

## Model Order Reduction for Maxwell Equations based on Moment Matching

Matthias Bollhöfer (joint work with André Bodendiek) INRIA Sophia Antipolis , July 28, 2015

## Outline

- Maxwell's Equations
- Model Order Reduction
- MOR for Maxwell Equations
- Model Order Reduction Based on Moment Matching
- Numerical Results
- Modified Adaptive-Order Rational Arnoldi Method
- Savings for the AORA method
- QMR Method
- Simplified QMR
- QMR with Subspace Recycling
- Numerical Results
- Conclusions


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## Model Problems



Left picture: branchline coupler on a substrate with PMC boundary conditions, two parallel microstriplines, coupled together in form of a transversal bridge, frequency range 1.0 to $10.0 \mathrm{GHz}, N=73^{\prime} 385$.

Right picture: coplanar waveguide with a dielectric overlay, PEC boundary conditions, frequency range 0.6 to $3.0 \mathrm{GHz}, N=32^{\prime} 924$.

## Model Problems



PCB circuit on a substrate within the frequency range from 7.5 to $10.0 \mathrm{GHz}, N=226^{\prime} 458$

PEC boundary condition for the conducting lines, PMC boundary condition for the rest

## Maxwell's Equations

$$
\begin{aligned}
\frac{\partial(\varepsilon \mathbf{E})}{\partial t} & =-\sigma \mathbf{E}+\nabla \times \mathbf{H} \\
\frac{\partial(\mu \mathbf{H})}{\partial t} & =-\nabla \times \mathbf{E} \\
(0 & =\nabla \cdot(\varepsilon \mathbf{E}), 0=\nabla \cdot(\mu \mathbf{H}))
\end{aligned}
$$

$\varepsilon$ electric permittivity, $\mu$ magnetic permeability, $\sigma$ electric conductivity.

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$\varepsilon$ electric permittivity, $\mu$ magnetic permeability, $\sigma$ electric conductivity.
Discrete equations:

$$
\begin{aligned}
M_{\varepsilon} \dot{E} & =-M_{\sigma} E+G H+B_{E} u \\
M_{\mu} \dot{H} & =-G^{T} E \quad+B_{H} u \\
(0 & \left.=D_{E} M_{\varepsilon} E, \quad 0=D_{H} M_{\mu} H\right) \\
y & =C_{E} E+C_{H} H
\end{aligned}
$$

- $u$ input, $y$ output
- $M_{\varepsilon}, M_{\mu}$ are sym. pos. def., $M_{\sigma}$ sym. pos. semidef. (mass matrices)
- $G$ highly singular! (curl operator)


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## Model Order Reduction

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\begin{aligned}
\mathcal{M} \dot{x} & =\mathcal{A} x+\mathcal{B} u \\
(0 & =\mathcal{D} x) \\
y & =\mathcal{C} x \\
\mathcal{M}=\left(\begin{array}{cc}
M_{\varepsilon} & 0 \\
0 & M_{\mu}
\end{array}\right), \mathcal{A} & =\left(\begin{array}{ll}
-M_{\sigma} & G \\
-G^{T} & 0
\end{array}\right), x=\binom{E}{H} .
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$$

Model Order Reduction: find full rank $S, T \in \mathbb{R}^{2 n, r}$ such that $r \ll 2 n$ and use instead

$$
\begin{aligned}
\left(S^{*} \mathcal{M} T\right) \dot{\tilde{x}} & =\left(S^{*} \mathcal{A} T\right) \tilde{x}+\left(S^{*} \mathcal{B}\right) u \\
(0 & =(\mathcal{D} T) \tilde{x}) \\
\tilde{y} & =(\mathcal{C} T) \tilde{x} \\
& \|y-\tilde{y}\| \text { small }
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## Structure-Preserving MOR for Maxwell's Equations

$$
\text { Here use } S=T=\left(\begin{array}{cc}
V & 0 \\
0 & W
\end{array}\right)
$$

Structured MOR: find full rank $V, W \in \mathbb{R}^{n, r}$ such that $r \ll n$ and use instead

$$
\begin{array}{rlr}
\left(V^{*} M_{\varepsilon} V\right) \dot{e} & =-\left(V^{*} M_{\sigma} V\right) e+\left(V^{*} G W\right) h+\left(V^{*} B_{E}\right) u \\
\left(W^{*} M_{\mu} W\right) \dot{h} & =-\left(W^{*} G^{T} V\right) e \quad+\left(W^{*} B_{H}\right) u \\
(0 & \left.=\left(D_{E} M_{\varepsilon} V\right) e, \quad 0=\left(D_{H} M_{\mu} W\right) h\right) \\
\tilde{y} & =\quad\left(C_{E} V\right) e+\left(C_{H} W\right) h
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\|y-\tilde{y}\| \text { small }
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## Moment-Matching — Basic Idea

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\mathcal{M}=\left(\begin{array}{cc}
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\end{array}\right), \mathcal{A}=\left(\begin{array}{cc}
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Transfer function

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\mathcal{H}(s)=\mathcal{C P}(s \mathcal{M}-\mathcal{A})^{-1} \mathcal{B} \quad(\mathcal{P} \text { projector to divergence-free part })
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Taylor/Laurent expansion at some expansion point $s_{j}$ :

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\mathcal{A}_{j}:=\left(s_{j} \mathcal{M}-\mathcal{A}\right)^{-1} \mathcal{M}, \mathcal{B}_{j}:=\left(s_{j} \mathcal{M}-\mathcal{A}\right)^{-1} \mathcal{B}, \mathcal{C}_{j}:=\mathcal{C P}\left(s_{j} \mathcal{M}-\mathcal{A}\right)^{-1}
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& \Rightarrow \mathcal{H}(s)=\sum_{p=0}^{\infty} \mathcal{C P} \overbrace{\left[-\mathcal{A}_{j}\right]^{p} \mathcal{B}_{j}}^{x_{j}^{(p)}}\left(s-s_{j}\right)^{p}=\sum_{p=0}^{\infty} \underbrace{\mathcal{C}_{j}\left[-\mathcal{A}_{j}\right]^{p}}_{Y_{j}^{(p)}} \mathcal{B}\left(s-s_{j}\right)^{p}
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\end{aligned}
$$

$X_{j}^{(p)}$ input moments, Taylor coefficients $Z_{j}^{(p)}=\mathcal{C P} X_{j}^{(p)}=Y_{j}^{(p)} \mathcal{B}$ output moments.

## Krylov Subspace Methods

Krylov subspace

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\left(Y_{j}^{(p)}\right)^{*} \in \mathcal{K}_{p}\left(\mathcal{A}_{j}^{*}, \mathcal{C}_{j}^{*}\right)
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- Lanczos-type methods [PVL,Gragg'74,Gutknecht'92,Feldmann,Freund'94,...] generate dual basis $T \equiv T_{r} \in \mathbb{R}^{2 n, r}$ and $S \equiv S_{r} \in \mathbb{R}^{2 n, r}$ of

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\mathcal{K}_{r}\left(\mathcal{A}_{j}, \mathcal{B}_{j}\right) \text { and } \mathcal{K}_{r}\left(\mathcal{A}_{j}^{*}, \mathcal{C}_{j}^{*}\right)
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such that $S^{*} T=I \rightarrow$ matches $2 r$ moments $Z_{j}^{(0)}, \ldots, Z_{j}^{(2 r-1)}$

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such that $S^{*} T=I \rightarrow$ matches $2 r$ moments $Z_{j}^{(0)}, \ldots, Z_{j}^{(2 r-1)}$

- Arnoldi-type methods [PRIMA,Odabasioglu et al.'96,'97], [SPRIM,Freund'04,'08] compute one orthonormal basis $S=T=Q \equiv Q_{r} \in \mathbb{R}^{2 n, r}$, say, from

$$
\mathcal{K}_{r}\left(\mathcal{A}_{j}, \mathcal{B}_{j}\right)
$$

using modified Gram-Schmidt $\rightarrow$ matches $r$ moments $Z_{j}^{(0)}, \ldots, Z_{j}^{(r-1)}$

## Structure-Preserving Moment Matching Methods

Structure Preservation for Maxwell Equations

$$
\begin{aligned}
& \mathcal{M}=\left(\begin{array}{cc}
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$$

- twice as big, but ...
- block-structure preserved, still $r$ moments matched
- if we are lucky, up to $2 r$ moments could be matched


## Structure-Preserving Moment Matching Methods

Problems

- (How to) select $s_{j} \in\left[f_{\text {min }}, f_{\text {max }}\right]$
- (No) error bounds!? Choice of $r, I$, accuracy of the reduced model

$$
\begin{gathered}
\mathcal{H}(s)=\mathcal{C}(s \mathcal{M}-\mathcal{A})^{-1} \mathcal{B}, \quad \mathcal{H}_{r}(s)=\hat{\mathcal{C}}(s \hat{\mathcal{M}}-\hat{\mathcal{A}})^{-1} \hat{\mathcal{B}} \\
\left\|\mathcal{H}(i \omega)-\mathcal{H}_{r}(i \omega)\right\| \leqslant \ldots
\end{gathered}
$$

- multiple expansion points $s_{1}, \ldots s_{l} \in\left[f_{\text {min }}, f_{\text {max }}\right]$
- Restarting Arnoldi and increasing $r$ or / whenever the "error estimate" is not accurate enough


## Rational Arnoldi Methods

- multiple expansion points $s_{1}, \ldots, s_{l}$
- multiple associated Taylor expansions

$$
\mathcal{H}(s)=\sum_{p=0}^{\infty} Z_{j}^{(p)}\left(s-s_{j}\right)^{p}, j=1, \ldots, l
$$

- Rational Krylov method: Compute basis $Q_{r}$ for the Krylov subspaces

$$
\sum_{j=1}^{1} \mathcal{K}_{r_{j}}\left(\mathcal{A}_{j}, \mathcal{B}_{j}\right)
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## Lemma (Key-Lemma, Grimme,Gallivan'98)

If $s_{j} \neq s_{k}$ then

$$
\mathcal{A}_{k} \cdot \mathcal{A}_{j}^{p-1} \mathcal{B}_{j} \in \mathcal{K}_{p}\left(\mathcal{A}_{j}, \mathcal{B}_{j}\right)+\mathcal{K}_{1}\left(\mathcal{A}_{k}, \mathcal{B}_{k}\right)
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$$

$\Rightarrow$ mixing inverses with different shifts leads to a separate sum of Krylov subspaces, no "mixed powers of inverses"

## Rational Arnoldi Methods

- Traditional Rational Arnoldi methods build $Q_{r}$ w.r.t. $\mathcal{K}_{r_{j}}\left(\mathcal{A}_{j}, \mathcal{B}_{j}\right), j=1, \ldots$, l one after another using modified Gram-Schmidt.


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$$
\mathcal{K}_{r_{1}}\left(\mathcal{A}_{1}, \mathcal{B}_{1}\right)
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$$
\mathcal{K}_{r_{1}}\left(\mathcal{A}_{1}, \mathcal{B}_{1}\right)+\mathcal{K}_{r_{2}}\left(\mathcal{A}_{2}, \mathcal{B}_{2}\right)
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$$

- At each step $s_{j}$ is selected w.r.t. the largest output moment error $Z_{j}^{(r)}-\tilde{Z}_{j}^{(r)}$ which can be computed cheaply


## Expansion Point Selection

- AORA called repeatedly $I=1,2,3, \ldots$ with increasing number of expansion points


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$$
s_{1}=2 \pi i f_{\min }, s_{2}=2 \pi i \sqrt{f_{\min } f_{\max }}, s_{3}=2 \pi i f_{\max }
$$

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$$
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$$
s_{1}, s_{2}, s_{3}, s_{4}, s_{5}, \mathbf{s}_{6}
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- global stopping criterion [Grimme,Gallivan'98] $\sum_{i=1}^{l} 2^{i-l} \frac{\left|\mathcal{H}_{r}^{(i)}(s)-\mathcal{H}_{r}^{(i-1)}(s)\right|}{\left|\mathcal{H}_{r}^{(i)}(s)\right|} \leqslant \varepsilon$


## Outline

- Maxwell's Equations
- Model Order Reduction
- MOR for Maxwell Equations
- Model Order Reduction Based on Moment Matching
- Numerical Results
- Modified Adaptive-Order Rational Arnoldi Method
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- QMR Method
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## Model Order Reduction - Numerical Results

- model problems have a frequency range in $\left[f_{\min }, f_{\max }\right]$
- computed reduced order models have fixed size $n=25(50)$
- SPRIM uses expansion point $s_{0}=\frac{f_{\min }+f_{\max }}{2}$
- expansion point selection based on relative error $10^{-9}$
- strategy finally leads to $I=8$ expansion points $s_{1}, \ldots, s_{8}$
- Rational Arnoldi (RA) and Adaptive Order Rational Arnoldi (AORA) repeated 5 times
- RA uses fixed sizes $j=n / l$ for each Krylov subspace $\mathcal{K}_{j}$
- AORA adaptively increases each $\mathcal{K}_{j_{i}}$


## Model Order Reduction - Numerical Results

## Branchline Coupler



- size $N=73385$, discretized using FIT
- frequency range $\left[f_{\text {min }}, f_{\text {max }}\right]=\left[10^{9}, 10^{10}\right]$


## Model Order Reduction - Numerical Results <br> Branchline Coupler

relative error

$$
\epsilon_{\text {rel }}(f)=\frac{|\mathcal{H}(s)-\tilde{\mathcal{H}}(s)|}{|\mathcal{H}(s)|}
$$

$s=2 \pi i f$ with $f \in\left[10^{9}, 10^{10}\right]$ displayed as reference



## Model Order Reduction - Numerical Results

## Coplanar Waveguide




- size $N=32924$, discretized using FIT
- frequency range $\left[f_{\min }, f_{\max }\right]=\left[0.6 \cdot 10^{9}, 3.0 \cdot 10^{9}\right]$


## Model Order Reduction - Numerical Results <br> Coplanar Waveguide

relative error

$$
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Repeatedly calling AORA method requires to recompute span $Q_{r}$ from scratch.

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## Lemma

Suppose that

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\operatorname{span} Q_{r}^{(I+1)}=\mathcal{K}_{r_{1}^{(I+1)}}\left(s_{1}\right)+\cdots+\mathcal{K}_{r_{l}^{(I+1)}}\left(s_{l}\right)+\mathcal{K}_{r_{l+1}^{(I+1)}}\left(s_{l+1}\right)
\end{gathered}
$$

such that

$$
r_{1}^{(I+1)} \leqslant r_{1}^{(I)} \cdots r_{I}^{(I+1)} \leqslant r_{l}^{(I)}
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Then $\mathcal{K}_{r_{1}^{(I+1)}}\left(s_{1}\right)+\cdots+\mathcal{K}_{r_{l}^{(I+1)}}\left(s_{l}\right)$ can be directly extracted from $\operatorname{span} Q_{r}^{(I)}$.

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- Lemma requires that shifts $s_{1}, \ldots, s_{l}$ are selected in the same order
- new shift $s_{I+1}$ can be injected at any time
- only new shift $s_{I+1}$ requires to solve systems with $s_{I+1} \mathcal{M}-\mathcal{A}$
- $\longrightarrow$ mAORA (modified AORA)


## Solving Systems in the Modified AORA Method

Recall

- $\mathcal{J}(s \mathcal{M}-\mathcal{A})$ is complex-symmetric, where $\mathcal{J}=\left(\begin{array}{cc}I & 0 \\ 0 & -I\end{array}\right)$
- Schur complement $\mathcal{S}(s)=s^{2} M_{\varepsilon}+s M_{\sigma}+G M_{\mu}^{-1} G^{T}$ is complex-symmetric


## Solving Systems in the Modified AORA Method

## Recall

- $\mathcal{J}(s \mathcal{M}-\mathcal{A})$ is complex-symmetric, where $\mathcal{J}=\left(\begin{array}{cc}1 & 0 \\ 0 & -1\end{array}\right)$
- Schur complement $\mathcal{S}(s)=s^{2} M_{\varepsilon}+s M_{\sigma}+G M_{\mu}^{-1} G^{T}$ is complex-symmetric
- (modified) AORA requires solving a sequence complex-symmetric systems with varying shifts and varying right hand sides

$$
A\left(s_{j}\right) x_{j p}=b_{j p}
$$

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- (modified) AORA requires solving a sequence complex-symmetric systems with varying shifts and varying right hand sides

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- Save memory and time: Only compute factorization of $A\left(s_{*}\right)$ for some characteristic shift $s_{*}=2 \pi i \sqrt{f_{\text {min }} f_{\max }}$.
- For all $s_{j} \neq s_{*}$ use (recycling) Krylov subspace method as wrapper.


## Outline

. Maxwell's Equations

- Model Order Reduction
- MOR for Maxwell Equations
- Model Order Reduction Based on Moment Matching
- Numerical Results
- Modified Adaptive-Order Rational Arnoldi Method
- Savings for the AORA method
- QMR Method
- Simplified QMR
- QMR with Subspace Recycling
. Numerical Results
- Conclusions


## QMR Method - Sketch

## Two-sided Lanczos Method

$$
A V_{k}=V_{k+1} \underline{T_{k}}, A^{T} \tilde{V}_{k}=\tilde{V}_{k+1} \underline{\tilde{T}_{k}}
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QMR for $A x=b$ based on two-sided Lanczos method:

- $r_{0}=b-A x_{0}, v_{1}=r_{0} /\left\|r_{0}\right\|$
- $x_{k}=x_{0}+V_{k} y$
- quasi-minimize

$$
\left\|b-A x_{k}\right\| \leqslant\left\|V_{k+1}\right\| \cdot\| \| r_{0}\left\|e_{1}-\underline{T_{k}} y\right\|
$$

$\longrightarrow y$ from least squares solution
$\longrightarrow x$

- [Freund et al. '91,'93,'94]


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## Simplified QMR Method for J-Symmetric Matrices

- $A^{T} J=J A$, where $J=J^{T}, J$ nonsingular
- use specific left initial guess $\tilde{v}_{1}=J v_{1}$


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- [Freund et al. '94,'95]
- Maxwell equations: complex-symmetric system $A$ with complex-symmetric preconditioner $P$ satisfy $J$-symmetry!


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## QMR with Subspace Recycling

Krylov subspace methods with subspace recycling

- GCROT, GCRO-DR [De Sturler et al.'99,'06]
- recycling BiCG, recycling BiCGStab [Ahuja et al. '12,'13]

Main idea: $U, \tilde{U} \in \mathbb{C}^{n, r}, C=A U, \tilde{C}=A^{T} \tilde{U}, r \ll n$ given subspaces, $D_{c}=\tilde{C}^{T} C$. Apply Lanczos (Arnoldi) to the systems

$$
\begin{aligned}
& A V_{k}=V_{k+1} \underline{T_{k}} \\
& A^{T} \tilde{V}_{k}=\tilde{V}_{k+1} \underline{\tilde{I}_{k}}
\end{aligned}
$$

with updated initial guess and projected initial residual

$$
x_{0}^{(\text {new })}=x_{0}+U D_{c}^{-1} \tilde{C}^{T} r_{0}, r_{0}^{(\text {new })}=\left(I-C D_{c}^{-1} \tilde{C}^{T}\right) r_{0}
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$$
\begin{aligned}
& \left(I-C D_{c}^{-1} \tilde{C}^{T}\right) A V_{k}=V_{k+1} \underline{T_{k}} \\
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## Recycling QMR/BiCG in a Nutshell

$$
U, \tilde{U}, C=A U, \tilde{C}=A^{\top} \tilde{U}, D_{c}=\tilde{C}^{\top} C .
$$

Changing the methods:

| BiCG/QMR |  | Recycling BiCG/QMR |
| :--- | :--- | :--- |
| init $x_{0}$ |  | $x_{0}^{(\text {new })}=x_{0}+U D_{c}^{-1} \tilde{C}^{T} r_{0}$ |
| init $r_{0}=b-A x_{0}$ |  | $r_{0}^{(\text {new })}=\left(I-C D_{c}^{-1} \tilde{C}^{T}\right) r_{0}$ |
| matvec $A x$ |  | $\left(I-C D_{c}^{-1} \tilde{C}^{T}\right) A x$ |
| matvec $A^{T} x$ |  | $\left(I-\tilde{C} D_{c}^{-T} C^{T}\right) A^{T} x$ |
| solution update $x=x+\alpha p$ | $x=x+\alpha\left(I-U D_{c}^{-1} \tilde{C}^{T}\right) p$ |  |

## Recycling QMR/BiCG in a Nutshell

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Preconditioning and simplified QMR: a little bit more tricky but similar

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## Modified AORA versus AORA





## Model Order Reduction <br> Numerical Results



Sampling 20 frequencies $s_{j}=2 \pi i f_{j}$
$f_{j} \in\left[f_{\text {min }}, f_{\text {max }}\right]=\left[0.6 \cdot 10^{9}, 3 \cdot 10^{9}\right]$
Use complex-symmetric Schur complement system
Preconditioner: $L U$-decomposition for $s_{*}=\sqrt{f_{\text {min }} f_{\text {max }}}$

Comparison

- simplified QMR without subspace recycling
- simplified QMR using subspace recycling


## Model Order Reduction

## Numerical Results

Preconditioner: LU decomposition for $\mathcal{S}\left(\boldsymbol{s}_{*}\right)$
Preconditioned SQMR Without Subspace Recycling



## Model Order Reduction

## Numerical Results

Preconditioner: LU decomposition for $\mathcal{S}\left(\boldsymbol{s}_{*}\right)$

## Preconditioned SQMR With Subspace Recycling




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## Conclusions

- (adaptive order) rational Arnoldi yields reduced order model for Maxwell equations
- main important structures can be preserved
- outer iteration: several calls of (modified) rational Arnoldi, inner iteration: preconditioned SQMR
- Recycling techniques $\rightarrow$ modified AORA, recycling SQMR

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