Multi-active Objects

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Outline

1. Active objects and their limitations

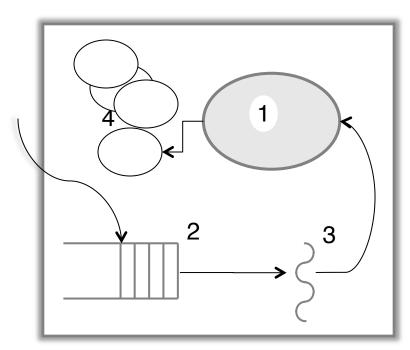
- 2. Related Works
- 3. Proposed solution
- 4. Experiment
- 5. Conclusion
- 6. Current and future work

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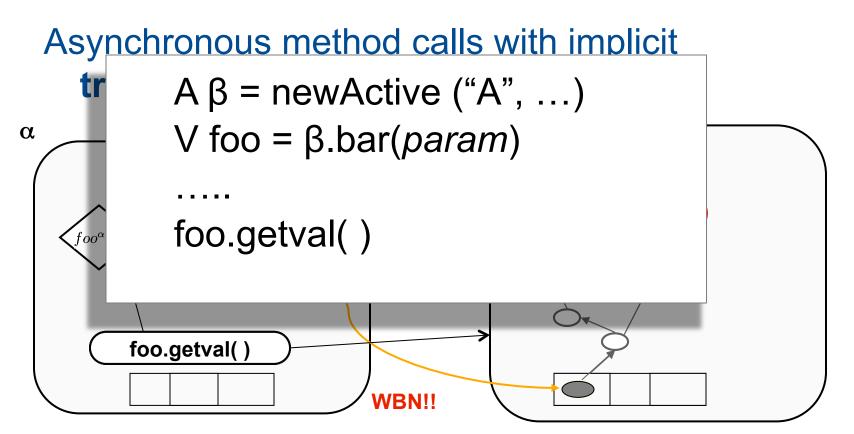
Active Objects

- Asynchronous communication with futures
- Location transparency
- Composition:
 - An active object (1)
 - a request queue (2)
 - one service thread (3)
 - Some passive objects (local state) (4)



ASP / ProActive

Active objects, asynchronous communication

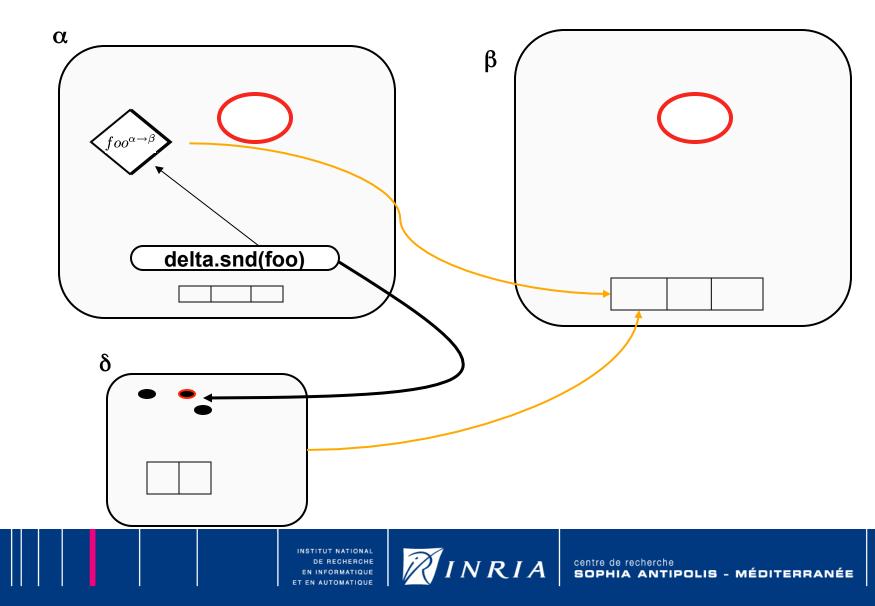


Caromel, D., Henrio, L.: A Theory of Distributed Object. Springer-Verlag (2005)

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First Class Futures



Active Objects – Limitations

No data sharing

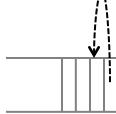
- Parameters of method calls are passed by value
 No data race-condition

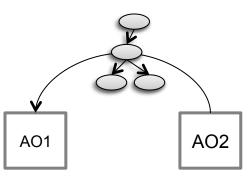
 >simpler programming + easy distribution
 >Slow local parallelism
 >Less efficient
- No re-entrant calls
 - Active object deadlocks by waiting on itself
 Modifications to the application logic

→ difficult to program

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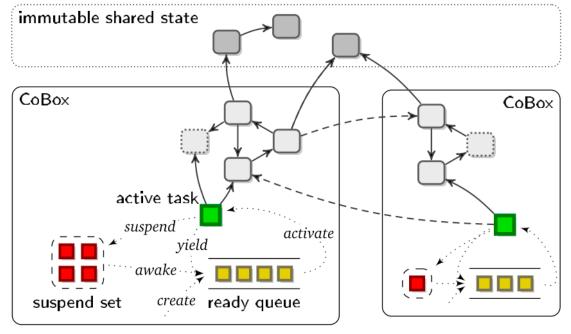
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Related Work (I)

Creol an JCobox:

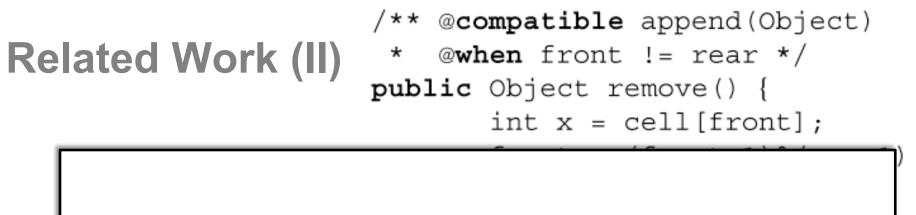
 Active object paradigm & futures



- Cooperative multithreading
 - -All requests are served at the same time
 - -But only one thread active at a time
 - Explicit release points in the code
- Less problem with re-entrance
- More difficult to program: less transparency

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Our "idea": adapt a simple version of JAC to simple active objects a la ASP to get efficient and easy to program multi-active objects

- Simulating active objects is possible but not trivial





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Proposal: Multi-active object

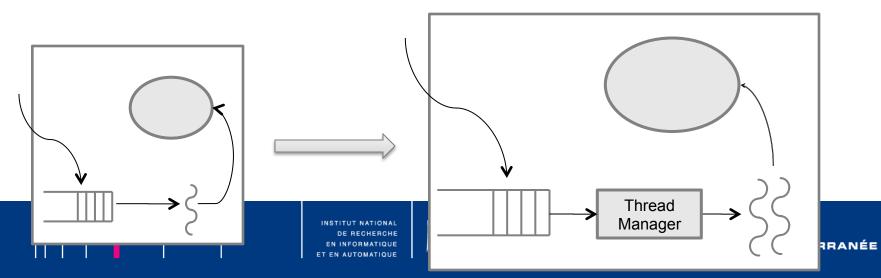
Starting from active objects à la ASP

Transparent multi-threading

Executing compatible requests in parallel

Safe parallel execution

A thread manager handles the queue



Proposal – Compatibility

Two methods are compatible if either: a) They do not access the same resources b) The user "protects" the locations of possible data races

How to express this information?

Annotate the code

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Proposal – Annotations (I)

@Group

- Group identifier
- Self compatible = methods of the group can run in parallel

```
@DefineGroups(
        @Group(name = "GroupFoo", selfCompatible = true),
        @Group(name = "GroupBar", selfCompatible = false)
        }
                                 NRIA
```

MÉDITERRANÉE

Proposal – Annotations (II)

@Compatible

A list of mutually compatible groups

```
@DefineRules(
    {
        Compatible({"GroupFoo", "GroupBar"})
    }
)
```

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Proposal – Annotations (III)

@MemberOf

- Refers to a group identifier
- Each method belongs to a single group
- Not annotated methods are incompatible with the others

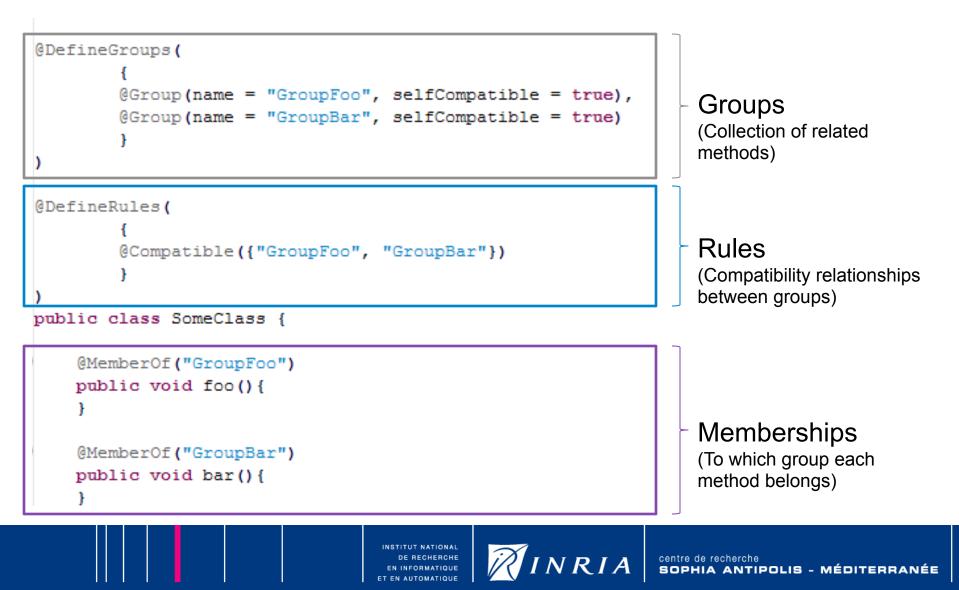
```
@MemberOf("GroupFoo")
public void foo(){
}
```

@MemberOf("GroupBar")
public void bar(){

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Proposal – Annotation Example



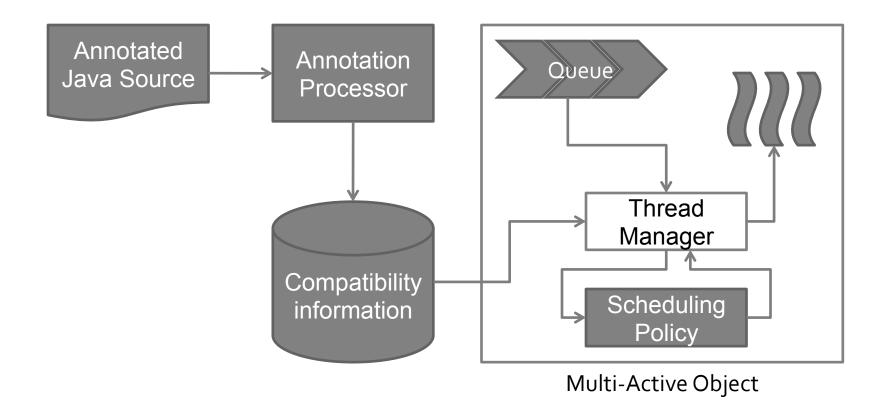
Proposal – Thread Manager

Scheduling

- Default: FIFO + compatibility
- Pluggable policies
- A policy is a function:
 - Input: info from scheduler
 - Output: list of requests to be started
- Scheduling Policy API provides:
 - Static compatibility information
 - Scheduler state (request queue, active requests)

```
method runActivity() {
  while (true) {
    if (compatible(requestQueue.peekFirst(),
        activeRequests)) {
        parallelServe(requestQueue.removeFirst());
}
```

Proposal – Main Elements





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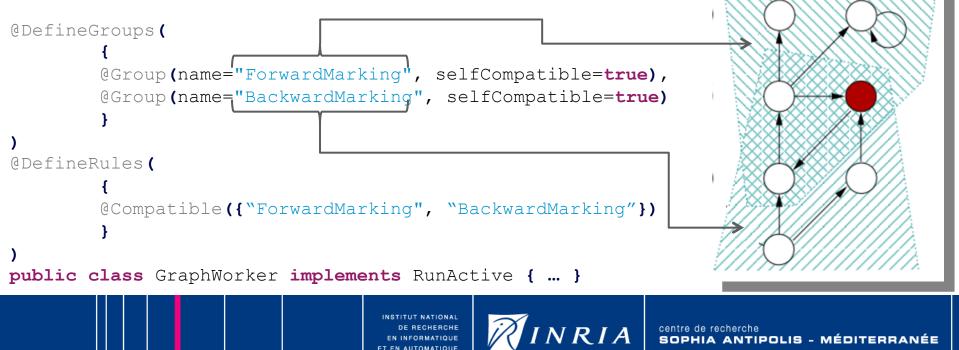
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Experiment – SCC Search

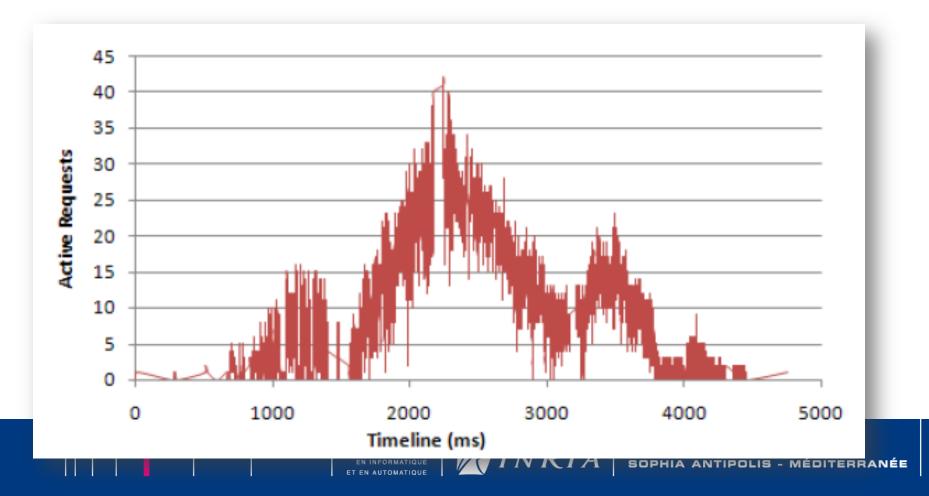
Strongly connected component search in a distributed graph

- Divide-and-conquer
- Local and distributed parallelism



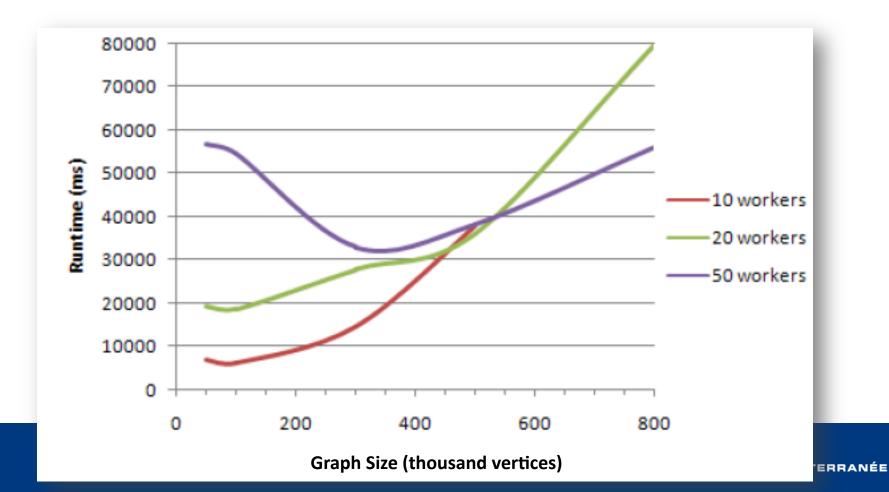
Experiment – Results (I)

Local parallelism



Experiment – Results (II)

Performance



Experiment – Results (III)

Multi-active objects are usable in large-scale applications:

- Easy distributed execution 80 workers
- Quite large graph sizes 800,000 vertices

No changes have to be done to the application logic

Acceptable speedup even without explicit optimization



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Conclusion

Multi-active objects feature:

- Active object model
 - Easy to program
 - Support for distribution
- Efficient utilization of multi-cores
 - Transparent multi-threading
 - Safe parallel execution
- Possibility to write re-entrant code
- Simple annotations

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Current and Future Work (I)

Compatibility

- Create dynamic rules use parameters
- Common parameter for methods in a group
- Comparison of parameters at runtime

```
@DefineRules({
    @Compatible(value={"data_access", "join"},
        condition="!this.inSplittingZone") })
```



Current and Future Work

Threads

- Too many threads can be harmful
- Limitation of threads without deadlocks

Formalisation of the multi-active object model

SERVELL

 $C = \mathcal{R}[Serve(M)] \mapsto f:[a_i \mapsto f_i] \| C'$ $ParallelSchedule(M, Futures(C), R) = ([m, f, \iota], R')$ $\alpha[F; C; R; \sigma] \| P \longrightarrow \alpha[F; \iota_0.m(\iota) \mapsto f' \| C; R'; \sigma] \| P$



Questions?

More information:

Adapting Active Objects to Multicore Architectures

Ludovic Henrio, Fabrice Huet, Zsolt István, and Gheorghen Sebestyén International Symposium on Parallel and Distributed Computing *(ISPDC 2011) - IEEE*

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