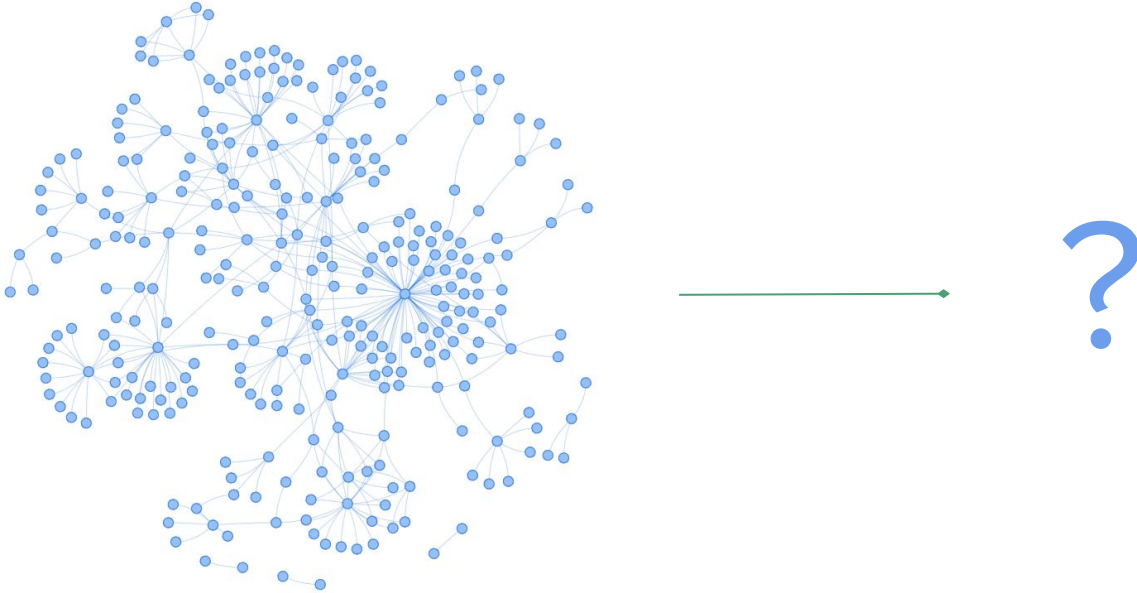


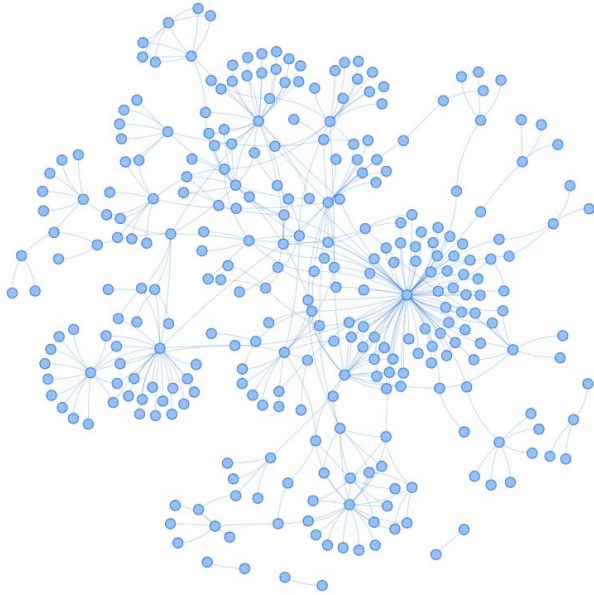
Explainable AI: Rule-aware Datalog Fact Explanation Using Group-SAT Solver

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Introduction: making sense of data

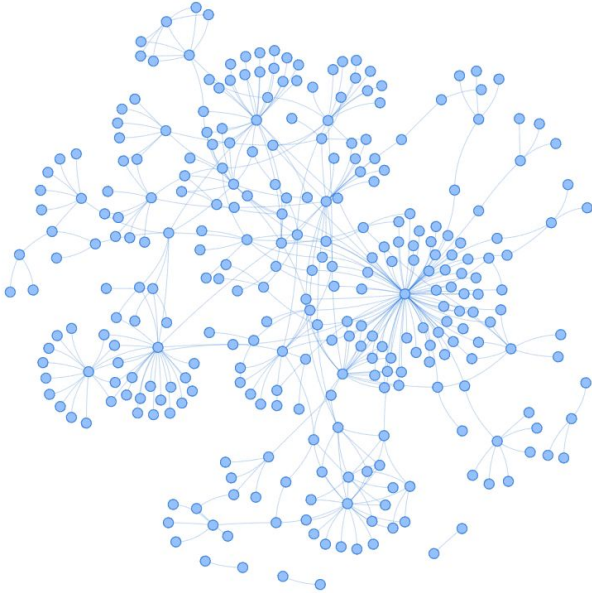


Introduction: making sense of data



? radioactive

Introduction: making sense of data

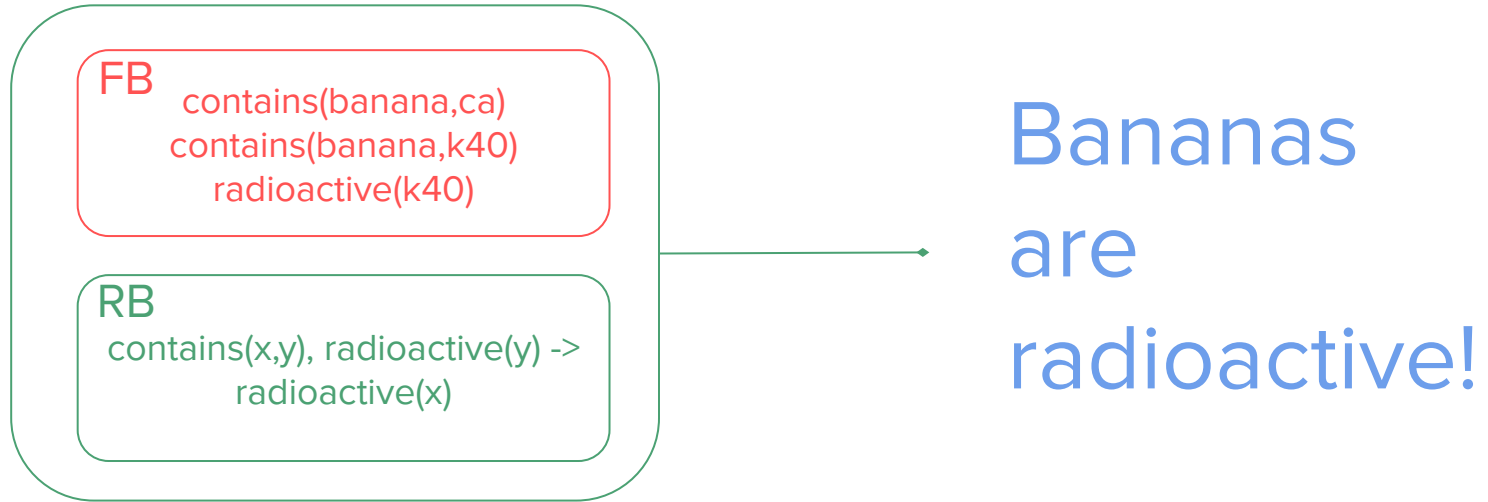


Bananas
are
radioactive!

Introduction: making sense of data



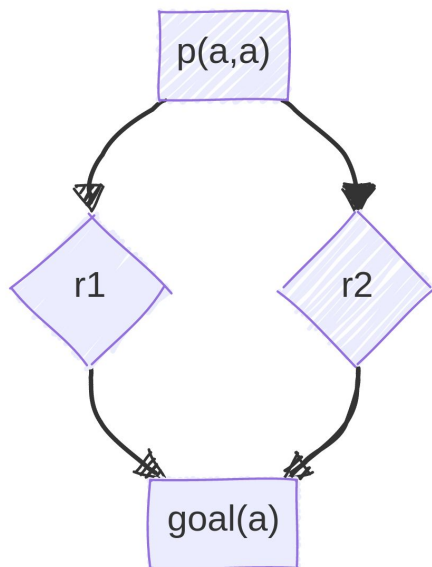
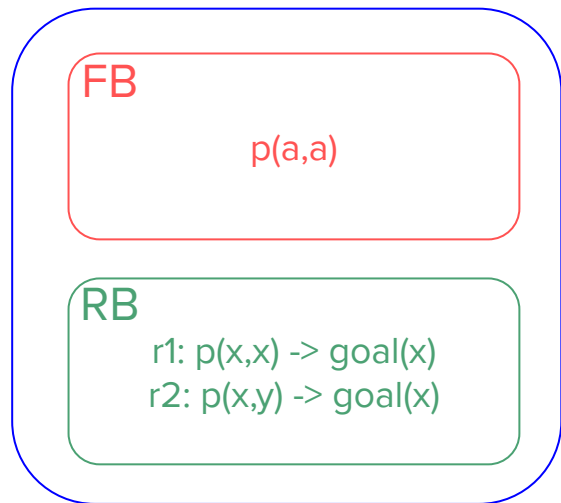
Introduction: making sense of data



Research Questions

- 1. What is a useful form of explanations?**
- 2. What is an efficient way to compute explanations?**

What is a 'good' explanation: fact-support vs KB-support



Explaining goal(a)

2 possible forms of explanations:

1. Fact-support (why-provenance)

1 explanation: p(a,a)

2. KB-support

2 explanations:

p(a,a), r1

p(a,a), r2

Research Questions

1. What is a useful form of explanations?

-> An explanation of a fact p with regard to KB K is a *minimal* subset of K that entails p .

2. What is an efficient way to compute KB-support explanations?

Research Questions

1. What is a useful form of explanations?

-> An explanation of a fact p is a *minimal* subset of KB that entails p .

2. What is an efficient way to compute KB-support explanations?

-> Hint: Explanations resemble MUS in SAT formula

Computing explanations: SAT Problem

Propositional formula:

$p \wedge (p \rightarrow q) \wedge (\text{NOT } q) \wedge s \wedge (q \vee \text{NOT } s) \wedge t$

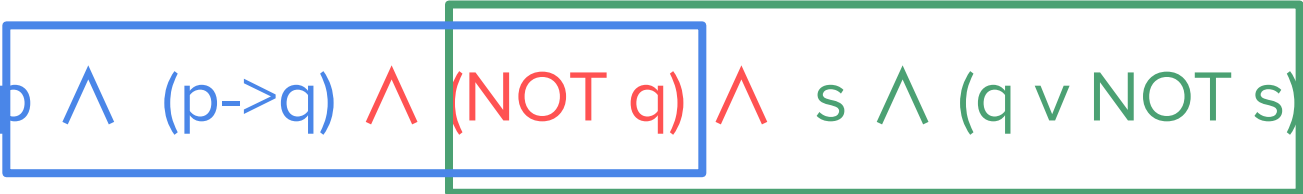
Computing explanations: MUS

MUS: minimal unsatisfiable subsets of clauses

$p \wedge (p \rightarrow q) \wedge (\text{NOT } q) \wedge s \wedge (q \vee \text{NOT } s) \wedge t$

Computing explanations: MUS

MUS: minimal unsatisfiable subsets of clauses

$$p \wedge (p \rightarrow q) \wedge (\text{NOT } q) \wedge s \wedge (q \vee \text{NOT } s) \wedge t$$


Research Questions

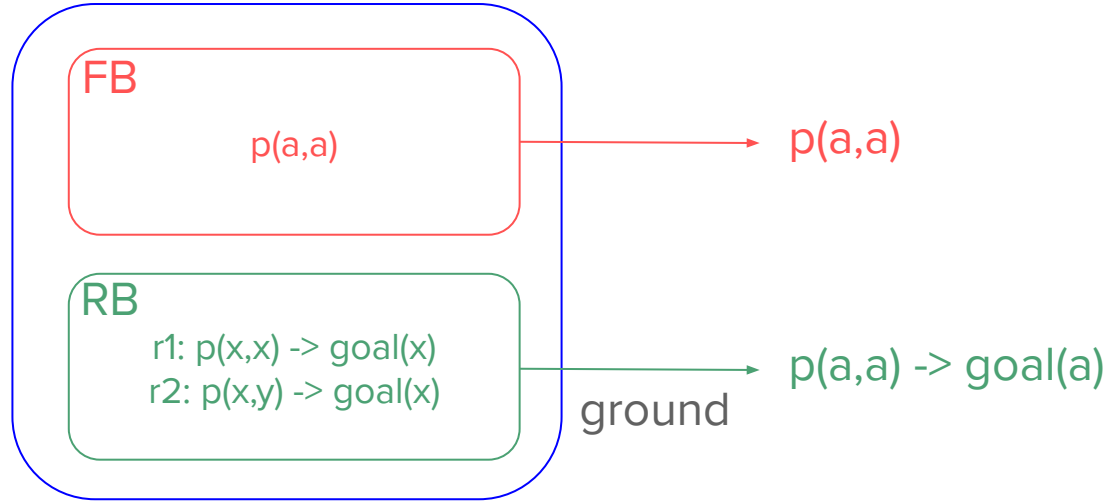
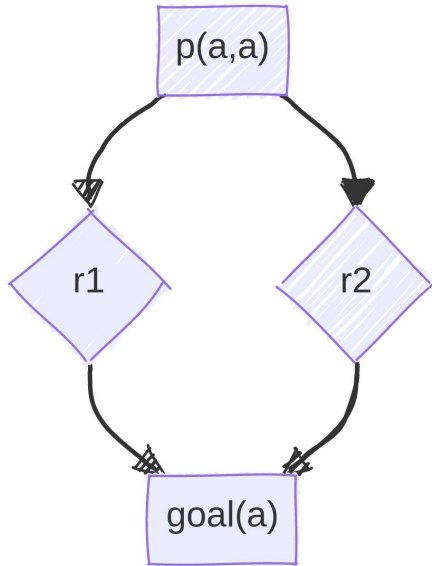
1. What is a useful form of explanations?

-> An explanation of a fact p is a *minimal* subset of KB that entails p .

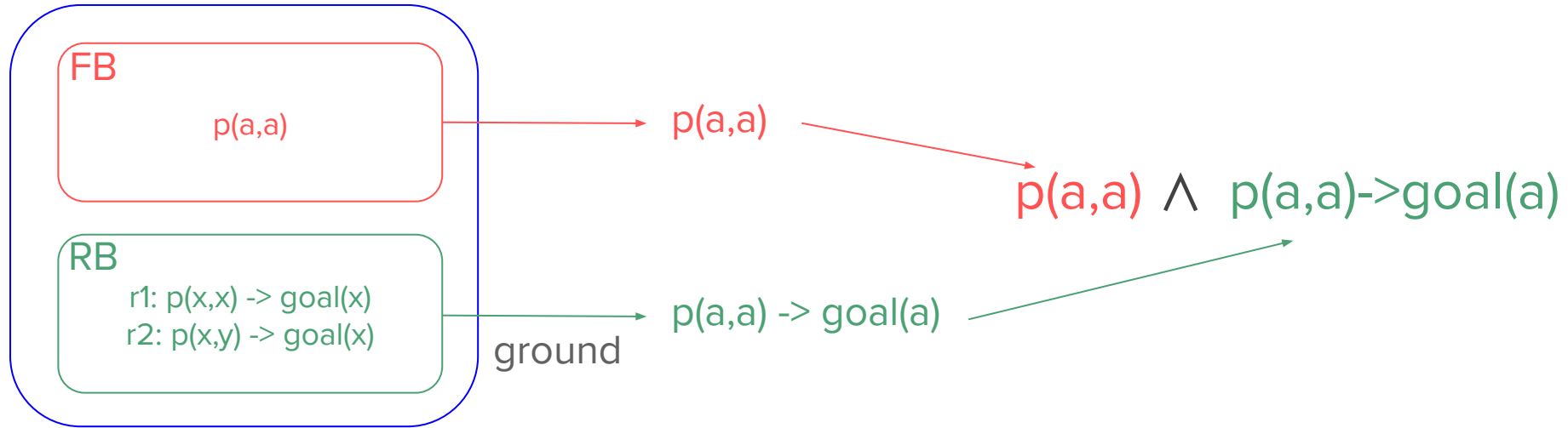
2. What is an efficient way to compute explanations?

-> Hint: Explanations resemble MUS in SAT formula

Building a SAT formula from a KB

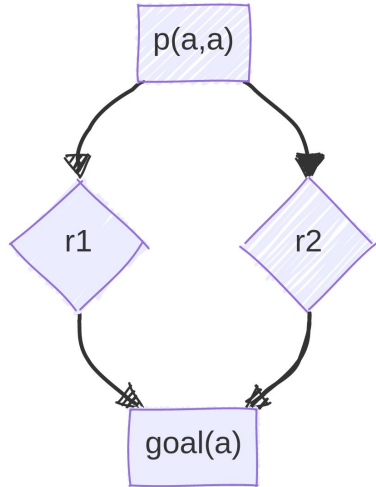


Building a SAT formula from a KB



Finally: reduction to MUS

r1: $p(x,x) \rightarrow \text{goal}(x)$
r2: $p(x,y) \rightarrow \text{goal}(x)$



Explaining $\text{goal}(a)$

$\text{NOT } \text{goal}(a)$

$\wedge p(a,a)$

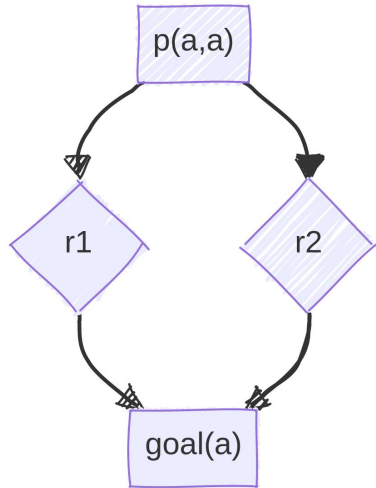
$\wedge p(a,a) \rightarrow \text{goal}(a)$

Computing explanations: reduction to MUS

r1: $p(x, \textcolor{blue}{x}) \rightarrow \text{goal}(x)$
r2: $p(x, \textcolor{red}{y}) \rightarrow \text{goal}(x)$

Encoding

MUS



Example 1

NOT $\text{goal}(a)$

$p(a, a)$

$p(a, a) \rightarrow \text{goal}(a)$

MUS
Solver

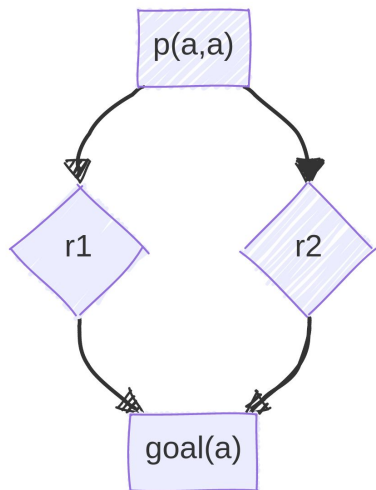
NOT $\text{goal}(a)$

$p(a, a)$

$p(a, a) \rightarrow \text{goal}(a)$

Computing explanations: reduction to (standard) MUS

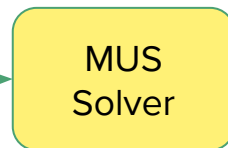
r1: $p(x, \textcolor{blue}{x}) \rightarrow \text{goal}(x)$
r2: $p(x, \textcolor{red}{y}) \rightarrow \text{goal}(x)$



Example 1

Encoding

NOT goal(a)
 $p(a,a)$
 $p(a,a) \rightarrow \text{goal}(a)$



MUS

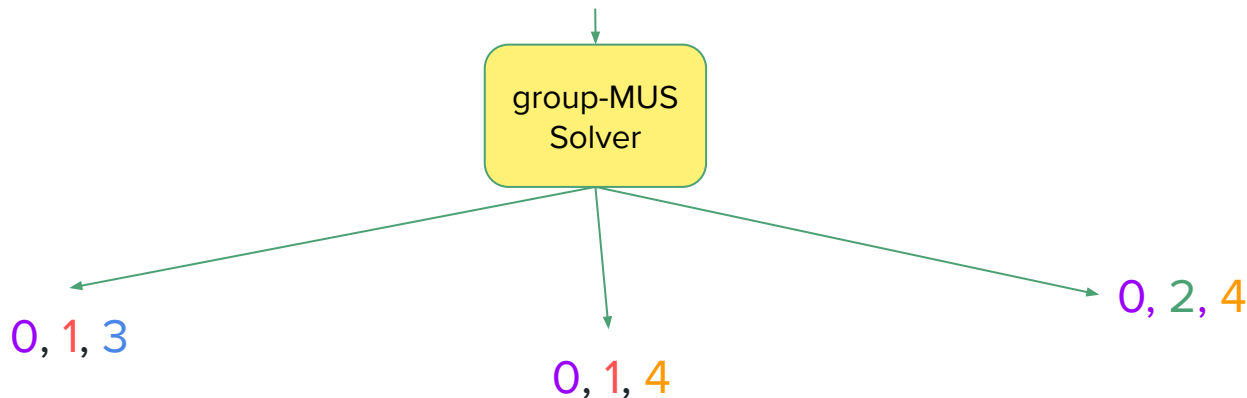
NOT goal(a)
 $p(a,a)$
 $p(a,a) \rightarrow \text{goal}(a)$

Issue with standard MUS:
1 MUS for 2 explanations

2, 4, 5 Computing explanations: group MUS

Group MUS: minimal unsatisfiable subsets of *groups of* clauses

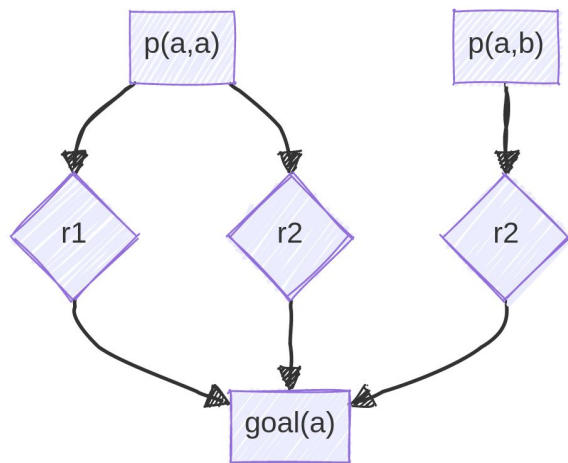
0 1 2 3 4
[NOT g(a)] $\wedge$ [p(a,a)] $\wedge$ [p(a,b)] $\wedge$ [p(a,a) $\rightarrow$ g(a)] $\wedge$ [p(a,a) $\rightarrow$ g(a), p(a,b) $\rightarrow$ g(a)]



Computing explanations: reduction to **group-MUS**

r1: $p(x,x) \rightarrow \text{goal}(x)$

r2: $p(x,y) \rightarrow \text{goal}(x)$



Example 1 (modified)

0

NOT goal(a)

1

p(a,a)

2

p(a,b)

3

p(a,a)→goal(a)

(group for r1)

4

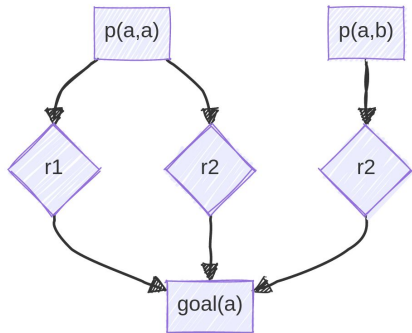
p(a,a)→goal(a)

p(a,b)→goal(a)

(group for r2)

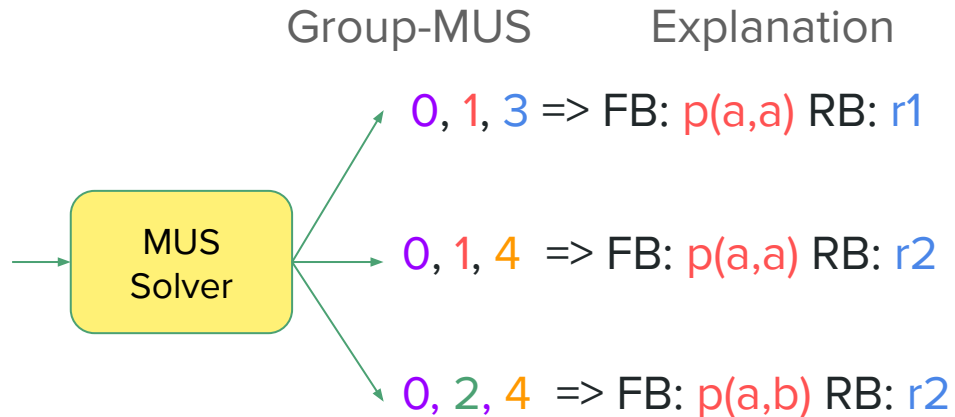
Computing explanations: reduction to **group-MUS**

r1: $p(x,x) \rightarrow \text{goal}(x)$
r2: $p(x,y) \rightarrow \text{goal}(x)$



Encoded KB

0 NOT goal(a)
1 $p(a,a)$
2 $p(a,b)$
3 $p(a,a) \rightarrow \text{goal}(a)$
4 $p(a,a) \rightarrow \text{goal}(a)$
4 $p(a,b) \rightarrow \text{goal}(a)$



Group MUS:
3 MUS for 3 explanations

Filtering: MUS solving is expensive

Computing explanation via MUS can be expensive

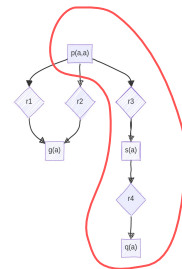
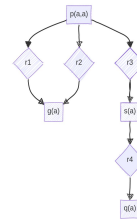
1. Complexity of group-MUS enumeration
2. KB can be very large

Filtering: MUS solving is expensive

Goal: Reduce the size of the encoded formula as much as possible while preserving soundness & completeness -> Find the relevant subset of the KB

2 steps:

1. Static step
 - entailment graph building
2. Dynamic step
 - fact relevance tracing



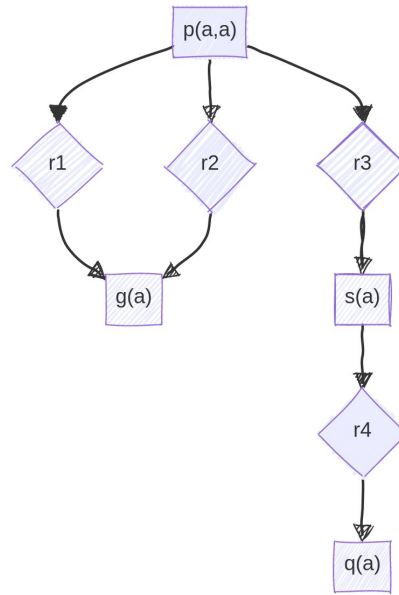
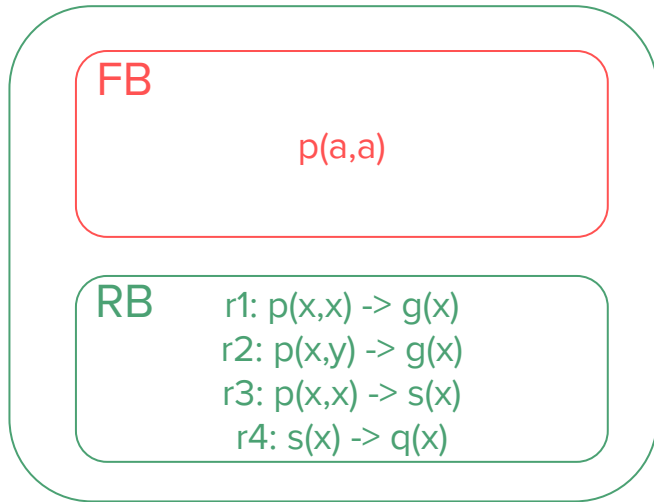
Static Step: Entailment Graph Building

Desiderata: a data structure to trace the lineage of an atom

Solution: use a rule-based encoding to build this data structure

Static Step: Entailment Graph Building

Desiderata: a data structure to trace the lineage of an atom



.dlgp graph

```
#FB
E r1(f p(a,a),f g(a))
E r2(f p(a,a),f g(a))
E r3(f p(a,a),f s(a))
E r4(f s(a),f q(a))

#RB
[REL(Y) ^ E r1(X1,Y) ->
REL(X1) ^ REL E r1(X1,Y) ,
...
REL(Y) ^ E r4(X1,Y) ->
REL(X1) ^ REL E r4(X1,Y) ]
```

Solution: use a rule-based encoding to build this data structure

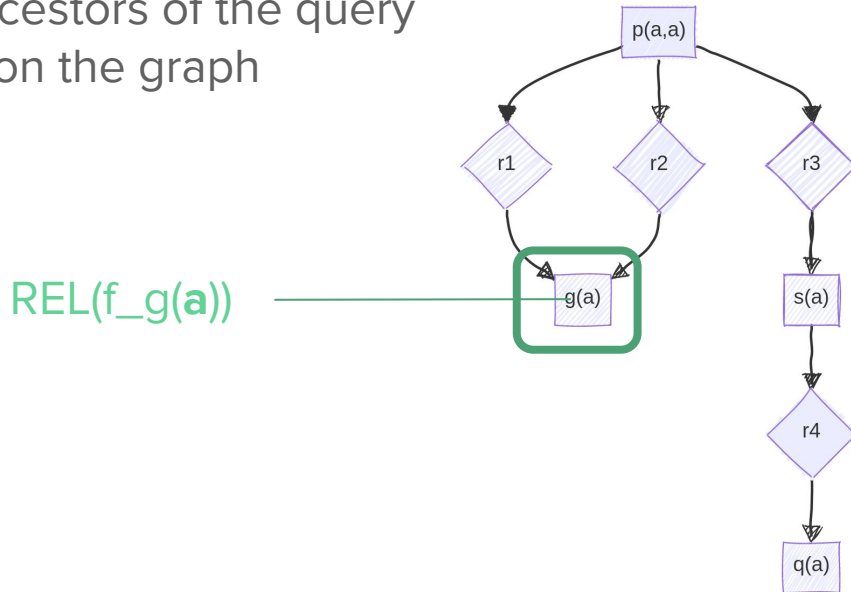
Dynamic Step: Fact Tracing

Desiderata: trace the
FB-ancestors of the query
atom on the graph

Solution: rule-based encoding of the tracing

Dynamic Step: Fact Tracing

Desiderata: trace the
FB-ancestors of the query
atom on the graph

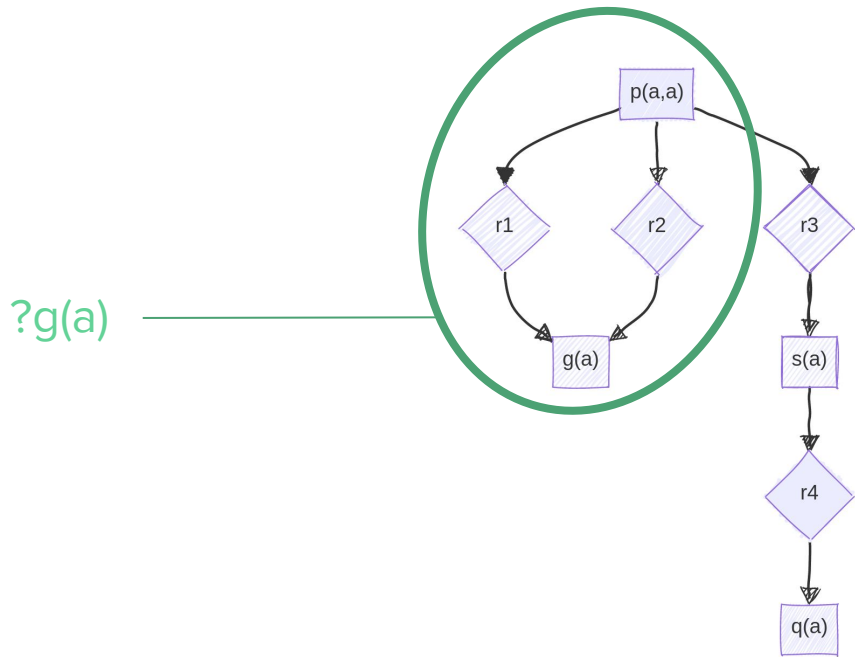


```
#FB  
E r1(f p(a,a), f g(a))  
E r2(f p(a,a), f g(a))  
E r3(f p(a,a), f s(a))  
E r4(f s(a), f q(a))  
REL(f g(a))
```

```
#RB  
[REL(Y) ^ E r1(X1,Y) ->  
REL(X1) ^ REL E r1(X1,Y),  
...  
REL(Y) ^ E r4(X1,Y) ->  
REL(X1) ^ REL E r4(X1,Y)]
```

Solution: rule-based encoding of the tracing

Dynamic Step: Fact Tracing



#FB

E r1(f p(a,a), f g(a))

E r2(f p(a,a), f g(a))

E r3(f p(a,a), f s(a))

E r4(f s(a), f q(a))

REL(f g(a))

#RB

[REL(Y) ^ E r1(X1,Y) ->

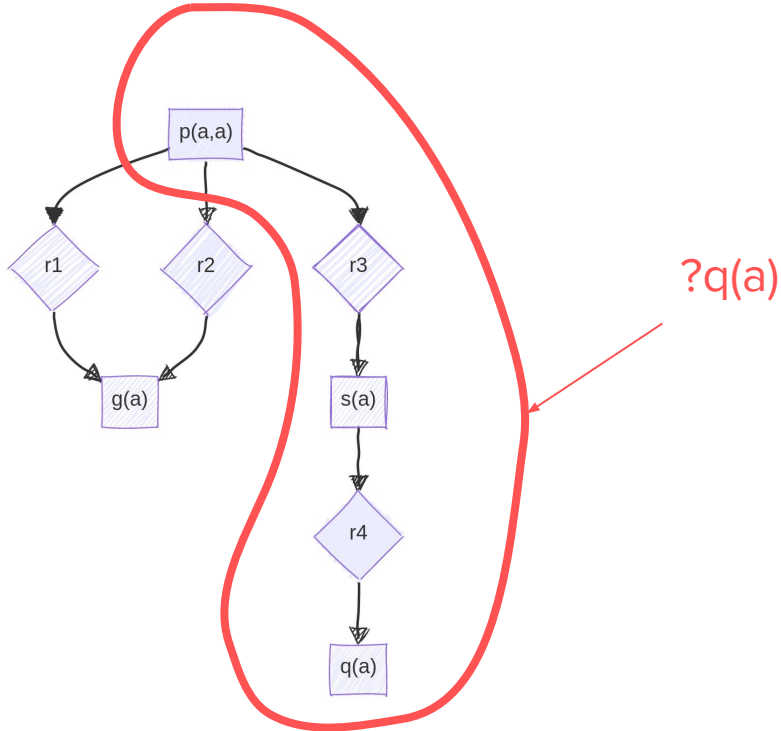
REL(X1) ^ REL E r1(X1,Y),

...

REL(Y) ^ E r4(X1,Y) ->

REL(X1) ^ REL E r4(X1,Y)]

Dynamic Step: Fact Tracing



#FB

E r1(f p(a,a), f g(a))

E r2(f p(a,a), f g(a))

E r3(f p(a,a), f s(a))

E r4(f s(a), f q(a))

REL(f q(a))

#RB

[REL(Y) ^ E r1(X1,Y) ->

REL(X1) ^ REL E r1(X1,Y) ,

...

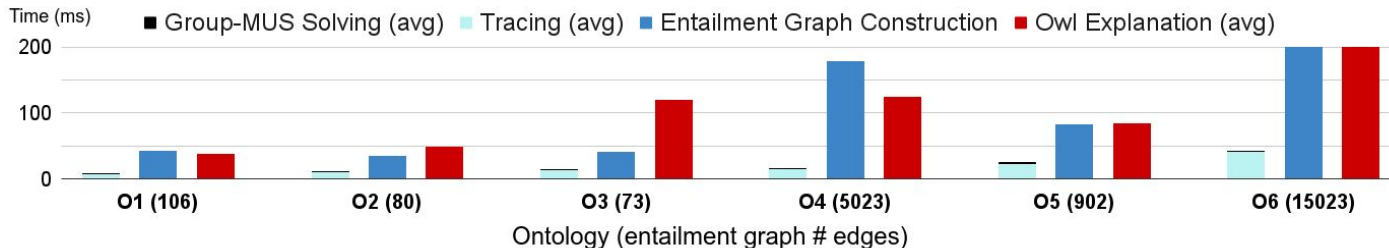
REL(Y) ^ E r4(X1,Y) ->

REL(X1) ^ REL E r4(X1,Y)]

Experiment

Preprocessing:

- 24 ontologies from MOWL corpus translated to .dlgp (via Carral et al.)
- Choose 5 facts from deepest reasoning level
- Run against OWL Explanation (owl api) tool



Conclusion: our approach is generally more competitive as the number of queries grows (and sometimes even for one query!)

Current Status & Future plans

- Formal/proof development
- Implementation/experiment
- Submission to RuleML24

2 future directions:

1. Extension
2. Optimisation & comparison with other approaches

Training

Past:

- RuleML Summer School (2023) - 34h
- French course - 30h

Ongoing:

- Research Integrity - 15h
- French course - 30h

Reference

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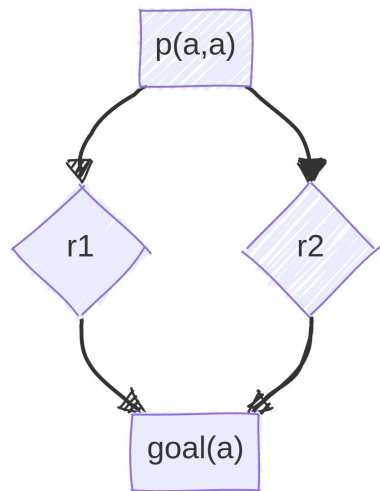
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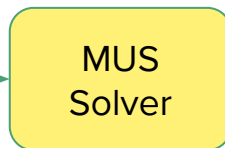
Computing explanations: reduction to **group-MUS**

r1: $p(x,x) \rightarrow \text{goal}(x)$
r2: $p(x,y) \rightarrow \text{goal}(x)$



Example 1

$\{\sim \text{goal}(a)\}$
 $\wedge \{p(a,a)\}_{f1}$
 $\wedge \{p(a,a) \rightarrow \text{goal}(a)\}_{r1}$
 $\wedge \{p(a,a) \rightarrow \text{goal}(a)\}_{r2}$



Group MUS:
2 MUS for 2 explanations

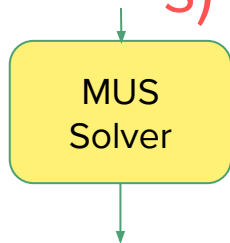
$\{\sim \text{goal}(a)\}$
 $\wedge \{p(a,a)\}_{f1}$
 $\wedge \{p(a,a) \rightarrow \text{goal}(a)\}_{r1}$

$\{\sim \text{goal}(a)\}$
 $\wedge \{p(a,a)\}_{f1}$
 $\wedge \{p(a,a) \rightarrow \text{goal}(a)\}_{r2}$

Computing explanations: MUS

MUS: minimal unsatisfiable subsets of clauses

$p \wedge (p \rightarrow q) \wedge (\text{NOT } q) \wedge s \wedge (q \vee \text{NOT } s) \wedge t$



MUS1: $(\text{NOT } q) \wedge p \wedge (p \rightarrow q)$

MUS2: $(\text{NOT } q) \wedge s \wedge (s \rightarrow q)$

Computing explanations: group MUS

Group MUS: minimal unsatisfiable subsets of *groups* of clauses

1 2 3 4 5 6
[p(a)] $\wedge$ [p(a)] $\wedge$ [p(b)] $\wedge$ [p(a),p(b)] $\wedge$ [p(a) $\wedge$ p(b) $\rightarrow$ q(a,b)] $\wedge$ [NOT

q(a,b)]

group-MUS
Solver

1, 3, 5, 6

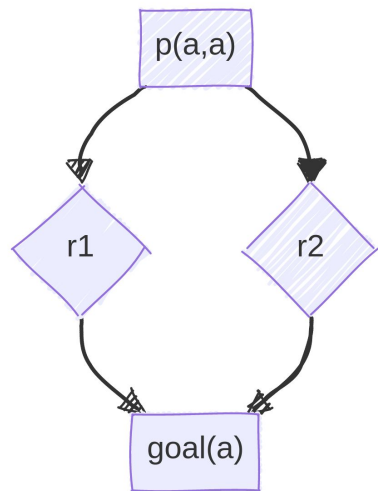
2, 3, 5, 6

4, 5, 6

Computing explanations: reduction to **group-MUS**

r1: $p(x,x) \rightarrow \text{goal}(x)$

r2: $p(x,y) \rightarrow \text{goal}(x)$



Example 1

$g1 : p(a,a) \rightarrow \text{goal}(a)$ (group for r1)

$g2 : p(a,a) \rightarrow \text{goal}(a)$ (group for r2)

$g3 : p(a,a)$

Computing explanations: reduction to **group-MUS**

r1: $p(x,x) \rightarrow \text{goal}(x)$
r2: $p(x,y) \rightarrow \text{goal}(x)$

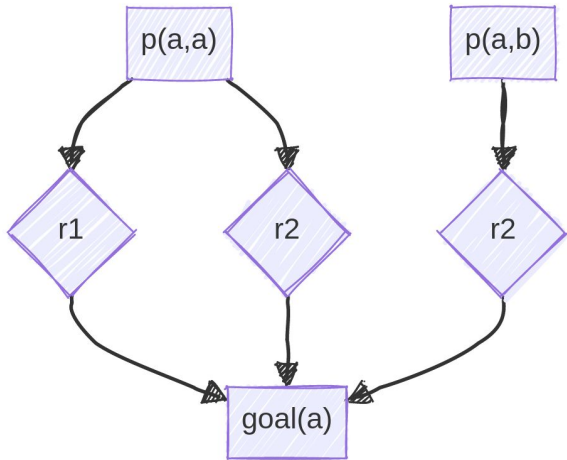
$g0 : \text{NOT goal}(a)$

$p(a,a)$

$p(a,b)$

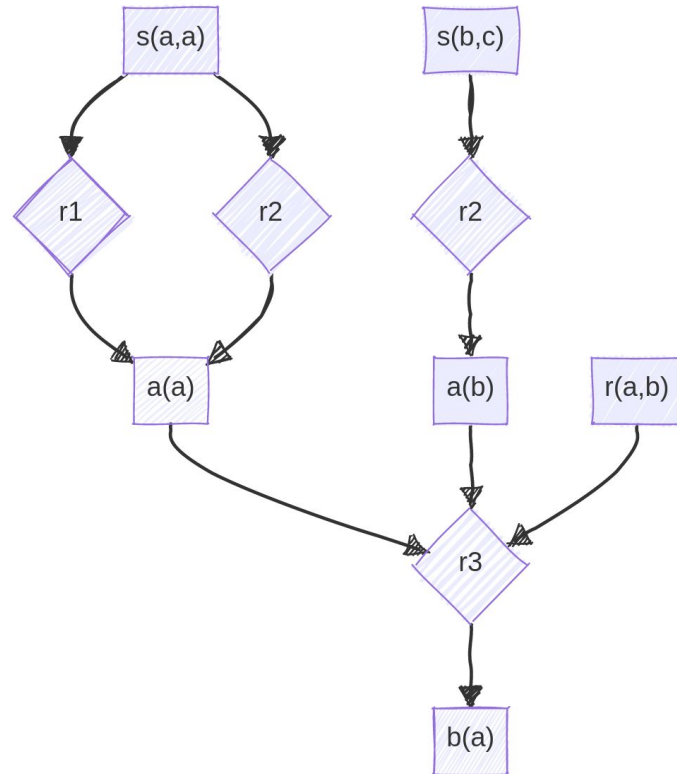
$p(a,a) \rightarrow \text{goal}(a)$ (group for r1)

$p(a,a) \rightarrow \text{goal}(a)$
 $p(a,b) \rightarrow \text{goal}(a)$ (group for r2)



Example 1
(modified)

Computing explanations: group MUS

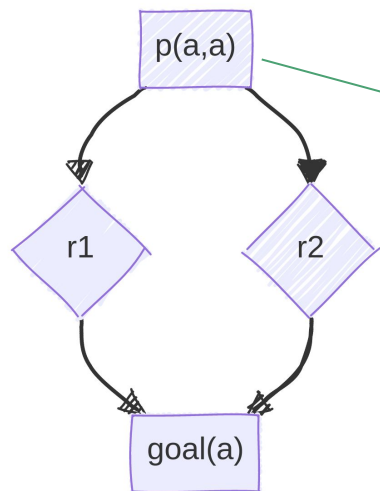


Computing explanations: reduction to MUS

r1: $p(x,x) \rightarrow \text{goal}(x)$

r2: $p(x,y) \rightarrow \text{goal}(x)$

How do we reduce a KB to SAT formula?



grounding

$p(a,a)$

$p(a,a) \rightarrow \text{goal}(a)$

Computing explanations: reduction to MUS

r1: $p(x, \mathbf{x}) \rightarrow \text{goal}(x)$
r2: $p(x, \mathbf{y}) \rightarrow \text{goal}(x)$

How do we reduce a KB to SAT formula?

contains(banana, ca)
contains(banana, k40)
radioactive(k40)

$\text{contains}(x, y), \text{radioactive}(y) \rightarrow$
 $\text{radioactive}(x)$

grounding

$p(a, a)$

$p(a, a) \rightarrow \text{goal}(a)$

