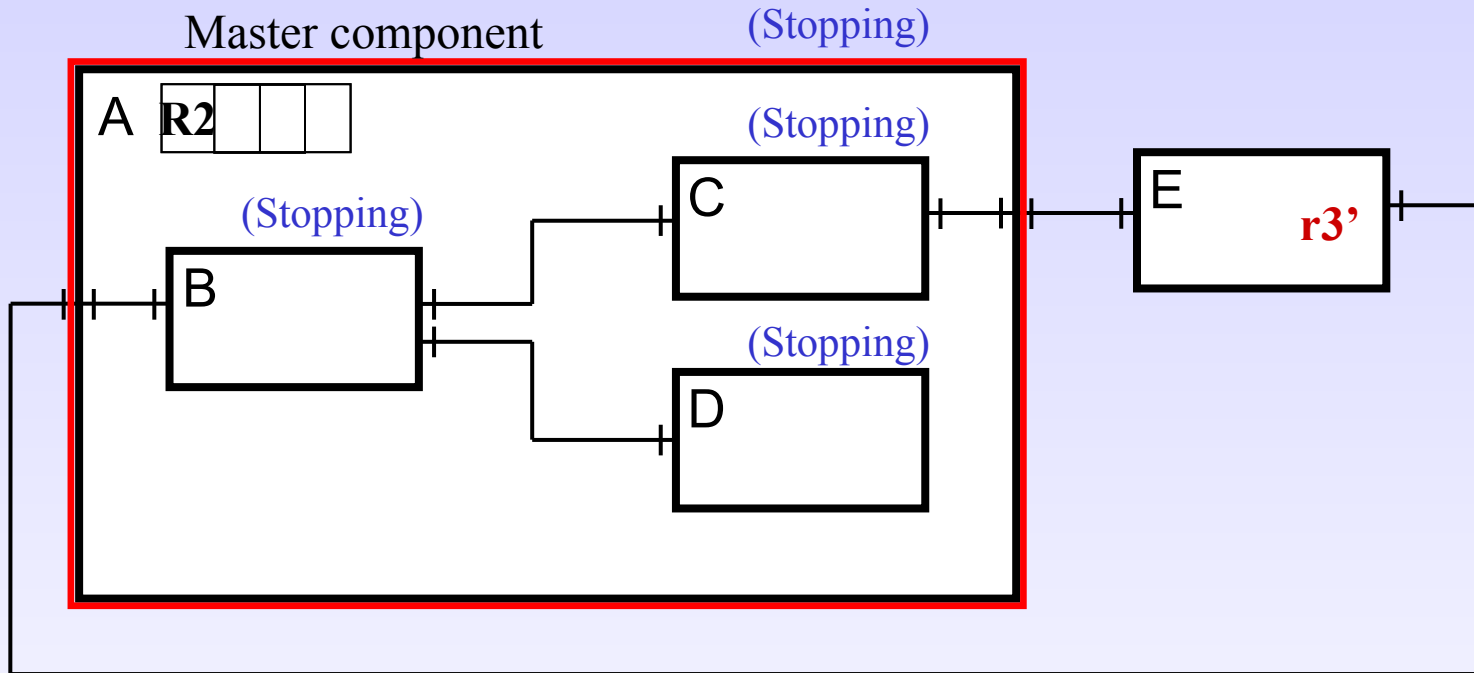
Example



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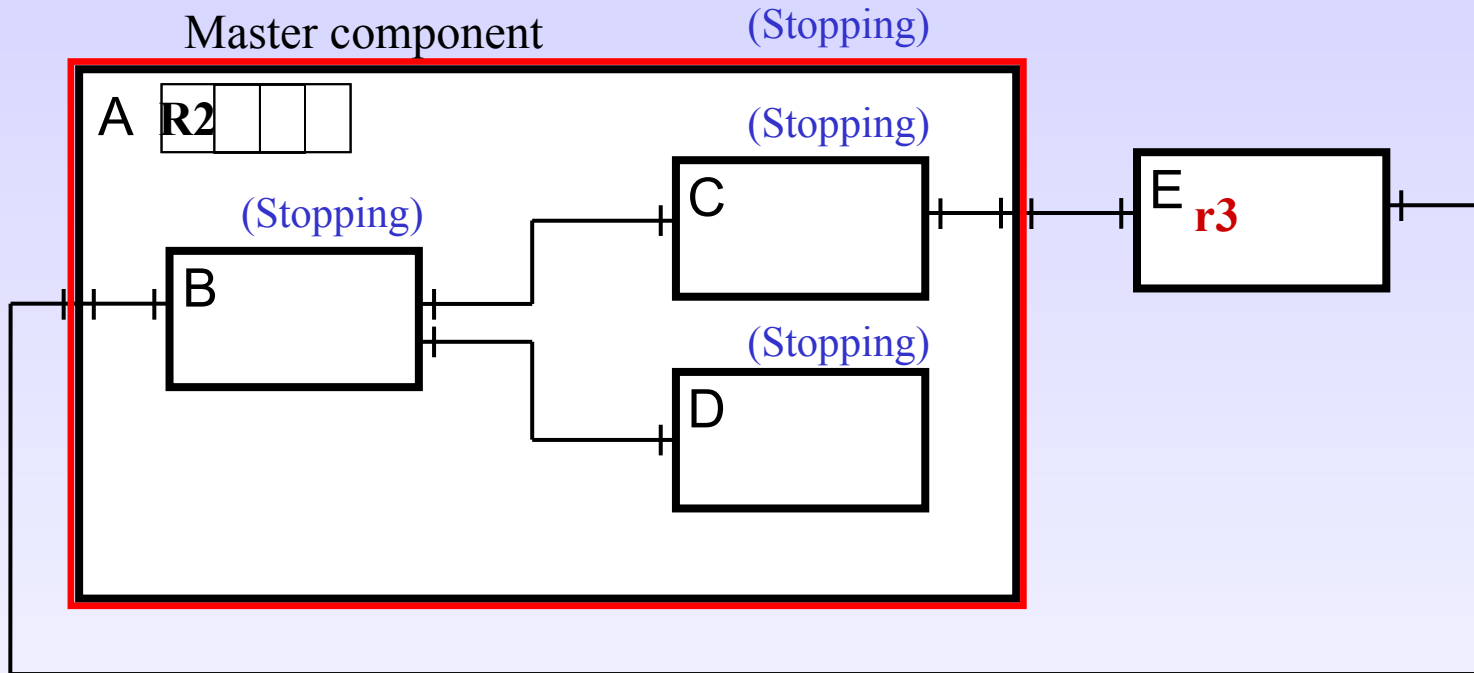


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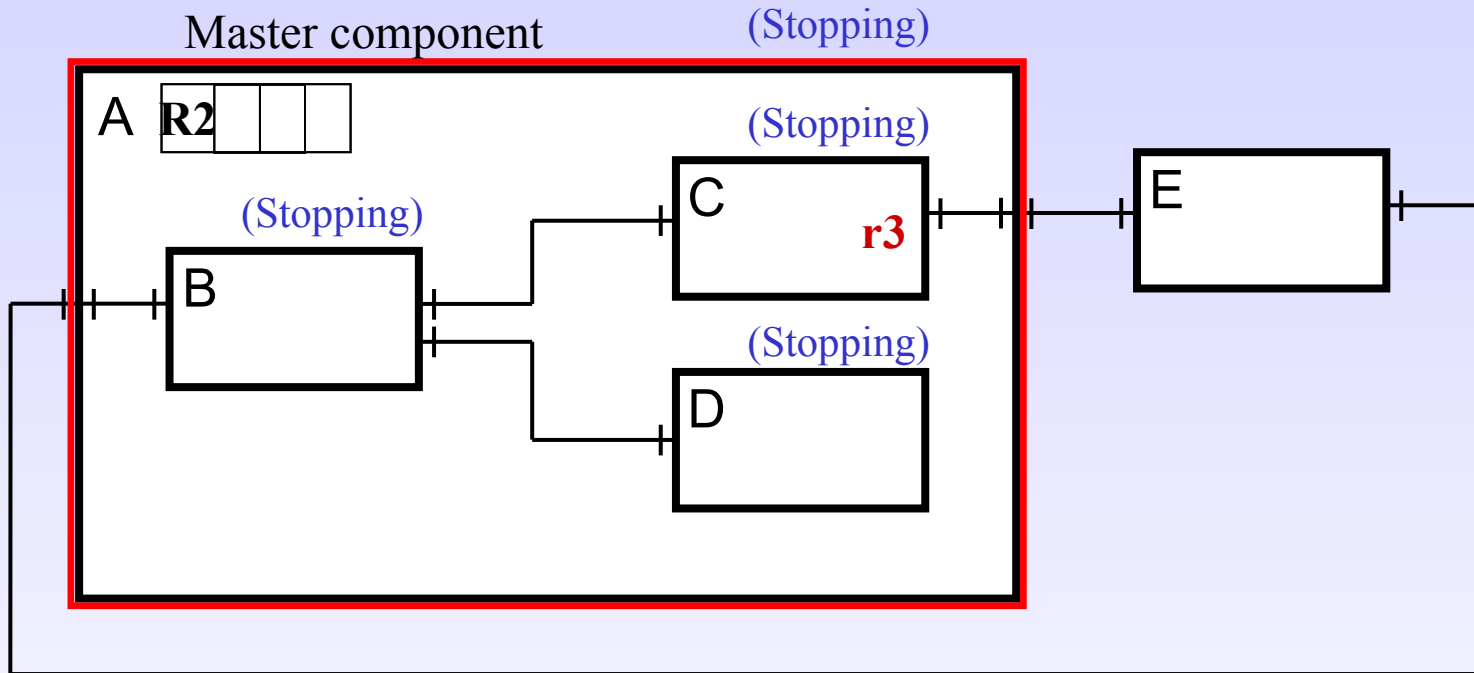


Example



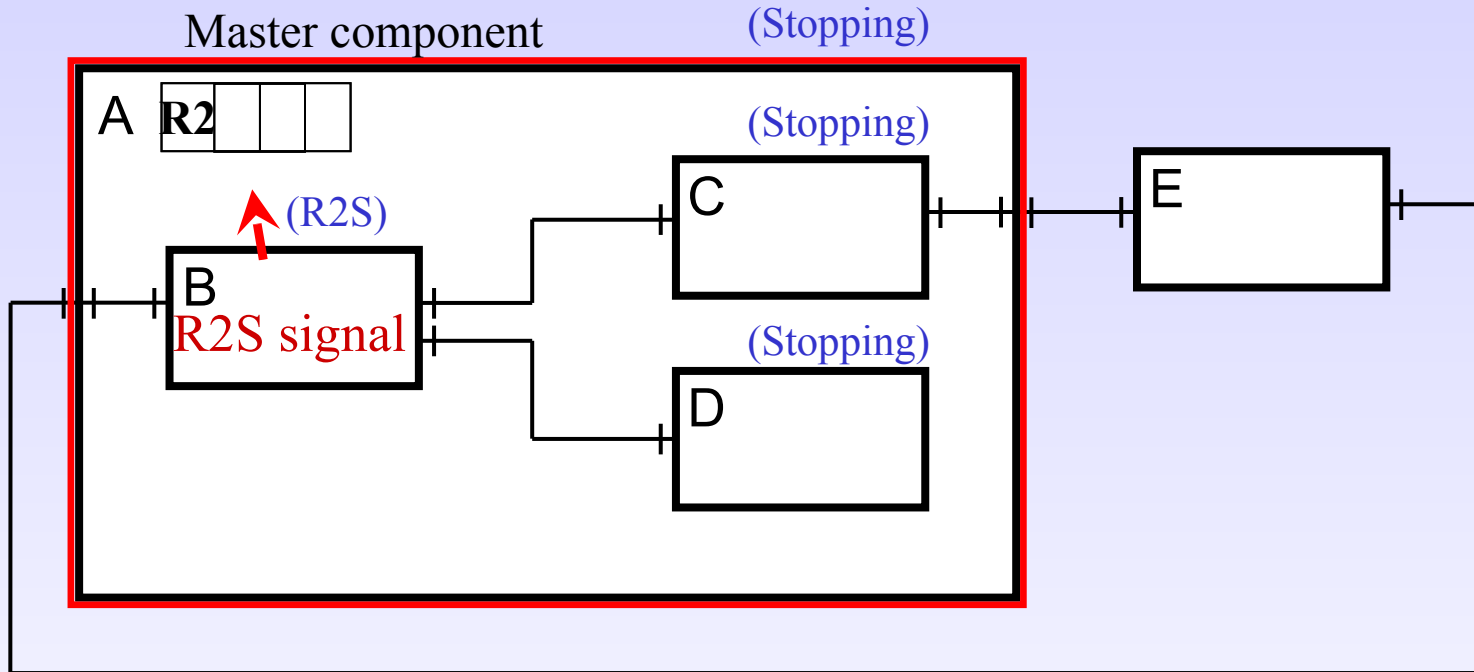


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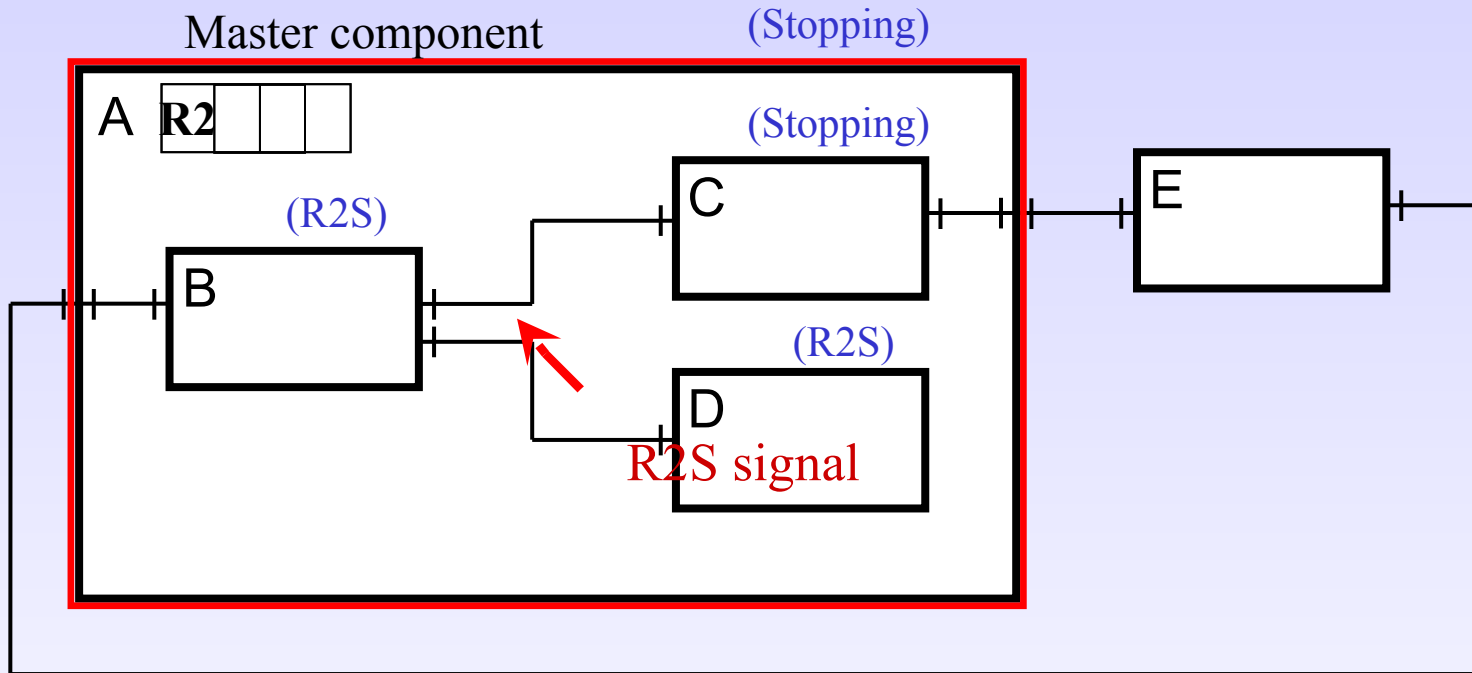


Example



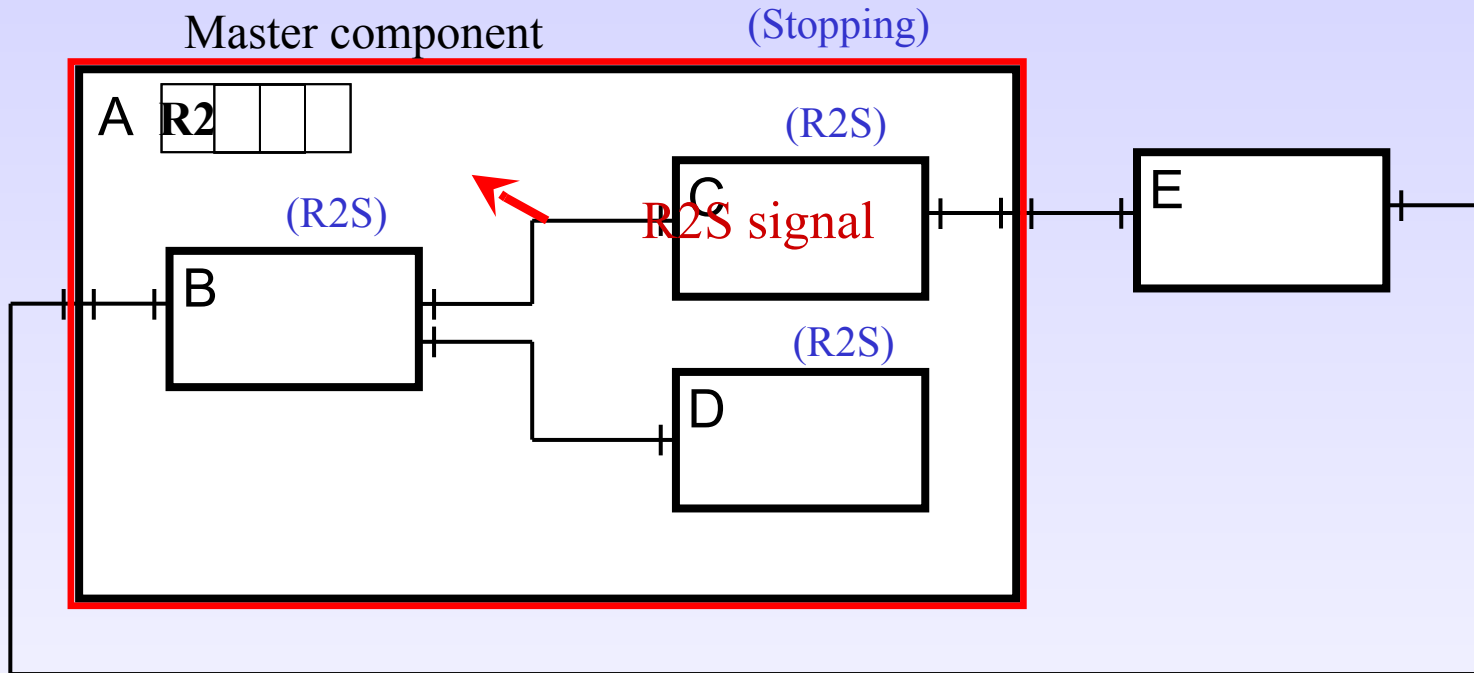


Example



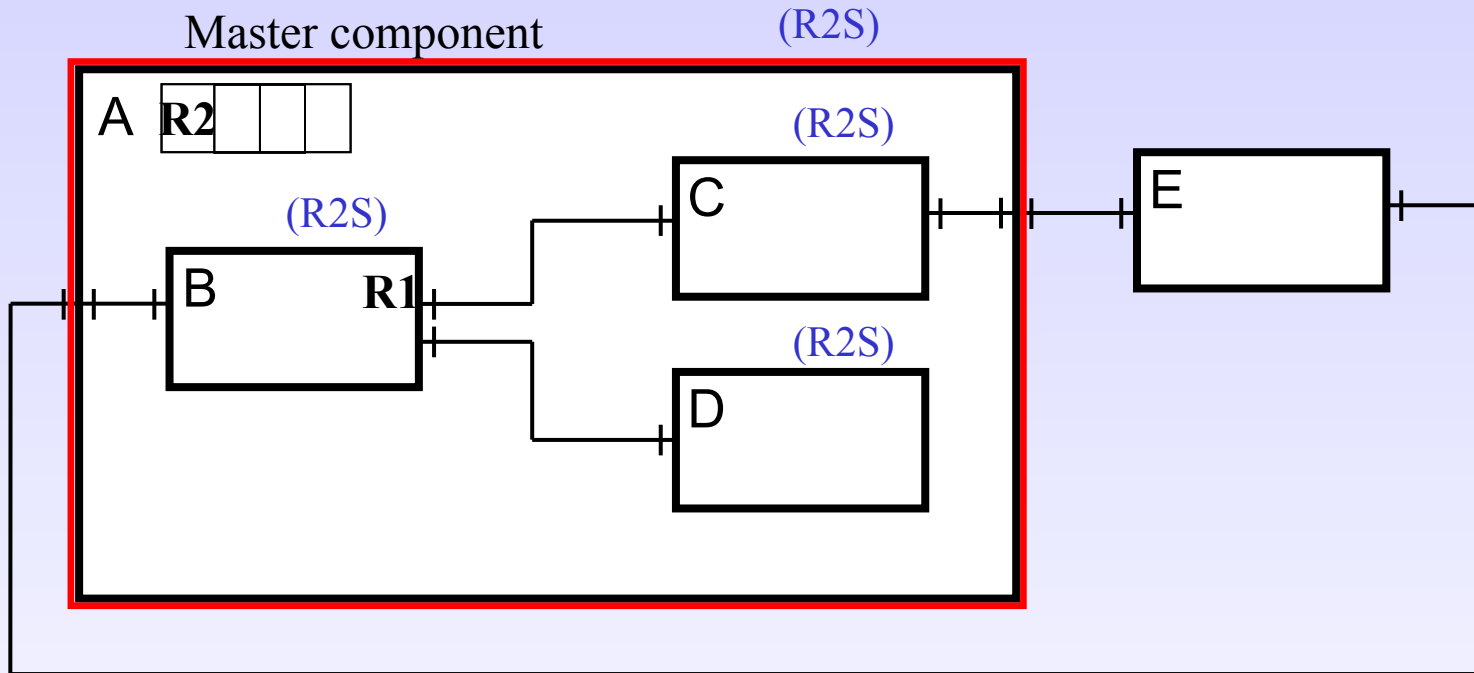


Example



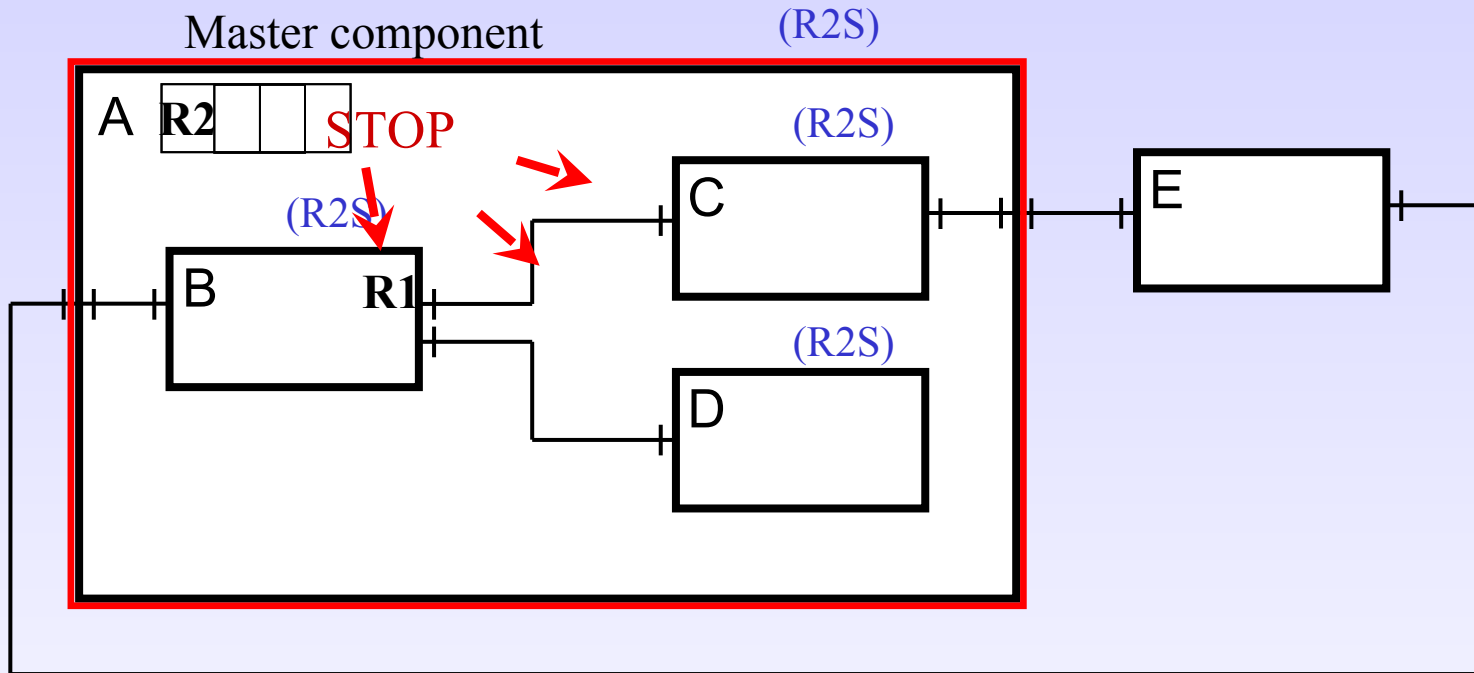


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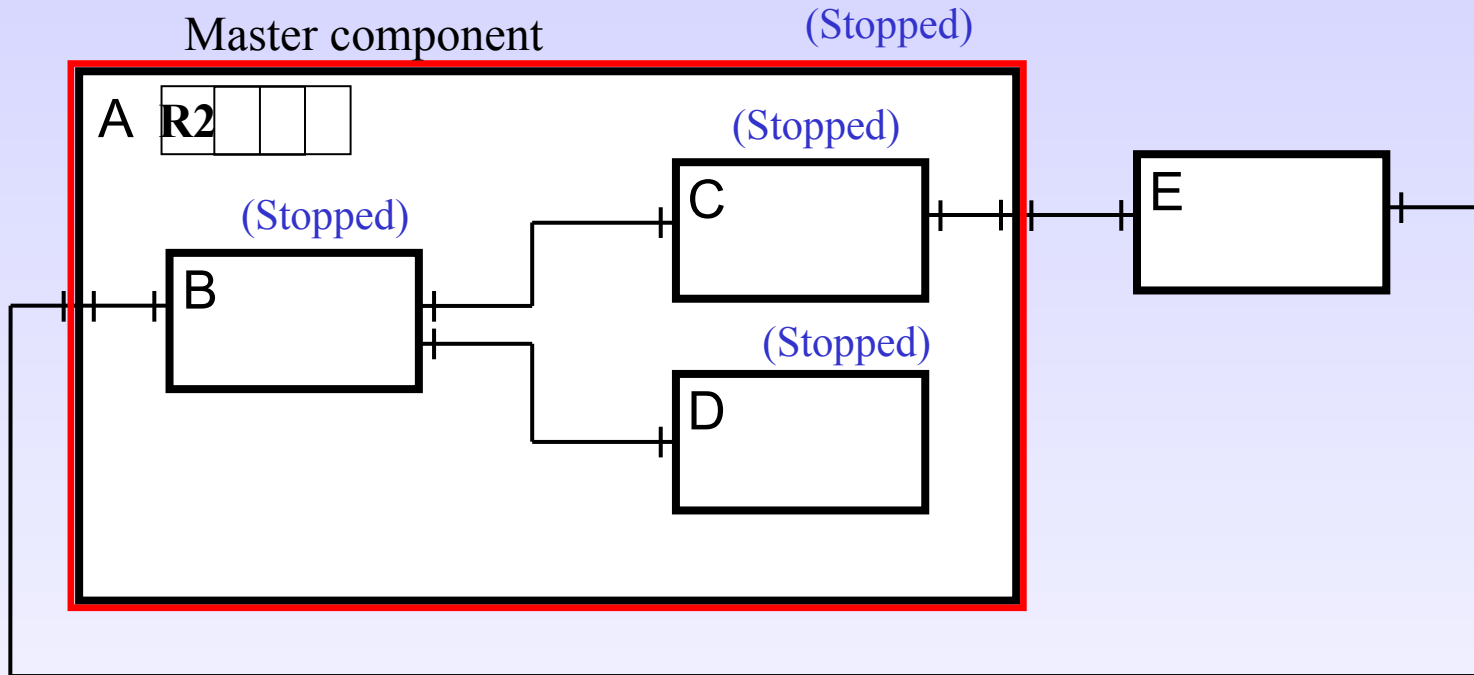


Example





Example





Algorithm description

- Initially all components are in *started state*.
- The *master component* changes to *waiting for stopping state* when it receives the stop request.
- While the master is in *waiting for stopping state* it marks all outgoing requests.
- When all requests are marked the *master component* changes to *stopping state* and it sends the *stopping signal* to its subcomponents.

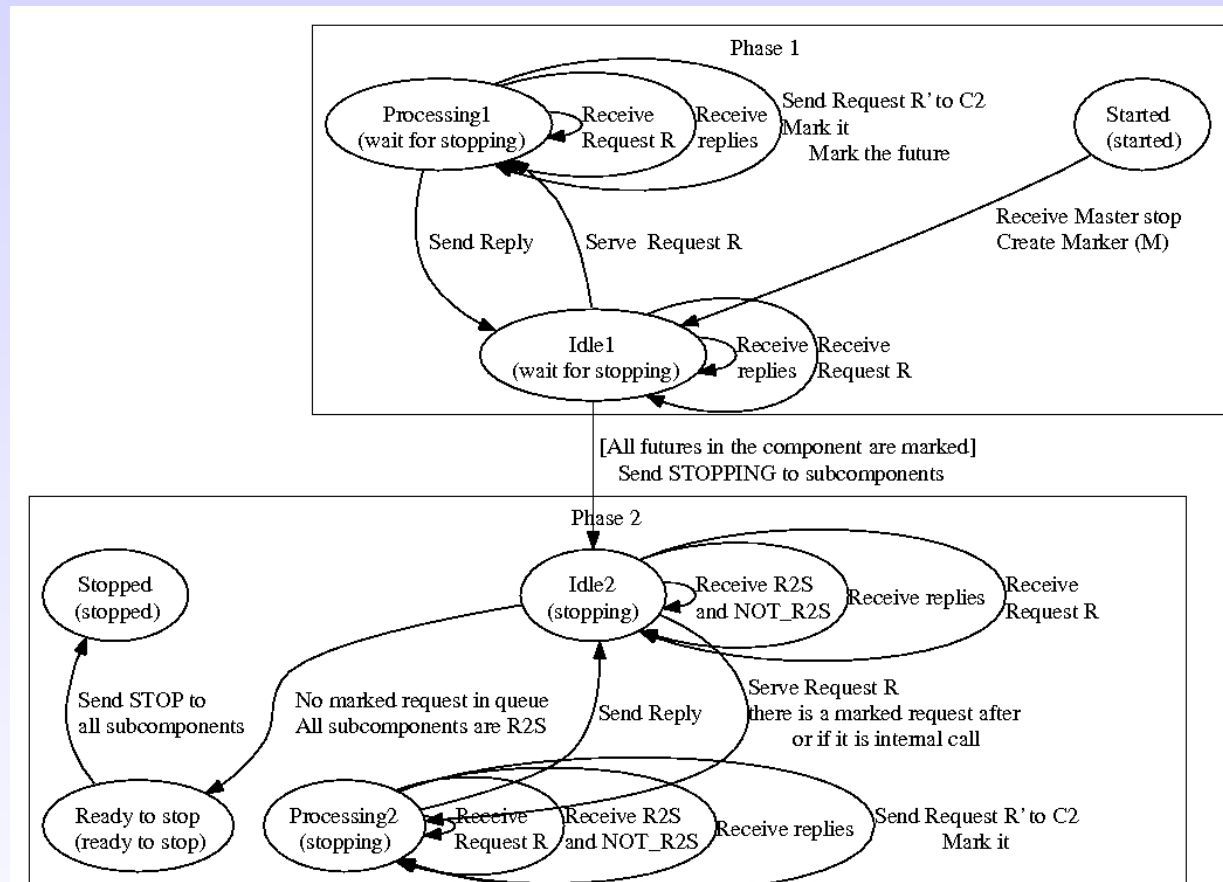


Algorithm description

- When the subcomponents receive the *stopping signal* they change to stopping state.
- When there are no more internal calls to process, each components send a *R2S signal* and change to *R2S state*.
- When the *master component* receives the *R2S signal* it sends the *stop signal* for stopped the system.



Algorithm for the master component





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Conclusion

- Any component without deadlock and livelock is safely stopped:
 - The master component is in *strongly idle state (quiescent)*.
 - Master queue only contains requests received after the last served request.
- The algorithm is independent of components topology or components behavior.
- Our contribution can be considered as an adaptation and an extension of the lifecycle controller of Fractal for GCM.
- Prototype implementation in GCM.



Extension and future works

- Stopping several components.
 - Marking states and requests with the identifier of the master.
- Non hierarchical systems
 - One central manager.
 - Intercepts calls outgoing the assembly.
 - Intercepts calls ingoing the assembly.
- Verify formally the stopping algorithm.
- Considering hierarchical systems with shared components.



Thanks