

Shape Reconstruction

Pierre Alliez

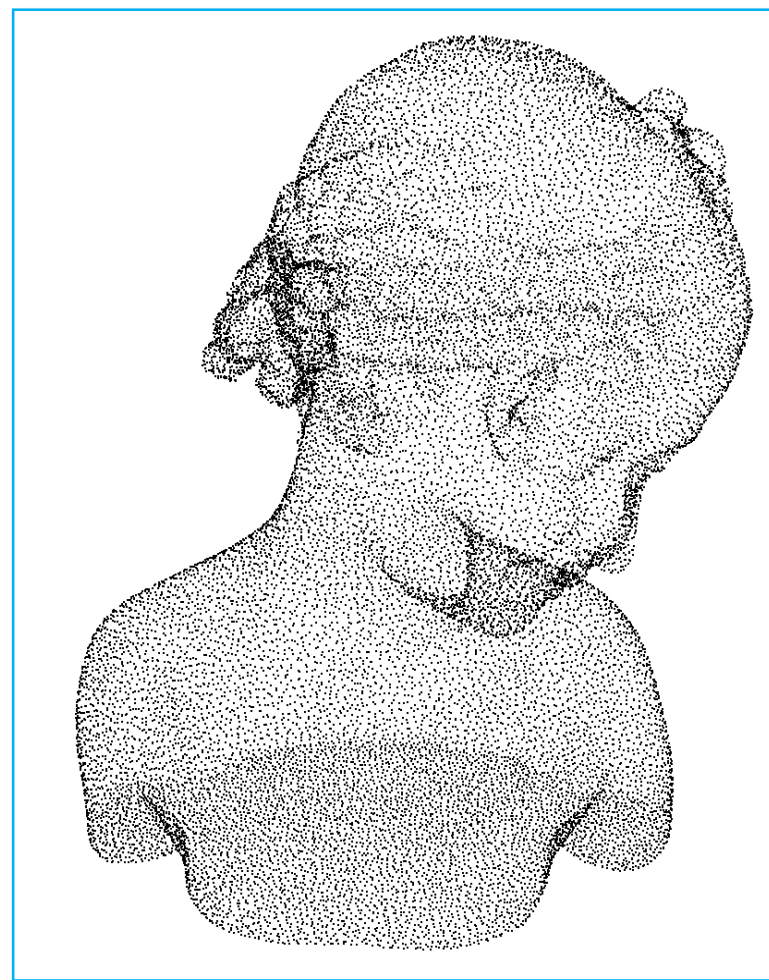
Inria

Outline

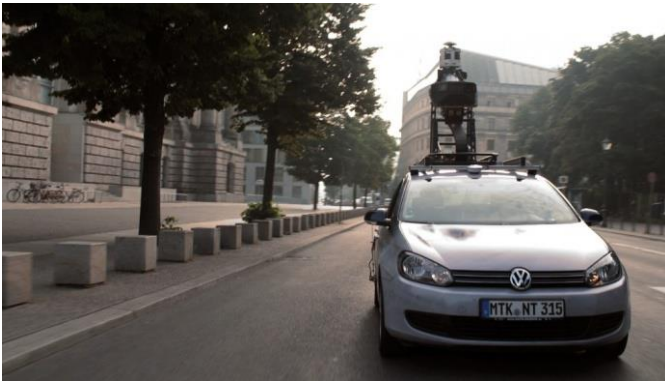
- Sensors
- Problem statement
- Computational Geometry
 - Convex hull, Voronoi/Delaunay, alpha-shapes

SENSORS

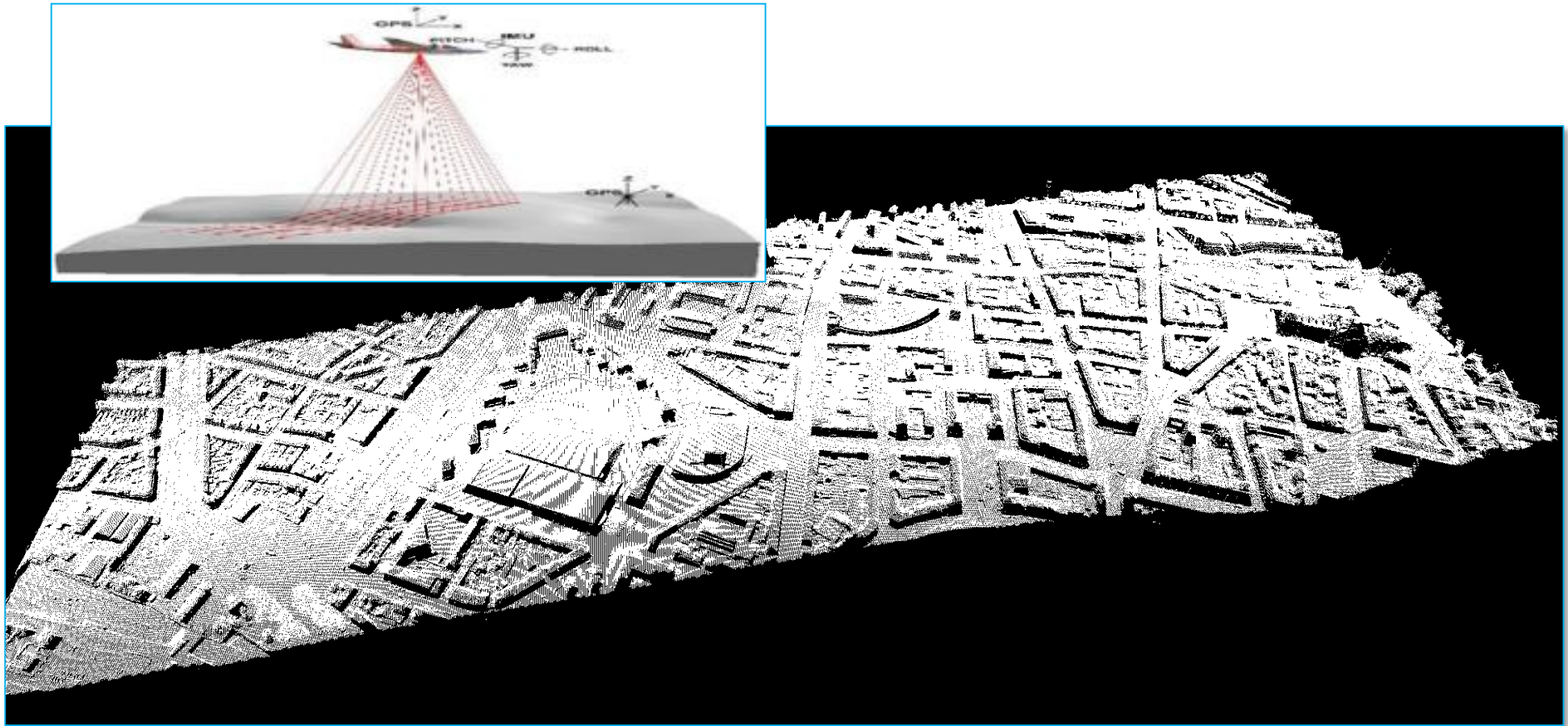
Laser scanning



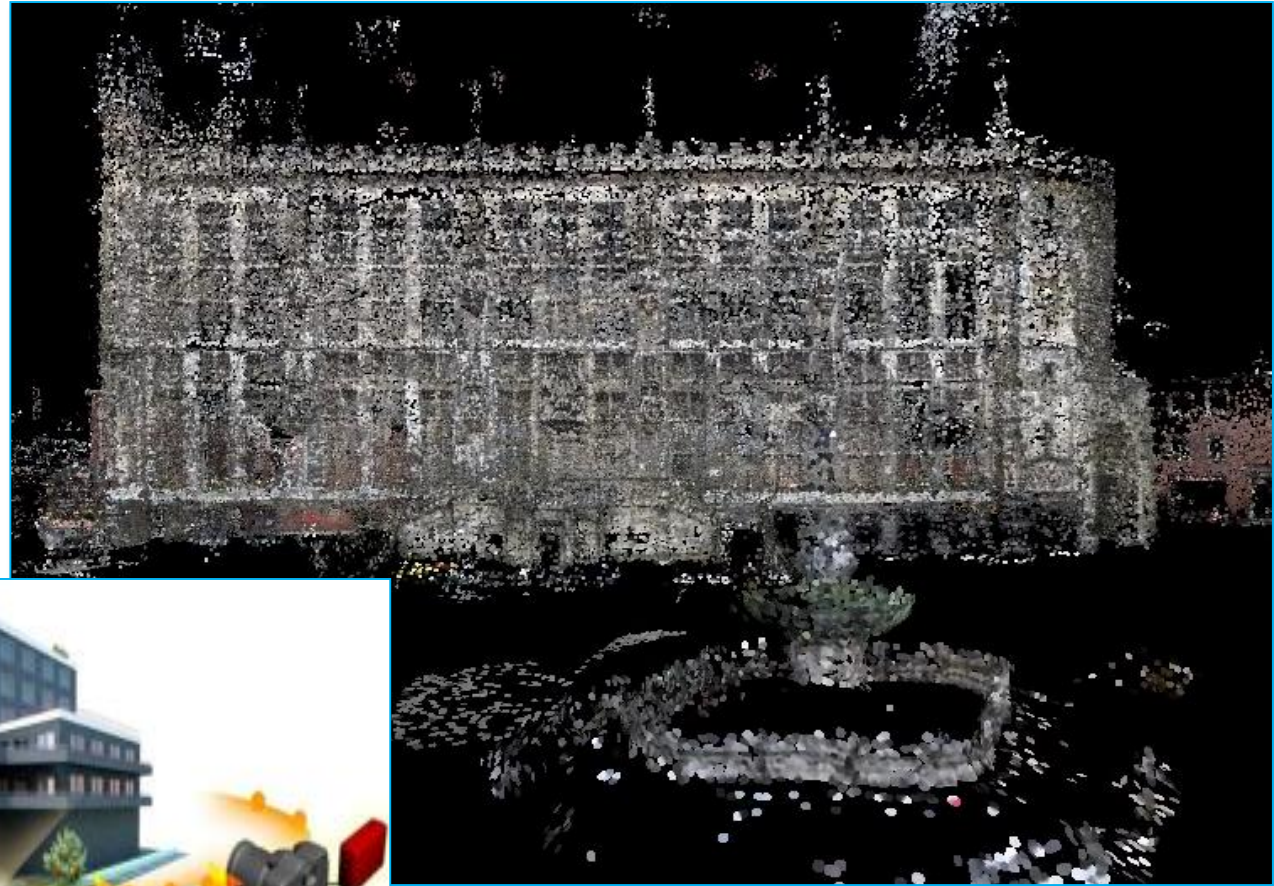
Car-based Laser



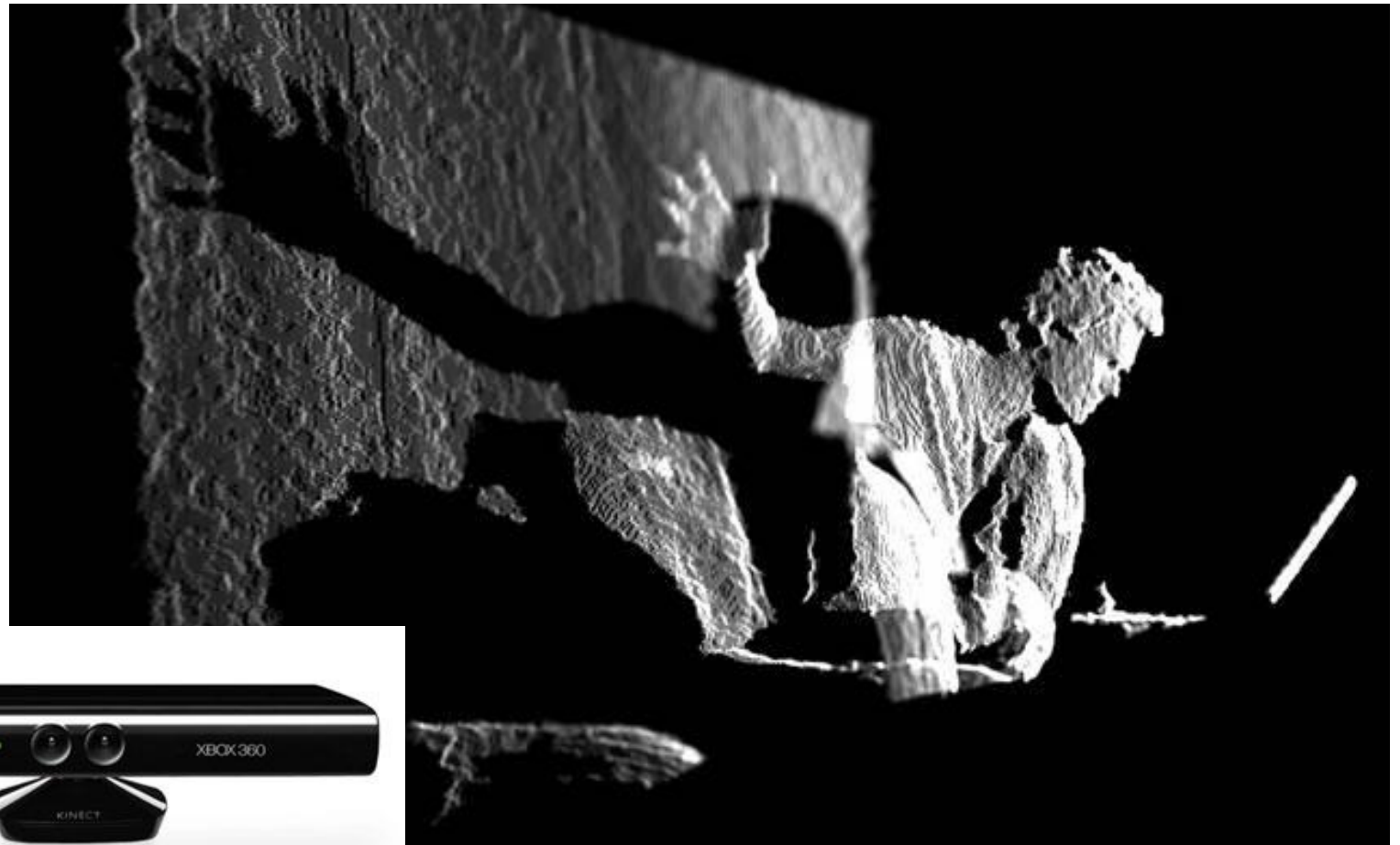
Airborne Lidar



Multi-View Stereo (MVS)



Depth Sensors



PROBLEM STATEMENT

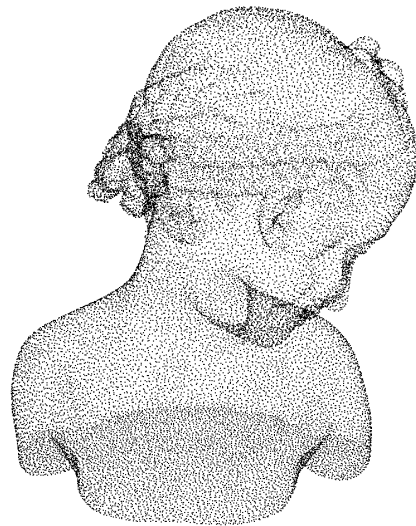
Reconstruction Problem

Input: point set P sampled over a surface S :

Non-uniform sampling

With holes

With uncertainty (noise)



point set

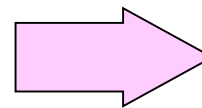
Output: surface

Approximation of S in terms of topology and geometry

Desired:

Watertight

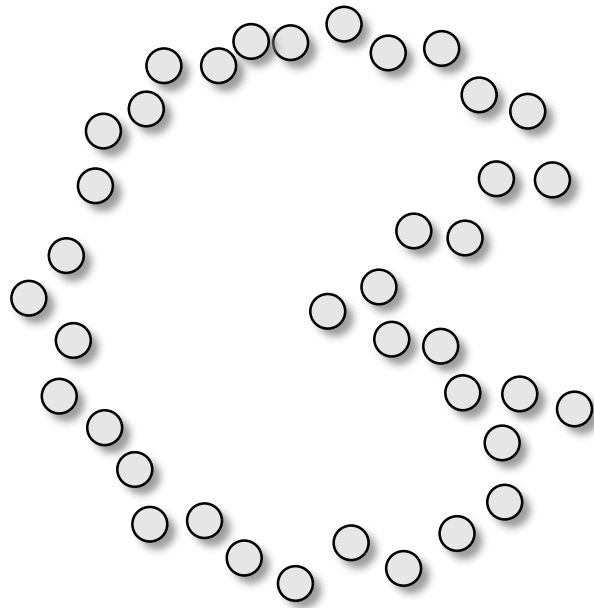
Intersection free



reconstruction

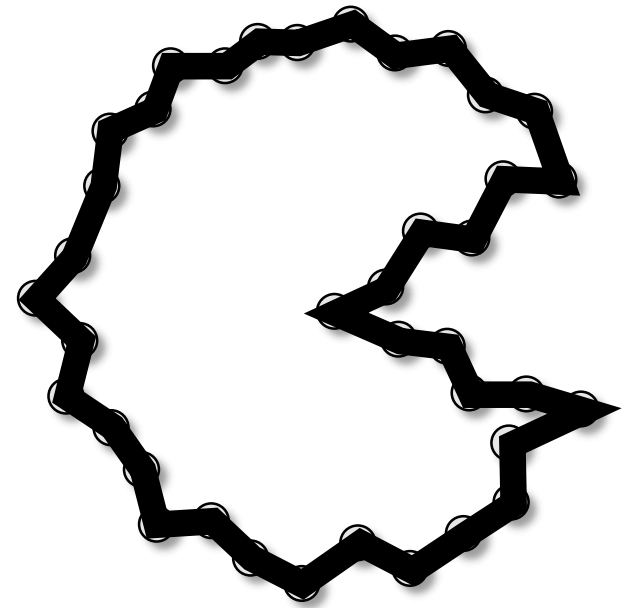
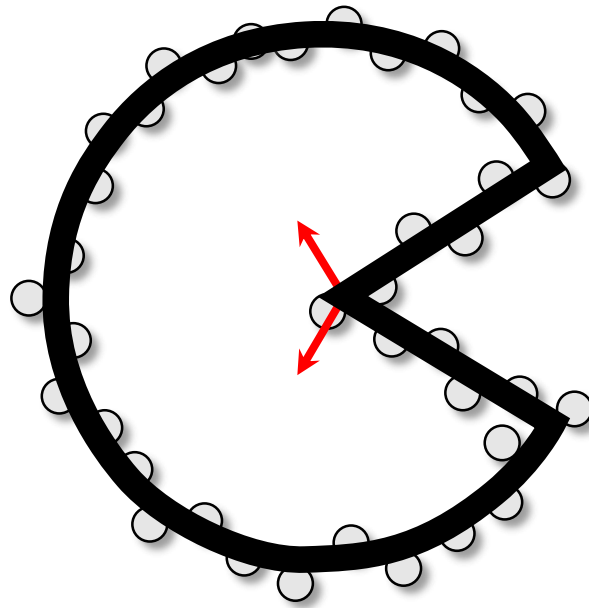
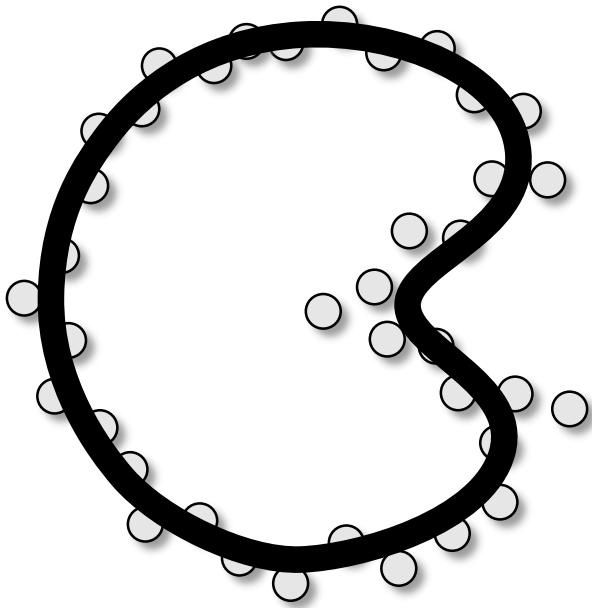
surface

Ill-posed Problem



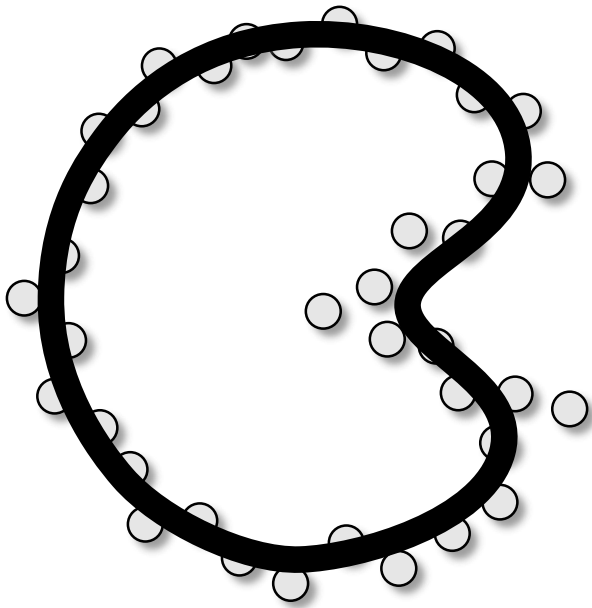
Many candidate surfaces for the reconstruction problem!

Ill-posed Problem

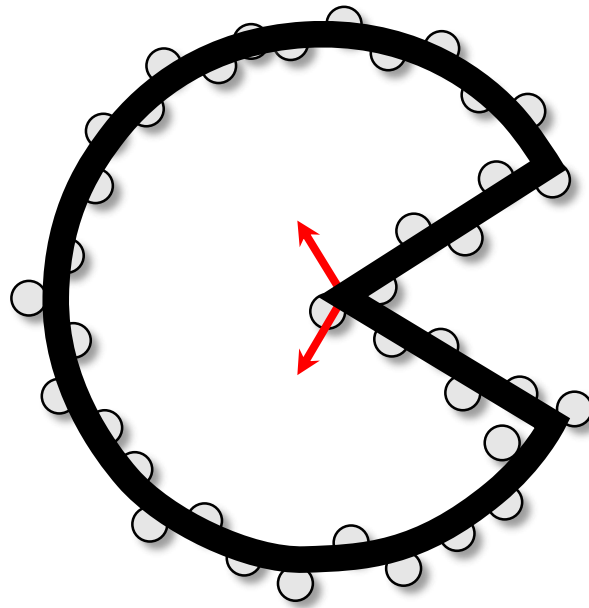


Many candidate surfaces for the reconstruction problem! How to pick?

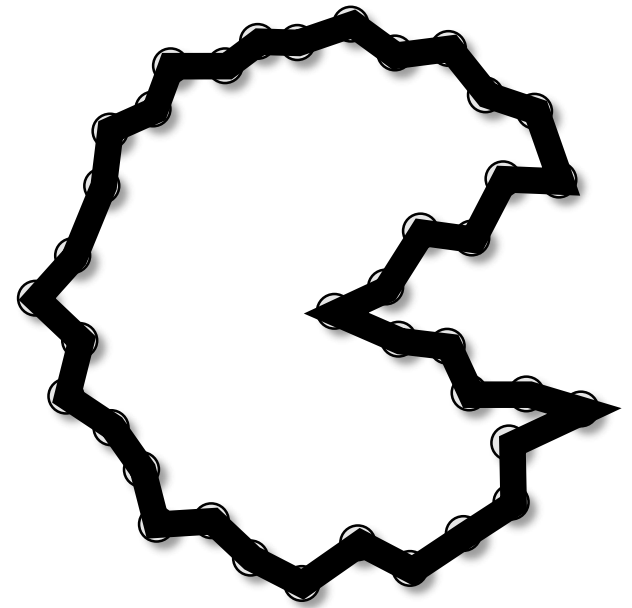
Priors



Smooth



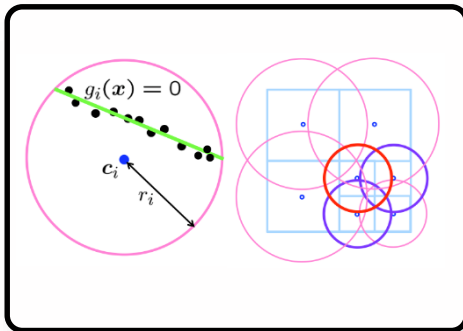
Piecewise Smooth



“Simple”

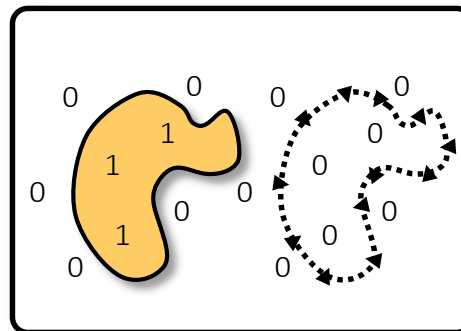
Surface Smoothness Priors

Local Smoothness



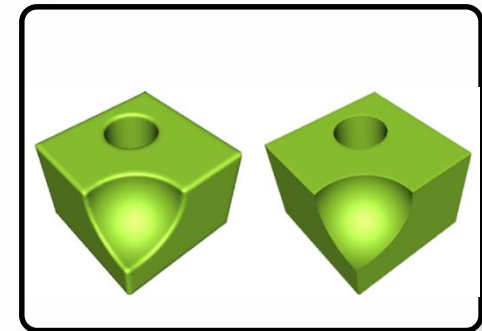
Local fitting
No control away from data
Solution by interpolation

Global Smoothness



Global: linear, eigen, graph cut, ...
Robustness to missing data

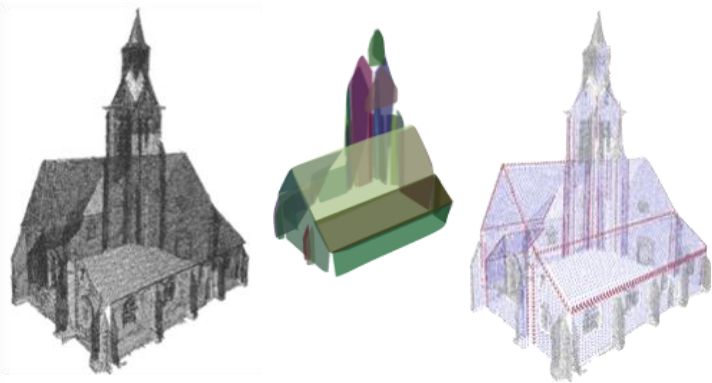
Piecewise Smoothness



Sharp near features
Smooth away from features

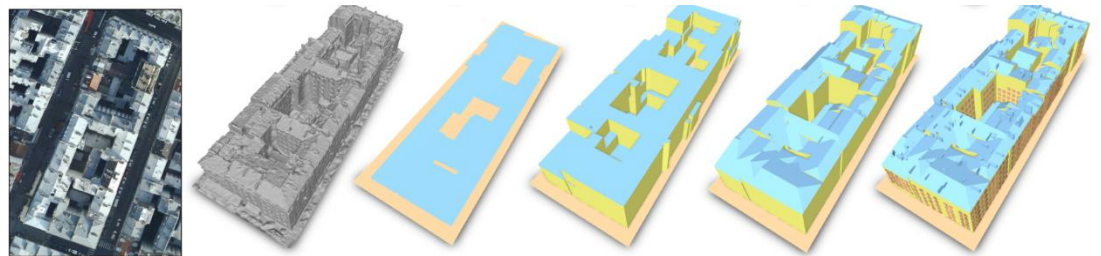
Domain-Specific Priors

Surface Reconstruction by Point Set Structuring



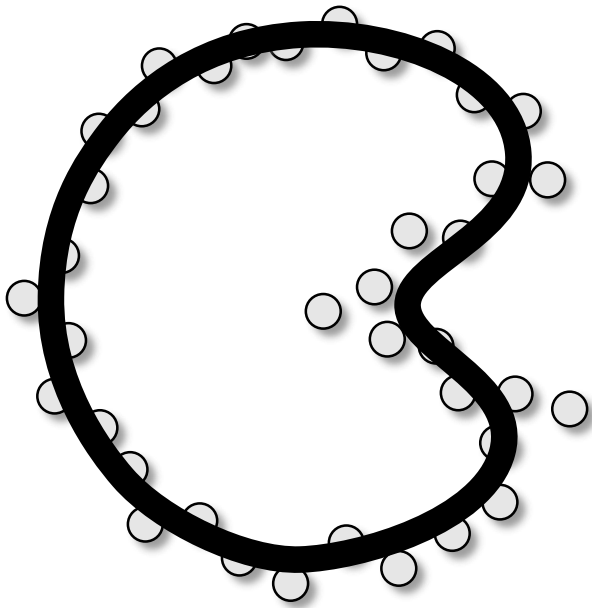
[Lafarge - A. EUROGRAPHICS 2013]

LOD Reconstruction for Urban Scenes

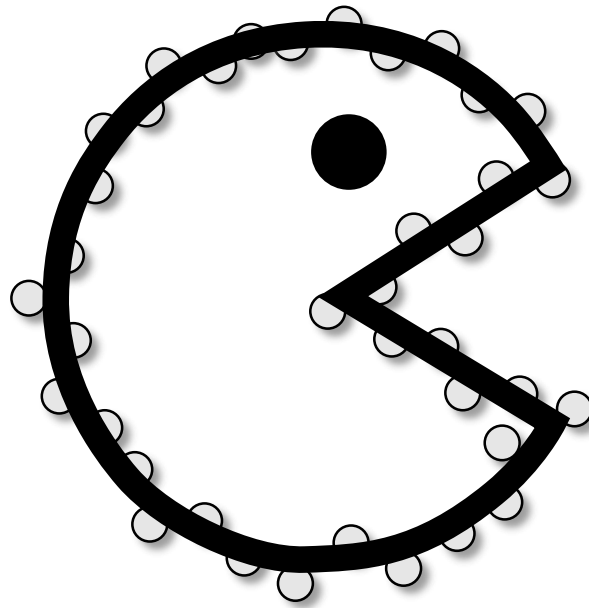


[Verdie, Lafarge - A. ACM Transactions on Graphics 2015]

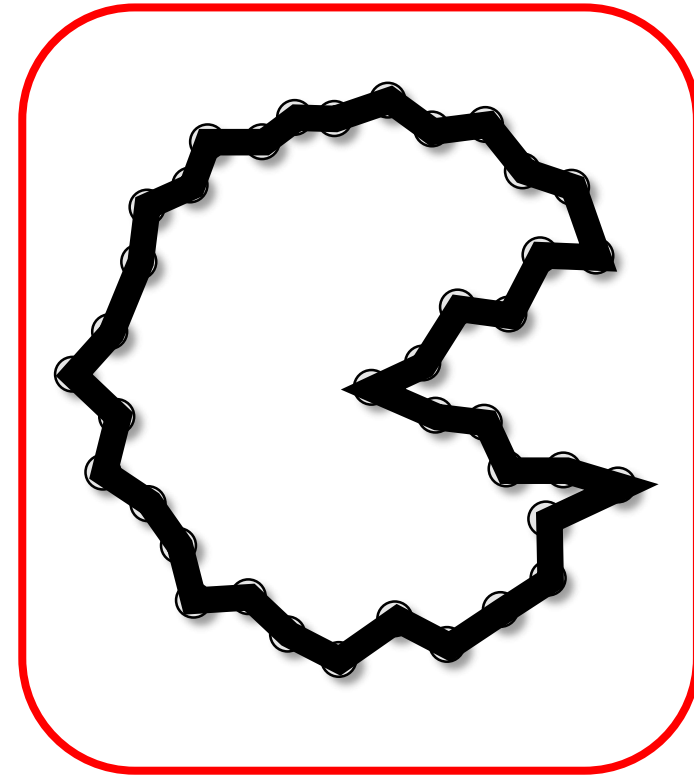
Warm-up



Smooth



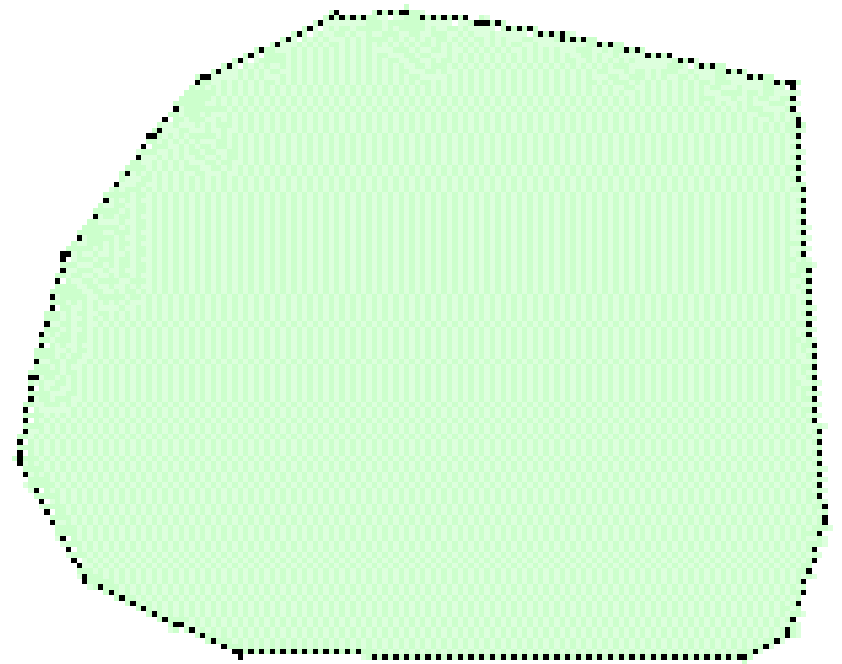
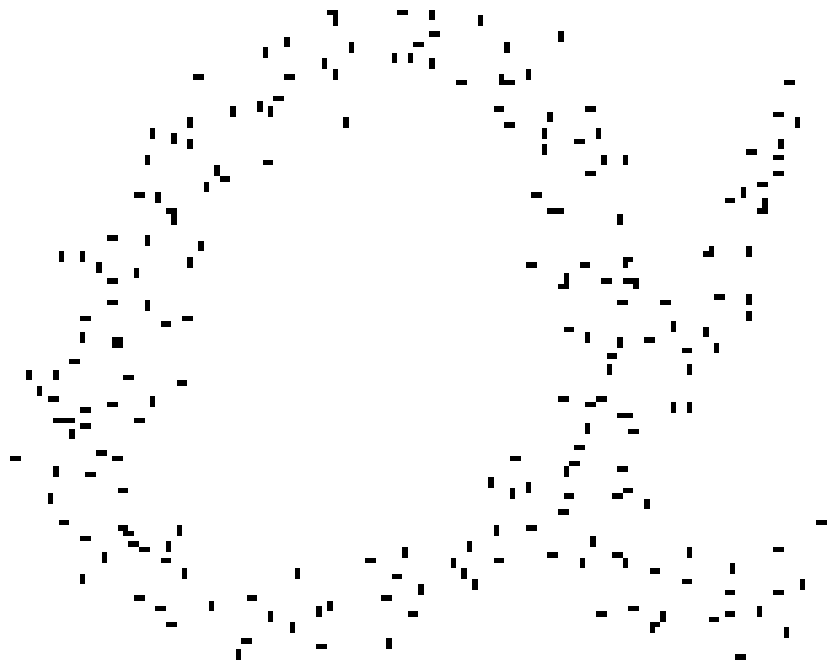
Piecewise Smooth



“Simple”

CONVEX HULL

Convex Hull

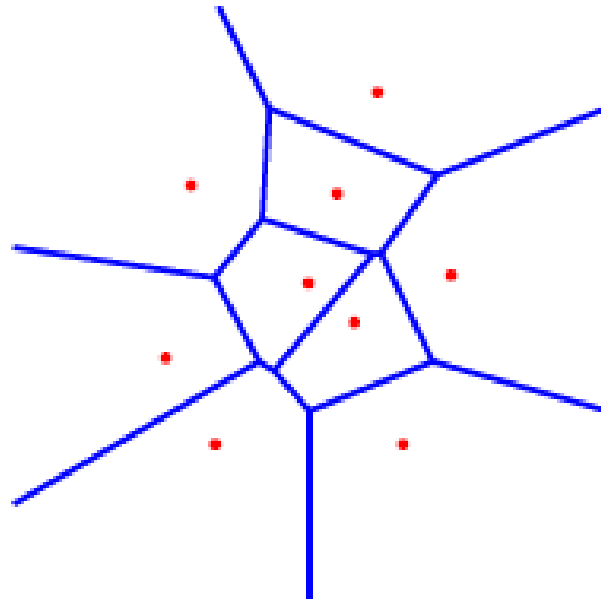


VORONOI / DELAUNAY

Voronoi Diagram

Let $\mathcal{E} = \{\mathbf{p}_1, \dots, \mathbf{p}_n\}$ be a set of points (so-called sites) in \mathbb{R}^d . We associate to each site \mathbf{p}_i its Voronoi region $V(\mathbf{p}_i)$ such that:

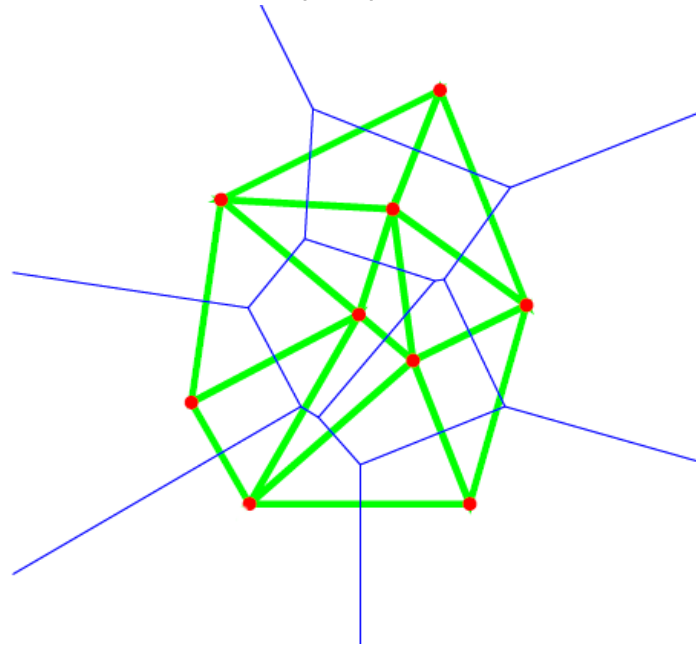
$$V(\mathbf{p}_i) = \{\mathbf{x} \in \mathbb{R}^d : \|\mathbf{x} - \mathbf{p}_i\| \leq \|\mathbf{x} - \mathbf{p}_j\|, \forall j \leq n\}.$$



Delaunay Triangulation

Dual structure of the Voronoi diagram.

The Delaunay triangulation of a set of sites E is a simplicial complex such that $k+1$ points in E form a Delaunay simplex if their Voronoi cells have nonempty intersection



Delaunay-based

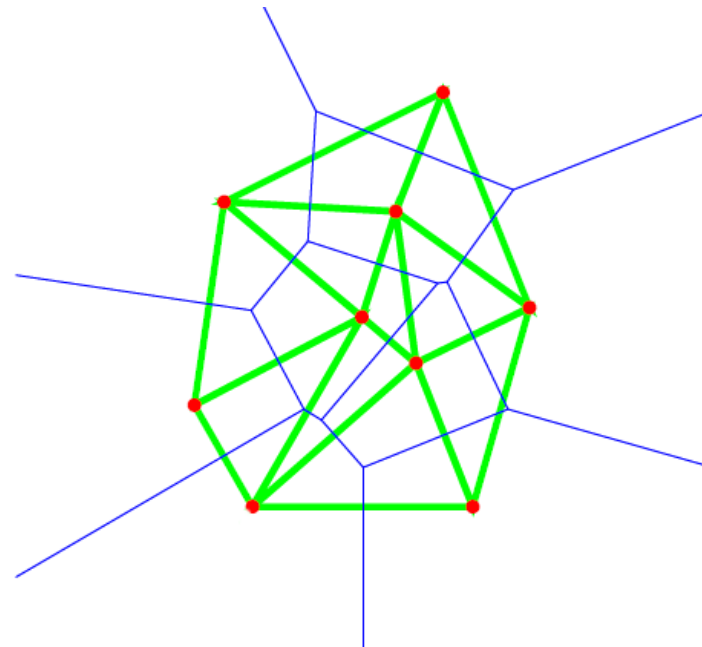
Key idea: assuming dense enough sampling, reconstructed triangles are Delaunay triangles.

First define

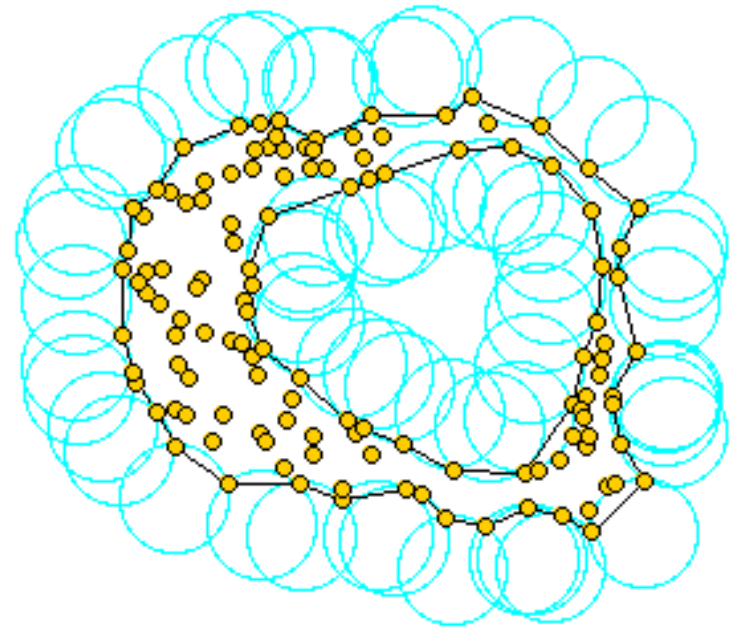
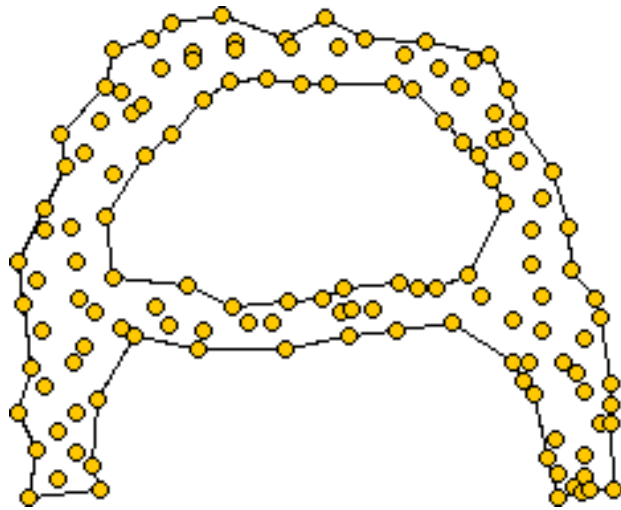
Medial axis

Local feature size

Epsilon-sampling



Alpha-Shapes [Edelsbrunner, Kirkpatrick, Seidel]



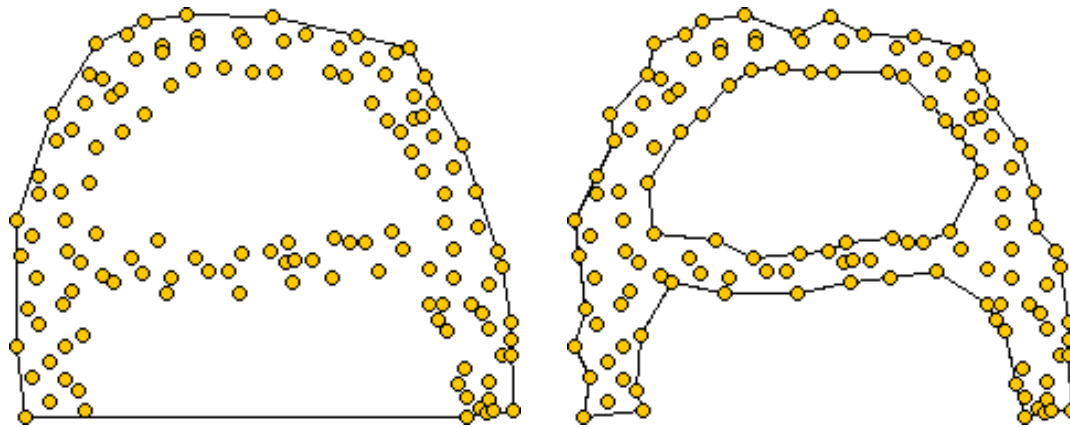
Segments: point pairs that can be touched by an empty disc of radius alpha.

Alpha-Shapes

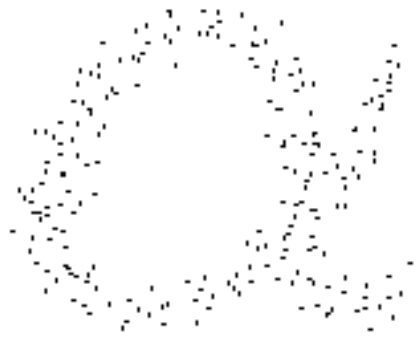
In 2D: family of piecewise linear simple curves constructed from a point set P .

Subcomplex of the Delaunay triangulation of P .

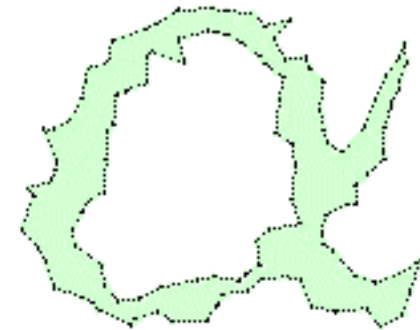
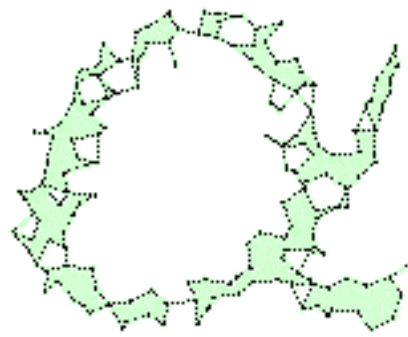
Generalization of the concept of the convex hull.



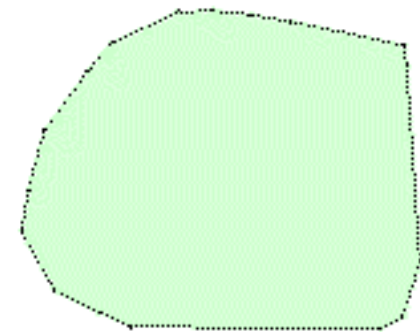
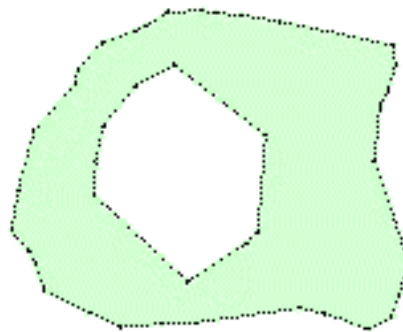
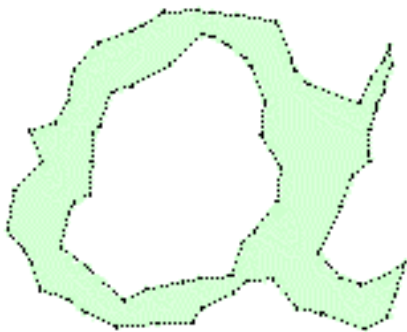
Alpha-Shapes



$$\alpha = 0$$

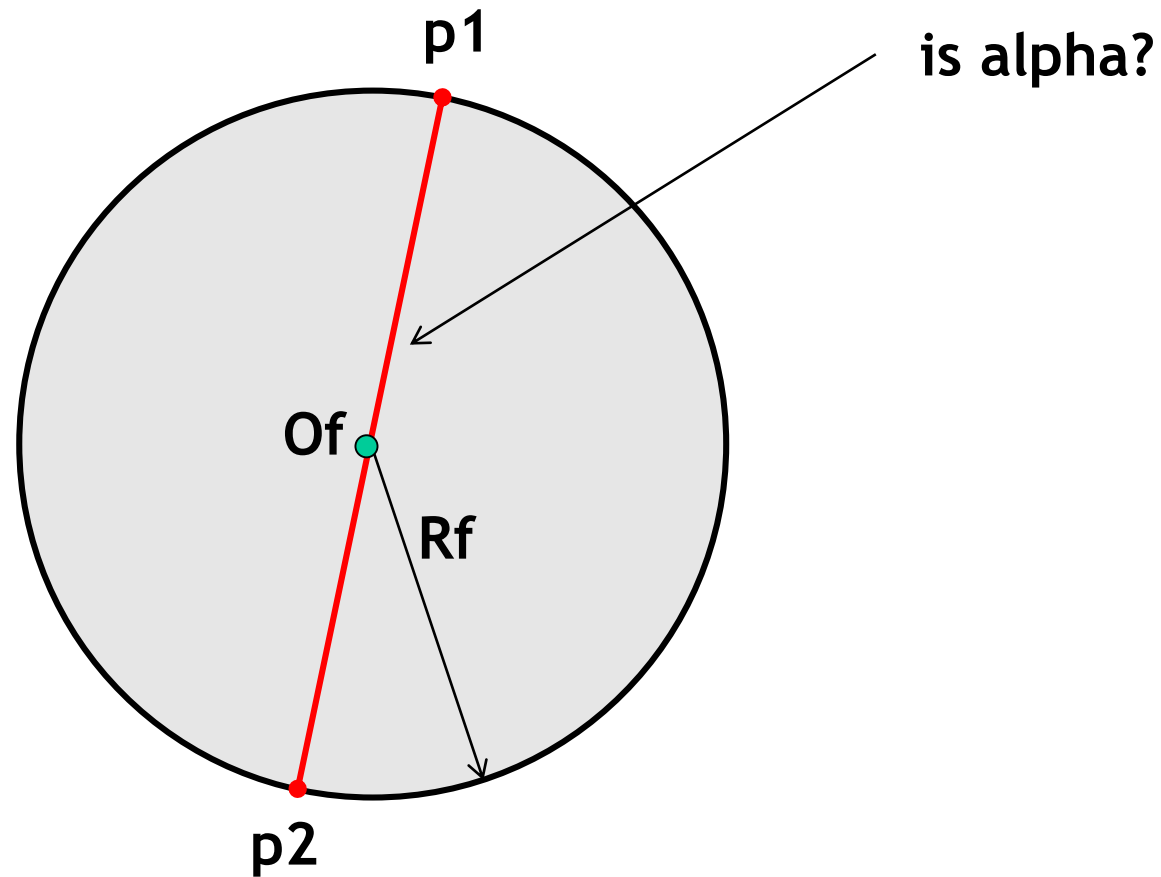


Alpha controls the desired level of detail.

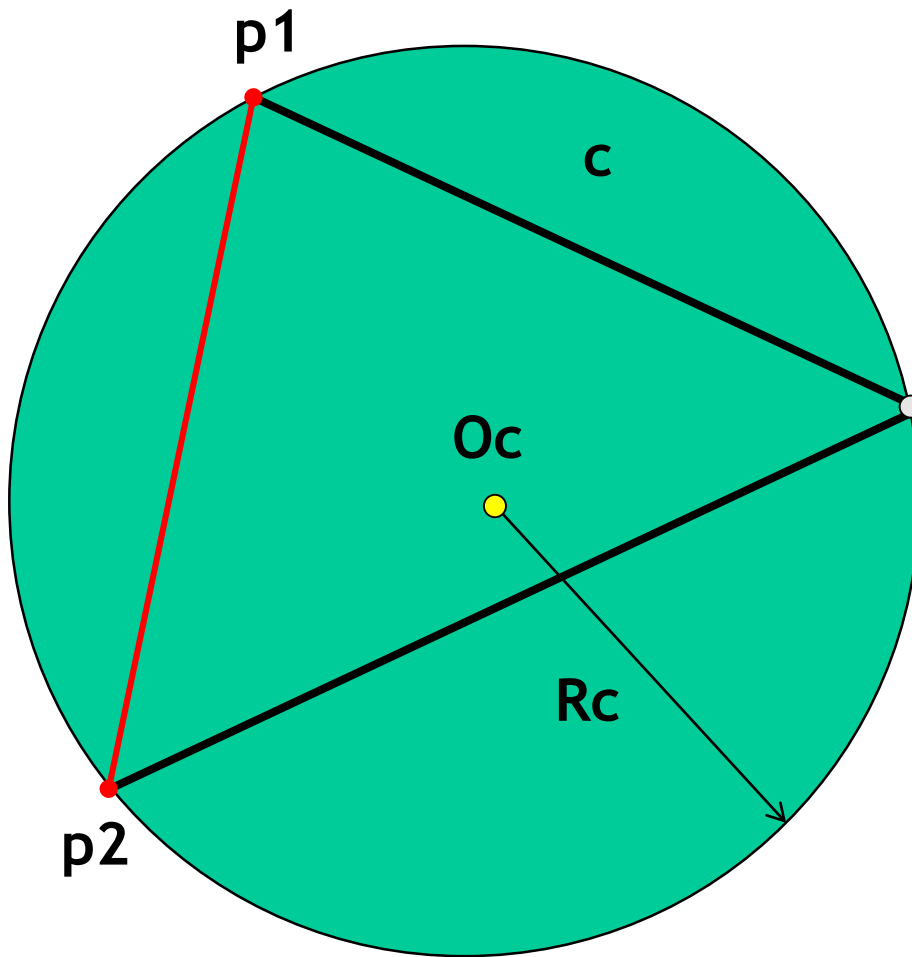


$$\alpha = \infty$$

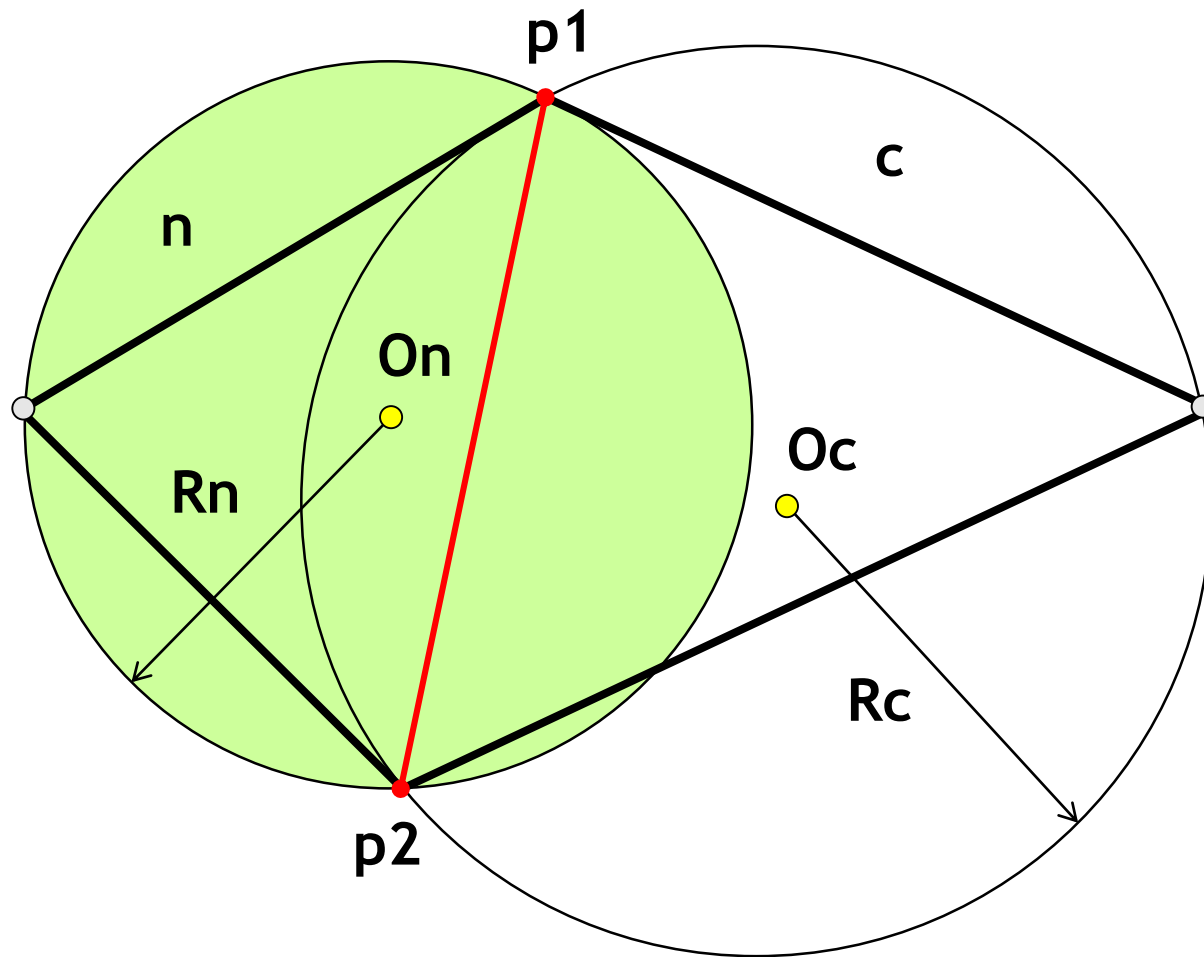
Computing Alpha-Shapes



Computing Alpha-Shapes

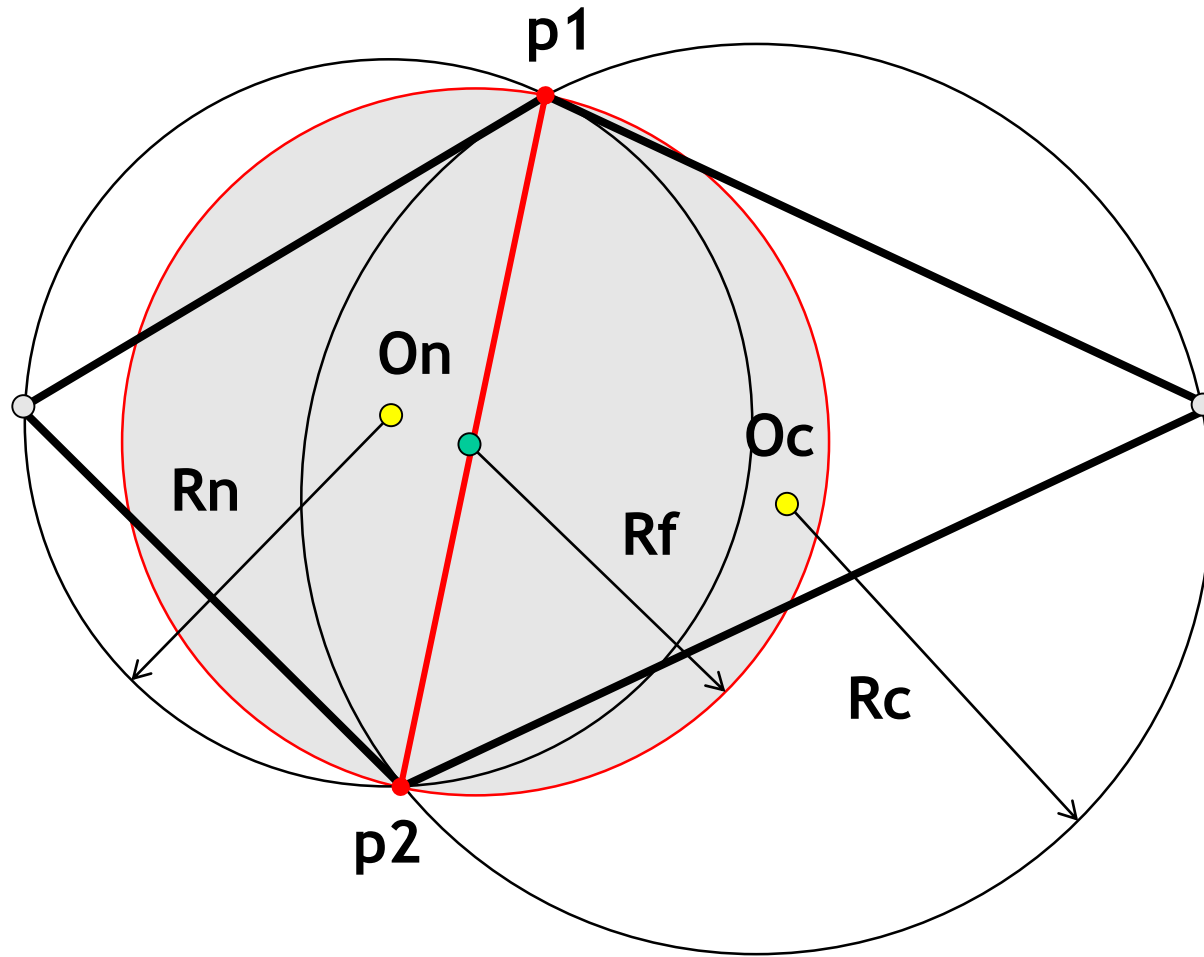
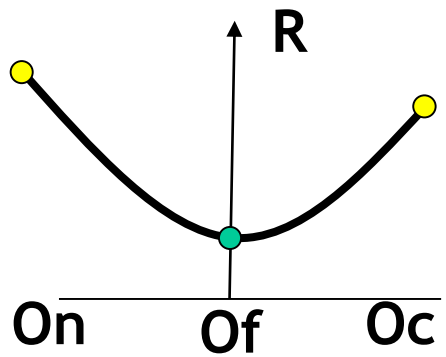


Computing Alpha-Shapes



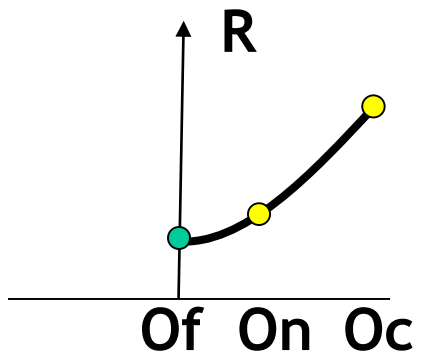
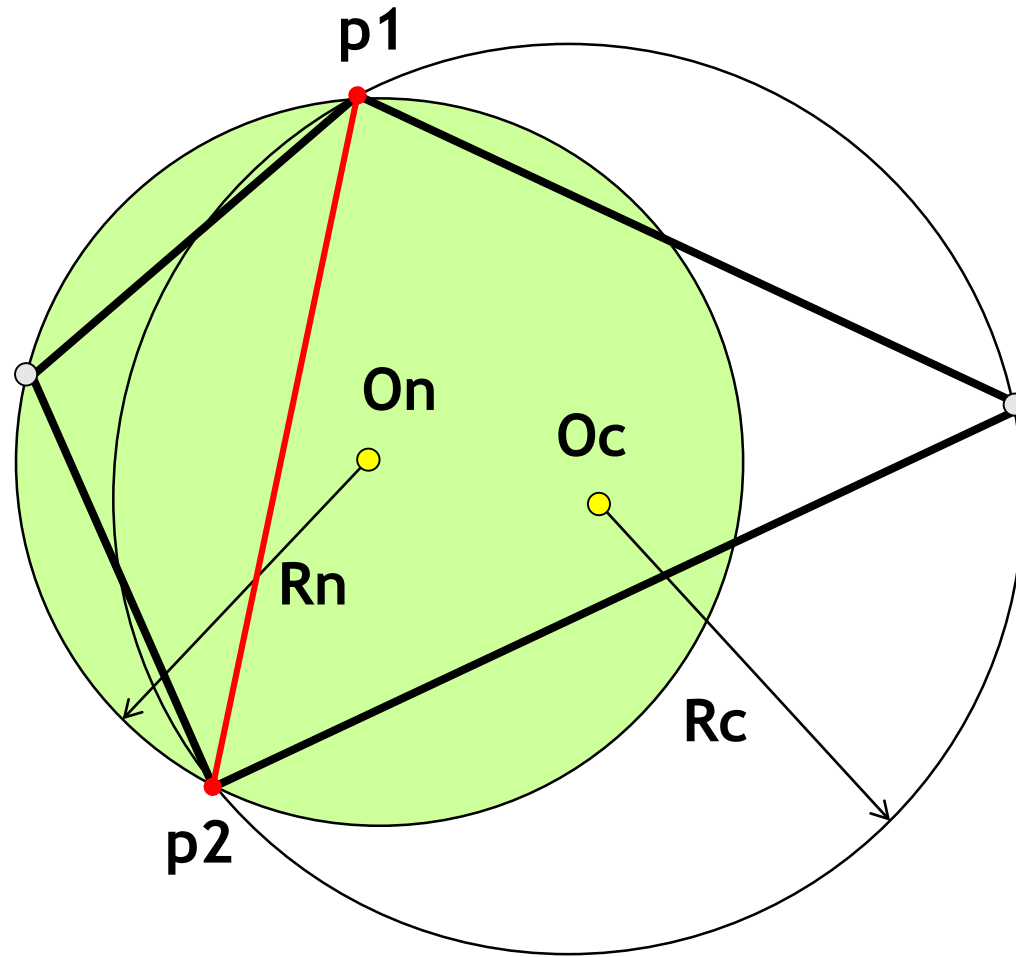
Computing Alpha-Shapes

is_alpha if
 $\alpha > R_f$
AND
 $\alpha < \max(R_c, R_n)$



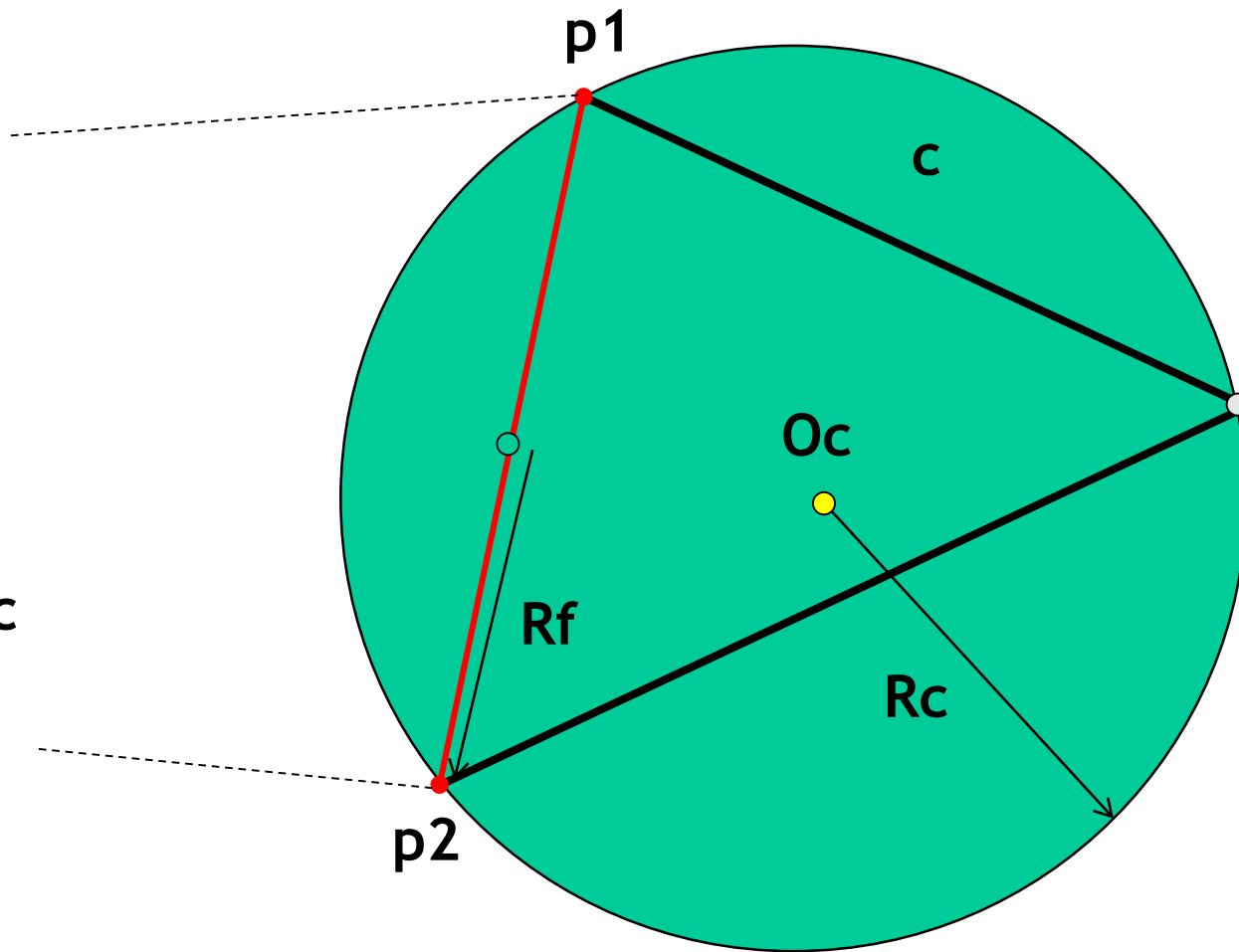
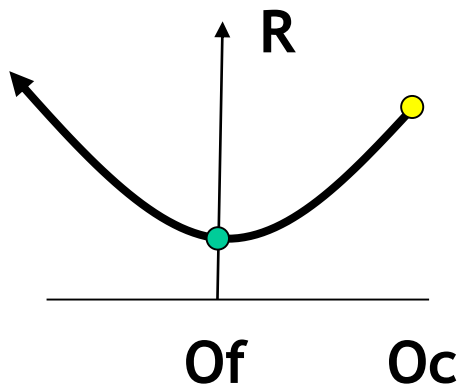
Computing Alpha-Shapes

is_alpha if
 $\alpha > \min(R_c, R_n)$
AND
 $\alpha < \max(R_c, R_n)$



Computing Alpha-Shapes

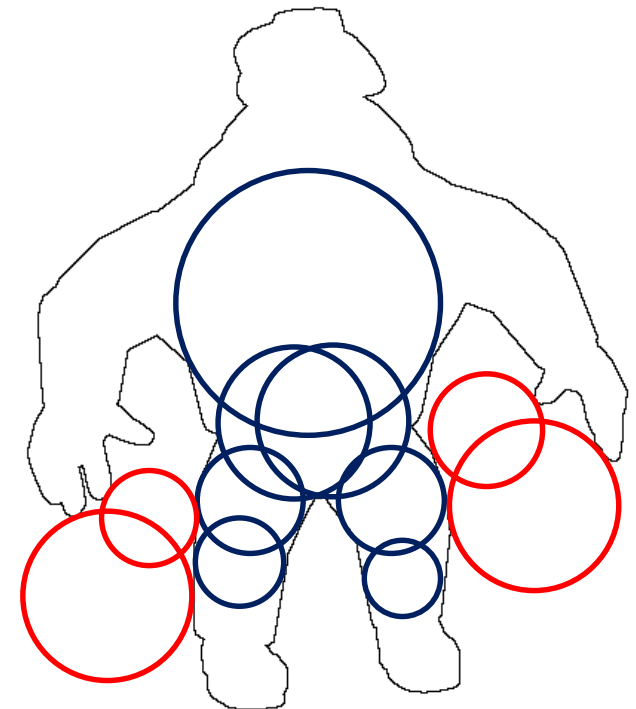
is_alpha if
 $\alpha > R_f$



MEDIAL AXIS

Medial Axis

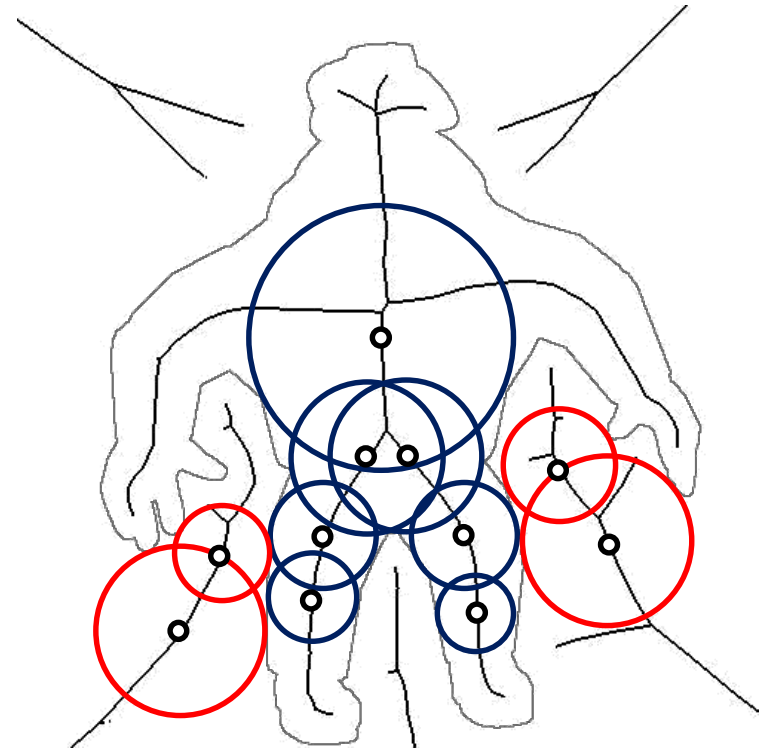
For a shape (curve/surface) a *Medial Ball* is a circle/sphere that only meets the shape tangentially, in at least two points.



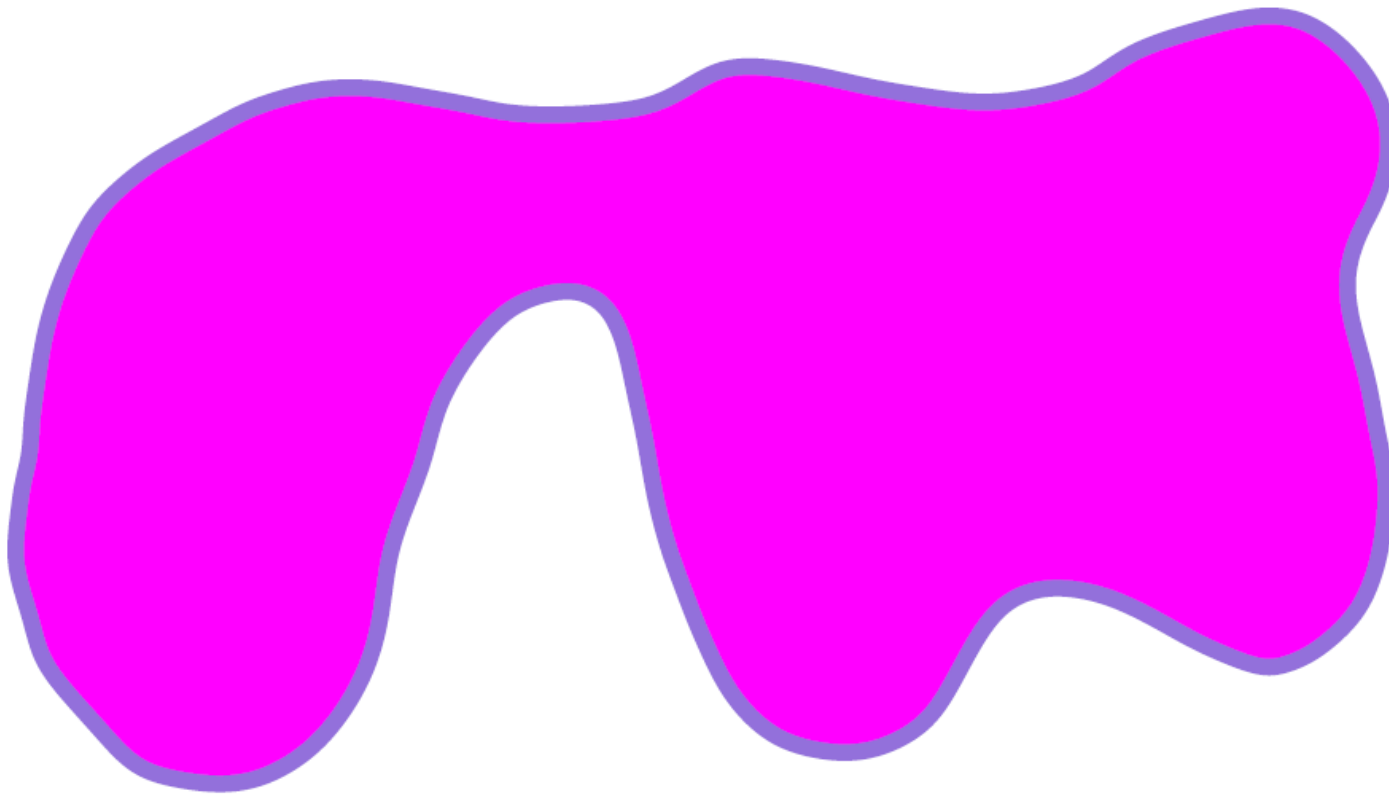
Medial Axis

For a shape (curve/surface) a *Medial Ball* is a circle/sphere that only meets the shape tangentially, in at least two points.

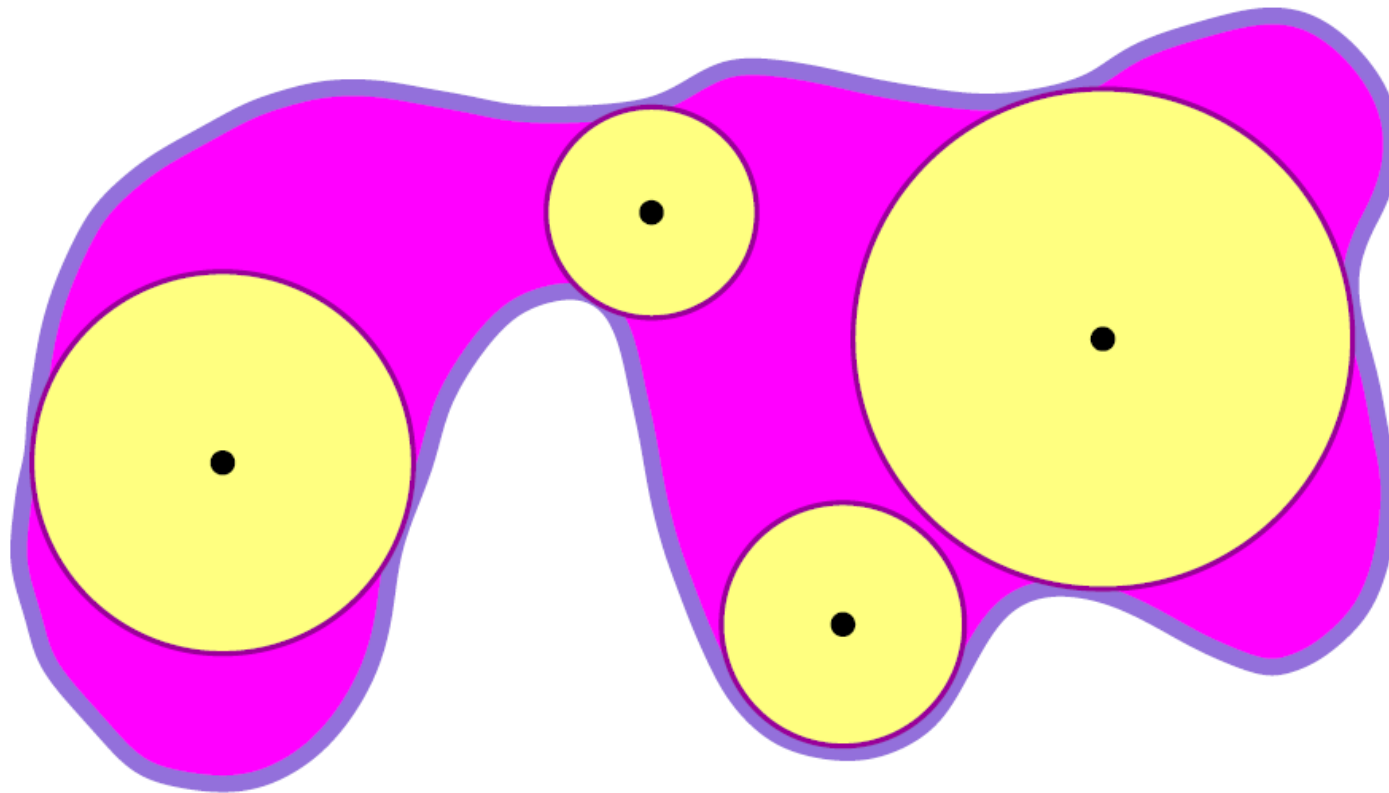
The centers of all such balls make up the *medial axis/skeleton*.



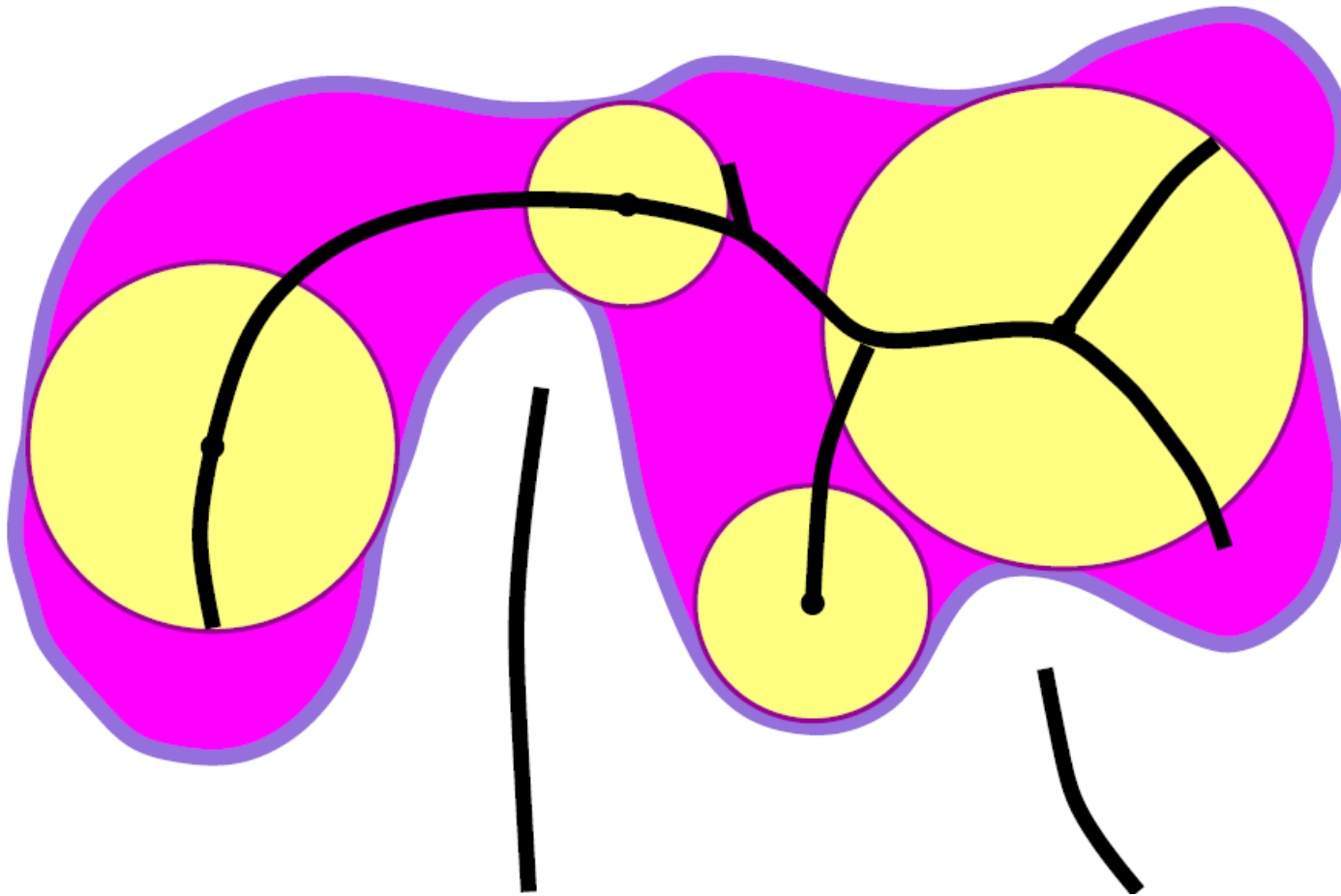
Medial Axis



Medial Axis



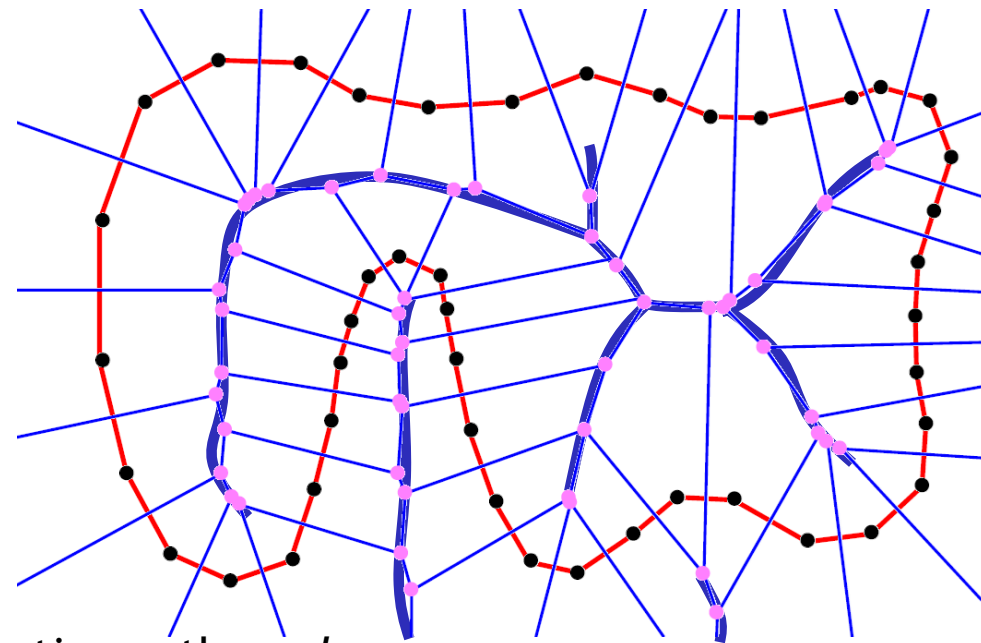
Medial Axis



Medial Axis

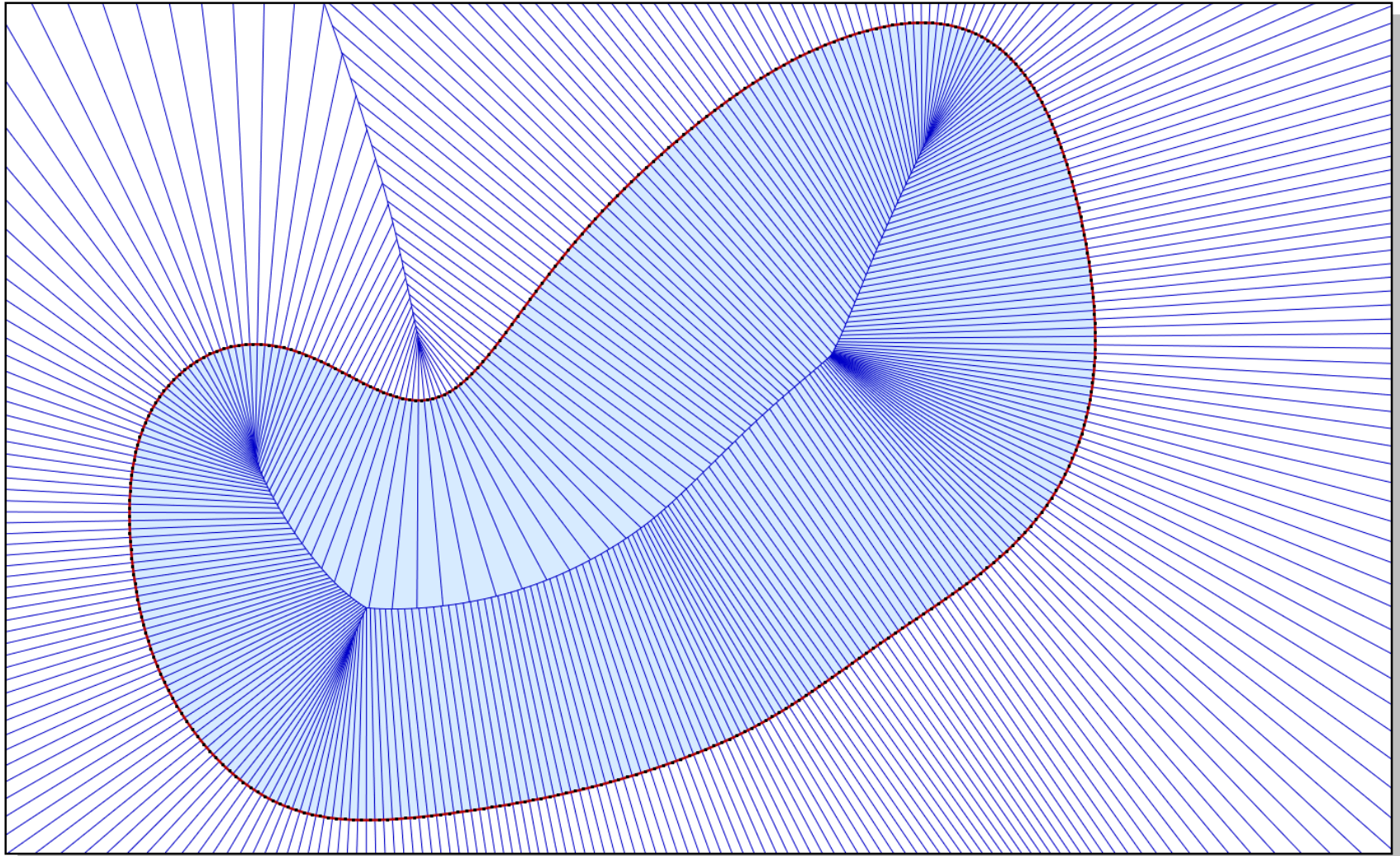
Observation*:

For a reasonable point sample, the medial axis is well-sampled by the Voronoi vertices.

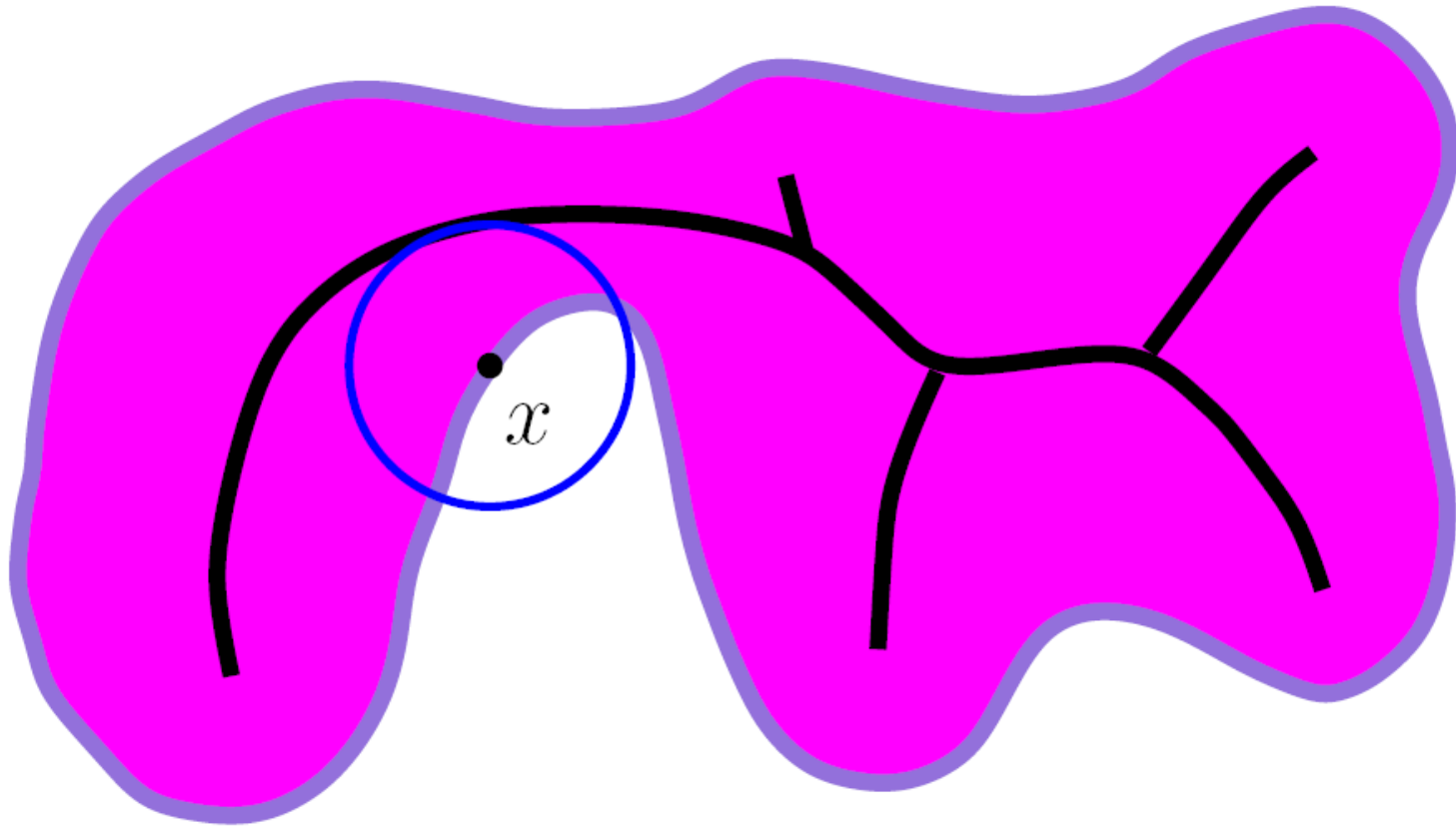


*In 3D, this is only true for a subset of the Voronoi vertices - the *poles*.

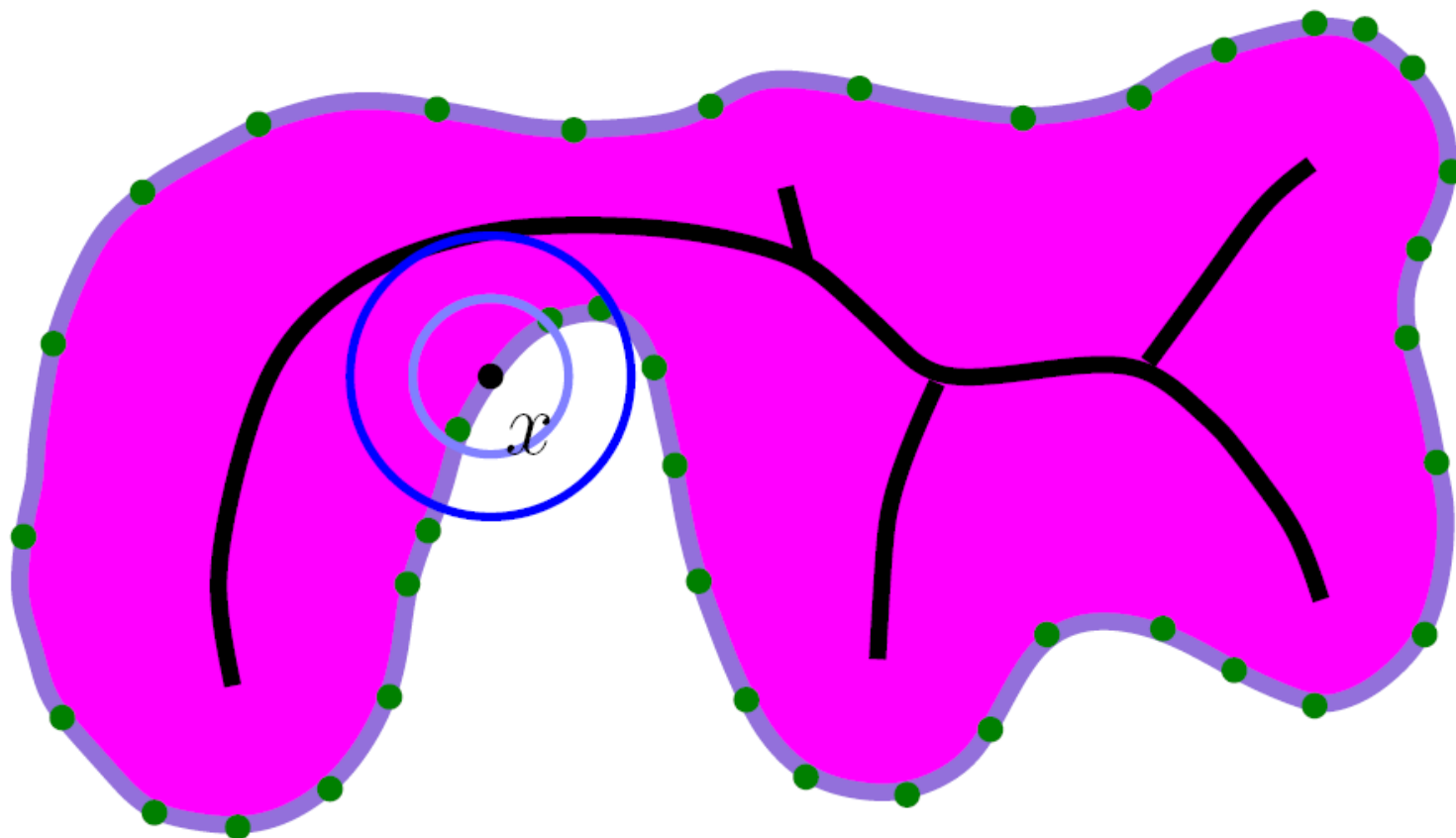
Voronoi & Medial Axis



Local Feature Size



Epsilon-Sampling



Crust [Amenta et al. 1998]

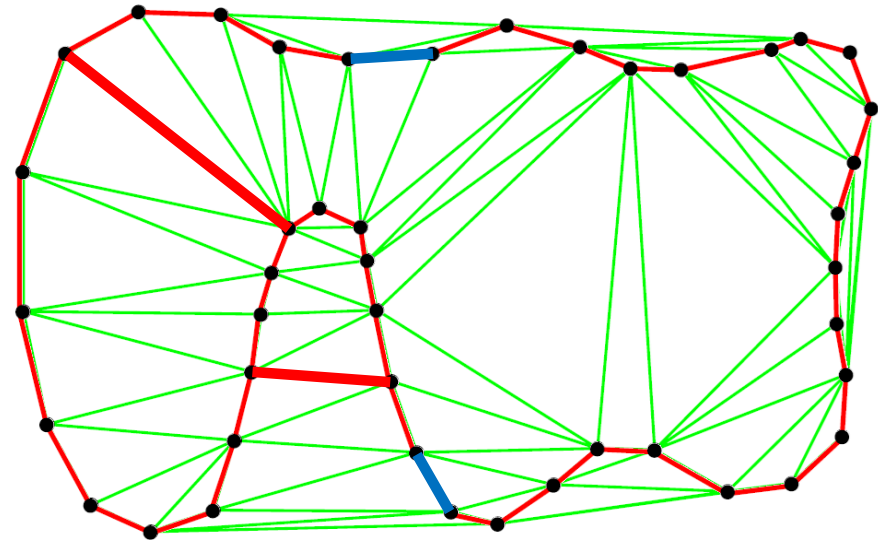
If we consider the Delaunay Triangulation of a point set, the shape boundary can be described as a subset of the Delaunay edges.

Q: How do we determine which edges to keep?

A: Two types of edges:

1. Those connecting adjacent points on the boundary
2. Those traversing the shape.

Discard those that traverse.



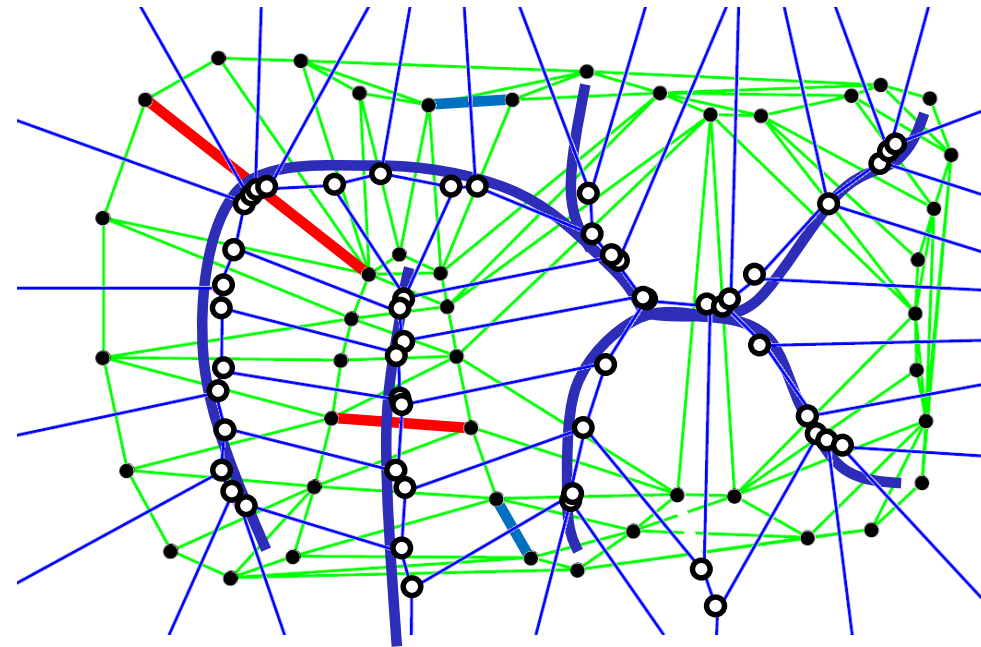
Crust [Amenta et al. 1998]

Observation:

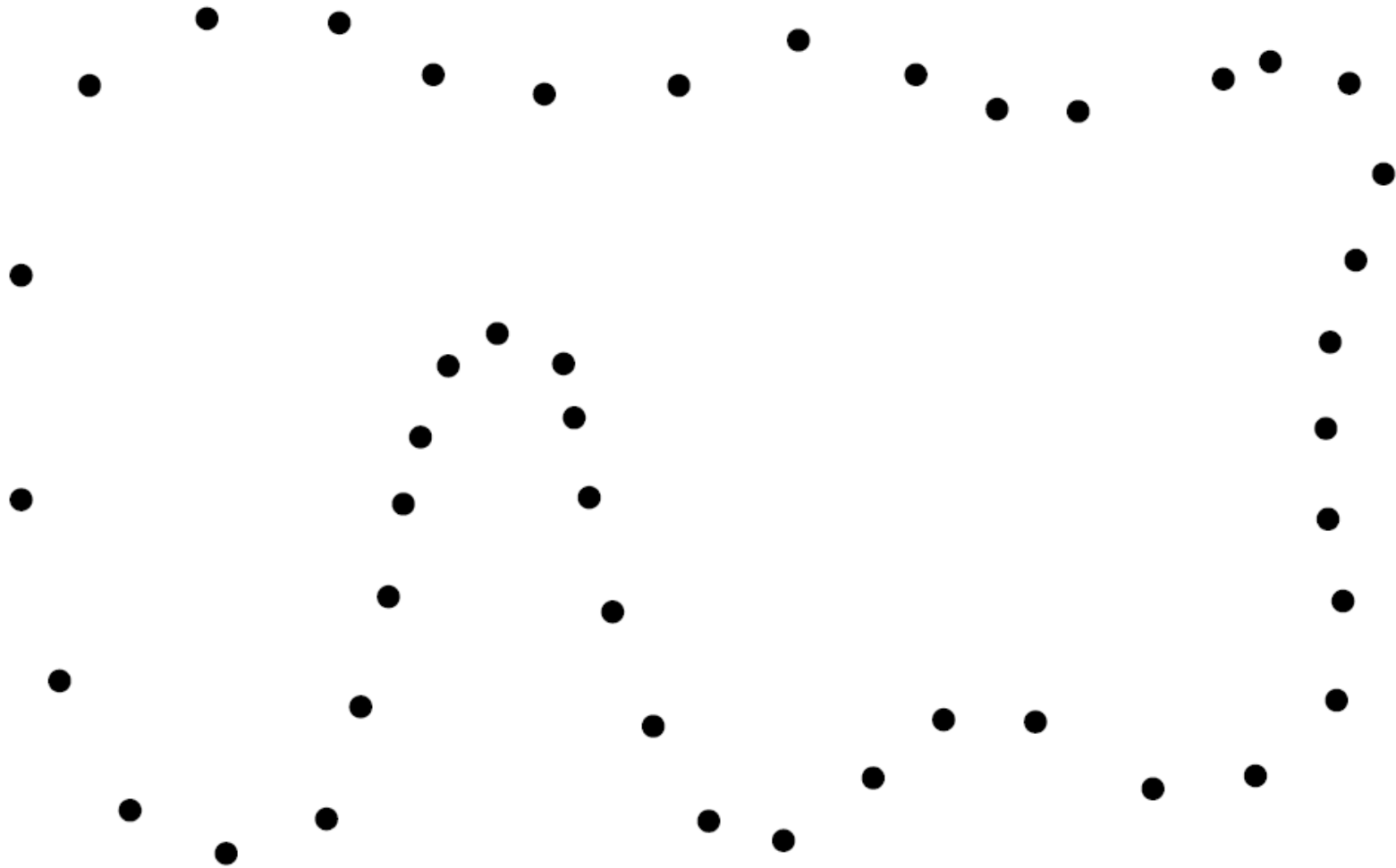
Edges that traverse cross the medial axis.

Although we don't know the axis, we can sample it with the Voronoi vertices.

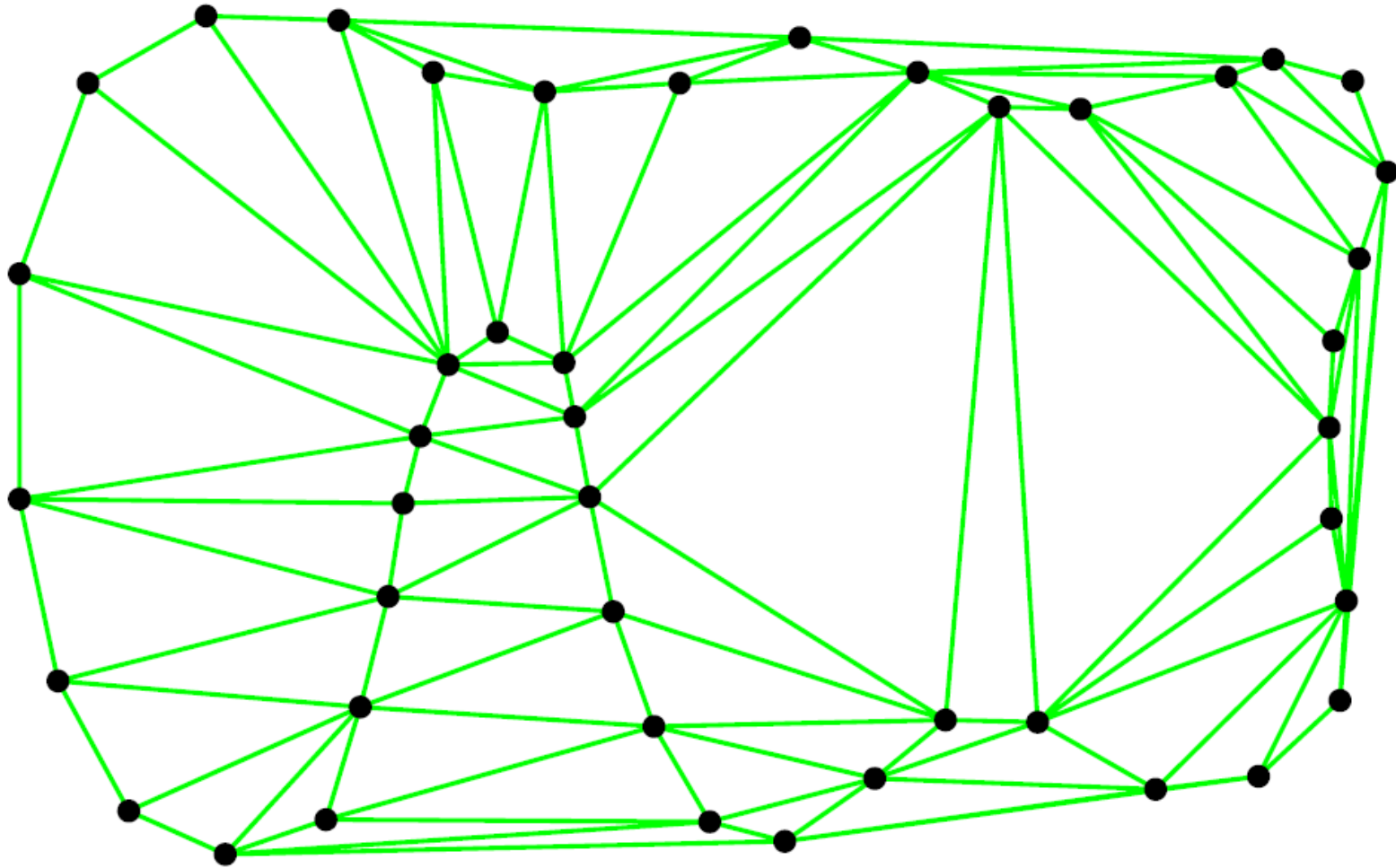
Edges that traverse must be near the Voronoi vertices.



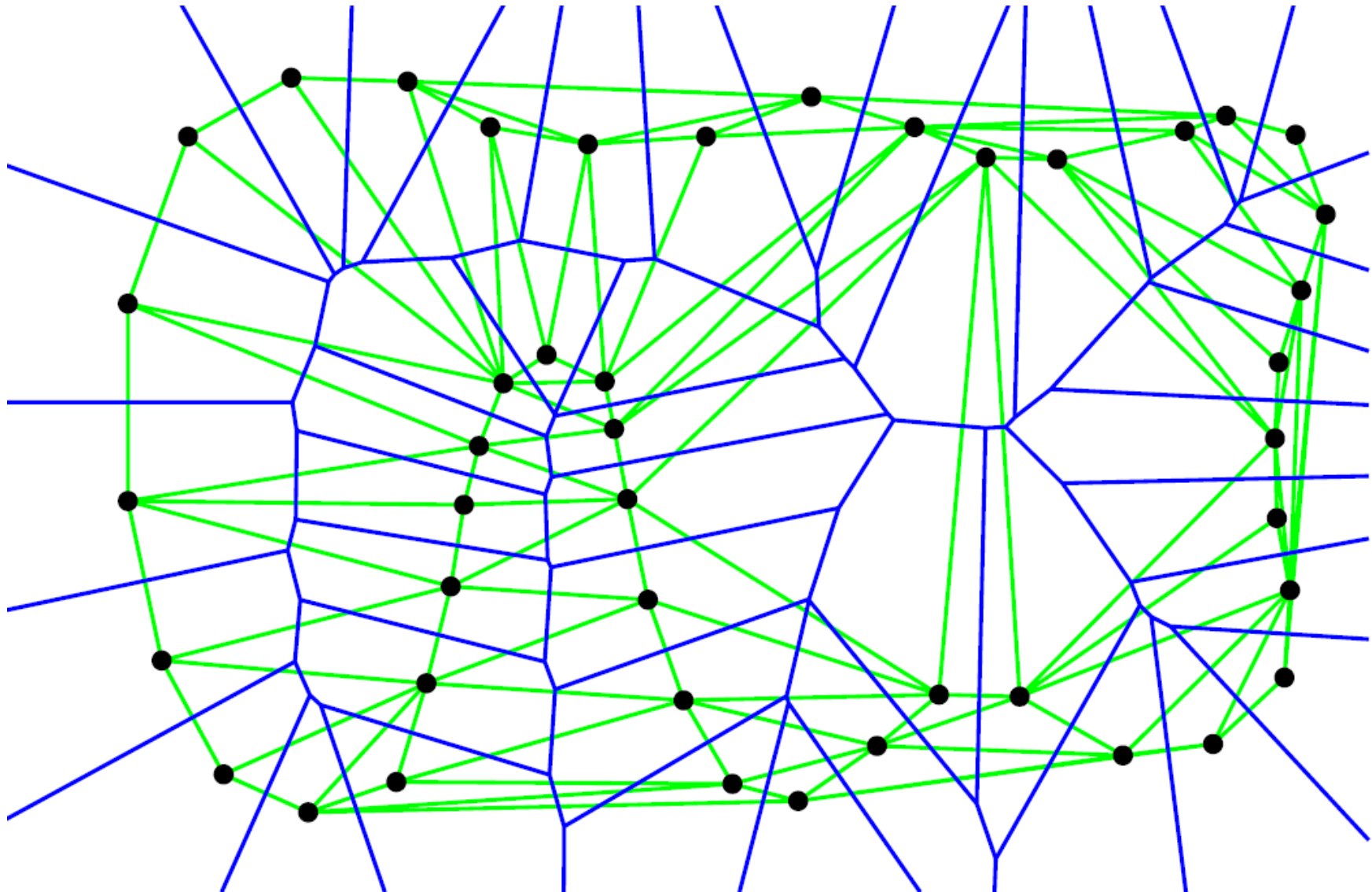
Crust [Amenta et al.]



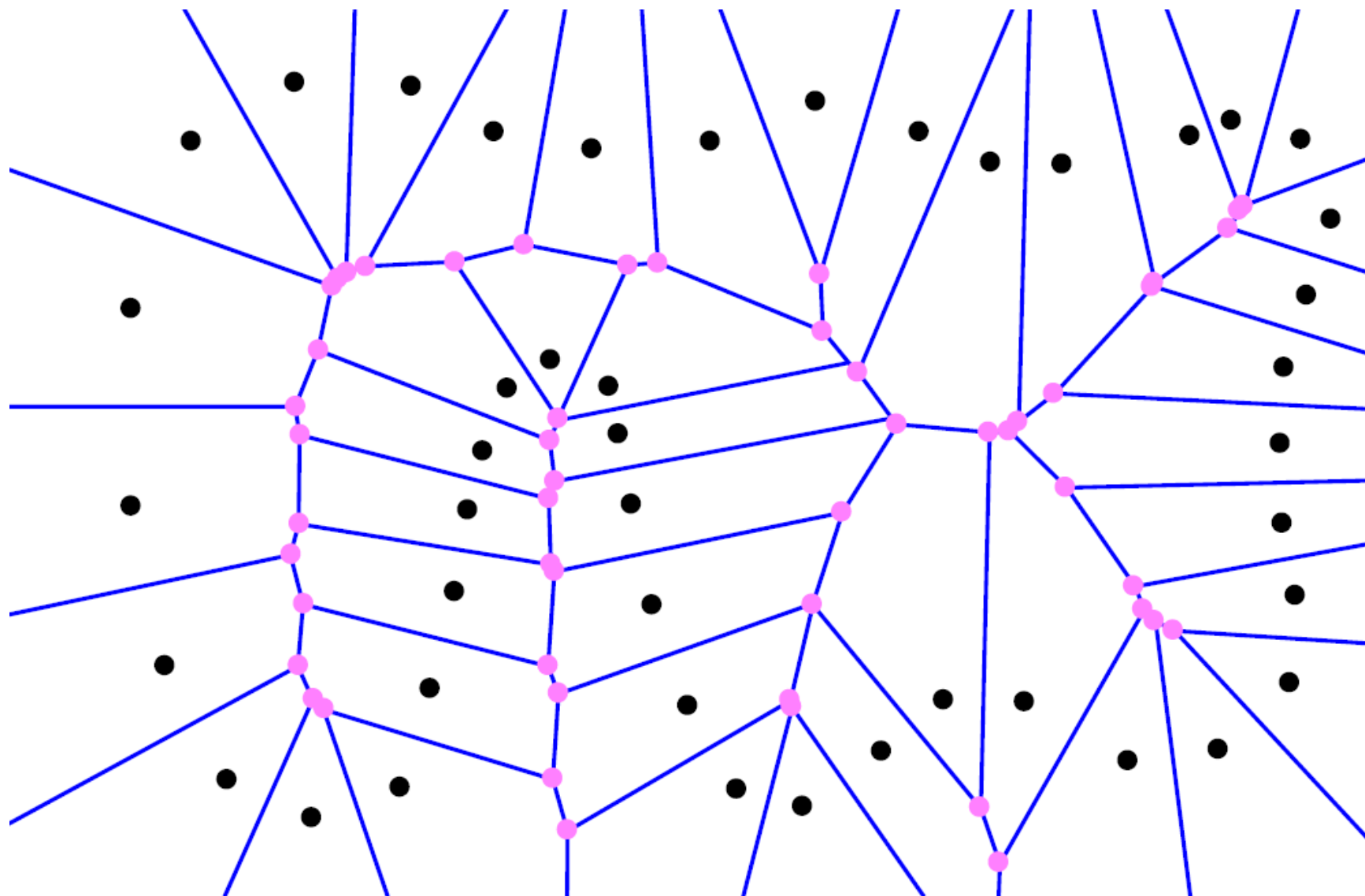
Delaunay Triangulation



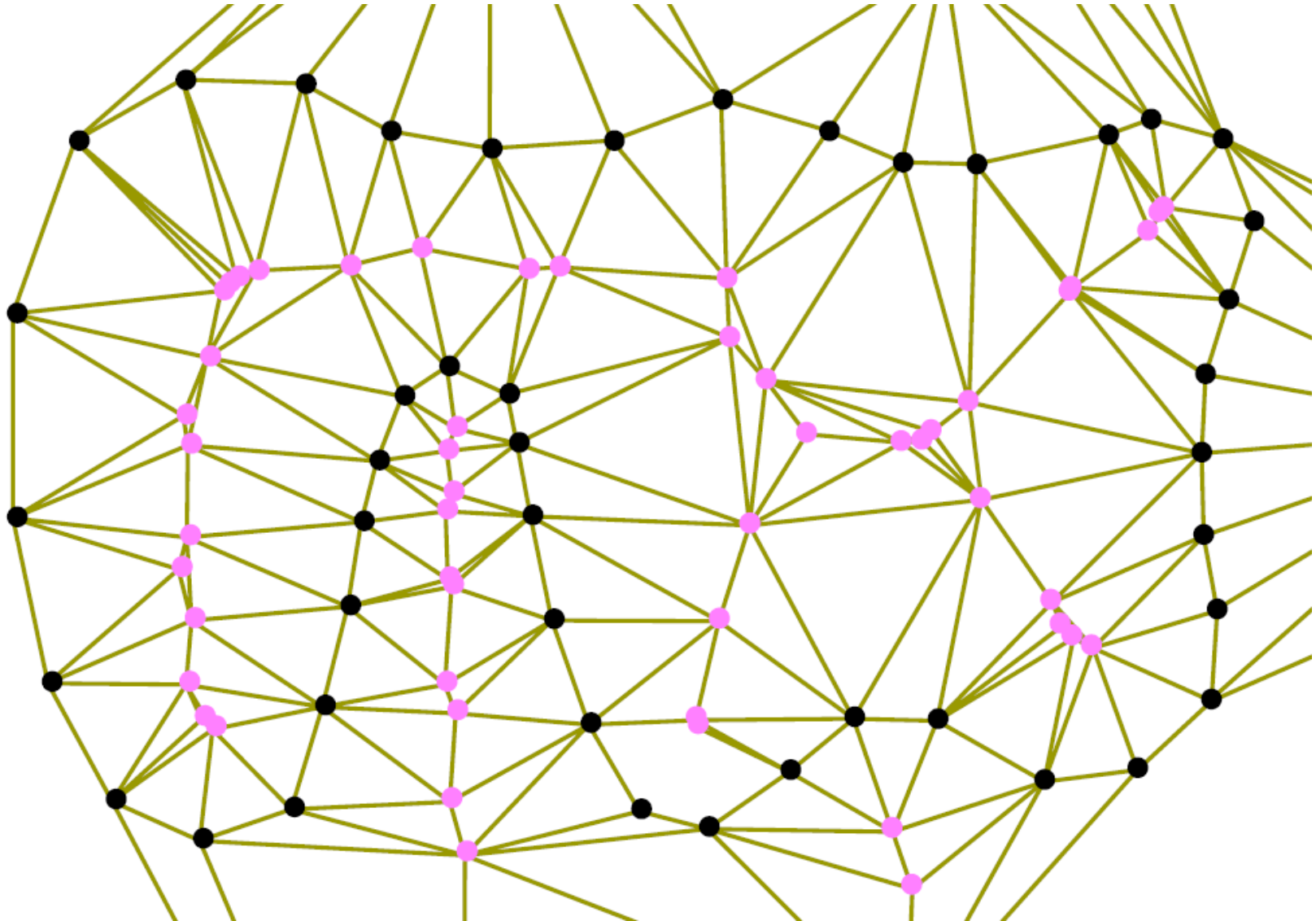
Delaunay Triangulation & Voronoi Diagram



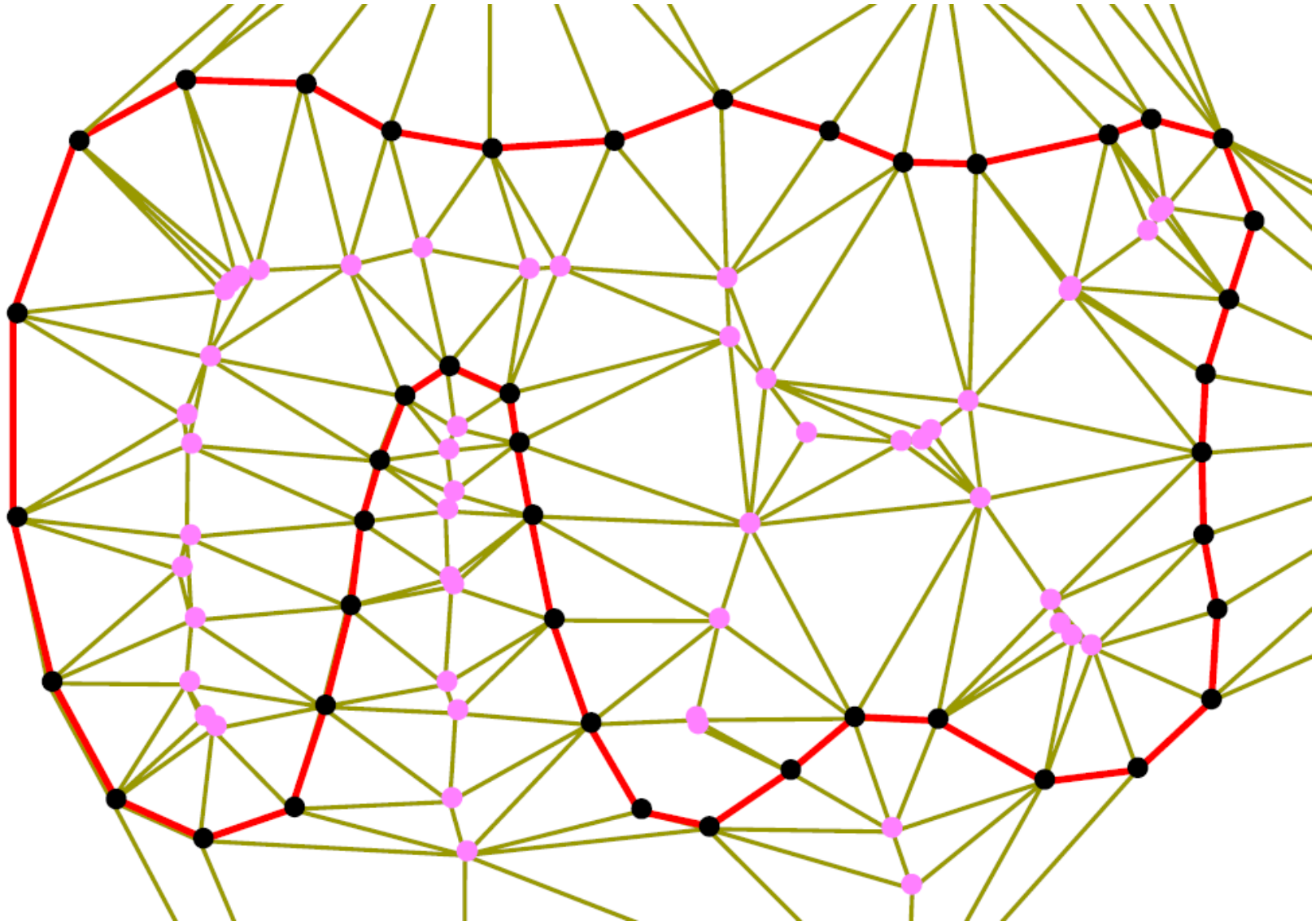
Voronoi Vertices



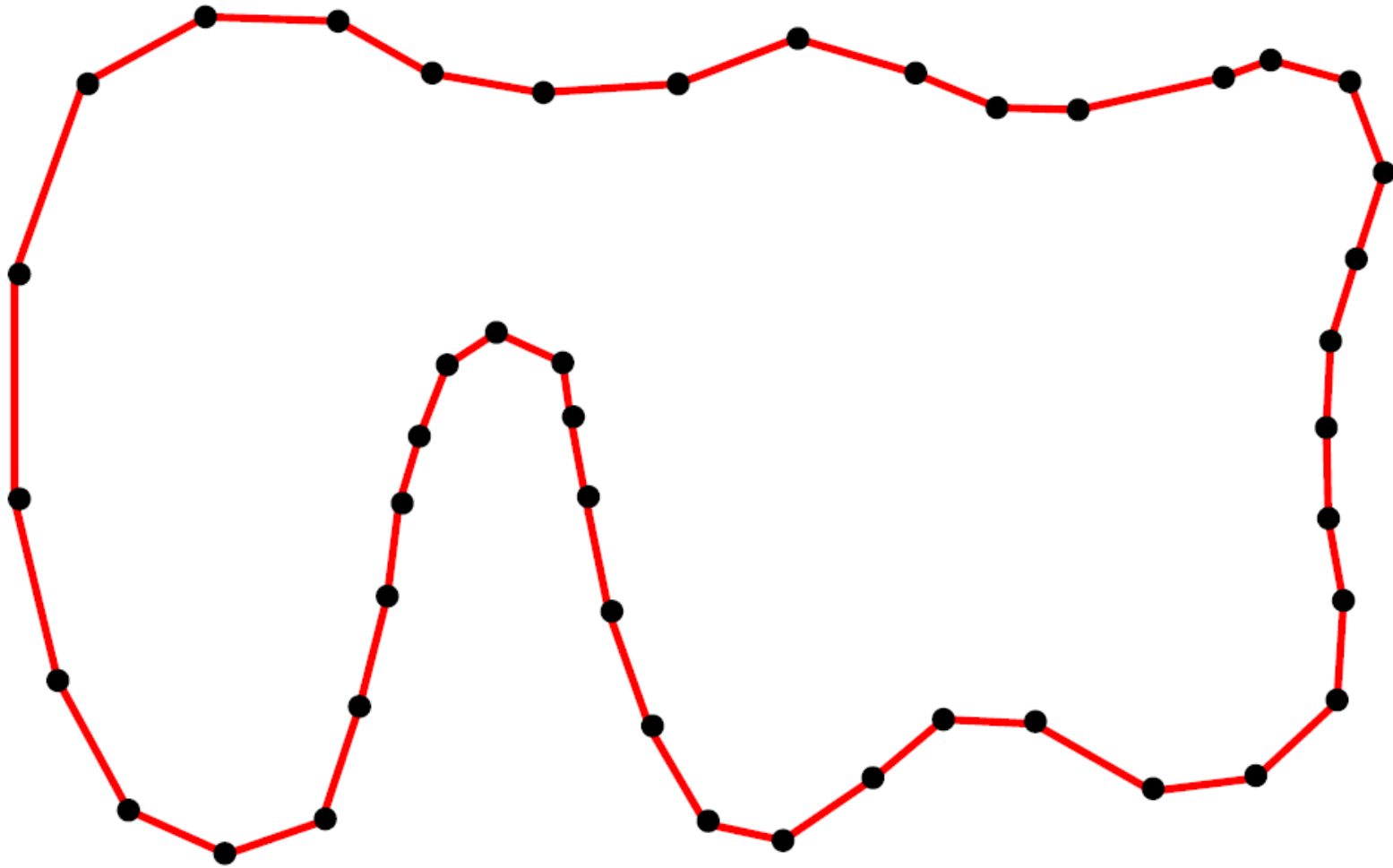
Refined Delaunay Triangulation



Crust



Crust



Crust (variant)

Algorithm:

1. Compute the Delaunay triangulation.
2. Compute the Voronoi vertices
3. Keep all edges for which there is a circle that contains the edge but no Voronoi vertices.

