

# CGL 4 - Homology Inference

Clément Maria

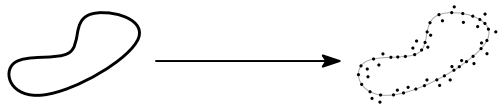
INRIA Sophia Antipolis-Méditerranée

MPRI 2020-2021

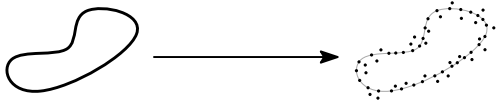
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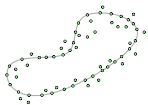
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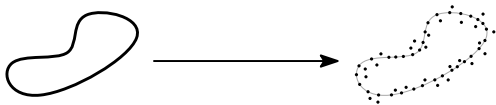
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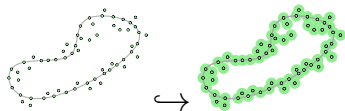
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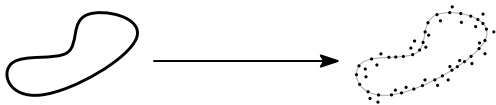
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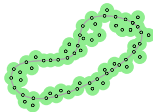
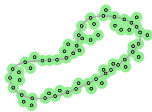
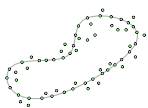
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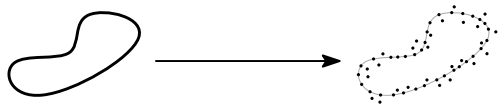
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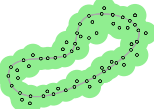
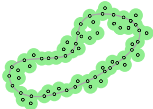
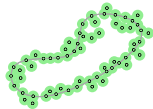
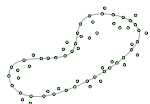
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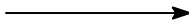
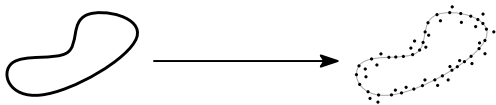
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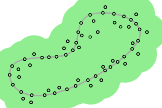
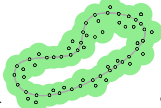
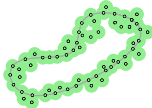
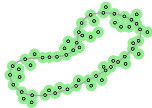
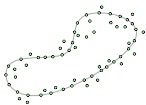
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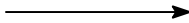
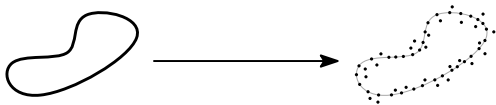


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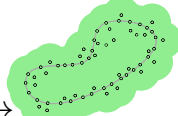
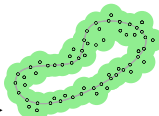
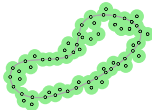
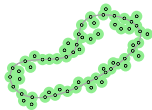
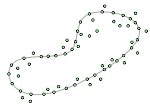




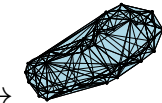
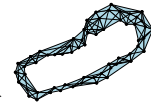
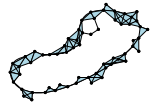
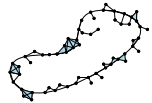
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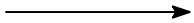
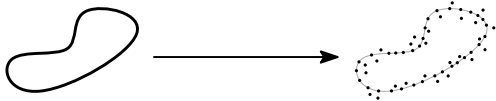
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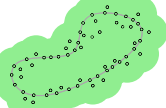
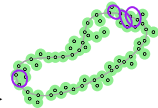
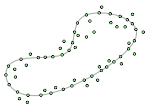
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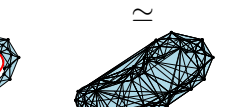
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$H_1(K_1)$



$H_1(K_2)$



$H_1(K_3)$

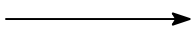
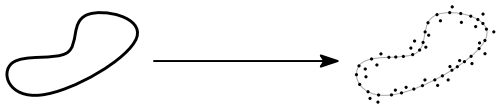


$H_1(K_4)$

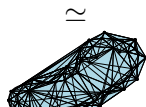
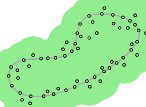
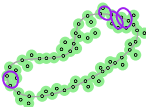
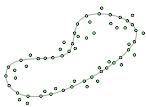


$H_1(K_5)$

# Homology Inference Problem



scale:



$H_1(K_1)$



$H_1(K_2)$



$H_1(K_3)$



$H_1(K_4)$



$H_1(K_5)$

barcode



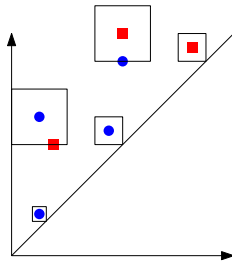
## Stability Theorems

Let  $D = \{(b_i, d_i)\}_{i \in I} \cup \{x = y\}$  and  $D' = \{(b'_j, d'_j)\}_{j \in J} \cup \{x = y\}$  be two persistence diagrams,  $b, d \in \mathbb{R}$ .

The **bottleneck distance**  $d_B(D, D')$  between  $D$  and  $D'$  is:

$$d_B(D, D') := \inf_{\substack{\Phi: D \rightarrow D' \\ \text{bijection}}} \sup_{p \in D} \|p - \Phi(p)\|_\infty$$

Defined even when  $|I| \neq |J|$  by sending points to the diagonal.



### Theorem (Stability on a simplicial complex)

Let  $f, g: \mathbf{K} \rightarrow \mathbb{R}$  be two functions on a same simplicial complex  $\mathbf{K}$ , inducing filtrations  $\mathbf{K}_\alpha = f^{-1}((-\infty; \alpha])$  and  $\mathbf{K}'_\gamma = g^{-1}((-\infty; \gamma])$ . Then

$$d_B(D(f), D(g)) \leq \|f - g\|_\infty.$$

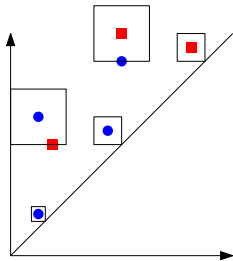
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### Theorem (Stability on general space)

Let  $f, g: M \rightarrow \mathbb{R}$  be two functions on a same metric space  $M$ , satisfying some “tameness” conditions. If  $\exists \varepsilon \geq 0$  s.t.  $\forall r \in \mathbb{R}$

$$f^{-1}(-\infty; r] \subseteq g^{-1}(-\infty; r + \varepsilon] \text{ and } g^{-1}(-\infty; r] \subseteq f^{-1}(-\infty; r + \varepsilon],$$

then  $d_B(D(f), D(g)) \leq \varepsilon$ .

# Distance between spaces

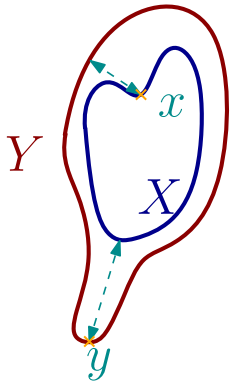
## Definition (Hausdorff distance)

The **Hausdorff distance** between two non-empty subsets  $X, Y$  of a metric space  $(M, d)$  is:

$$d_H(X, Y) := \max \left\{ \sup_{x \in X} d(x, Y), \sup_{y \in Y} d(X, y) \right\},$$

where,

$$d(x, Y) = \inf_{y \in Y} d(x, y) \text{ and } d(X, y) = \inf_{x \in X} d(x, y).$$



## Definition ( $\varepsilon$ -sample)

Let  $K \subset \mathbb{R}^D$  be a compact set. An  **$\varepsilon$ -sample** of  $K$ , for some  $\varepsilon \geq 0$ , is a finite set of points  $P$  such that  $d_H(P, K) \leq \varepsilon$ .



## Sampling compacts

Now consider compacts  $K$  in Euclidean space  $\mathbb{R}^D$ . Any compact  $K$  defines a function:

$$d_K: \mathbb{R}^D \rightarrow \mathbb{R}, \quad x \mapsto d_K(x) = d(x, K).$$

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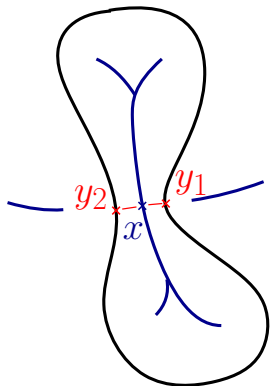
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### Definition (Critical point)

A point  $x \in \mathbb{R}^D$  is a **critical point** for  $d_K$  if there exist distinct points  $y_1, y_2 \in \mathbf{K}$  such that:

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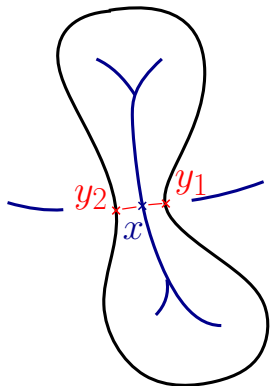
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### Definition (Reach of a compact)

The **reach** of a compact  $K$  is:

$$\text{reach } \mathbf{K} := \inf\{d_K(x) : x \text{ critical point for } d_K\}.$$



# Sampling, distance function approach

## Lemma

*Let  $P$  be an  $\varepsilon$ -sample of a compact  $K$ . Then,*

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Consequently, by triangle inequality:

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And vice versa, exchanging  $K$  and  $P$ .

# Reconstruction theorem

## Theorem (Reconstruction)

Let  $K \subset \mathbb{R}^D$  be a compact set, and  $\varepsilon > 0$  be such that:

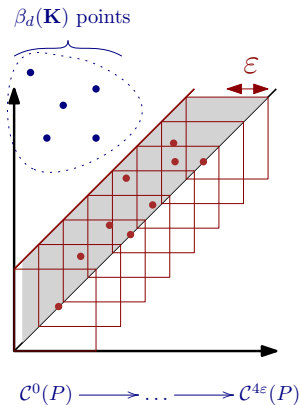
$$0 < 4\varepsilon < \text{reach } \mathbf{K}.$$

Let  $P$  be an  $\varepsilon$ -sample of  $K$ . Then,

- there is a *one-to-one correspondence* between the points of the persistence diagram of the Čech filtration:

$$\check{C}^0(P) \longrightarrow \dots \longrightarrow \check{C}^{4\varepsilon}(P)$$

that are *at least  $\varepsilon$ -away* from the diagonal  $\Delta$  in  $\infty$  norm, and the homology features of  $K$ .



# Reconstruction theorem

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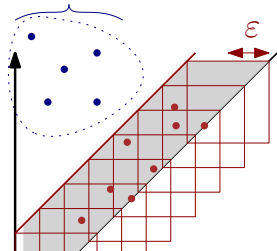
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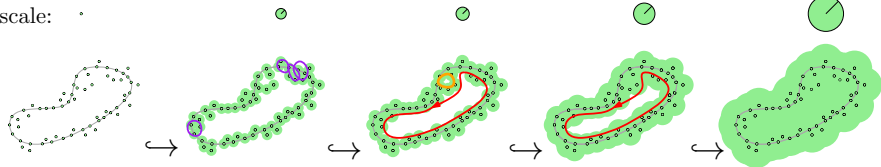
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$\beta_d(\mathbf{K})$  points



scale:  $\cdot$



barcodes





## Reconstruction theorem

*Sketch of proof:*

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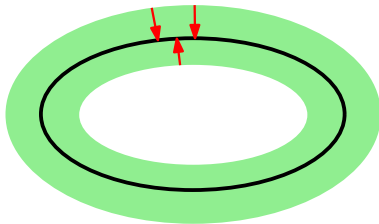
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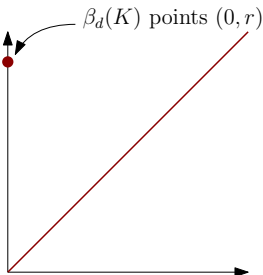
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—→ by projecting each  $x \in d_K^{-1}(-\infty; r] \setminus K$  to its **unique** nearest neighbor on  $K$ .

Consequently, for all  $0 < r_1 \leq r_2 < \text{reach}(K)$  we have

$$\mathbf{H}(d_K^{-1}(-\infty; r_1]) \xrightarrow{\text{id}} \mathbf{H}(d_K^{-1}(-\infty; r_2])$$



## Reconstruction theorem

*Sketch of proof:*

(2) Consider the persistent homology of

$$d_p^{-1}(-\infty; r] = \bigcup_{p \in P} \mathcal{B}_r(p) \simeq |\check{\mathcal{C}}^r(P)| \text{ (Nerve lemma).}$$

## Reconstruction theorem

*Sketch of proof:*

(2) Consider the persistent homology of

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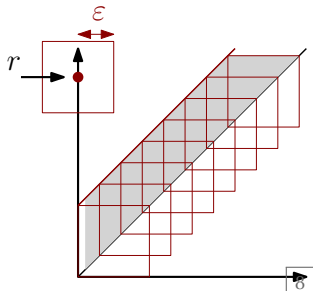
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For  $\text{reach}(K) > r > 4\varepsilon > 0$ , the radius  $r$   $\|\cdot\|_\infty$ -ball in  $\mathbb{R}^2$  around  $(0, r)$ , and the  $\|\cdot\|_\infty$ -band of width  $\varepsilon$  around the diagonal are disjoint.





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### Definition (Flag complex)

Let  $G = (V, E)$  be a graph. The **flag complex** induced by  $G$  is the abstract simplicial complex  $\mathcal{F}(G)$  satisfying:

- the vertices of  $\mathcal{F}(G)$  are  $V$ ,
- $\sigma = \{v_0, \dots, v_d\} \in \mathcal{F}(G)$  iff  $v_0, \dots, v_d$  form a clique in  $G$  (i.e., an induced complete subgraph of  $G$ ).

It is the maximal simplicial complex with  $G$  as 1-skeleton.

# Rips complex

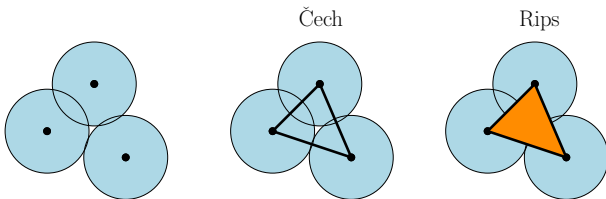
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## Definition (Rips complex)

Let  $P = \{x_1, \dots, x_n\}$  be a set of points in  $\mathbb{R}^N$ . The **neighbor graph of threshold  $r$**  of  $P$  is the graph  $N_r(P)$  with:

- vertices the points in  $P$ , and
- any two points  $x_i, x_j, i \neq j$  are connected by an edge iff  $\|x_i - x_j\| \leq r$ .

The **Rips complex** of threshold  $r$ , denoted by  $\mathcal{R}^r(P)$ , is the flag complex  $\mathcal{F}(N_{2r}(P))$ .



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**Computation:** Simply compute pairwise distances (low arithmetic complexity predicate), then expand the flag complex combinatorial

$\sim O(\#\text{points}^2) + \text{linear in the number of cliques/size of the complex.}$

## Rips vs Čech

### Lemma

*For a point cloud  $P$  in Euclidean space, and any  $r > 0$ :*

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Consequently, for all  $i$ :

$$\|z - x_i\| \leq \text{diam}(\text{CH}(x_0, \dots, x_k)) \leq \max_{i,j} \|x_i - x_j\| \leq 2r.$$

In conclusion,  $z \in \bigcap_{i=0 \dots k} \mathcal{B}_{2r}(x_i) \neq \emptyset$  and  $\sigma \in \check{C}^{2r}(P)$ .



## Ranks in sequence of morphisms

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Consider the following sequence of vector spaces and morphisms:

$$A \xrightarrow{\alpha} B \xrightarrow{\beta} C \xrightarrow{\gamma} D \xrightarrow{\delta} E \xrightarrow{\eta} F$$

such that  $\text{rk} ( A \xrightarrow{\eta \circ \dots \circ \alpha} F ) = \text{rk} ( C \xrightarrow{\gamma} D ) = k$ , then:

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## Cech persistence, a closer look

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Let  $P$  be an  $\varepsilon$ -sample of a compact  $K$  of  $\text{reach}(K) > 0$  in  $\mathbb{R}^N$ . Then, for any  $\varepsilon_0 > \varepsilon$ , and  $\alpha$ , such that  $\varepsilon \leq \alpha < \text{reach}(K) - 2\varepsilon_0$ , we have:

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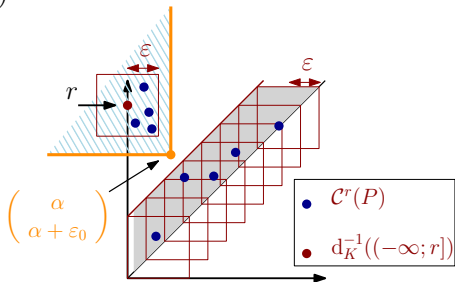
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## Theorem

*Let  $P$  be an  $\varepsilon$ -sample of a compact  $K$  s.t.  $0 < \varepsilon < \frac{\text{reach}(K)}{8}$ . The homology of  $K$  can be read off the persistence diagram of  $\mathcal{R}^r(P)$  for  $r \geq \varepsilon + 2\varepsilon_0$ , for any  $\varepsilon < \varepsilon_0 < \frac{\text{reach}(K)}{8}$ .*



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Now, with  $2\varepsilon + 4\varepsilon_0 < \text{reach}(K)$ ,

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In conclusion,  $\dim \mathbf{H}_d(K)$  is exactly the number of points in the persistence diagram of the Rips filtration (up to  $r = \varepsilon + 2\varepsilon_0$ ) in the upper left quadrant based at  $(\varepsilon, \varepsilon + 2\varepsilon_0)$ .

# Morse filtration

## Definition

Let  $\mathbf{K}$  be a complex, equipped with a filtration,

$$\emptyset := \mathbf{K}_0 \xrightarrow{\mathfrak{S}_1} \mathbf{K}_1 \xrightarrow{\mathfrak{S}_2} \mathbf{K}_2 \xrightarrow{\mathfrak{S}_3} \dots \xrightarrow{\mathfrak{S}_m} \mathbf{K}_m = \mathbf{K}$$

where  $\mathfrak{S}_i$  contains **at least** one simplex.

A **Morse filtration** for the filtration of  $\mathbf{K}$  is a collection of Morse matchings  $(X_i \sqcup T_i \sqcup S_i, \omega_i: T_i \rightarrow S_i)$  for each  $\mathbf{K}_i$ ,  $i = 0 \dots m$ , satisfying:

$$X_i \subseteq X_{i+1}, \quad T_i \subseteq T_{i+1}, \quad S_i \subseteq S_{i+1},$$

$$\omega_{i+1}|_{T_i} = \omega_i, \quad \partial^{X_{i+1}}|_{X_i} = \partial^{X_i}, \quad \forall i = 0 \dots m - 1.$$

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and a Morse filtration  $(X_i \sqcup T_i \sqcup S_i, \omega_i: T_i \rightarrow S_i)$ ,  $i = 0 \dots m$ :

$$\begin{aligned} X_i &\subseteq X_{i+1}, & T_i &\subseteq T_{i+1}, & S_i &\subseteq S_{i+1}, \\ \omega_{i+1}|_{T_i} &= \omega_i, & \partial^{X_{i+1}}|_{X_i} &= \partial^{X_i}, & \forall i &= 0 \dots m-1. \end{aligned}$$

Then the following persistence modules:

$$\mathbf{H}(\mathbf{K}_0, \partial) \xrightarrow{\subseteq} \mathbf{H}(\mathbf{K}_1, \partial) \xrightarrow{\subseteq} \mathbf{H}(\mathbf{K}_2, \partial) \xrightarrow{\subseteq} \dots \xrightarrow{\subseteq} \mathbf{H}(\mathbf{K}_m, \partial)$$

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is **well-defined** filtration of complexes inducing the persist. module:

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*Sketch of proof:*

For each  $i$  there is a chain equivalence:

$$(\mathbf{C}(\mathbf{K}_i), \partial) \xrightarrow{\phi_i} (\mathbf{C}(X_i), \partial^{X_i}) \xrightarrow{\psi_i} (\mathbf{C}(\mathbf{K}_i), \partial^{X_i})$$

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### Theorem

The following persistence modules are isomorphic:

$$\mathbf{H}(\mathbf{K}_0, \partial) \xrightarrow{\subseteq} \mathbf{H}(\mathbf{K}_1, \partial) \xrightarrow{\subseteq} \mathbf{H}(\mathbf{K}_2, \partial) \xrightarrow{\subseteq} \dots \xrightarrow{\subseteq} \mathbf{H}(\mathbf{K}_m, \partial)$$
$$\mathbf{H}(X_0, \partial^{X_0}) \xrightarrow{\subseteq} \mathbf{H}(X_1, \partial^{X_1}) \xrightarrow{\subseteq} \mathbf{H}(X_2, \partial^{X_2}) \xrightarrow{\subseteq} \dots \xrightarrow{\subseteq} \mathbf{H}(X_m, \partial^{X_m})$$

Sketch of proof:

For each  $i$  there is a chain equivalence:

$$(\mathbf{C}(\mathbf{K}_i), \partial) \xrightarrow{\phi_i} (\mathbf{C}(X_i), \partial^{X_i}) \xrightarrow{\psi_i} (\mathbf{C}(\mathbf{K}_i), \partial^{X_i})$$

where each chain map **commutes with inclusion**:

$$\begin{array}{ccc} \mathbf{C}(\mathbf{K}_i) & \xrightarrow{\subseteq} & \mathbf{C}(\mathbf{K}_{i+1}) \\ \phi_i \downarrow & & \downarrow \phi_{i+1} \\ \mathbf{C}(X_i) & \xrightarrow{\subseteq} & \mathbf{C}(X_{i+1}) \end{array}$$

$$\begin{array}{ccc} \mathbf{C}(\mathbf{K}_i) & \xrightarrow{\subseteq} & \mathbf{C}(\mathbf{K}_{i+1}) \\ \psi_i \uparrow & & \uparrow \psi_{i+1} \\ \mathbf{C}(X_i) & \xrightarrow{\subseteq} & \mathbf{C}(X_{i+1}) \end{array}$$

## Morse filtration

### Theorem

The following persistence modules are isomorphic:

$$\mathbf{H}(\mathbf{K}_0, \partial) \xrightarrow{\subseteq} \mathbf{H}(\mathbf{K}_1, \partial) \xrightarrow{\subseteq} \mathbf{H}(\mathbf{K}_2, \partial) \xrightarrow{\subseteq} \dots \xrightarrow{\subseteq} \mathbf{H}(\mathbf{K}_m, \partial)$$
$$\mathbf{H}(X_0, \partial^{X_0}) \xrightarrow{\subseteq} \mathbf{H}(X_1, \partial^{X_1}) \xrightarrow{\subseteq} \mathbf{H}(X_2, \partial^{X_2}) \xrightarrow{\subseteq} \dots \xrightarrow{\subseteq} \mathbf{H}(X_m, \partial^{X_m})$$

*Sketch of proof:*

For each  $i$  there is a chain equivalence:

$$(\mathbf{C}(\mathbf{K}_i), \partial) \xrightarrow{\phi_i} (\mathbf{C}(X_i), \partial^{X_i}) \xrightarrow{\psi_i} (\mathbf{C}(\mathbf{K}_i), \partial^{X_i})$$

→ enough prove commutativity for addition of a Morse pair  $(\tau, \sigma)$ :

$$(\mathbf{C}(X), \partial^X) \xrightarrow{\phi} (\mathbf{C}(X'), \partial^{X'}) \xrightarrow{\psi} (\mathbf{C}(X), \partial^X)$$

with  $X' = X - \{\sigma, \tau\}$  as used previously (same chain maps).

# Topology Inference Problem

compact  $K$  in Euclidean space



$P$   $\epsilon$ -sample with

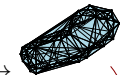
$$0 < \epsilon < \frac{\text{reach}(K)}{8}$$



scale:



Rips filtration



Morse filtration

$(X_1, \partial^{X_1})$

$(X_2, \partial^{X_2})$

$(X_3, \partial^{X_3})$

$(X_4, \partial^{X_4})$

$(X_5, \partial^{X_5})$

Persistent homology  $H_1(K_1)$

$\longrightarrow H_1(K_2)$

$\longrightarrow H_1(K_3)$

$\longrightarrow H_1(K_4)$

$\longrightarrow H_1(K_5)$

barcodes



$O(|P|^2) + \tilde{O}(n)$

$\mathcal{R}^{3.1\epsilon}$   
 $n$  simplices

$\tilde{O}(n)$

$m$  cells

$O(m^3)$