

TAPENADE: a tool for Automatic Differentiation of programs

Laurent Hascoët

`Laurent.Hascoet@sophia.inria.fr`

TROPICS Project, INRIA Sophia-Antipolis

ECCOMAS 2004, Jyväskylä, July 25-28, 2004

PLAN:

- AD: principles of Tangent and Reverse
- Tapenade: technology from Compilation and Parallelization
- Tapenade: differentiation model on examples
- Tapenade: an AD tool on the web
- Conclusion and Further Developments

AD: Principles of Tangent and Reverse

AD rewrites **source programs** to make them compute derivatives.

consider:

$P : \{I_1; I_2; \dots; I_p;\}$ implementing $f : \mathbf{IR}^m \rightarrow \mathbf{IR}^n$

AD: Principles of Tangent and Reverse

AD rewrites **source programs** to make them compute derivatives.

consider: $P : \{I_1; I_2; \dots; I_p;\}$ implementing $f : \mathbf{IR}^m \rightarrow \mathbf{IR}^n$

identify with: $f = f_p \circ f_{p-1} \circ \dots \circ f_1$

name: $x_0 = x$ and $x_k = f_k(x_{k-1})$

AD: Principles of Tangent and Reverse

AD rewrites **source programs** to make them compute derivatives.

consider: $P : \{I_1; I_2; \dots; I_p;\}$ implementing $f : \mathbf{IR}^m \rightarrow \mathbf{IR}^n$

identify with: $f = f_p \circ f_{p-1} \circ \dots \circ f_1$

name: $x_0 = x$ and $x_k = f_k(x_{k-1})$

chain rule: $f'(x) = f'_p(x_{p-1}) \cdot f'_{p-1}(x_{p-2}) \cdot \dots \cdot f'_1(x_0)$

AD: Principles of Tangent and Reverse

AD rewrites **source programs** to make them compute derivatives.

consider: $P : \{I_1; I_2; \dots; I_p;\}$ implementing $f : \mathbf{IR}^m \rightarrow \mathbf{IR}^n$

identify with: $f = f_p \circ f_{p-1} \circ \dots \circ f_1$

name: $x_0 = x$ and $x_k = f_k(x_{k-1})$

chain rule: $f'(x) = f'_p(x_{p-1}) \cdot f'_{p-1}(x_{p-2}) \cdot \dots \cdot f'_1(x_0)$

$f'(x)$ generally too large and expensive \Rightarrow take useful views!

$\dot{y} = f'(x) \cdot \dot{x} = f'_p(x_{p-1}) \cdot f'_{p-1}(x_{p-2}) \cdot \dots \cdot f'_1(x_0) \cdot \dot{x}$ **tangent AD**

$\bar{x} = f'^*(x) \cdot \bar{y} = f'^*_1(x_0) \cdot \dots \cdot f'^*_{p-1}(x_{p-2}) \cdot f'^*_p(x_{p-1}) \cdot \bar{y}$ **reverse AD**

Evaluate both **from right to left** !

AD: Example

...

$$\mathbf{v}_2 = 2 * \mathbf{v}_1 + 5$$

$$\mathbf{v}_4 = \mathbf{v}_2 + p_1 * \mathbf{v}_3 / \mathbf{v}_2$$

...

AD: Example

$$\begin{aligned}
 & \dots \\
 \mathbf{v}_2 &= 2 * \mathbf{v}_1 + 5 \\
 \mathbf{v}_4 &= \mathbf{v}_2 + p_1 * \mathbf{v}_3 / \mathbf{v}_2 \\
 & \dots
 \end{aligned}$$

The corresponding (fragment of) Jacobian is:

$$f'(x) = \dots \begin{pmatrix} 1 & & & \\ & 1 & & \\ & & 1 & \\ 0 & 1 - \frac{p_1 * v_3}{v_2^2} & \frac{p_1}{v_2} & 0 \end{pmatrix} \begin{pmatrix} 1 & & & \\ 2 & 0 & & \\ & & 1 & \\ & & & 1 \end{pmatrix} \dots$$

Tangent AD keeps the structure of P :

$$\dot{y} = f'(x) \cdot \dot{x} = f'_p(x_{p-1}) \cdot f'_{p-1}(x_{p-2}) \cdot \dots \cdot f'_1(x_0) \cdot \dot{x}$$

...

$$\mathbf{v}_2 = 2 * \mathbf{v}_1 + 5$$

$$\mathbf{v}_4 = \mathbf{v}_2 + p_1 * \mathbf{v}_3 / \mathbf{v}_2$$

...

Tangent AD keeps the structure of P :

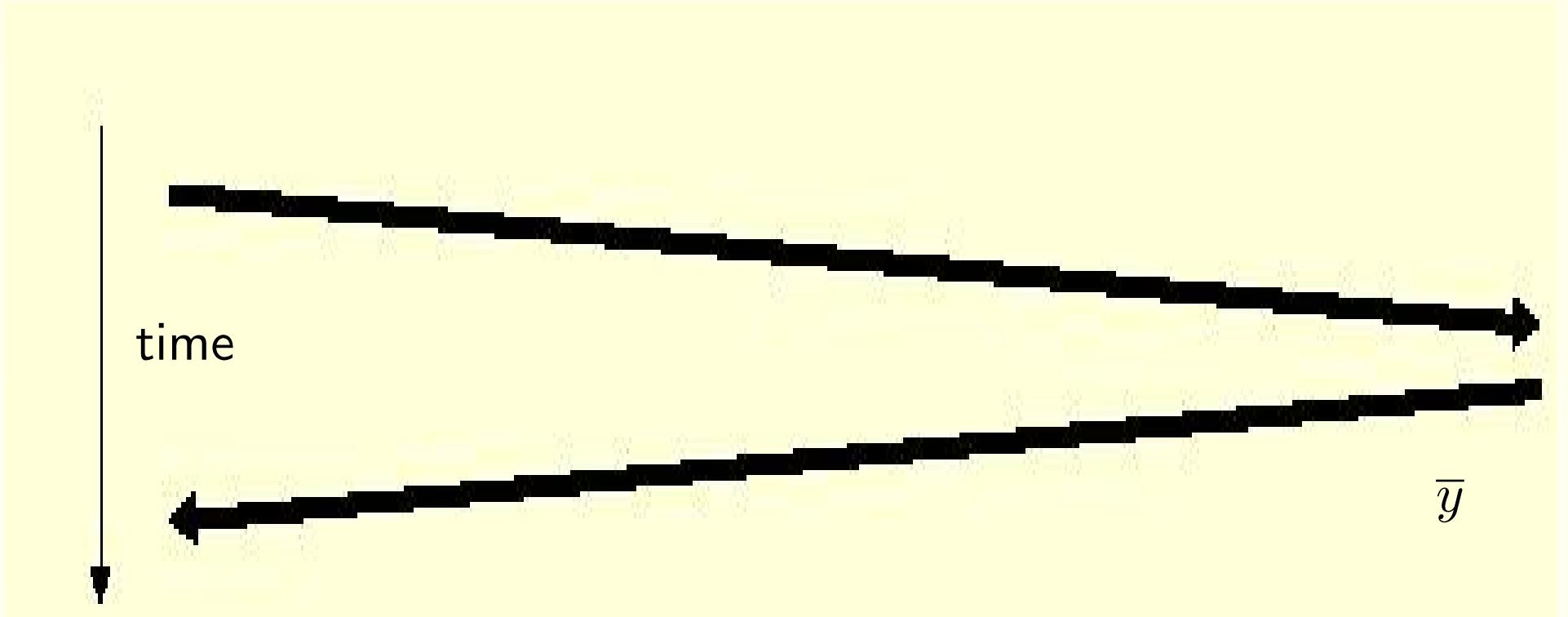
$$\dot{y} = f'(x) \cdot \dot{x} = f'_p(x_{p-1}) \cdot f'_{p-1}(x_{p-2}) \cdot \dots \cdot f'_1(x_0) \cdot \dot{x}$$

$$\begin{aligned}\dot{\mathbf{v}}_2 &= \mathbf{2} * \dot{\mathbf{v}}_1 \\ \mathbf{v}_2 &= \mathbf{2} * \mathbf{v}_1 + \mathbf{5} \\ \dot{\mathbf{v}}_4 &= \dot{\mathbf{v}}_2 * (1 - \mathbf{p}_1 * \mathbf{v}_3 / \mathbf{v}_2^2) + \dot{\mathbf{v}}_3 * \mathbf{p}_1 / \mathbf{v}_2 \\ \mathbf{v}_4 &= \mathbf{v}_2 + \mathbf{p}_1 * \mathbf{v}_3 / \mathbf{v}_2 \\ &\dots\end{aligned}$$

just inserts the products $\dot{x}_k = f'_k(x_{k-1})$ for $k = 1$ to p .

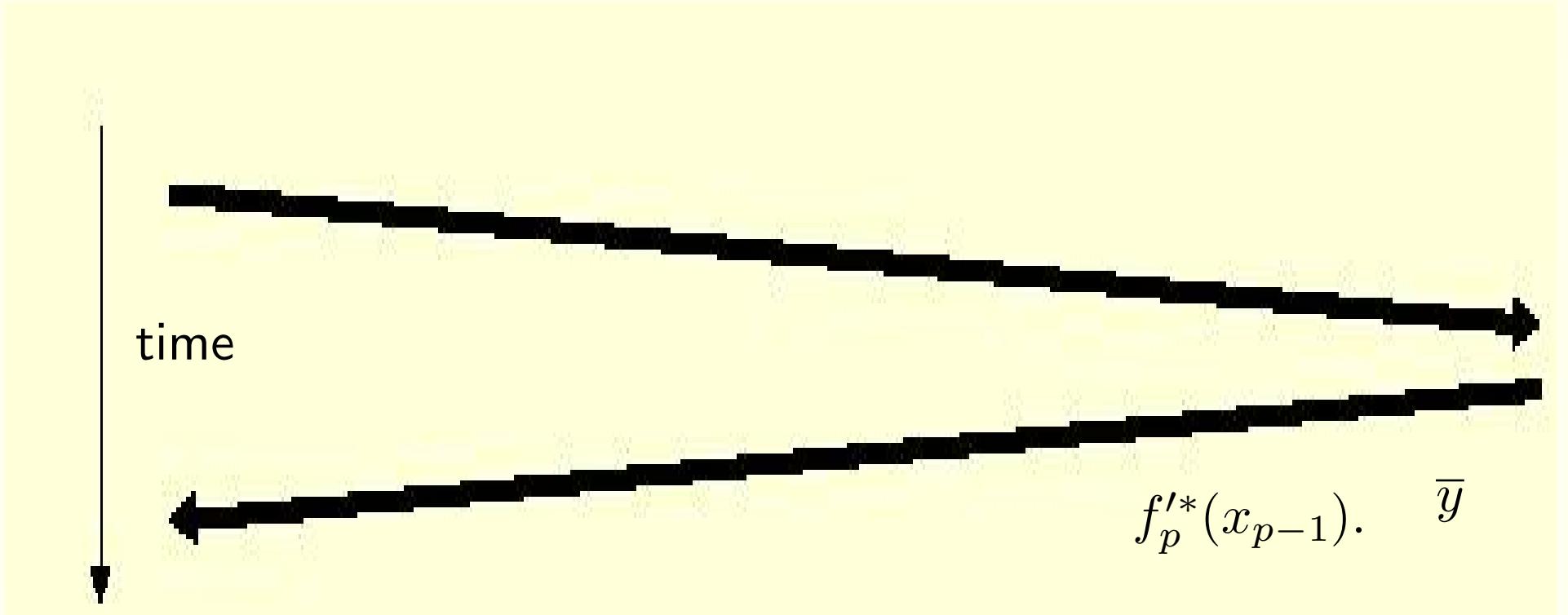
AD: Reverse is more tricky than Tangent

$$\bar{x} = f'^*(x). \bar{y} = f_1'^*(x_0) \dots f_{p-1}'^*(x_{p-2}). f_p'^*(x_{p-1}). \bar{y}$$



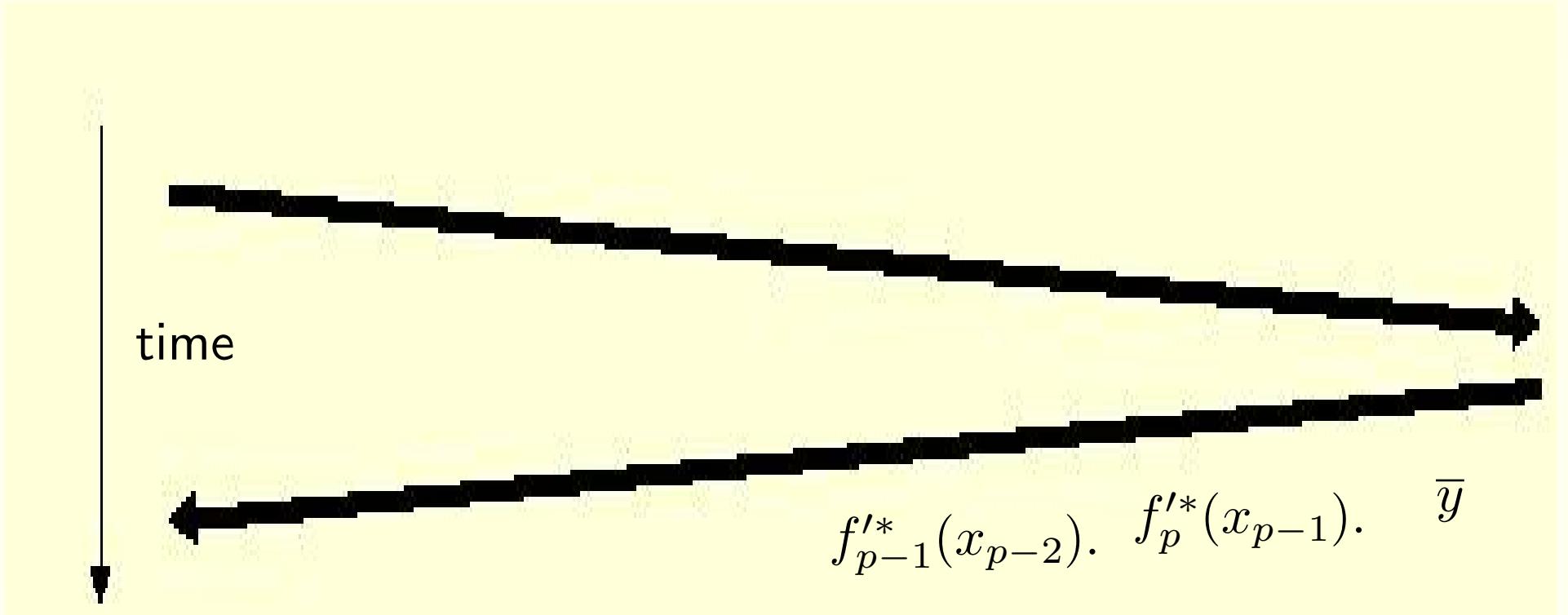
AD: Reverse is more tricky than Tangent

$$\bar{x} = f'^*(x). \bar{y} = f_1'^*(x_0) \dots f_{p-1}'^*(x_{p-2}). f_p'^*(x_{p-1}). \bar{y}$$



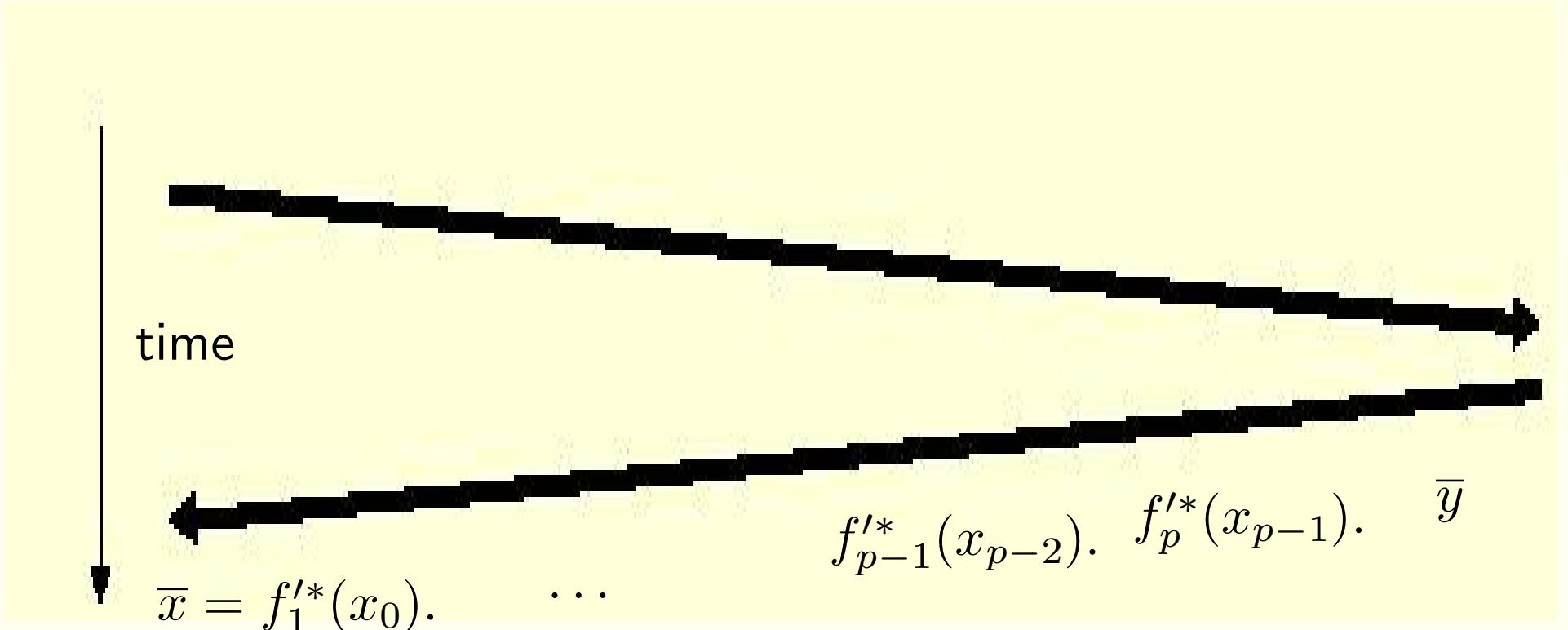
AD: Reverse is more tricky than Tangent

$$\bar{x} = f'^*(x). \bar{y} = f_1'^*(x_0) \dots f_{p-1}'^*(x_{p-2}) \cdot f_p'^*(x_{p-1}) \cdot \bar{y}$$



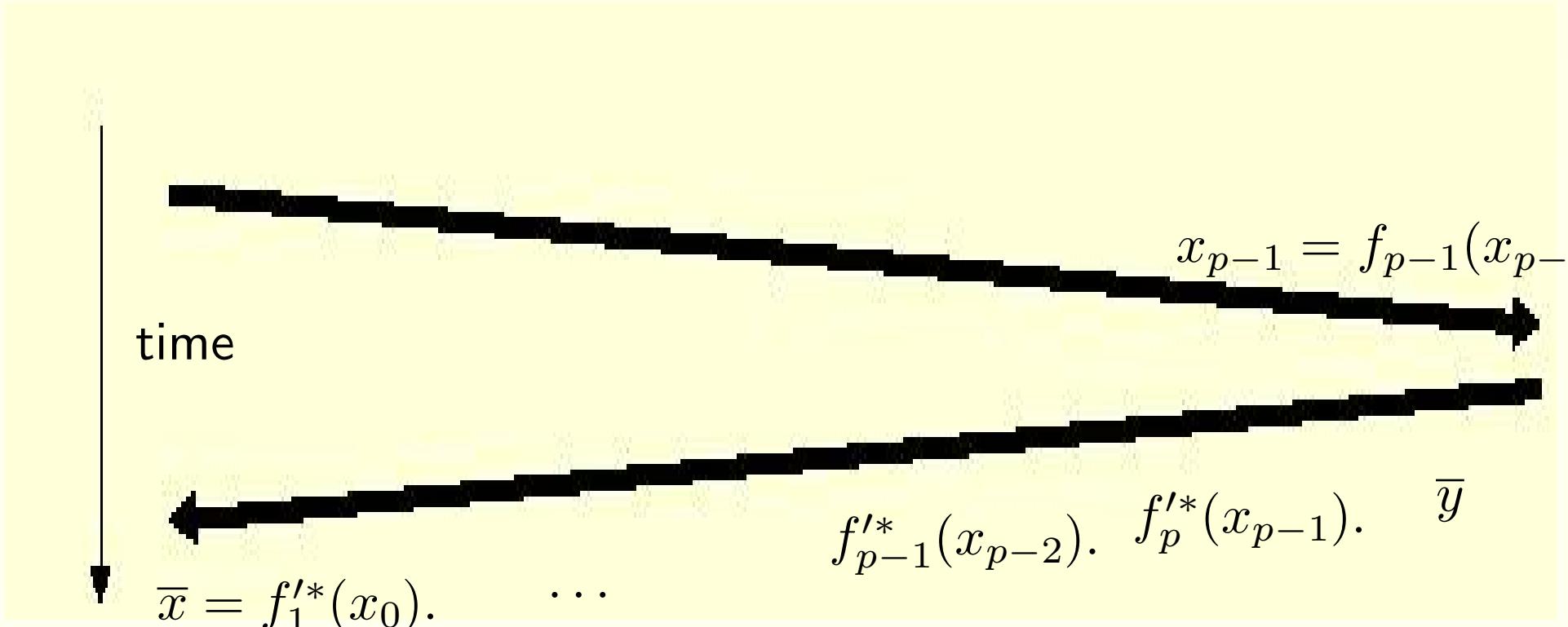
AD: Reverse is more tricky than Tangent

$$\bar{x} = f'^*(x). \bar{y} = f_1'^*(x_0) \dots f_{p-1}'^*(x_{p-2}) \cdot f_p'^*(x_{p-1}) \cdot \bar{y}$$



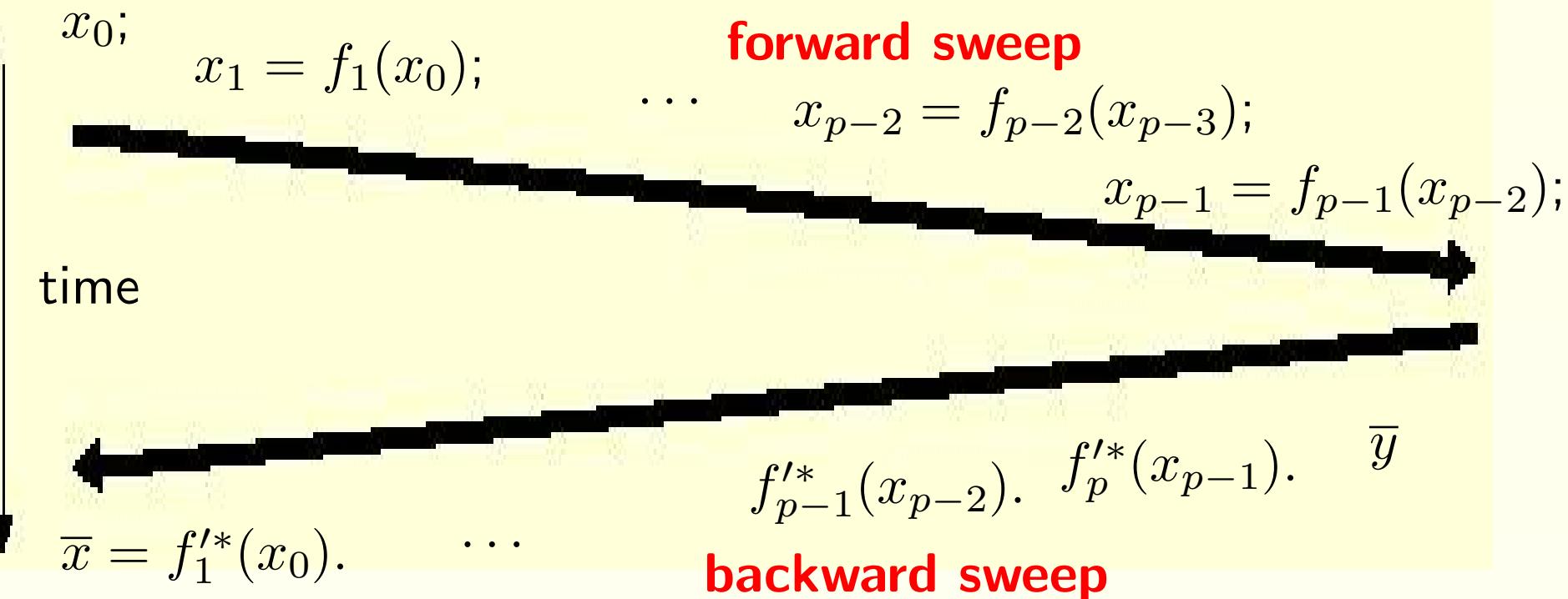
AD: Reverse is more tricky than Tangent

$$\bar{x} = f'^*(x). \bar{y} = f_1'^*(x_0) \dots f_{p-1}'^*(x_{p-2}) \cdot f_p'^*(x_{p-1}) \cdot \bar{y}$$



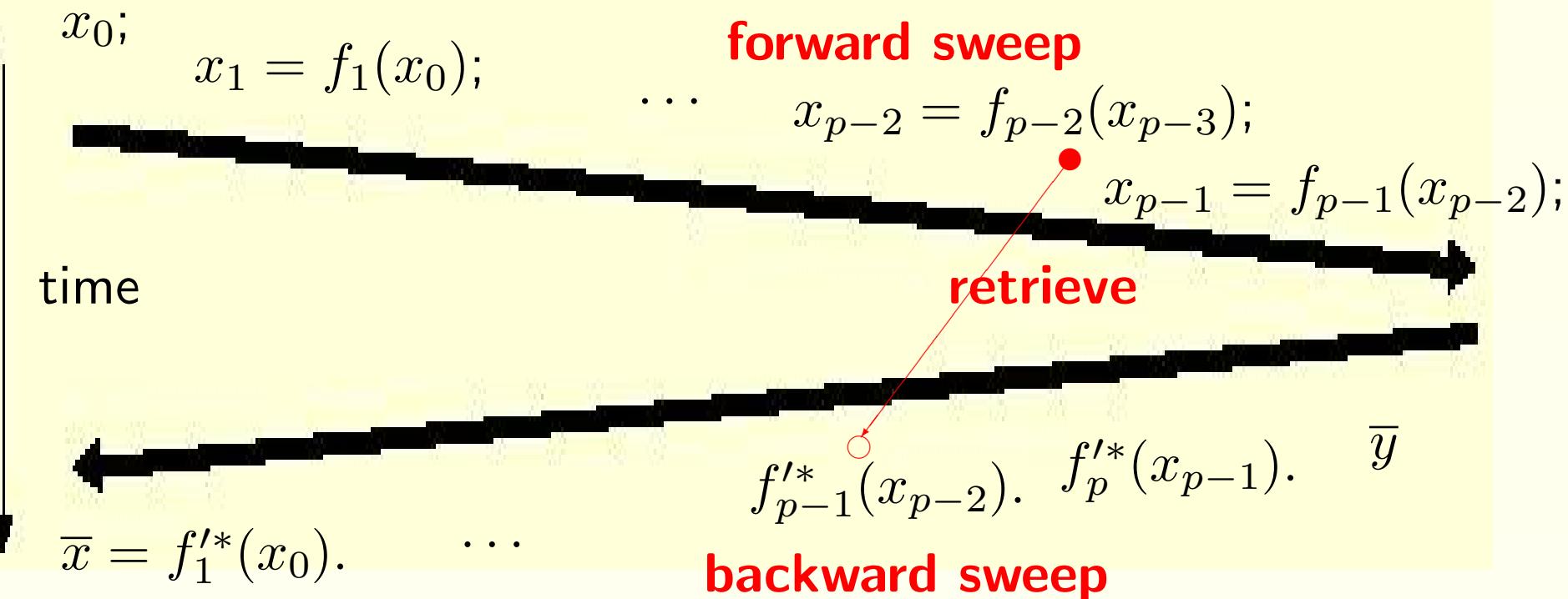
AD: Reverse is more tricky than Tangent

$$\bar{x} = f'^*(x). \bar{y} = f_1'^*(x_0) \dots f_{p-1}'^*(x_{p-2}) \cdot f_p'^*(x_{p-1}) \cdot \bar{y}$$



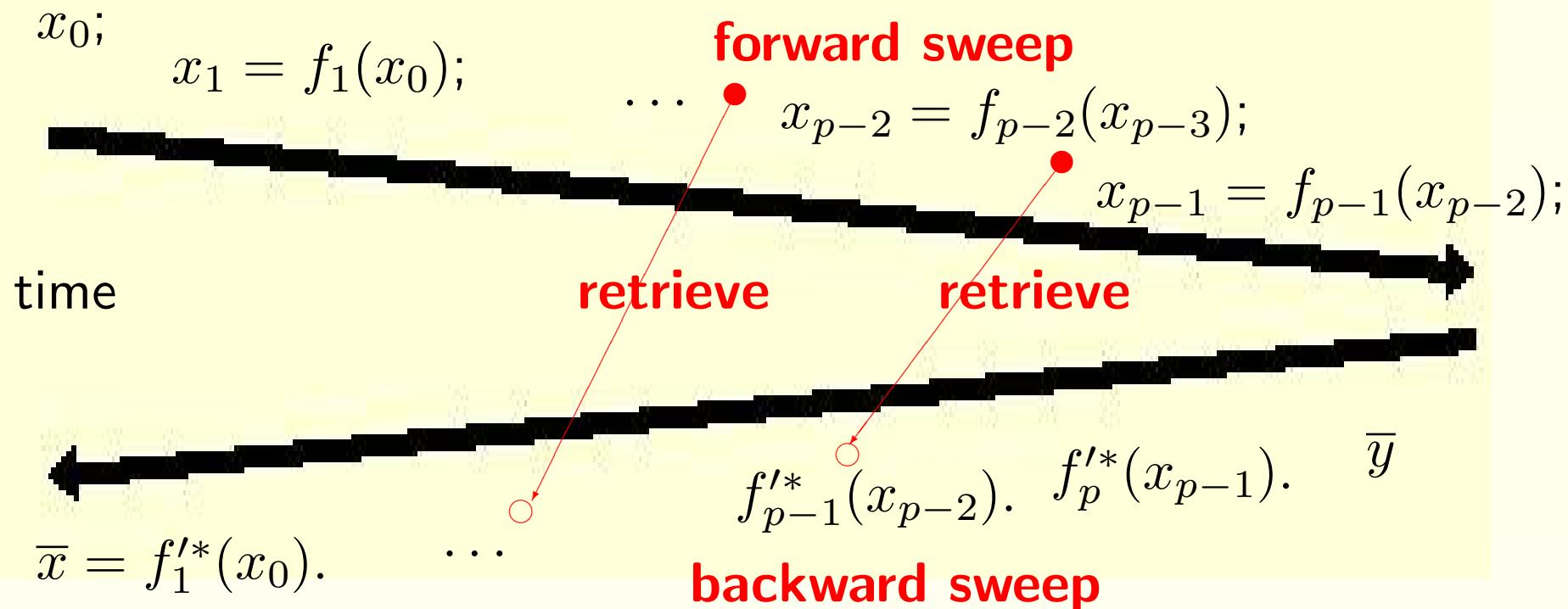
AD: Reverse is more tricky than Tangent

$$\bar{x} = f'^*(x). \bar{y} = f_1'^*(x_0) \dots f_{p-1}'^*(x_{p-2}) \cdot f_p'^*(x_{p-1}) \cdot \bar{y}$$



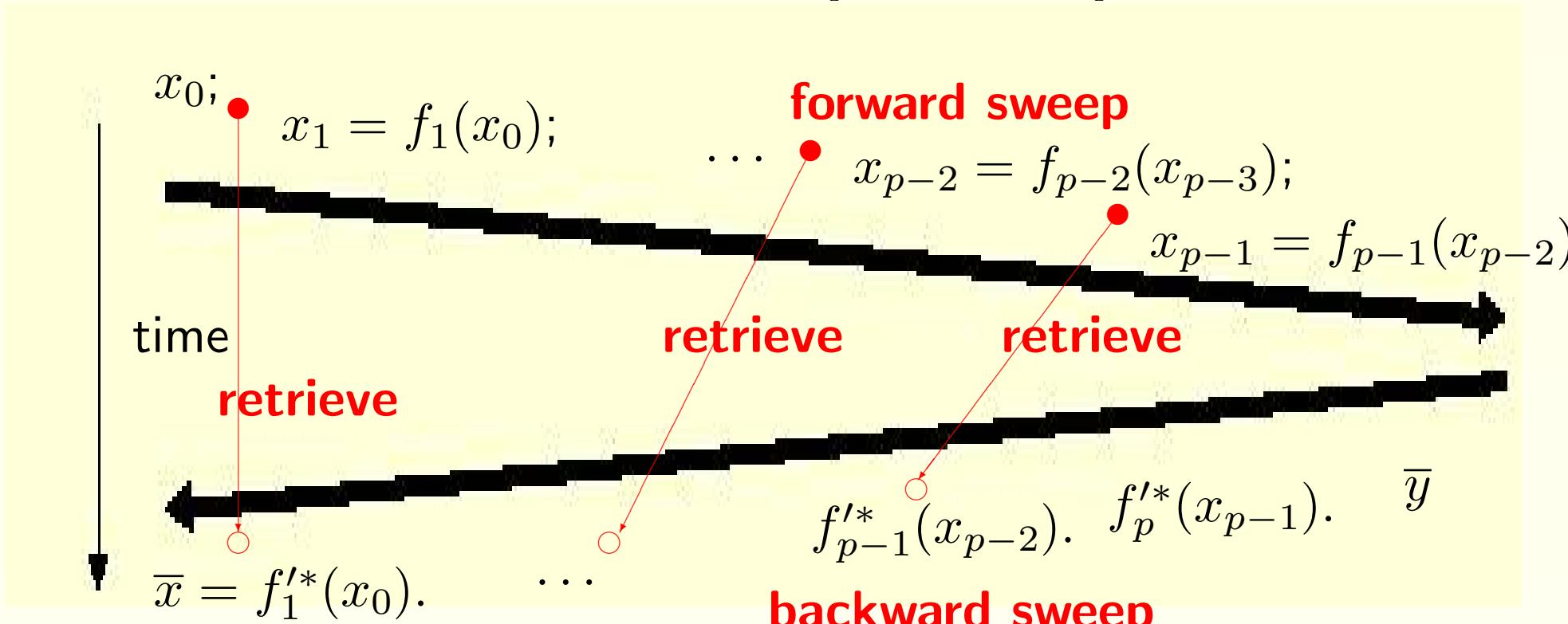
AD: Reverse is more tricky than Tangent

$$\bar{x} = f'^*(x). \bar{y} = f_1'^*(x_0) \dots f_{p-1}'^*(x_{p-2}) \cdot f_p'^*(x_{p-1}) \cdot \bar{y}$$



AD: Reverse is more tricky than Tangent

$$\bar{x} = f'^*(x). \bar{y} = f_1'^*(x_0) \dots f_{p-1}'^*(x_{p-2}). f_p'^*(x_{p-1}). \bar{y}$$



Memory usage (“Tape”) is the bottleneck!

AD: Continued Example

Program fragment:

...

$$\mathbf{v}_2 = \mathbf{2} * \mathbf{v}_1 + \mathbf{5}$$

$$\mathbf{v}_4 = \mathbf{v}_2 + \mathbf{p}_1 * \mathbf{v}_3 / \mathbf{v}_2$$

...

AD: Continued Example

Program fragment:

```

...
v2 = 2 * v1 + 5
v4 = v2 + p1 * v3/v2
...

```

Corresponding **transposed** Partial Jacobians:

$$f'^*(x) = \dots \begin{pmatrix} 1 & 2 & & \\ & 0 & & \\ & & 1 & \\ & & & 1 \end{pmatrix} \begin{pmatrix} 1 & & 0 & \\ & 1 & 1 - \frac{p_1 * v_3}{v_2^2} & \\ & & 1 & \frac{p_1}{v_2} \\ & & & 0 \end{pmatrix} \dots$$

AD: Reverse mode on the example

...

$$\begin{aligned}\bar{v}_2 &= \bar{v}_2 + \bar{v}_4 * (1 - p_1 * v_3/v_2^2) \\ \bar{v}_3 &= \bar{v}_3 + \bar{v}_4 * p_1/v_2 \\ \bar{v}_4 &= 0\end{aligned}$$

$$\bar{v}_1 = \bar{v}_1 + 2 * \bar{v}_2$$

$$\bar{v}_2 = 0$$

...

AD: Reverse mode on the example

...

$$\mathbf{v}_2 = 2 * \mathbf{v}_1 + 5$$

$$\mathbf{v}_4 = \mathbf{v}_2 + p_1 * \mathbf{v}_3 / \mathbf{v}_2$$

...

...

$$\bar{\mathbf{v}}_2 = \bar{\mathbf{v}}_2 + \bar{\mathbf{v}}_4 * (1 - p_1 * \mathbf{v}_3 / \mathbf{v}_2^2)$$

$$\bar{\mathbf{v}}_3 = \bar{\mathbf{v}}_3 + \bar{\mathbf{v}}_4 * p_1 / \mathbf{v}_2$$

$$\bar{\mathbf{v}}_4 = 0$$

$$\bar{\mathbf{v}}_1 = \bar{\mathbf{v}}_1 + 2 * \bar{\mathbf{v}}_2$$

$$\bar{\mathbf{v}}_2 = 0$$

...

AD: Reverse mode on the example

$\text{Push}(\overline{v}_2)$

$$\mathbf{v}_2 = 2 * \mathbf{v}_1 + 5$$

$\text{Push}(\overline{v}_4)$

$$\mathbf{v}_4 = \mathbf{v}_2 + p_1 * \mathbf{v}_3 / \mathbf{v}_2$$

...

$\text{Pop}(\overline{v}_4)$

$$\bar{\mathbf{v}}_2 = \bar{\mathbf{v}}_2 + \bar{\mathbf{v}}_4 * (1 - p_1 * \mathbf{v}_3 / \mathbf{v}_2^2)$$

$$\bar{\mathbf{v}}_3 = \bar{\mathbf{v}}_3 + \bar{\mathbf{v}}_4 * p_1 / \mathbf{v}_2$$

$$\bar{\mathbf{v}}_4 = 0$$

$\text{Pop}(\overline{v}_2)$

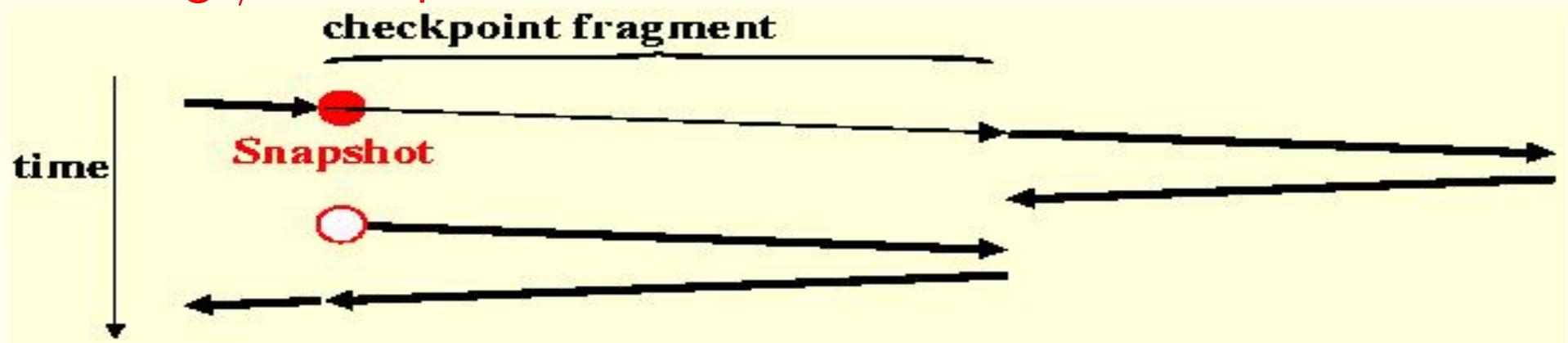
$$\bar{\mathbf{v}}_1 = \bar{\mathbf{v}}_1 + 2 * \bar{\mathbf{v}}_2$$

$$\bar{\mathbf{v}}_2 = 0$$

...

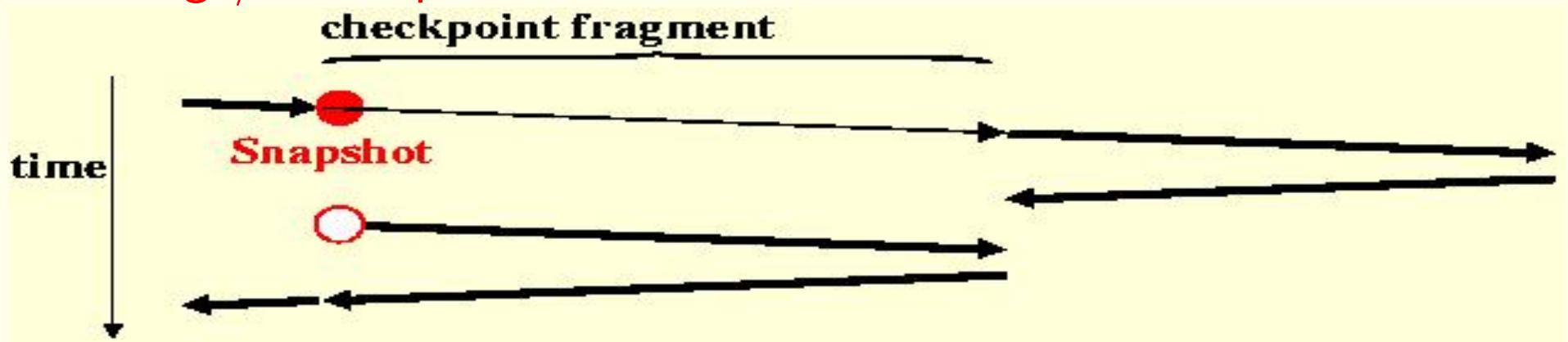
AD: The Checkpointing tactic

A Storage/Recomputation tradeoff:

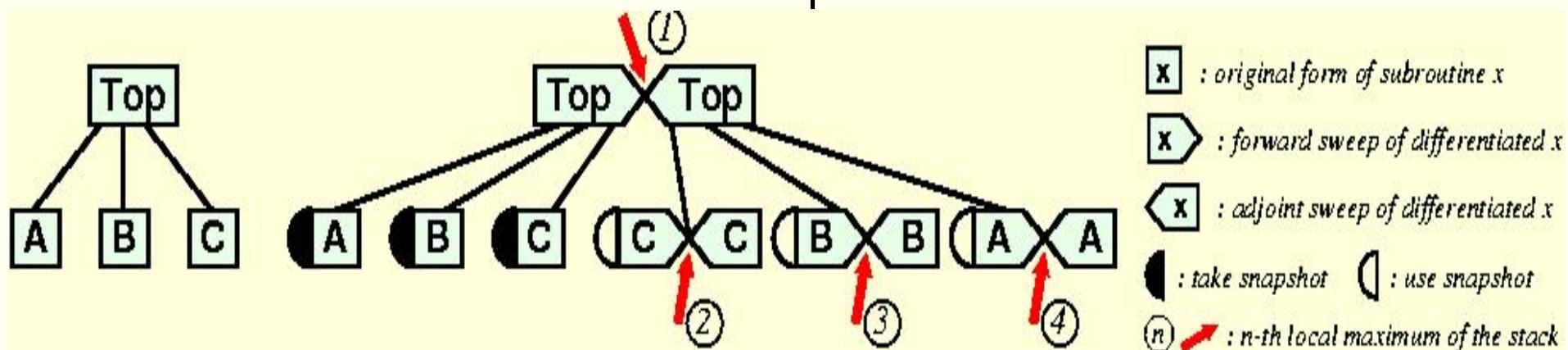


AD: The Checkpointing tactic

A Storage/Recomputation tradeoff:



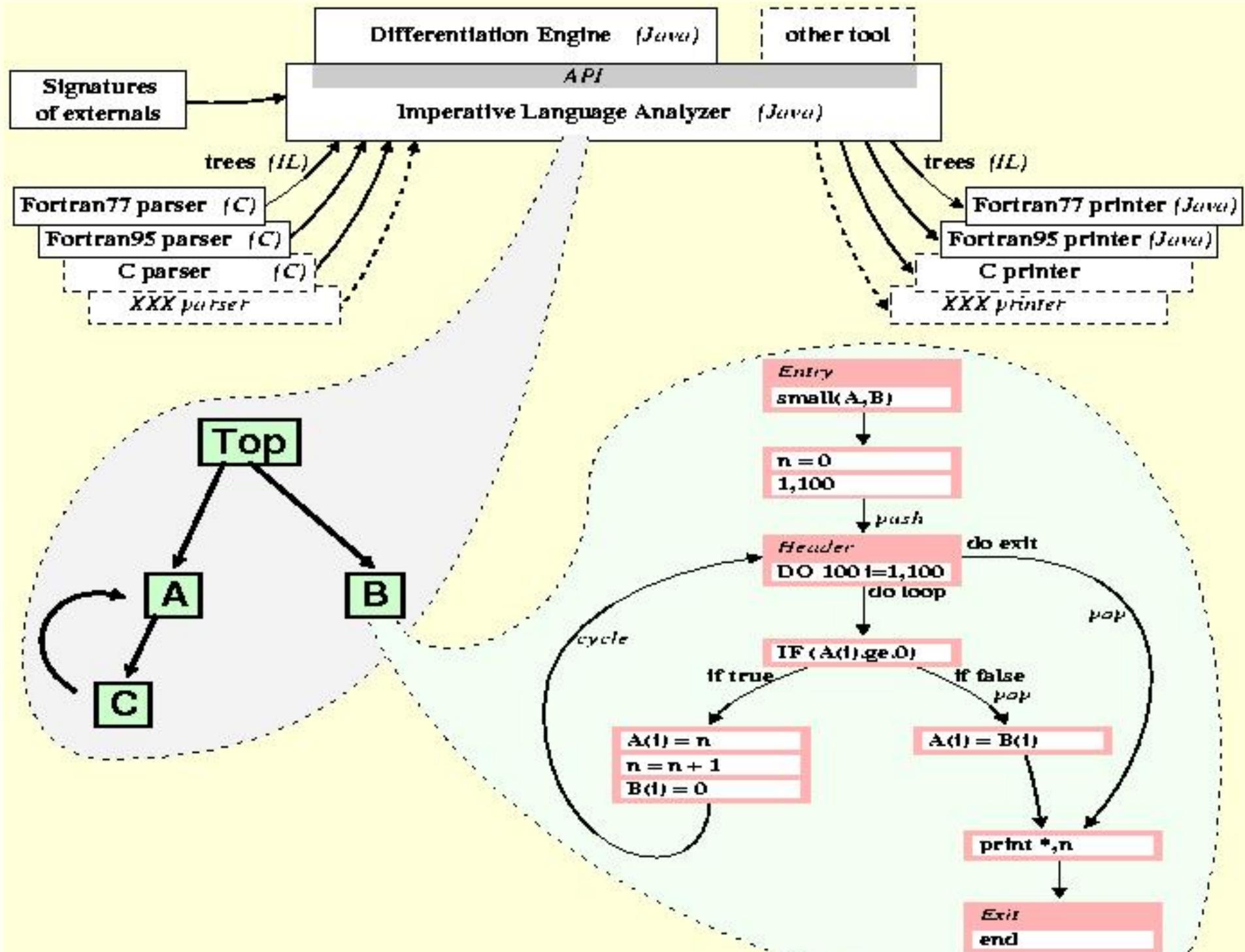
TAPENADE does it on the Call Graph :



Tapenade: Internal Representation

Take profit of well-known techniques
from Compilation and Parallelization:

- Use a general abstract *Imperative Language (IL)*
- Represent programs as *Call Graphs* or *Flow Graphs*
- Store symbol declarations in nested *Symbol Tables*



Tapenade Modes

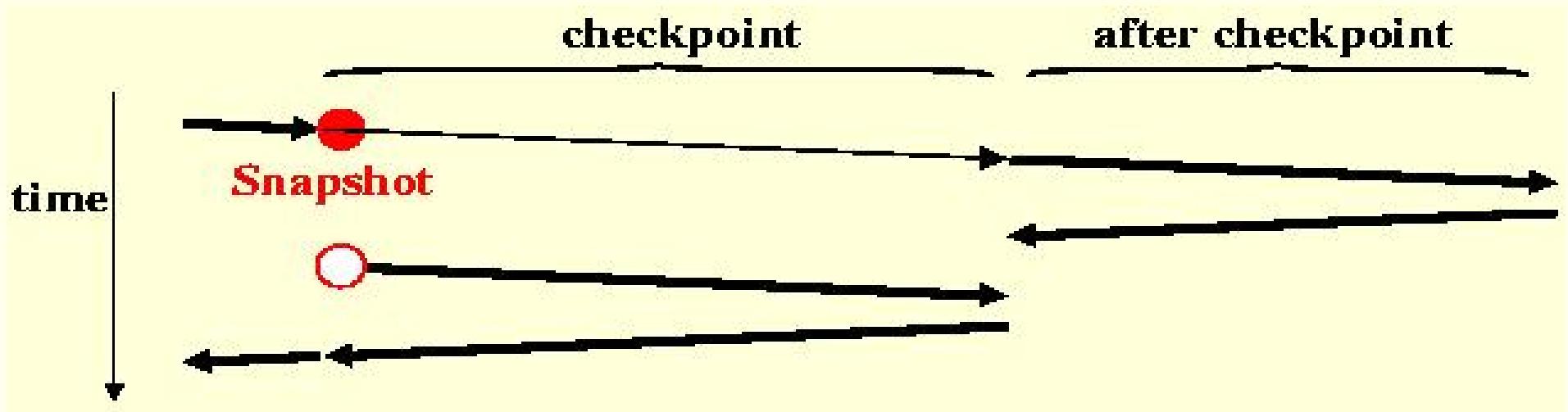
- normalize
- tangent
- multi-directional tangent
- reverse

Tapenade: Global Static Analyses on Flow Graphs

- forward (*resp.* backward) dependence wrt *independent inputs* (*resp.* *dependent inputs*)
- classical *IN-OUT* analysis (e.g. for snapshots)
- specific for the reverse mode: "*To Be Restored*", *Adjoint Dead Code*
- data-dependency analysis to reorder instructions.
- . . . *pointer analysis* . . .

Usual restrictions: conservative assumptions, arrays . . .

Example: reduced snapshots



Snapshot = $\overline{\text{IN}(\text{checkpoint})} \cap \text{OUT}(\text{checkpoint and after})$

Tapenade Differentiation model on examples

- Control structures
- Procedure calls and checkpointing
- "To Be Restored" analysis
- Instructions Reordering
- Dead Adjoint Code

Control structures

original program	Tapenade reverse: fwd sweep
<pre> SUBROUTINE S1(a, n, x) ... DO i=2,n,7 IF (a(i).GT.1.0) THEN a(i) = LOG(a(i)) + a(i-1) IF (a(i).LT.0.0) a(i)=2*a(i) END IF ENDDO </pre>	<pre> DO i=2,n,7 IF (a(i).GT.1.0) THEN CALL PUSHREAL4(a(i)) a(i) = LOG(a(i))+a(i-1) IF (a(i).LT.0.0) THEN CALL PUSHREAL4(a(i)) a(i) = 2*a(i) CALL PUSHINTEGER4(3) ELSE ... </pre>
Tapenade tangent	Tapenade reverse: bwd sweep
<pre> SUBROUTINE S1_D(a, ad, n, x) ... DO i=2,n,7 IF (a(i).GT.1.0) THEN ad(i)=ad(i)/a(i)+ad(i-1) a(i) = LOG(a(i)) + a(i-1) IF (a(i).LT.0.0) THEN ad(i) = 2*ad(i) a(i) = 2*a(i) END IF </pre>	<pre> CALL POPINTEGER4(adTo) DO i=adTo,2,-7 CALL POPINTEGER4(branch) IF (branch .GE. 2) THEN IF (branch .GE. 3) THEN CALL POPREAL4(a(i)) ab(i) = 2*ab(i) END IF CALL POPREAL4(a(i)) ab(j-1) = ab(j-1) + ab(j) ELSE ... </pre>

Procedure calls and checkpointing

original program	Tapenade reverse: fwd sweep
<pre>x = x**3 CALL SUB(a, x, 1.5, z) x = x*y</pre>	<pre>CALL PUSHREAL4(x) x = x**3 CALL PUSHREAL4(x) CALL SUB(a, x, 1.5, z) x = x*y</pre>
Tapenade tangent	Tapenade reverse: bwd sweep
<pre>xd = 3*x**2*xd x = x**3 CALL SUB_D(a, ad, x, xd, 1.5, 0.0, z) xd = y*xd x = x*y</pre>	<pre>xb = y*xb CALL POPREAL4(x) CALL SUB_B(a, ab, x, xb, 1.5, arg2b, z) CALL POPREAL4(x) xb = 3*x**2*xb</pre>

"To Be Restored" analysis

original program	reverse mode: naive bwd sweep	reverse mode: bwd sweep with TBR
<pre>x = x + EXP(a) y = x + a**2 a = 3*z</pre>	<pre>CALL POPREAL4(a) zb = zb + 3*ab ab = 0.0 CALL POPREAL4(y) ab = ab + 2*a*yb xb = xb + yb yb = 0.0 CALL POPREAL4(x) ab = ab + EXP(a)*xb</pre>	<pre>CALL POPREAL4(a) zb = zb + 3*ab ab = 0.0 ab = ab + 2*a*yb xb = xb + yb yb = 0.0 ab = ab + EXP(a)*xb</pre>

Instructions Reordering

original program	reverse mode: backward sweep with TBR	Tapenade reverse: non-incremental backward sweep
<pre>x = x + EXP(a) y = x + a**2 a = 3*z</pre>	<pre>CALL POPREAL4(a) zb = zb + 3*ab ab = 0.0 ab = ab + 2*a*yb xb = xb + yb yb = 0.0 ab = ab + EXP(a)*xb</pre>	<pre>CALL POPREAL4(a) zb = zb + 3*ab xb = xb + yb ab = 2*a*yb+EXP(a)*xb yb = 0.0</pre>

Dead Adjoint Code

original program	reverse mode:	Tapenade reverse: ajd. dead code removed
<pre> IF (a.GT.0.0) THEN a = LOG(a) ELSE a = LOG(c) CALL SUB(a) ENDIF END </pre>	<pre> IF (a .GT. 0.0) THEN CALL PUSHREAL4(a) a = LOG(a) CALL POPREAL4(a) ab = ab/a ELSE a = LOG(c) CALL PUSHREAL4(a) CALL SUB(a) CALL POPREAL4(a) CALL SUB_B(a, ab) cb = cb + ab/c ab = 0.0 END IF </pre>	<pre> IF (a .GT. 0.0) THEN ab = ab/a ELSE a = LOG(c) CALL SUB_B(a, ab) cb = cb + ab/c ab = 0.0 END IF </pre>

Tapenade: an AD tool on the web

Original call graph

- adj
 - sub2
 - sub1
 - maxx

```

C
SUBROUTINE ADJ(u, z, t)
REAL t, u, z
REAL x(14), y
COMMON /cc/ x, y
INTEGER i, MAXX
REAL v
EXTERNAL MAXX

i = 5
x(1) = y * u + t
z = MAXX(z, t)
u = 0.0
CALL SUB1(u, x(i), z, v)
t = t + x(1) * z + 3 * v
y = 0.0
i = 6
CALL SUB2(u, x(3), z, v)

```

Differentiated call graph

- adj_dv
 - maxx_dv
 - sub1_dv
 - sub2_dv

```

x(1) = y * u + t
CALL MAXX_DV(z, zd, t, td, z)
u = 0.0
CALL SUB1_DV(u, ud, x(i), xd(1)
DO nd=1,nbdirs
  td(nd) = td(nd) + z * xd(nd, :
ENDDO
t = t + x(1) * z + 3 * v
y = 0.0
i = 6
CALL SUB2_DV(u, ud, x(3), xd(1)
DO nd=1,nbdirs
  td(nd) = td(nd) + z * xd(nd, :
ENDDO
t = t + x(1) * z + 3 * u
DO nd=1,nbdirs
  zd(nd) = 0.0
ENDDO

```

2 adj: undeclared external routine: maxx
 3 adj: Return type of maxx set by implicit rule to INTEGER
 4 adj: argument type mismatch in call of sub1, REAL(0:6) expected, receives I
 5 adj: argument type mismatch in call of sub2, REAL(0:12) expected, receives

6 maxx: Tool: Please provide a differentiated function for unit maxx for arg

- Servlet on <http://www-sop.inria.fr/tropics> or batch
- Uploads your Files and Includes
- Displays results and messages with links to source

Conclusion and future work

Tapenade now 3 years old.

Several applications on industrial and academic codes:
Aeronautics, Hydrology, Chemistry, Biology, Agronomy...

Future developments:

- In progress: FORTRAN95, and then C (\Rightarrow pointers!)
- User Directives: active I-O, checkpoints, special loops
- Validity domain for derivatives