

Performance Evaluation

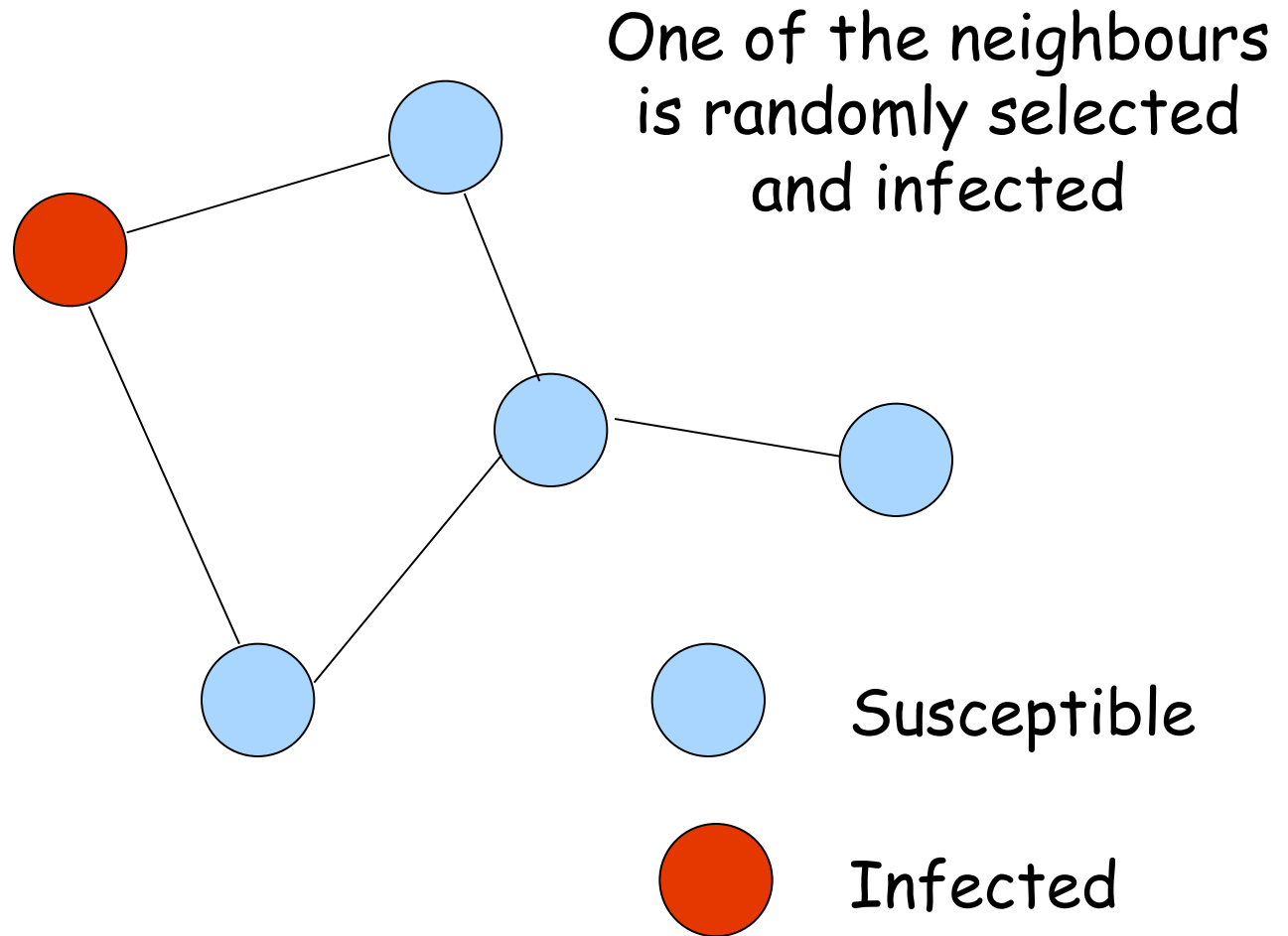
Lecture 2: Epidemics

Giovanni Neglia

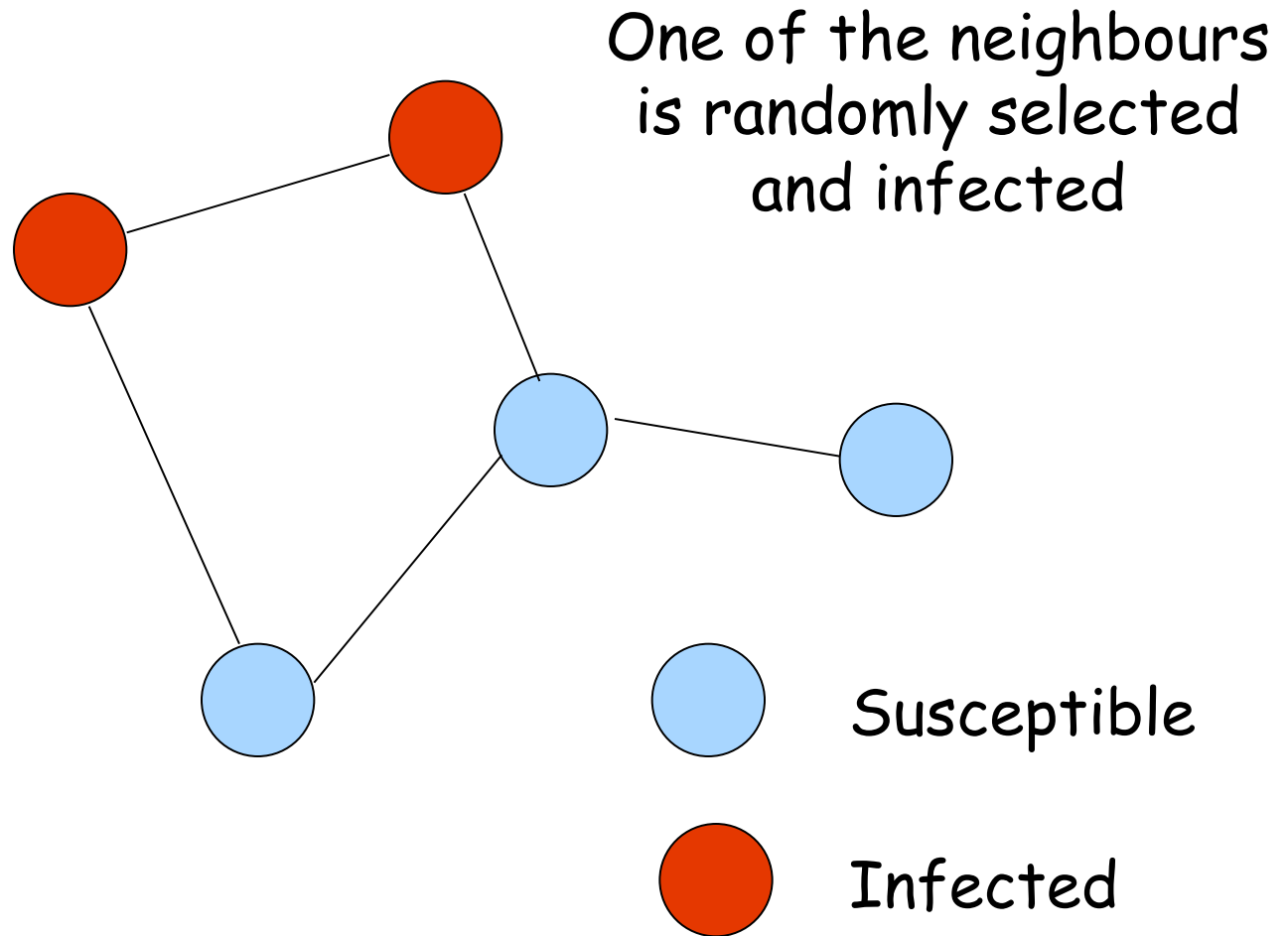
INRIA – EPI Maestro

12 December 2012

Epidemics on a graph: SI model



Epidemics on a graph: SI model

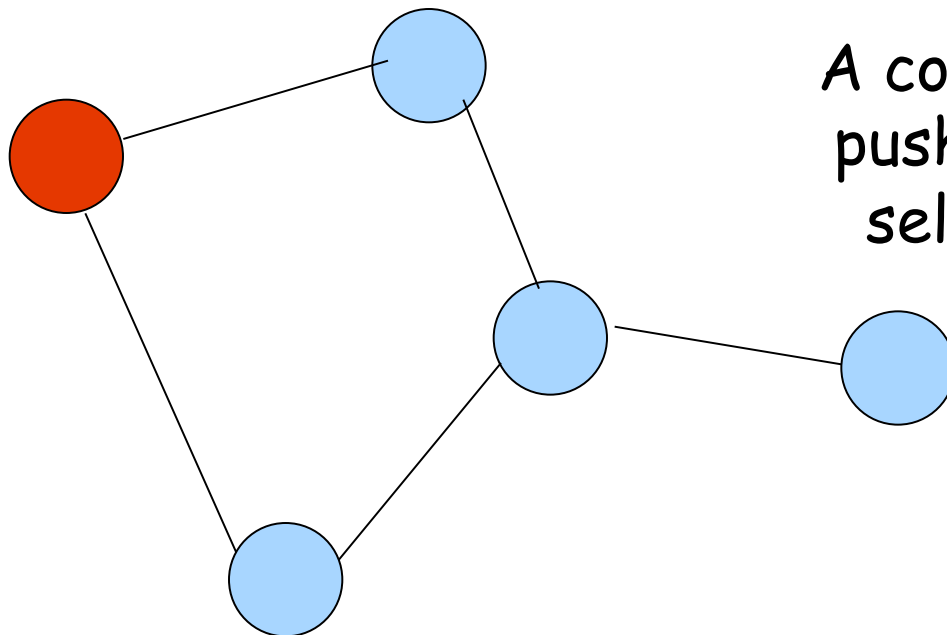


Any interest for Computer Networks?

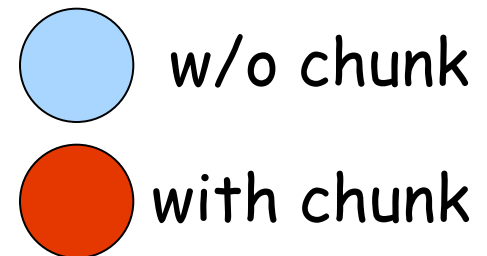
□ Flooding

- Epidemic Routing in Delay Tolerant Networks

□ Chunk distribution in a P2P streaming system (push algorithms)

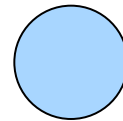
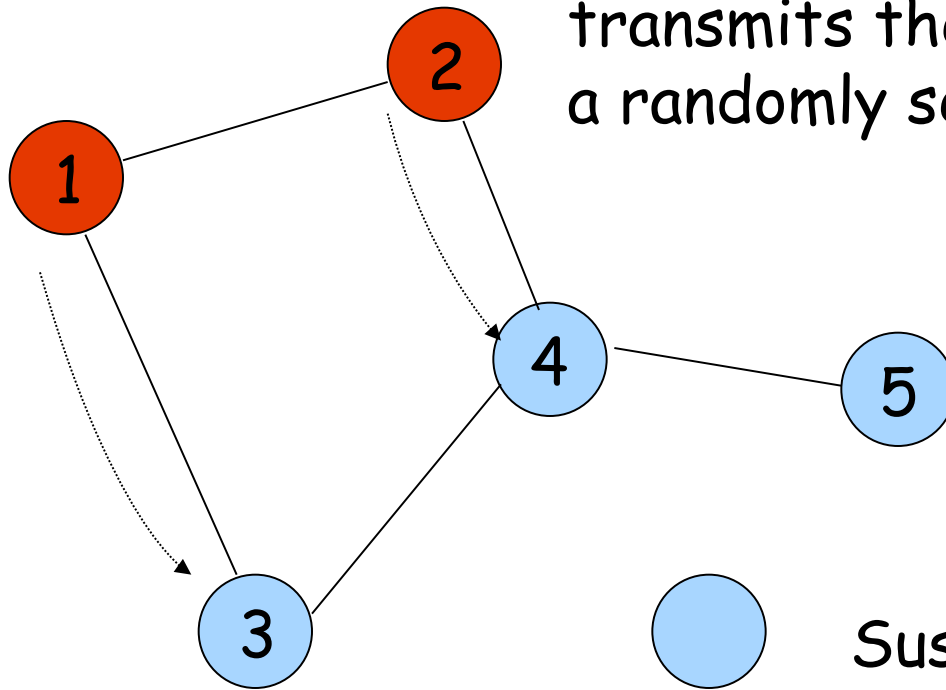


A copy of the chunk is pushed to a randomly selected neighbour

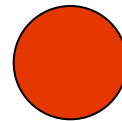


Time-slotted synchronous behaviour

Every T an infected node transmits the disease to a randomly selected neighbour



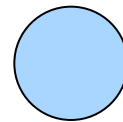
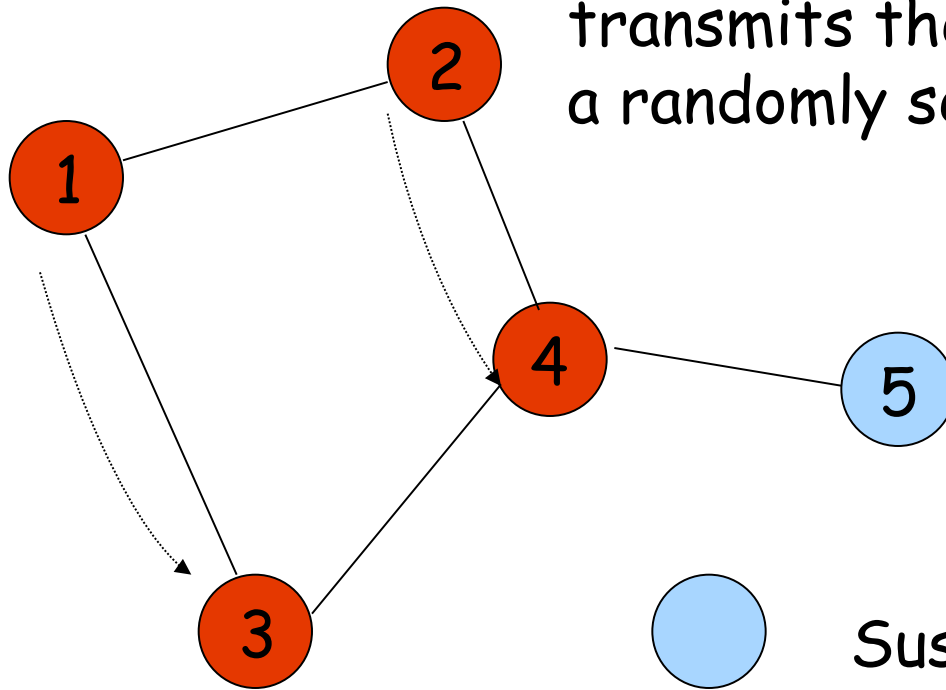
Susceptible



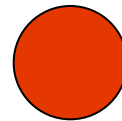
Infected

Time-slotted synchronous behaviour

Every T an infected node transmits the disease to a randomly selected neighbour



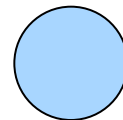
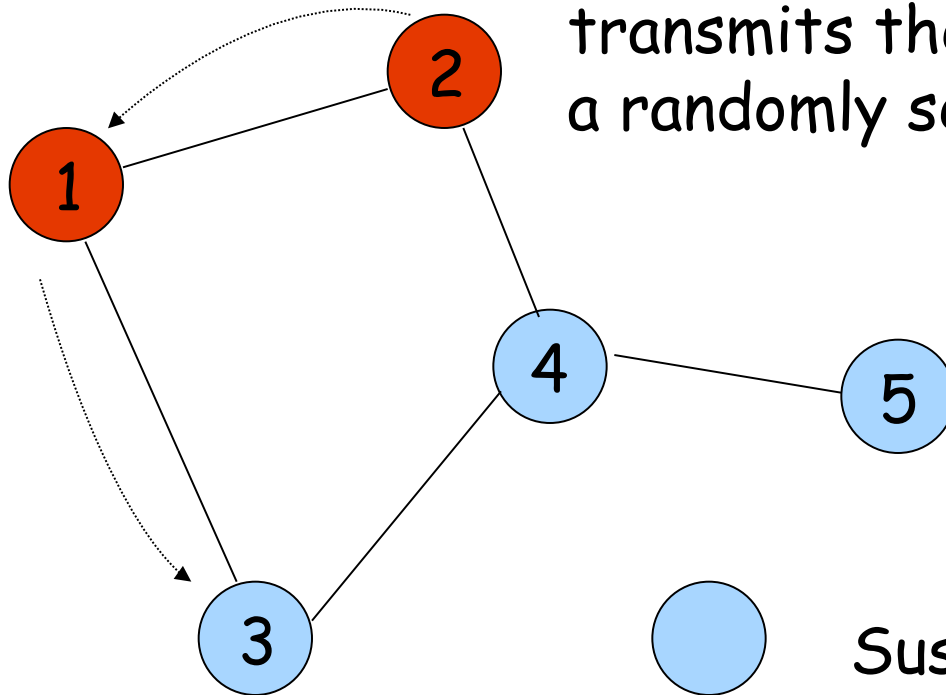
Susceptible



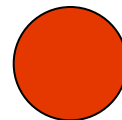
Infected

Time-slotted synchronous behaviour

Every T an infected node transmits the disease to a randomly selected neighbour



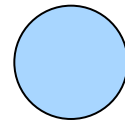
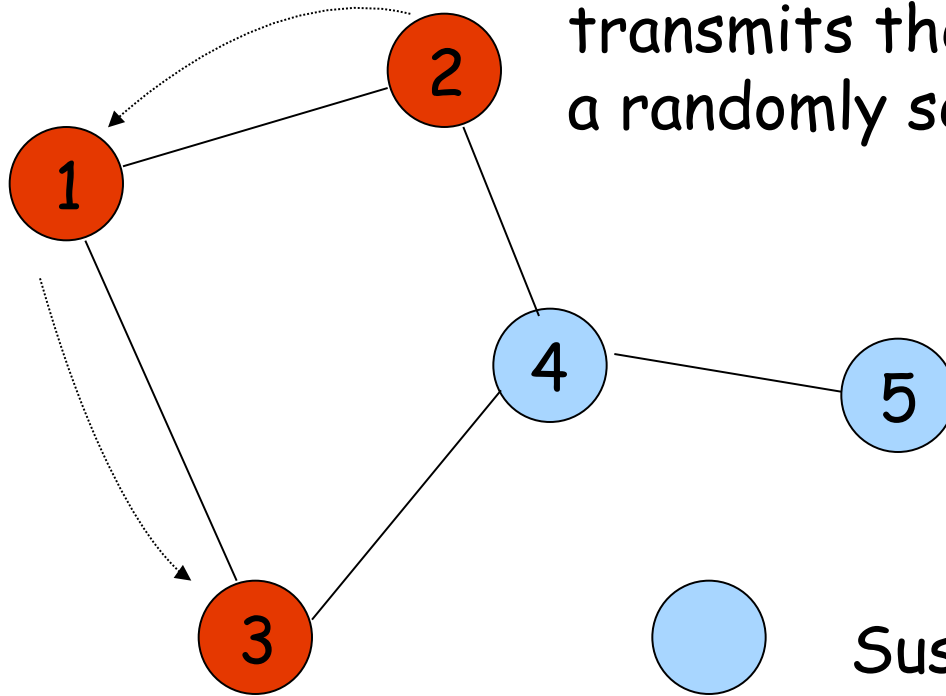
Susceptible



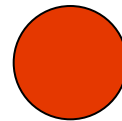
Infected

Time-slotted synchronous behaviour

Every T an infected node transmits the disease to a randomly selected neighbour



Susceptible

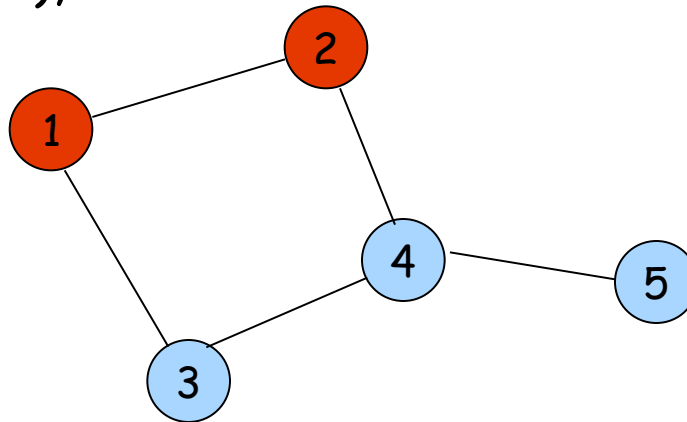


Infected

How do you model it?

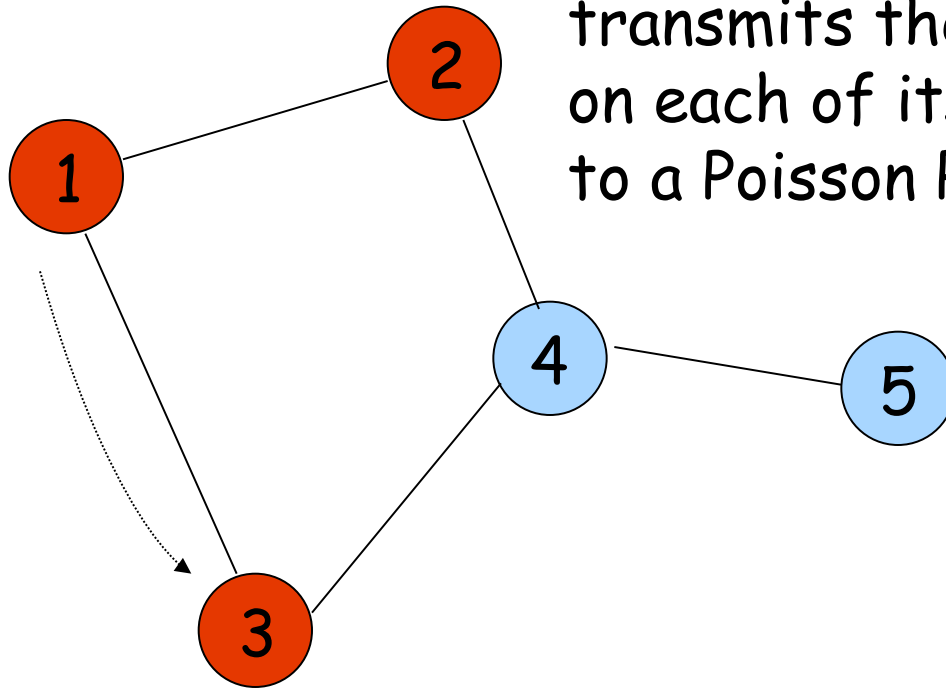
□ A Markov Chain

- System state at time k is a vector specifying if every node is infected (1) or not (0)
 - e.g. $(1,1,0,0,0)$, size: 2^5



- Probability transitions among states
 - e.g. $\text{Prob}((1,1,0,0,0) \rightarrow (1,1,1,0,0)) = 1/4$

Asynchronous behaviour

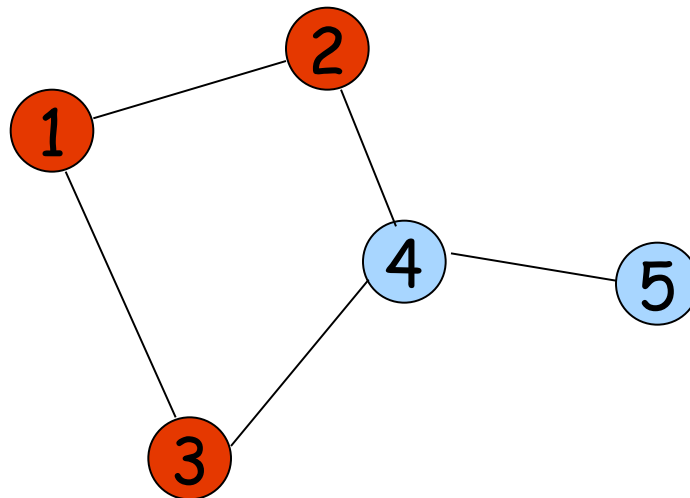


Every infected node transmits the disease on each of its links according to a Poisson Process with rate β

How to model it?

□ A Continuous-time Markov process/Chain (C-MC)

- System state at time t is a vector specifying if every node is infected (1) or not (0)
- Rate transitions between state pairs
 - e.g, $q((1,1,1,0,0) \rightarrow (1,1,1,1,0)) = 2\beta$



What to study and how

- P the transition matrix ($2^N \times 2^N$)
- Transient analysis
 - $\pi(k+1) = \pi(k)P$,
 - $\pi(k+1) = \pi(0)P^{k+1}$,
- Stationary distribution (equilibrium)
 - $\pi = \pi P$
 - If the Markov chain is irreducible and aperiodic
 - Computational cost:
 - $O((2^N)^3)$ if we solve the system
 - $O(K M)$ where M is the number of non-null entries in P if we adopt the iterative procedure (K is the number of iterations), in our case $M = O((2^N)^2)$

Similar for C-MC

□ Stationary distribution (equilibrium)

- $\pi = \pi P$, D-MC

- $\pi Q = 0$, C-MC

□ Transient analysis

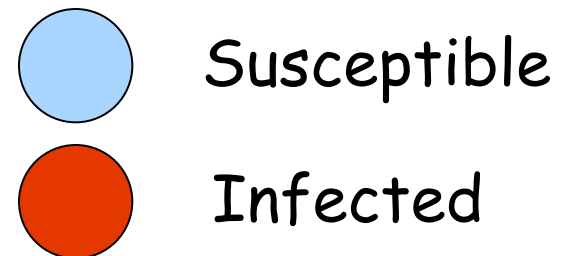
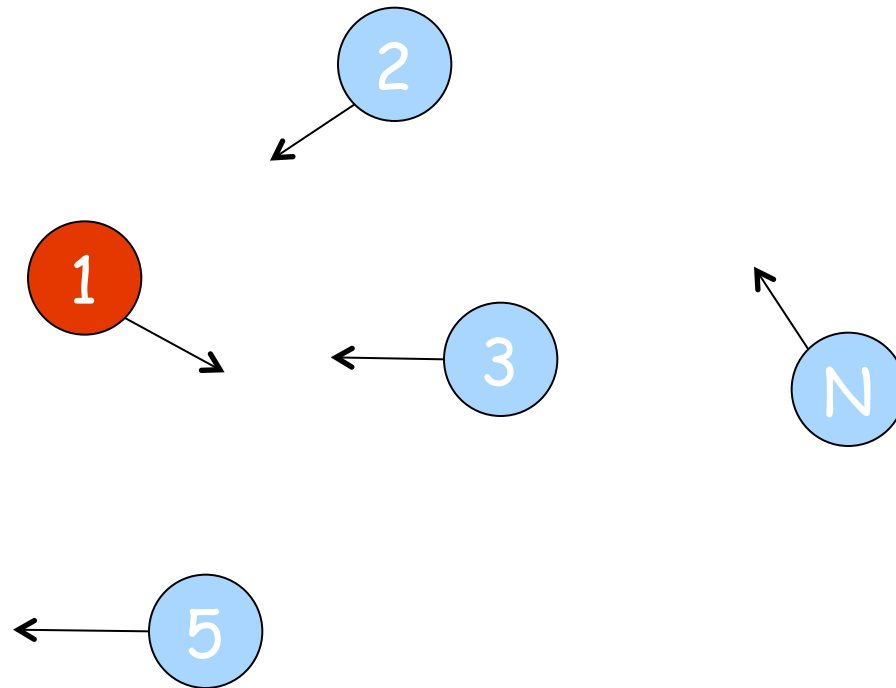
- $\pi(k+1) = \pi(k)P$, D-MC

- $d\pi(t)/dt = \pi(t)Q$, C-MC

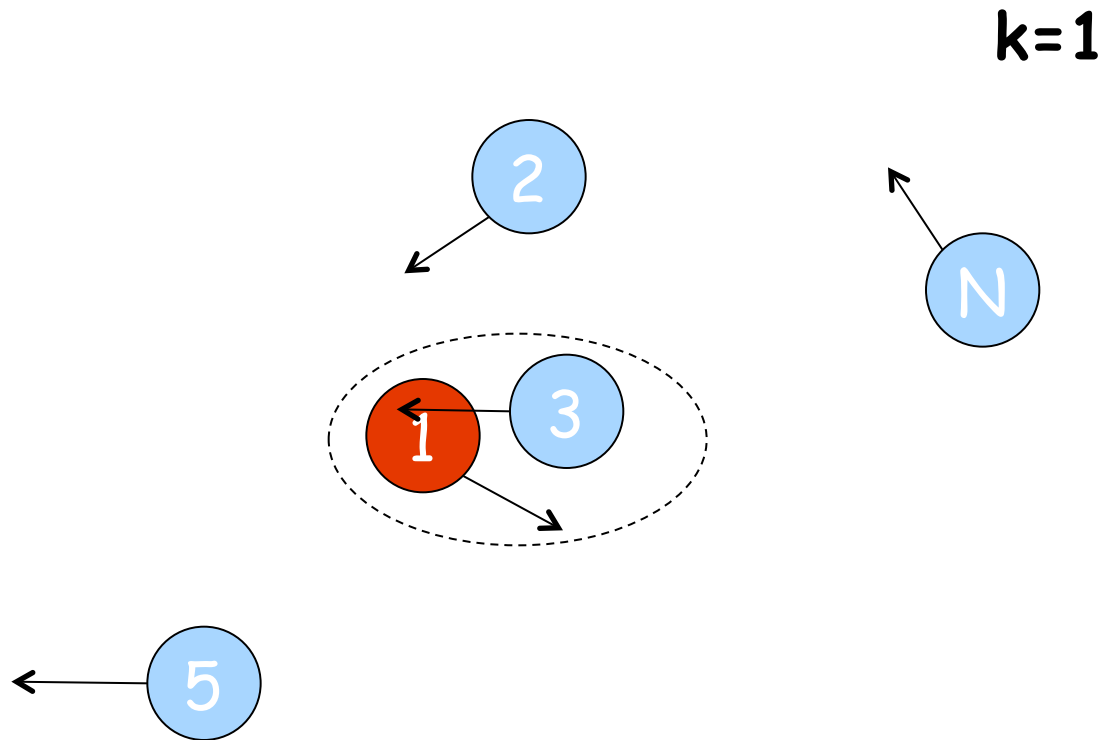
Outline

- Limit of Markovian models
- Mean Field (or Fluid) models
 - exact results
 - extensions to graphs
 - applications

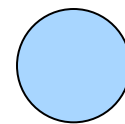
A motivating example: epidemics



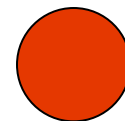
A motivating example: epidemics



At each slot there is a probability p
that two given nodes meet.
Assume meetings to be independent.

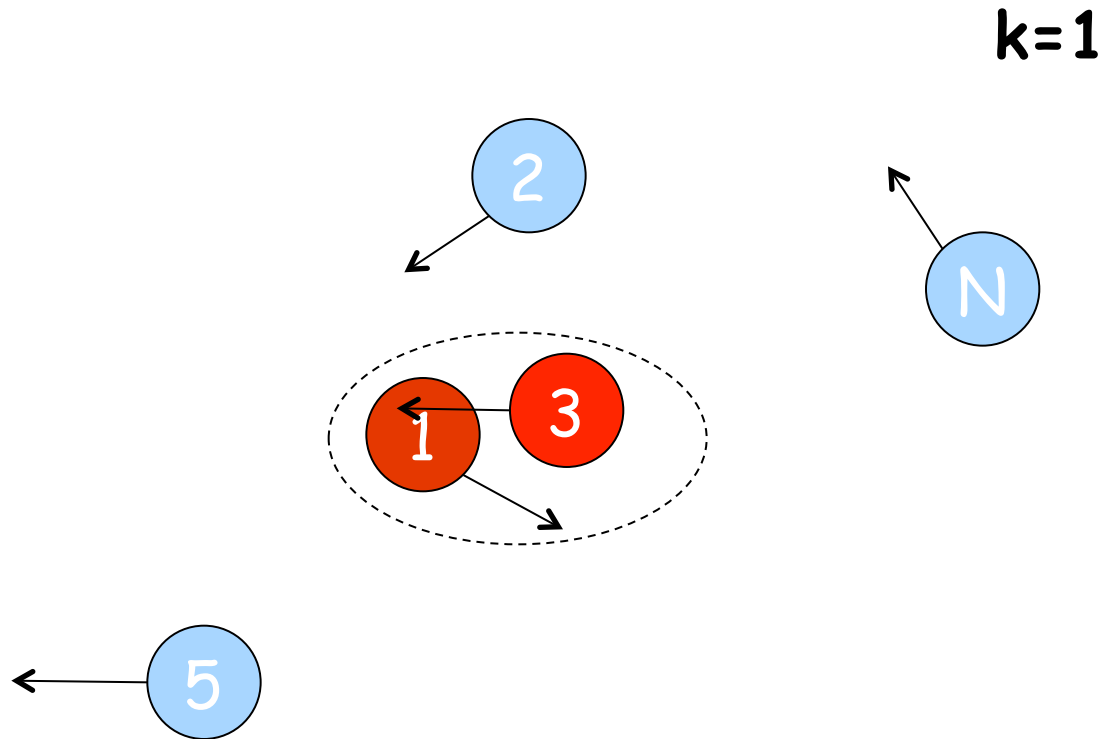


Susceptible

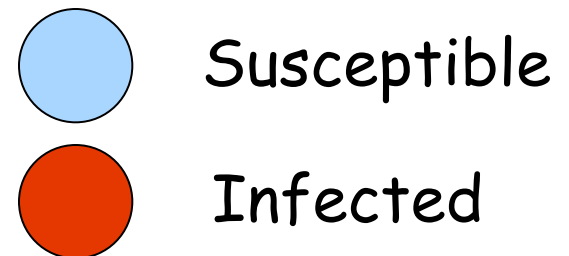


Infected

A motivating example: epidemics

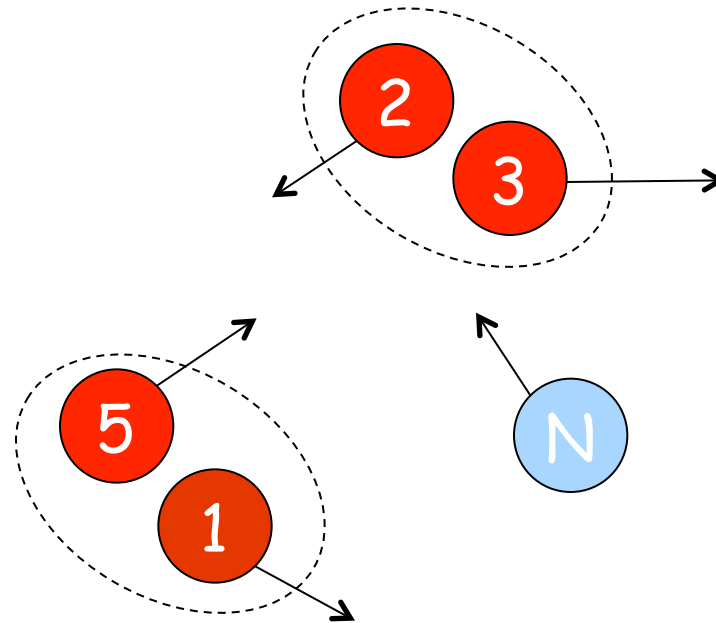


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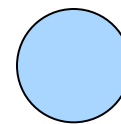


A motivating example: epidemics

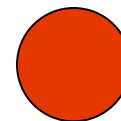
$k=2$



At each slot there is a probability p
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Assume meetings to be independent.



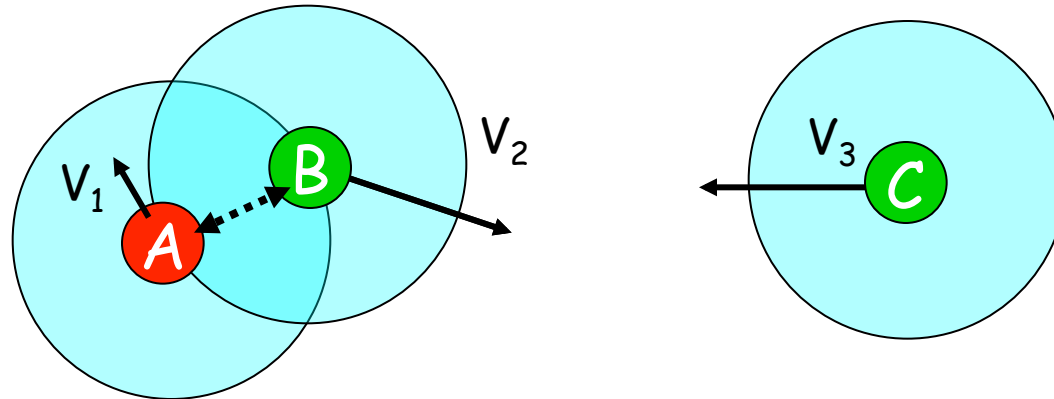
Susceptible



Infected

Delay Tolerant Networks

(a.k.a. Intermittently Connected Networks)



mobile wireless networks

no path at a given time instant between two nodes
because of power constraint, fast mobility dynamics
maintain capacity, when number of nodes (N) diverges

Fixed wireless networks: $C = \Theta(\sqrt{1/N})$ [Gupta99]

Mobile wireless networks: $C = \Theta(1)$, [Grossglauser01]

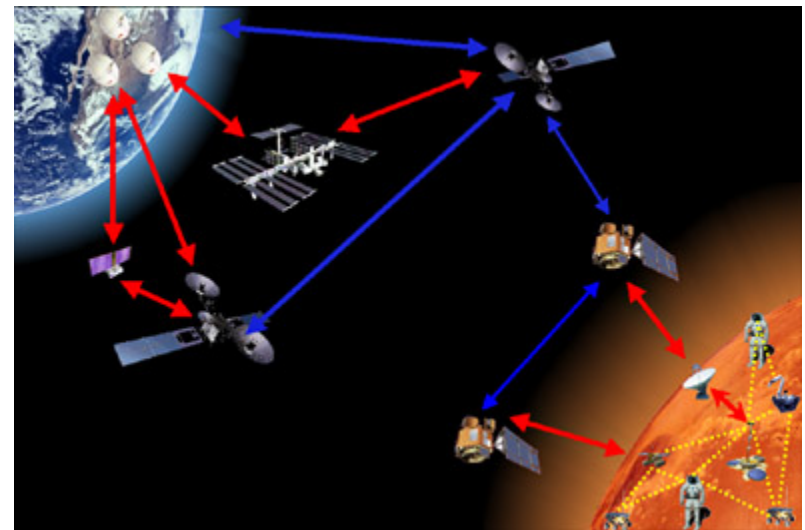
a really challenging network scenario

No traditional protocol works

Some examples

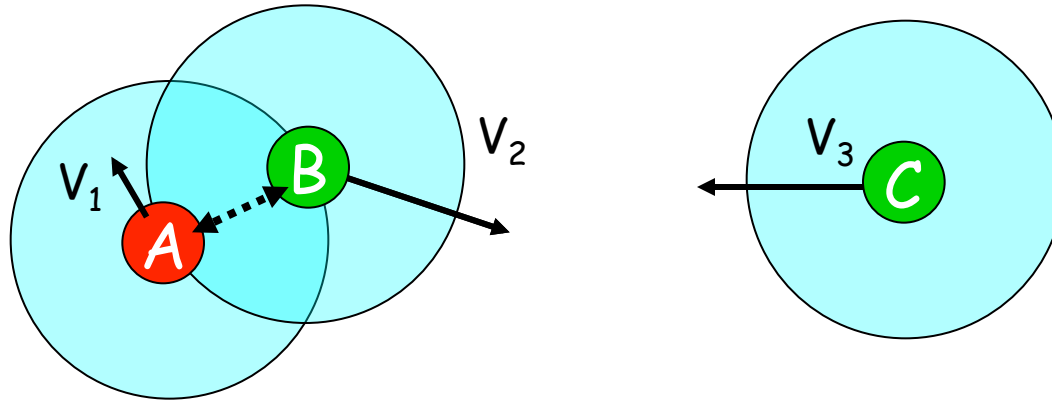


- Network for disaster relief team
- Military battle-field network
- ...



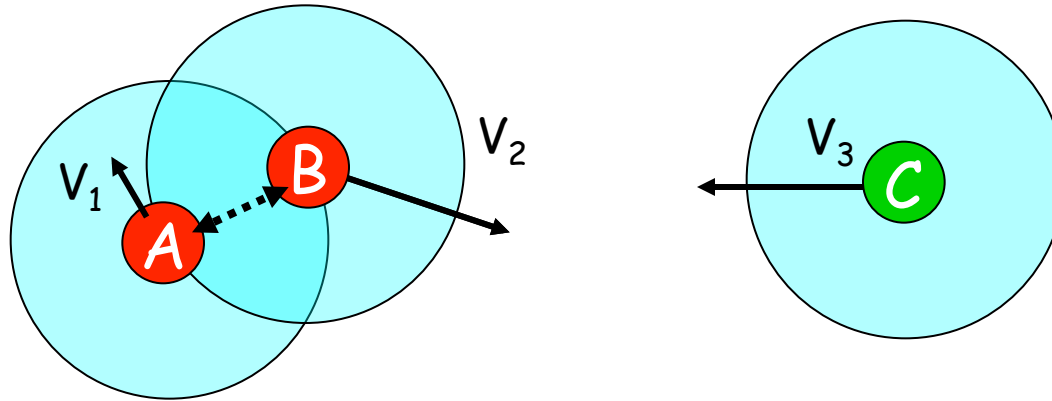
Inter-planetary backbone

Epidemic Routing



- Message as a disease, carried around and transmitted

Epidemic Routing

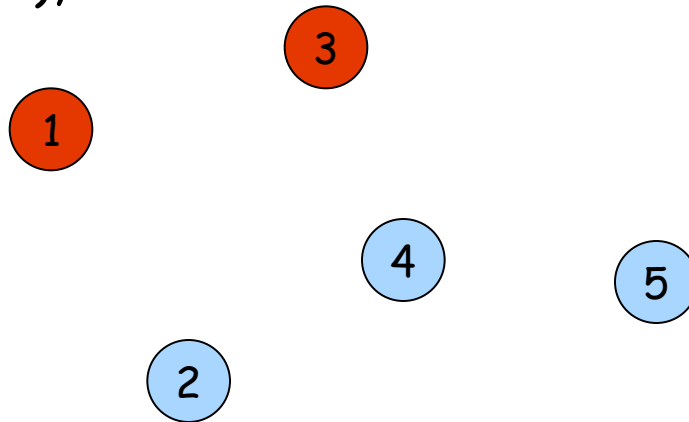


- Message as a disease, carried around and transmitted
 - Store, Carry and Forward paradigm

How do you model it?

□ A Markov Chain

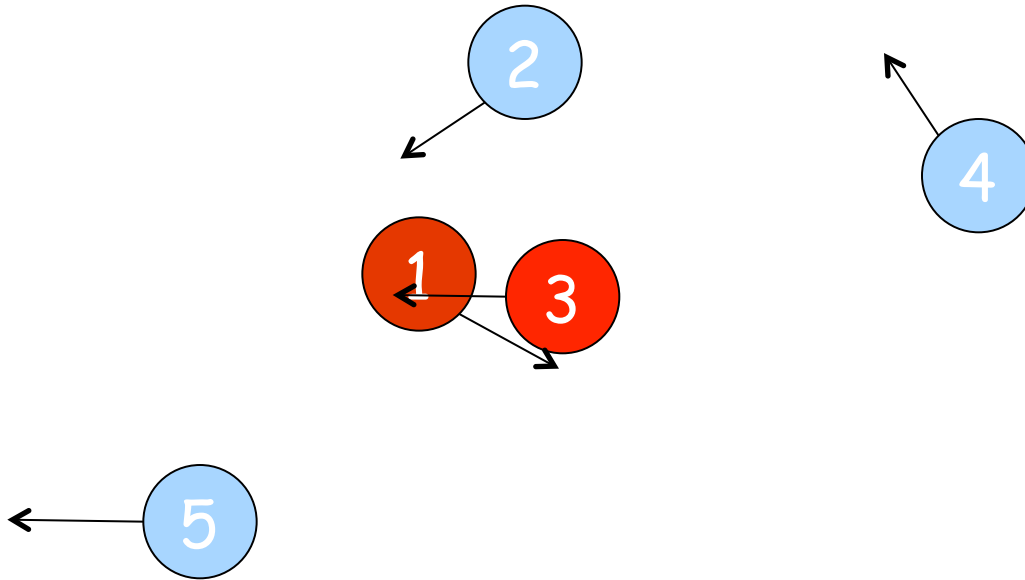
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- Probability transitions among states
 - e.g. $\text{Prob}((1,0,1,0,0) \rightarrow (1,1,1,0,0)) = ?$

Transition probabilities

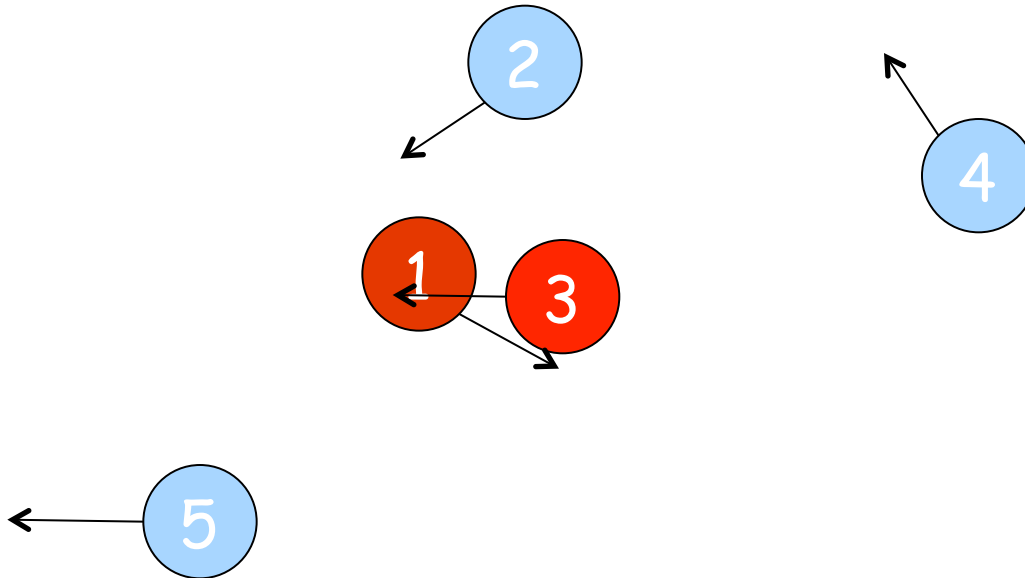
$$\text{Prob}((1,0,1,0,0) \rightarrow (1,1,1,0,0)) = ?$$



At slot k , when there are $I(=I(k))$ infected nodes, the prob. that node 2 gets infected is: $q_I = 1 - (1-p)^I$

Transition probabilities

Prob((1,0,1,0,0)->(1,1,1,0,0))=?



Prob((1,0,1,0,0)->(1,1,1,0,0))= $q_2(1-q_2)^2$
Where $q_I=1-(1-p)^I$

What to study and how

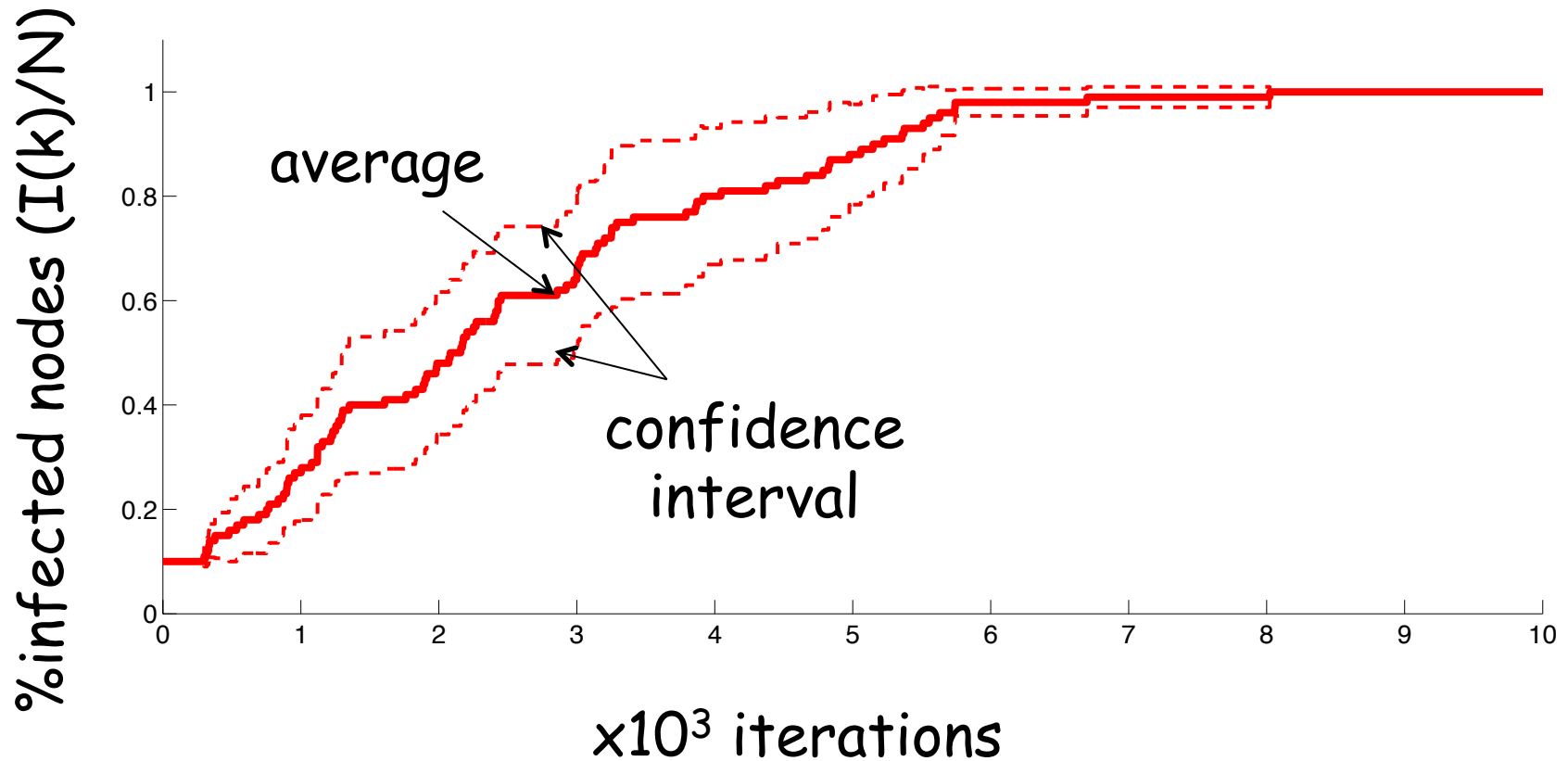
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Can we simplify the problem?

- all the nodes in the same state (infected or susceptibles) are **equivalent**
- If we are interested only in the number of nodes in a given status, we can have a more succinct model
 - state of the system at slot k : $I(k)$
 - it is still a MC
 - $\text{Prob}(I(k+1)=I+n \mid I(k)=I) = C_{N-I}^n q_I^n (1-q_I)^{N-n-I}$
 - $(I(k+1)-I(k) \mid I(k)=I) \sim \text{Bin}(N-I, q_I)$
 - $q_I = 1 - (1-p)^I$

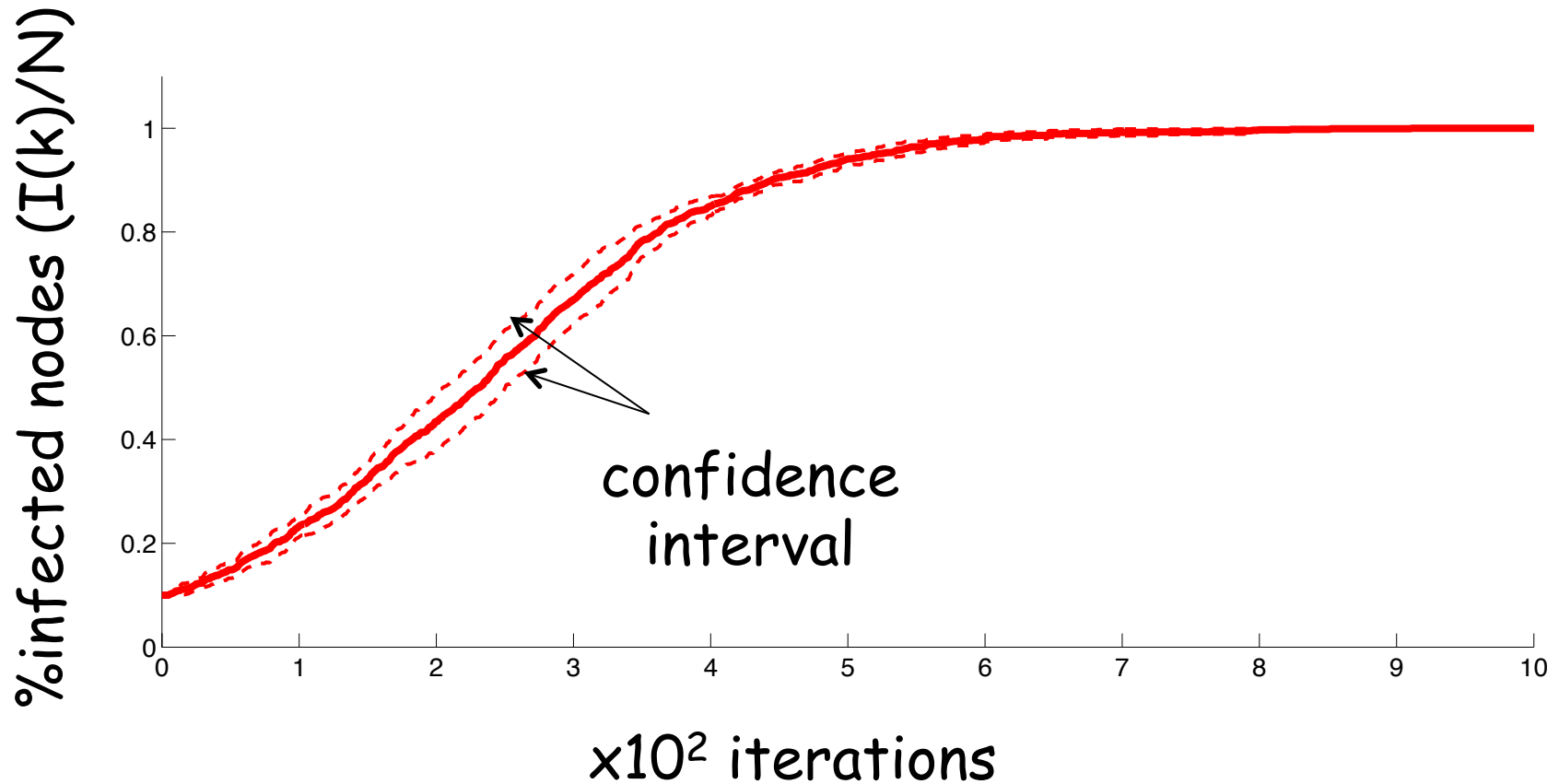
Some numerical examples

$p=10^{-4}$, $N=10$, $I(0)=N/10$, 10 runs



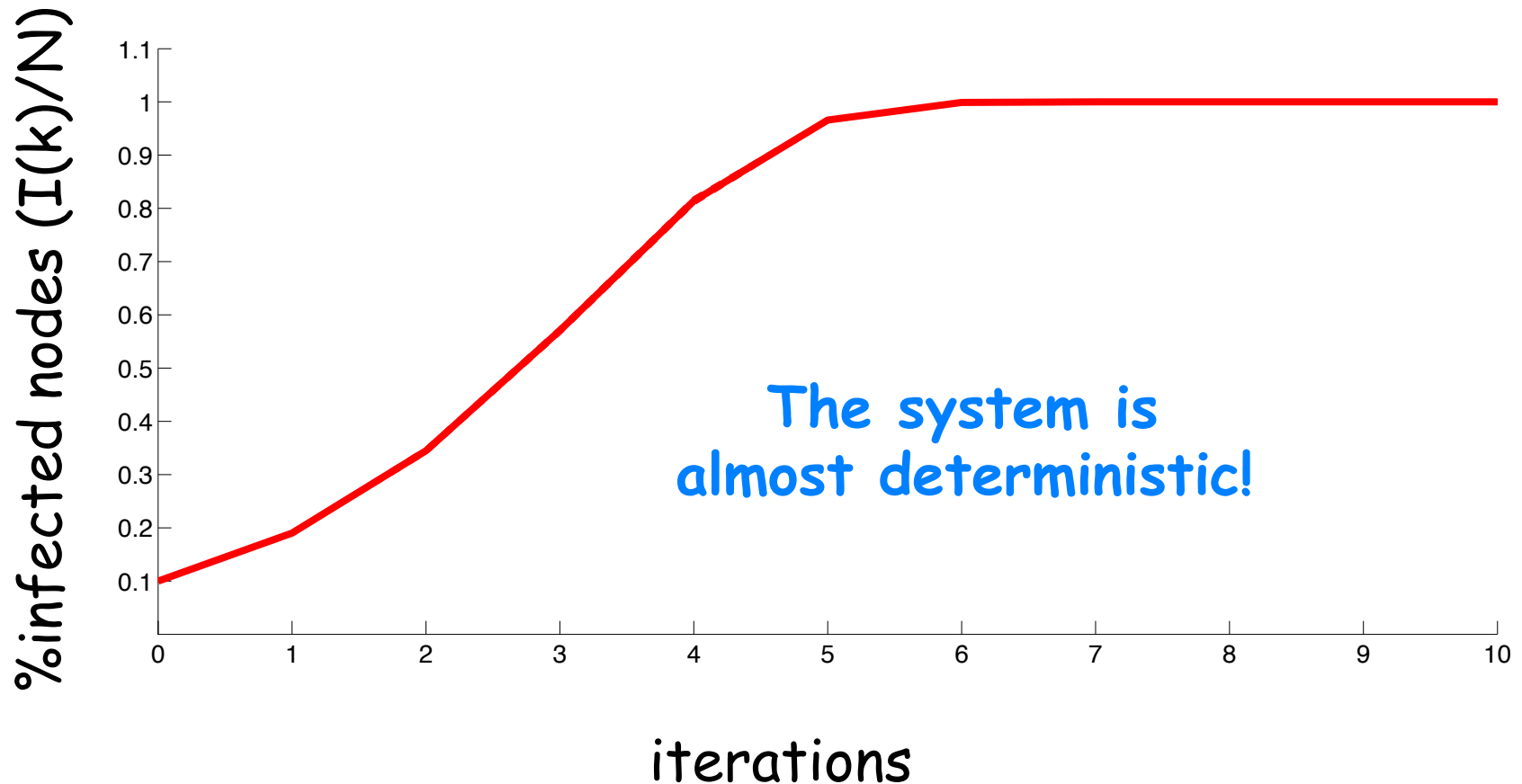
Some numerical examples

$p=10^{-4}$, $N=100$, $I(0)=N/10$, 10 runs



Some numerical examples

$p=10^{-4}$, $N=10000$, $I(0)=N/10$, 10 runs



Summary

- ❑ For a large system of interacting equivalent objects, the Markov model can be untractable...
- ❑ but a deterministic description of the system seems feasible in terms of the empirical measure (% of objects in each status)
 - intuition: kind of law of large numbers
- ❑ Mean field models describe the deterministic limit of Markov models when the number of objects diverges

Spoiler

□ $i^{(N)}(k)$, fraction of infected nodes at time k

□ Solve

$$di(t)/dt = i(t)(1-i(t)),$$

with $i=i_0$

- Solution: $i(t) = 1 / ((1/i_0 - 1) e^{-t} + 1)$

□ If $i^{(N)}(0) = i_0$,

$$i^{(N)}(k) \approx i(k p_0 / N) = 1 / ((1/i_0 - 1) \exp(-k p_0 / N) + 1) \\ = 1 / ((1/i_0 - 1) \exp(-k N p) + 1)$$

Outline

- Limit of Markovian models
- Mean Field (or Fluid) models
 - exact results
 - extensions to graphs
 - applications

References

- Results here for discrete time Markov Chains
 - Benaim, Le Boudec "A Class of Mean Field Interaction Models for Computer and Communication Systems", LCA-Report-2008-010
- A survey with pointers to continuous time Markov processes and links to stochastic approximation and propagation of chaos
 - Ch. 2 of Nicolas Gast's PhD thesis "Optimization and Control of Large Systems, Fighting the Curse of Dimensionality"

Necessary hypothesis: Objects' Equivalence

- $\pi(k+1) = \pi(k)P$
- A state $\sigma = (v_1, v_2, \dots, v_N)$, $v_j \in V$ ($|V|=V$, finite)
 - E.g. in our example $V = \{0, 1\}$
- P is invariant under any label permutation φ :
 - $$P_{\sigma, \sigma'} = \text{Prob}((v_1, v_2, \dots, v_N) \rightarrow (u_1, u_2, \dots, u_N)) =$$
$$\text{Prob}((v_{\varphi(1)}, v_{\varphi(2)}, \dots, v_{\varphi(N)}) \rightarrow (u_{\varphi(1)}, u_{\varphi(2)}, \dots, u_{\varphi(N)}))$$

Some notation and definitions

- $X_n^{(N)}(k)$: state of node n at slot k
- $M_v^{(N)}(k)$: occupancy measure of state v at slot k
 - $M_v^{(N)}(k) = \sum_n \mathbf{1}(X_n^{(N)}(k) = v) / N$
 - SI model: $M_2^{(N)}(k) = I^{(N)}(k) / N = i^{(N)}(k)$,
 $M_1^{(N)}(k) = S^{(N)}(k) / N = s^{(N)}(k) = 1 - i^{(N)}(k)$
- $\mathbf{M}^{(N)}(k) = (M_1^{(N)}(k), M_2^{(N)}(k), \dots, M_v^{(N)}(k))$
 - SI model: $(1 - i^{(N)}(k), i^{(N)}(k))$
- $\mathbf{f}^{(N)}(\mathbf{m}) = E[\mathbf{M}^{(N)}(k+1) - \mathbf{M}^{(N)}(k) | \mathbf{M}^{(N)}(k) = \mathbf{m}]$
 - Drift or intensity, it is the **mean field**

Other hypotheses

- Intensity vanishes at a rate $\varepsilon(N)$
 - $\lim_{N \rightarrow \infty} \mathbf{f}^{(N)}(\mathbf{m})/\varepsilon(N) = \mathbf{f}(\mathbf{m})$
- Second moment of number of object transitions per slot is bounded
 - #transitions $< W^N(k)$,
 - $E[W^N(k)^2 | \mathbf{M}^{(N)}(k) = \mathbf{m}] < cN^2\varepsilon(N)^2$
- Drift is a smooth function of \mathbf{m} and $1/N$
 - $\mathbf{f}^{(N)}(\mathbf{m})/\varepsilon(N)$ has continuous derivatives in \mathbf{m} and in $1/N$ on $[0,1]^V \times [0,\beta]$, with $\beta > 0$

Convergence Result

- Define $\underline{\mathbf{M}}^{(N)}(t)$ with t real, such that
 - $\underline{\mathbf{M}}^{(N)}(k \varepsilon(N)) = \mathbf{M}^{(N)}(k)$ for k integer
 - $\underline{\mathbf{M}}^{(N)}(t)$ is affine on $[k \varepsilon(N), (k+1)\varepsilon(N)]$
- Consider the Differential Equation
 - $d\boldsymbol{\mu}(t)/dt = \mathbf{f}(\boldsymbol{\mu})$, with $\boldsymbol{\mu}(0) = \mathbf{m}_0$
- Theorem
 - For all $T > 0$, if $\mathbf{M}^{(N)}(0) \rightarrow \mathbf{m}_0$ in probability (/mean square) as $N \rightarrow \infty$, then
$$\sup_{0 \leq t \leq T} \|\underline{\mathbf{M}}^{(N)}(t) - \boldsymbol{\mu}(t)\| \rightarrow 0$$
in probability (/mean square)

Convergence of random variables

- The sequence of random variables $X^{(N)}$ converges to X in probability if
 - for all $\delta > 0$ $\lim_{N \rightarrow \infty} \text{Prob}(|X^{(N)} - X| > \delta) = 0$
- The sequence of random variables X^N converges to X in mean square if
 - $\lim_{N \rightarrow \infty} E[|X^{(N)} - X|^2] = 0$
- Convergence in mean square implies convergence in probability

Application to the SI model

□ Assumptions' check

✓ Nodes are equivalent

- Intensity vanishes at a rate $\varepsilon(N)$

$$f^{(N)}(\mathbf{m}) = E[\mathbf{M}^{(N)}(k+1) - \mathbf{M}^{(N)}(k) \mid \mathbf{M}^{(N)}(k) = \mathbf{m}]$$

$$M_2^{(N)}(k) = I^{(N)}(k)/N = i^{(N)}(k), M_1^{(N)}(k) = 1 - M_2^{(N)}(k)$$

$$(I^{(N)}(k+1) - I^{(N)}(k) \mid I^{(N)}(k) = I) \sim \text{Bin}(N - I, q_I) \Rightarrow$$

$$E[I^{(N)}(k+1) - I^{(N)}(k) \mid I^{(N)}(k) = I] = q_I (N - I)$$

$$E[i^{(N)}(k+1) - i^{(N)}(k) \mid i^{(N)}(k) = i] = (1 - i) q_I$$

$$= (1 - i)(1 - (1 - p)^i)^N \rightarrow (1 - i) \text{ when } N \text{ diverges!}$$

Application to the SI model

- Out of the impasse: introduce a scaling for p
 - If $p^{(N)} = p_0/N^a$ $a > 1 \Rightarrow (1-i)(1-(1-p^{(N)})^i)^N \rightarrow 0$
 - Consider $a=2$
 - $(1-i)(1-(1-p^{(N)})^i)^N \sim (1-i) i p_0/N$ (for N large)
 - $\varepsilon(N) = p_0/N$
 - $f_2(\mathbf{m}) = f_2((s,i)) = s i = i(1-i)$
- Lesson to keep: often we need to introduce some parameter scaling

Application to the SI model

□ Assumptions' check

- ✓ Nodes are equivalent
- ✓ Intensity vanishes at a rate $\varepsilon(N)=p_0/N$
- Second moment of number of object transitions per slot is bounded

#transitions $\leq W^N(k)$,

$$E[W^N(k)^2 | \mathbf{M}^{(N)}(k)=\mathbf{m}] \leq cN^2 \varepsilon(N)^2$$

$$W^N(k) = \#trans. \sim \text{Bin}(N-I(k), q_I)$$

$$E[W^N(k)^2] = ((N-I(k))q_I)^2 + (N-I(k))q_I(1-q_I)$$

is in $O(N^2 \varepsilon(N)^2)$

Application to the SI model

□ Assumptions' check

- ✓ Nodes are equivalent
- ✓ Intensity vanishes at a rate $\varepsilon(N)=p_0/N$
- ✓ Second moment of number of object transitions per slot is bounded
- ✓ Drift is a smooth function of \mathbf{m} and $1/N$
 - $f_2^{(N)}(\mathbf{m})/\varepsilon(N) =$
 $= (1-i) (1 - (1-(p_0/N^2))^i)^N / (p_0/N)$
 - continuous derivatives in i and in $1/N$
(not evident)

Practical use of the convergence result

□ Theorem

- For all $T > 0$, if $\mathbf{M}^{(N)}(0) \rightarrow \mathbf{m}_0$ in probability (/mean square) as $N \rightarrow \infty$, then $\sup_{0 \leq t \leq T} \|\underline{\mathbf{M}}^{(N)}(t) - \boldsymbol{\mu}(t)\| \rightarrow 0$ in probability (/mean square)
- Where $\boldsymbol{\mu}(t)$ is the solution of $d\boldsymbol{\mu}(t)/dt = \mathbf{f}(\boldsymbol{\mu})$, with $\boldsymbol{\mu}(0) = \mathbf{m}_0$

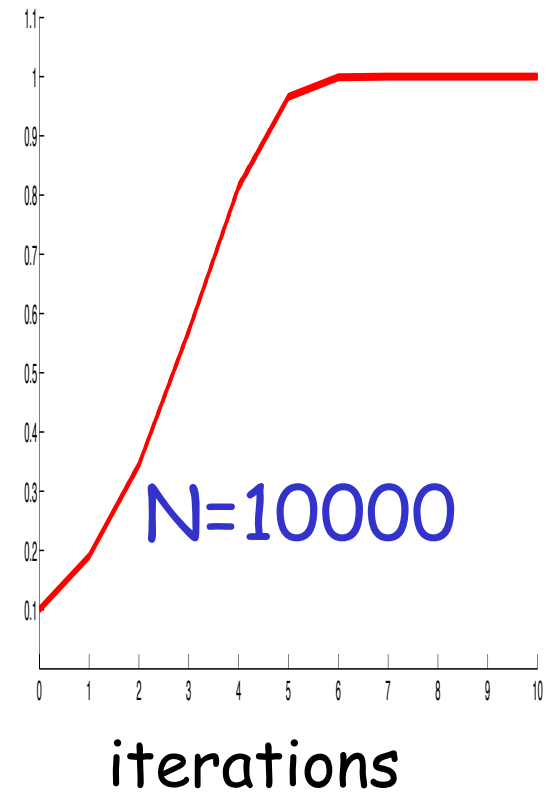
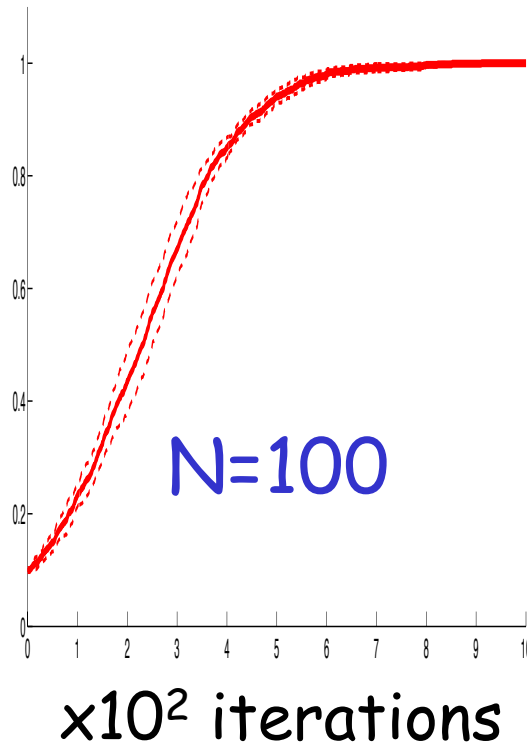
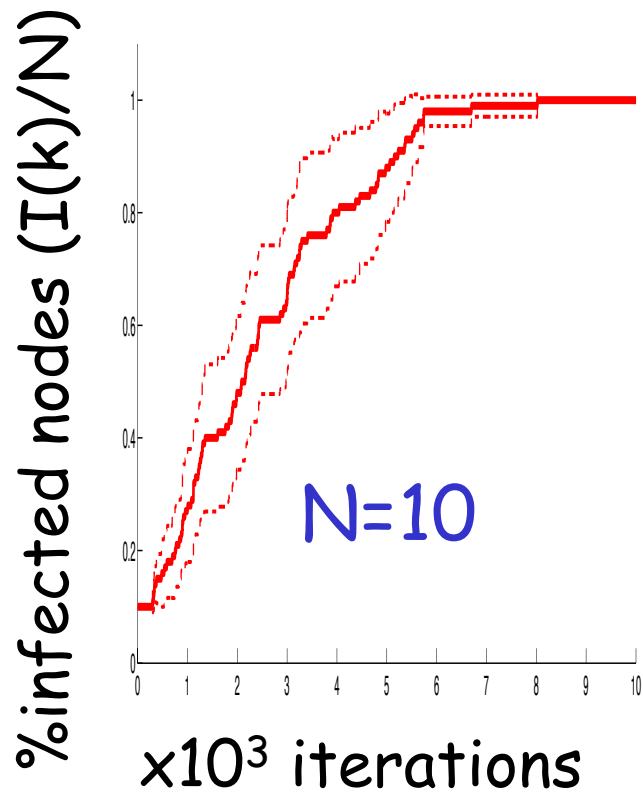
□ $\mathbf{M}^{(N)}(0) = \mathbf{m}_0$, $\mathbf{M}^{(N)}(k) = \underline{\mathbf{M}}^{(N)}(k\varepsilon(N)) \approx \boldsymbol{\mu}(k\varepsilon(N))$

Application to the SI model

- $f_2(\mathbf{m}) = f_2((s, i)) = i(1-i)$
- $d\mu_2(t)/dt = f_2(\mu_2(t)) = \mu_2(t)(1-\mu_2(t))$,
with $\mu_2(0) = \mu_{0,2}$
 - Solution: $\mu_2(t) = 1 / ((1/\mu_{0,2} - 1) e^{-t} + 1)$
- If $i^{(N)}(0) = i_0$,
 $i^{(N)}(k) \approx \mu_2(k\varepsilon(N)) = 1 / ((1/i_0 - 1) \exp(-k p_0/N) + 1)$
 $= 1 / ((1/i_0 - 1) \exp(-k N p) + 1)$

Back to the numerical examples

$p=10^{-4}$, $I(0)=N/10$, 10 runs



Advantage of Mean Field

- If $i^{(N)}(0) = i_0$,
$$i^{(N)}(k) \approx \mu_2(k\varepsilon(N)) = 1 / ((1/i_0 - 1) \exp(-k p_0/N) + 1)$$
$$= 1 / ((1/i_0 - 1) \exp(-k N p) + 1)$$
 - solved for each N with negligible computational cost
- In general: solve numerically the solution of a system of ordinary differential equations (size = #of possible status)
 - simpler than solving the Markov chain