Statistical analysis: complementary topics

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Statistical analysis: complementary topics

Statistical analysis: complementary topics

Beyond Two-Sample-Tests
Problem
Jensen-Shannon divergence and discrepancy
Density based clustering

Comparing clusterings
Motivation
Problem statement
Previous work
D-family matching: problem
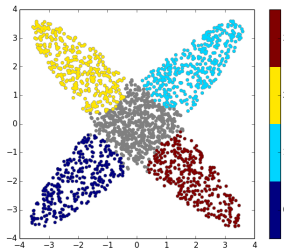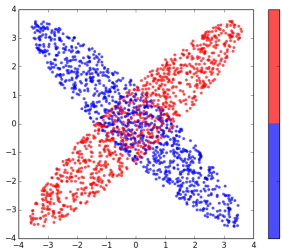Hardness
Algorithms
On the choice of D
Experiments

Maximum Information Coefficient

Beyond Two-sample-tests:
Localizing Data Discrepancies
in High-dimensional Spaces

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- ▶ `http://team.inria.fr/abs`
- ▶ `http://sbl.inria.fr`

Statistical analysis: complementary topics
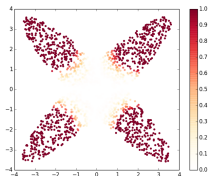
Localizing data discrepancies

▷ Problem: two populations differ in parameter/feature space: where are the differences?

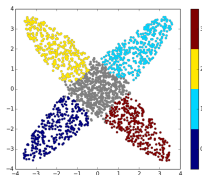▷ Contribution: density difference clustering based method
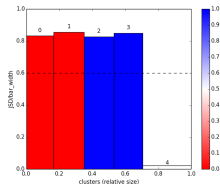
Given two point clouds,

we localize the discrepancy,

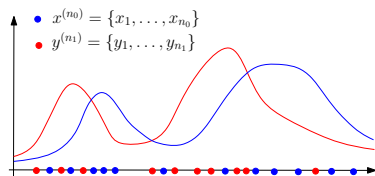to find spatially coherent regions of high discrepancy,

and provide a cluster based decomposed effect size.

Data discrepancies: two-sample problem and effect size

▷ **The two-sample test (TST) approach**

- ► Two datasets i.i.d. samples from two unknown densities f_X and f_Y:
  $x^{(n_0)} \equiv \{x_1, \ldots, x_{n_0}\}$ and $y^{(n_1)} \equiv \{y_1, \ldots, y_{n_1}\}$ in $\mathbb{R}^d$

- $x^{(n_1)} = \{x_1, \ldots, x_{n_0}\}$
- $y^{(n_1)} = \{y_1, \ldots, y_{n_1}\}$

$$\begin{cases} \mathrm{H}_0 : f_X = f_Y, \\ \mathrm{H}_1 : f_X \text{ and } f_Y \text{ differ in some way} \end{cases} \tag{1}$$

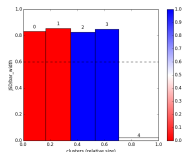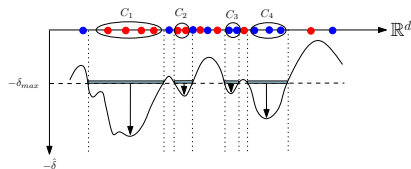▷ **Classical TST**

- ► p-value gives magnitude of the statistical significance, but
- ► (i) accept/reject: summarizes difference in a single bit
- ► (ii) the statistic of TST reflects the global discrepancy / effect size

▷ **Goal:** towards a nonparametric multivariate effect size

- ► (i') localize discrepancies accounting for the differences
- ► (ii') provide standardized (normalized) effect size

The three steps of the method

▷ **Step 1:** Estimate a measure of local discrepancy on each given point
Using $f \equiv (f_X + f_Y)/2$, define the Jensen-Shannon divergence:

$$JS(f_X \| f_Y) \equiv \frac{1}{2}\left(D_{\mathrm{KL}}(f_X \| f) + D_{\mathrm{KL}}(f_Y \| f)\right)$$

▷ **Step 2:** Aggregate local discrepancy in a spatial coherent way, using topological persistence analysis to spot stable features, and produce clusters by removing low discrepancy points

▷ **Step 3:** Produce an effect size bar plot to summarize the discrepancy profile

Statistical analysis: complementary topics

Pre-requisite: Jensen-Shannon divergence

▷ Kullback-Leibler divergence (KLD):

$$\begin{cases} D_{\mathrm{KL}}\left(f\|g\right) \equiv \int_{-\infty}^{\infty} f(x)\log\frac{f(x)}{g(x)}\,dx \\[2mm] D_{\mathrm{KL}}\left(P\|Q\right) \equiv \sum_{l\in\mathcal{A}} P(l)\log\frac{P(l)}{Q(l)} \end{cases}$$

▷ The Jensen-Shannon divergence (JSD): symmetrizes and smoothes the KLD:
Consider $f \equiv (f_X + f_Y)/2$, then

$$JS\left(f_X\|f_Y\right) \equiv \frac{1}{2}\left(D_{\mathrm{KL}}\left(f_X\|f\right) + D_{\mathrm{KL}}\left(f_Y\|f\right)\right)$$

▷ Main properties of JSD:

– JSD is symmetric
– JSD is bounded between 0 and 1
– Its square root yields a metric

▷Ref: Endres and Schindelin; IEEE Trans. Info. Theory, 2003

Step 1: Jensen-Shannon divergence and its decomposition

▷ Notations: two unknown densities f_X and f_Y, and the associated samples $x^{(n_0)}$ and $y^{(n_1)}$

▷ Two random variables are implicitly defined:
  – a position variable Z with density $f_Z \equiv f = (f_X + f_Y)/2$
  – a binary label $L \in \{0, 1\}$ with pmf $P(0) = 1/2$,
      indicating from which density (f_X or f_Y) an instance of Z is obtained.

▷ Equivalently, one defines the following pair of random variables:

$$(L, Z) = \begin{cases} (0, X) & \text{with prob. } \frac{1}{2} \\ (1, Y) & \text{with prob. } \frac{1}{2} \end{cases}$$

▷ Associated conditional and unconditional probability mass functions:

$$\begin{cases} P(l|z) = \mathbb{P}\left(L = l | Z = z\right) \\ P(l) = \mathbb{P}\left(L = l\right) = \frac{1}{2} \end{cases}$$

▷ Lemma: the JSD can be expressed as:

$$JS\left(f_X \| f_Y\right) = \int_{\mathbb{R}^d} f_Z(z) D_{\mathrm{KL}}\left(P(\cdot|z) \| P(\cdot)\right) dz$$

Step 1: the local discrepancy

▷ From

$$JS\left(f_X \| f_Y\right) = \int_{\mathbb{R}^d} f_Z(z) D_{\mathrm{KL}}\left(P(\cdot|z) \| P(\cdot)\right) dz$$

▷ We define the *discrepancy* at location z as

$$\delta(z) \equiv D_{\mathrm{KL}}\left(P(\cdot|z) \| P(\cdot)\right).$$

▷ Remarks:

    – $\delta(z) \in [0, 1]$ and $\delta(z) = 0 \Leftrightarrow f_X(z) = f_Y(z)$.

    – $P(I)$ is known but $P(I|z)$ is not:
        we need to estimate $P(I|z)$ at each given location z.

Step 1: random design nonparametric regression

▷ **Consider random variables:** location $Z \in \mathbb{R}^d$, and response variable $R \in \mathbb{R}$

▷ Associated regression function:

$$m(z) \equiv \mathbb{E}\left[R | Z = z\right].$$

▷ **Consider data:** $\{(Z_i, R_i)\}_{i=1,\ldots,n}$

▷ k_n-**nearest neighbors regressor:** upon sorting samples by increasing distance to z:

$$m_n(z) = \frac{1}{k_n} \sum_{i=1,\ldots,k_n} R_{(i,n)}(z)$$

▷ **NB:** $m_n(z)$ is a random variables: some convergence assessment is in order.

▷Ref: L. Györfi and A. Krzyzak; A distribution-free theory of nonparametric regression; 2002

Step 1: estimation via k-nearest neighbors

$\triangleright$ Using the labels as reponse variable $R \equiv L$

$\triangleright$ Estimate $P(\cdot|z)$ via random design nonparametric regression :
– build an estimator $m_n(z)$ using n i.i.d. realizations of (L, Z) for:

$$m(z) = \mathbb{E}\left[L|Z = z\right] = P(1|z).$$

– Then, if $0 \leq m_n(z) \leq 1$, we can use the following estimator for $P(l|z)$:

$$\hat{P}_n\left(l|z\right) \equiv |1 - l - m_n(z)|.$$

$\triangleright$ Thm: Using a k_n-nearest neighbors regressor, s.t. $\frac{k_n}{\log n} \to \infty$ and $\frac{k_n}{n} \to 0$:

$$\hat{\delta}_n(z) \equiv D_{\mathrm{KL}}\left(\hat{P}_n\left(\cdot|z\right)\|P(\cdot)\right) \xrightarrow{n\to\infty} \delta(z)\text{a.s.}$$

for f-almost all $z \in \mathbb{R}^d$.

The random multiplexer to obtain i.i.d. realizations of (L, Z)

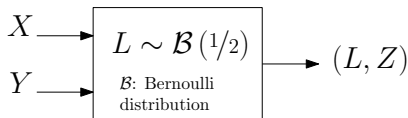▷ A random sampler produces i.i.d. realizations of (Z, L) from $x^{(n_0)}$ and $y^{(n_1)}$:

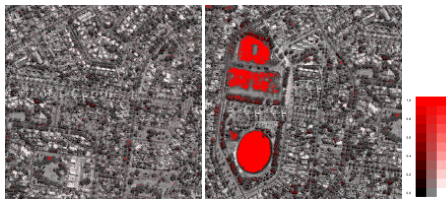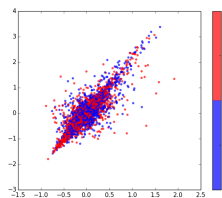Figure: **Random multiplexer generating pairs (label, position).**

▷ The case of populations of uneven sizes:

  – the multiplexer will consume faster the *small* population, and halt
  – unused samples of the large population: detrimental since information loss
  – resample B times and take the median of estimates, on a per sample basis
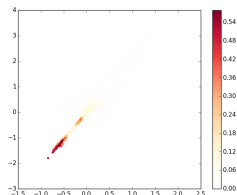
Step 1: Illustration: statistical image comparison

▷ **Images:** taking 2×2 blocks in each color channel (R,G,B) yields points in $\mathbb{R}^{12}$.

▷ **Interpolate** gray scale pixel color with red scale representing discrepancy at each pixel (upper left corner of the corresponding block) estimated with $k_n = n^{1/3}$

▷ **Multidimensional Scaling of parameter space:**

The two populations. . .

. . . colored with $\hat{\delta}$:

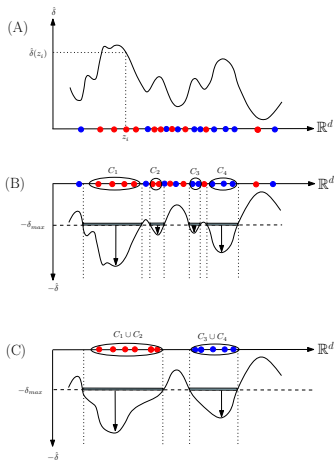# Statistical analysis: complementary topics

Step 2: Building the clusters from sublevel sets of $-\bar{\bar{\delta}}_z(z)$

▷ Ingredients:

- ▶ Height function / landscape: estimated discrepancy $\bar{\bar{\delta}}_z(z)$

- ▶ Parameter: significance threshold δ_{max}

▷ Construction:

- ▶ Idea: one cluster $\sim$ one connected component of the sublevel set of $-\bar{\bar{\delta}}_z(z)$ defined by δ_{max}

- ▶ Extra ingredient: smoothing the landscape to get rid of small clusters : smoothing using topological persistence at threshold ρ
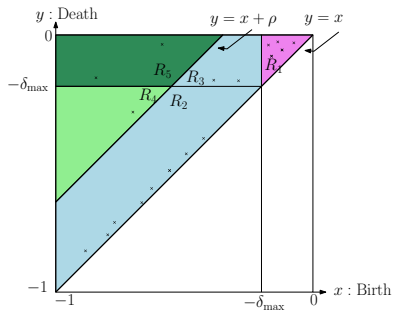

▷ NB: spurions samples removed from clusters due to filtering wrt δ_{max}.

Step 2: Building the clusters: persistence diagram

▷ Partition of the PD induced by:

- Significance threshold δ_{max}
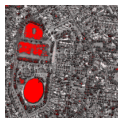- Persistence threshold ρ

▷ Local minimum m of $-\bar{\delta}_z(z)$:
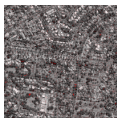
- Selected/rejected: m was born before $-\delta_{max}$.
- Persistent/canceled: persistence$(m) \geq \rho$
- Filtered (un-filtered): the catchment basin of m dies after (before) $-\delta_{max}$.

▷ Observation:
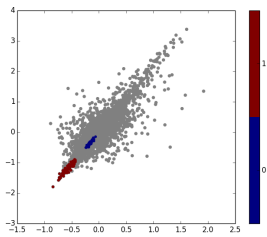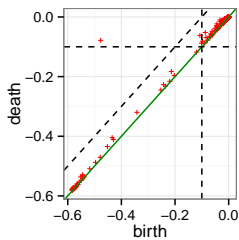
- # clusters : $1 +$ # points in region R_5 of the PD.
- # persistent local minima : 1 + num points in the region $R_4 \cup R_5$ of the PD.

Step 2: Illustration: statistical image comparison

▷ Images again:

▷ Parameters: $k = 10$ (NNG), $\rho = 0.1, \delta_{max} = 0.1$

Step 3: Effect size: discrepancy profile

▷ Global estimated JSD: area under dashed line
▷ Maximum JSD: area under continuous line ($=1$)
▷ Contribution of each cluster C to JSD: area of bar

$$JS_C\left(f_X\|f_Y\right) \equiv \frac{1}{n_0 + n_1} \sum_{z \in (x^{(n_0)} \cup y^{(n_1)}) \cap C} \hat{\delta}(z).$$

▷ Mass of each cluster: bar width
▷ Population balance in each cluster: bar color

▷ Ellipses:
 – Large global JSD (dashed line)
 – Contributed by $\mathbf{2+2}$ balanced clusters

▷ Images:
 – Smaller global JSD (dashed line)
 – Contributed by $\mathbf{2}$ clusters

Wrapping-up: workflow

▷ Compulsory parameters:
  k_n: regression parameter
  δ_{max}: discrepancy significance threshold
  ρ: persistence threshold
  k: num. of nearest neighbors for the persistence based clustering
▷ Optional parameter:
  B: num. repetition in case of unbalanced populations

Outlook: about regression

- k-NN based regressors: adapt to local intrinsic dimension: convergence results proved (L_2 sense) for marginals μ which are doubling measures.

- random projection tree based regressors: convergence results proved (L_2 sense) when $\mathcal{X}$ has Assouad dimension d. NB: more efficient than k-NN since cells of RPT have constant size.

- Open problem (AFAIK): strong pointwise consistency using RPTrees.

▷Ref: Kpotufe; k-NN regression adapts to local intrinsic dimension; NIPS 2011
▷Ref: Kpotufe and Dasgupta; A tree-based regressor that adapts to intrinsic dimension; J. of Computer and System Sciences, 2012

Outlook: general

- About p-values:
    - Use a classical test, possibly Maximum Mean Discrepancy (Gretton et al).
    - Also: the k-NN estimator used in a sequential way can be used to compute a p-value in a flexible way—the number of samples to process need not be known in advance.
- More applications:
    - Finding clusters with low discrepancy: study $\hat{\delta}$.
    - Goodness-of-fit analysis: sampling from a given model, then comparing data to spot discrepancies
- Feedback versus feature based selection: Compare to NIPS 2015 paper *Principal differences analysis*: feature based identification in the context of TST

Try me: http://sbl.inria.fr

Structural Bioinformatics Library
Template C++ / Python API for developping structural bioinformatics applications.

Structural Bioinformatic

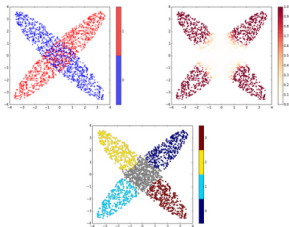## User Manual

Density_difference_based_clustering

Authors: *A. Lheritier and F. Cazals*

1. Goals

Comparing two sets of multivariate samples is a central problem in data analysis. From a statistical standpoint, one way to perform such a comparison is to resort to a non-parametric two-sample test (TST), which checks whether the two sets can be seen as i.i.d. samples of an identical unknown distribution (the null hypothesis, denoted H0).

Consistency of sequence regression estimates $\{m_n\}$ Based on L_2 norm

▷ Consider the following RV–induced by the data D_n:

$$\int \mid m_n(x) - m(x) \mid^2 \mu(dx). \tag{2}$$

▷ Def: The sequence $\{m_n\}$ is **weakly consistent** for a certain distribution of (X, Y) if

$$\lim_{n\to\infty} \mathsf{E}\Big[\int (m_n(x) - m(x))^2 \mu(dx)\Big] = 0. \tag{3}$$

▷ Def: The sequence $\{m_n\}$ is **strongly consistent** for a certain distribution of (X, Y) if

$$\lim_{n\to\infty} \int (m_n(x) - m(x))^2 \mu(dx) = 0 \text{ with proba. one.} \tag{4}$$

▷ Def: The sequence $\{m_n\}$ is **weakly universally consistent** if it is weakly consistent for all distributions of (X, Y) with $\mathbb{E}\left[Y^2\right] < \infty$.

▷ Def: The sequence $\{m_n\}$ is **strongly universally consistent** if it is strongly consistent for all distributions of (X, Y) with $\mathbb{E}\left[Y^2\right] < \infty$.

▷Ref: book
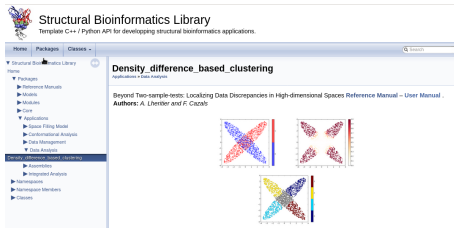
Consistency of sequence regression estimates $\{m_n\}$ based on pointwise convergence

▷ Def: The sequence $\{m_n(x)\}$ is called **strongly pointwise consistent** is $m_n(x) \to m(x)$ a.s.

▷ Def: The sequence is called **strongly universal pointwise consistent** if it is strongly pointwise consistent for all distributions of (X, Y) with $\mathbb{E}\left[Y^2\right] < \infty$.

▷Ref: book

Structural Bioinformatics Library

Package   Density Difference Based Clustering @ `http://sbl.inria.fr`

- ▶ User manual `https://sbl.inria.fr/doc/Density_difference_based_clustering-user-manual.html`

- ▶ General entry: `http://sbl.inria.fr`

▷Ref: Cazals and A. Lhéritier, IEEE/ACM DSAA, 2015
▷Ref: Kim, Lee, Lei, Electronic Journal of Statistic, 2019

Statistical analysis: complementary topics

Comparing two clusterings using matchings between clusters of clusters

F. Cazals, D. Mazauric, R. Tetley, and R. Watrigant
ACM Trans. Exp. Algorithms, 2019
`https://sbl.inria.fr/doc/D_family_matching-user-manual.html`

Statistical analysis: complementary topics

Clustering algorithms

- Many algorithms: which one?
- Many parameters: which ones?
- Many clustering: are they consistent? A problem of scale...

Statistical analysis: complementary topics

Grouping clusters into **metaclusters:**
problem formalization in terms of intersection graph

▷ Goal: recovering some coherence between groups of clusters

  ▶ as a function of a scale parameter D

▷ Rationale: many-to-many

  ▶ Aggregating many clusters, map to many clusters

  ▶ Characterize the *scale* at which clusters merge

Structurally conserved motifs in protein structures
Many-to-many correspondence between clusters

▷ Handling small and conserved structural motifs in proteins

Merging clusters: a matter of scale

On the role of the scale parameter D

(A) Two clusterings (kmeans++, Tomato, etc)   (B) Meta-clusters as union of clusters

Statistical analysis: complementary topics

Comparing clusterings: previous work

REVISE REVISE REVISE REVISE

▷ 1-1 mapping of clusters: equivalent to the problem of computing a maximum weighted matching in weighted bipartite graph.

▷ Solution: solved in $O(n^2 \log n + nm)$

▷ Particular case of the D-family-matching problem for $D = 1$ – see later

Comparing clusterings: the Variation of Information

- A set Z of t items
- A clustering F of size r for Z: $F = \{F_1, \ldots, F_r\}$; $n_k = |F_k|$; $p_k = n_k/t$.
- A clustering F of size r' for Z: $F = \{F_1, \ldots, F_r\}$; $n'_k = |F'_{k'}|$;
- Overlap between two clusters: $p(k, k') = |F_k \cap F'_{k'}|/t$.
- Entropy of clustering: $H(F) = -\sum_{k=1,\ldots,r} p(k) \ln p(k)$

- Mutual information between F and F':

$$I(F, F') = \sum_k \sum_{k'} p(k, k') \ln \frac{p(k, k')}{p(k)p(k')}.$$

- Variation of information (VI):

$$VI(F, F') = H(F) + H(F') - 2I(F, F').$$

- Main properties:
  - ▶ VI is a metric
  - ▶ $VI(F, F') \leq \ln t$
  - ▷Ref: M. Meila, Journal of Multivariate Analysis, 2007

Statistical analysis: complementary topics

Intersection graph

- Data: $Z = \{z_1, \ldots, z_t\}$
- Clustering F of size r: $F = \{F_1, \ldots, F_r\}$

  $$F_i \subseteq Z, F_i \neq \emptyset \text{ and } F_i \cap F_j = \emptyset \text{ for every } i, j \in \{1, \ldots, r\}, i \neq j.$$

- Clustering F' of size r': $F' = \{F'_1, \ldots, F'_{r'}\}$

  $$F'_i \subseteq Z, F'_i \neq \emptyset, \text{ and } F'_i \cap F'_j = \emptyset \text{ for every } i, j \in \{1, \ldots, r'\}, i \neq j.$$

NB: a clustering may not contain all t items

Definition 1 (Intersection graph $G = (U, U', E, w)$ for F and F').

The set $U = \{u_1, \ldots, u_r\}$: vertices of F
The set $U' = \{u'_1, \ldots, u'_{r'}\}$: vertices of F'
Edges $E = \{\{u_i, u'_j\} \mid F_i \cap F'_j \neq \emptyset, 1 \leq i \leq r, 1 \leq j \leq r'\}$.
Edge weight of edge $e = \{u_i, u'_j\} \in E$ is $w_e = |F_i \cap F'_j|$.

D-family matching

a constraint on the diameter of certain subgraph of the intersection graph

Definition 2. [D-family-matching for an intersection graph]

A family $\mathcal{S} = \{S_1, \ldots, S_k\}$, $k \geq 1$, such that

- for every $i, j \in \{1, \ldots, k\}$, if $i \neq j$, then: $S_i \subseteq V$, $S_i \neq \emptyset$, $S_i \cap S_j = \emptyset$,
- and the graph $G[S_i]$ induced by the set of nodes S_i has diameter at most D.

▷ Comments:

- $D = 1$: matching
- $D = 2$: clusters as stars

▷ Notations:

- Set of all D-family matchings of a graph G: $\mathcal{S}_D(G)$

D-family matching problem

$$\Phi(\mathcal{S}) = \sum_{i=1}^{k} \sum_{e \in E(G[S_i])} w_e. \tag{5}$$

▷ Remarks:

- ▶ The sum runs over **all** edges of a connected component. (Later: see algorithms based on spanning trees.)
- ▶ We wish to compute a D-family-matching which minimizes the inconsistencies.

Definition 3 (D-family-matching problem). Let $D \in \mathbb{N}^+$. Given an intersection graph G, the D-family-matching problem consists in computing

$$(\text{Opt score for a given } D) \quad \Phi_D(G) = \max_{\mathcal{S} \in \mathcal{S}_D(G)} \Phi(\mathcal{S}). \tag{6}$$

NB: Score with the diameter D stressed: $\Phi(\mathcal{S}^{D=d})$

D-family matching: role of the diameter, illustration

Figure: **Simple instance of the D-family-matching problem and solutions: panels (c,d,e,f) represent optimal solutions for different values of D. (a)** Simple instance of the D-family-matching problem with $t = 12$, $r = 5$, $r' = 4$, and so $n = 9$. The family F contains five sets and the family F' contains four sets. **(b)** Intersection graph G. **(c)** Optimal solution $\mathcal{S}$ for $D \geq 7$ with $\Phi(\mathcal{S}) = \Phi_D(G) = 12$. **(d)** Optimal solution $\mathcal{S}$ for $D = 3$ with $\Phi(\mathcal{S}) = \Phi_3(G) = 11$. **(e)** Optimal solution $\mathcal{S}$ for $D = 2$ with $\Phi(\mathcal{S}) = \Phi_2(G) = 9$. **(f)** Optimal solution $\mathcal{S}$ for $D = 1$ with $\Phi(\mathcal{S}) = \Phi_1(G) = 8$.

Notations, recap

Notation	Definition
$Z = \{z_1, \ldots, z_t\}$	Set of $t \geq 1$ elements
$F = \{F_1, \ldots, F_r\}$	Family of $r \geq 1$ disjoint subsets of Z
$F' = \{F'_1, \ldots, F'_{r'}\}$	Family of $r' \geq 1$ disjoint subsets of Z
$G = (V, E, w)$	Intersection graph of $n \geq 1$ nodes and $m \geq 1$ edges
$N_G(v) = \{v' \mid \{v, v'\} \in E\}$	Set of neighbors of node $v \in V$
$\Delta = \max_{v \in V} |N_G(v)|$	Maximum degree of G
$cc(G)$	Set of maximal connected components of G
$\mathcal{S} = \{S_1, \ldots, S_k\}$	D-family-matching
$\Phi(\mathcal{S}) = \sum\limits_{i=1}^{k} \sum\limits_{e \in E(G[S_i])} w_e$	Score of a D-family-matching $\mathcal{S}$
$\mathcal{S}_D(G)$	Set of all D-family-matching for G
$\Phi_D(G) = \max_{\mathcal{S} \in \mathcal{S}(G,D)} \Phi(\mathcal{S})$	Optimal score for the D-family-matching problem
$\mathcal{S}_D(G, T_r)$	Set of all D-family-matching constrained by T_r
$\Phi_D(G, T_r) = \max_{\mathcal{S} \in \mathcal{S}_D(G, T_r)} \Phi(\mathcal{S})$	Optimal score for the D-family-matching problem constrained by T_r

Statistical analysis: complementary topics

Main result

Theorem 4. Let $D \geq 2$ be any integer. The decision version of the D-family-matching problem is NP-complete for :

- bipartite graphs of maximum degree 3;
- bipartite graphs of maximum degree 4 even if the maximum weight is constant.

Moreover, the 2-family-matching problem is *APX*-hard for bipartite graphs of maximum degree 3 with unary weights.

▷ Open pb.: Is the D-family-matching problem in APX or not (constant factor approximation)?

- Nb: P $\neq$ NP: APX-hard pb. not in PTAS, i.e. no $(1 + \varepsilon)$ approx

Greedy strategy on the diameter is not an option

Lemma 5. For any integer $n \geq 1$, then there exists an intersection graph $G = (V, E, w)$ composed of n nodes such that $\Phi_2(G)/\Phi_1(G) \geq n - 1$.

▷ One has:

- ▶ $\Phi(\mathcal{S}^{D=1}) = 1$ (one edge)
- ▶ $\Phi(\mathcal{S}^{D=2}) = t = n - 1$ (all edges)

Statistical analysis: complementary topics

Trees: theorems

Theorem 6 (Computation of $\Phi_D(G)$ for trees). Let $D \in \mathbb{N}^+$. Consider any intersection tree $T = (V, E, w)$ of maximum degree $\Delta \geq 0$. Then, there exists an $O(D^2 \Delta^2 n)$-time complexity algorithm for the D-family-matching problem for T.

Proof.
See black board. $\qquad\qquad\qquad\qquad\qquad\qquad\qquad\qquad\qquad\qquad\qquad\qquad \square$

Theorem 7. For any $D \in \mathbb{N}^+$, the D-family-matching problem can be solved:

- in $O(Dn)$ time if G is a path;
- in $O(D^2 n)$ time if G is a cycle(s) or a graph of maximum degree 2.

Proof.
See paper. $\qquad\qquad\qquad\qquad\qquad\qquad\qquad\qquad\qquad\qquad\qquad\qquad\qquad\quad \square$

Generic approach on spanning trees

Definition 8 (D-family-matching constrained by a tree). Let $G = (V, E, w)$ be an intersection graph and T be a spanning tree of G. A D-family-matching for G *constrained by* T is a D-family-matching $\mathcal{S}$ for G such that all $S_i \in \mathcal{S}$ induces a connected subtree in T. The set of all D-family-matching constrained by T is denoted $\mathcal{S}_D(G, T)$.

With this Def., we obtain the following sub-problem of D-family-matching:

Definition 9 (D-family-matching problem constrained by a tree). The D-family-matching problem consists in computing

$$\Phi_D(G, T) = \max_{\mathcal{S} \in \mathcal{S}_D(G, T)} \Phi(\mathcal{S}) \tag{7}$$

Generic algorithm for the D-family-matching problem

▷ **Three ingredients:**

- ▶ **A property** $\Pi(\mathcal{M})$, depending on the set $\mathcal{M}$ of already computed D-family-matchings, represents the halting condition of the algorithm.

- ▶ **A spanning tree generator** $\mathcal{R}(G, \lambda)$ computes the rooted spanning tree T^λ of G that is used at step $\lambda \geq 1$ by Algorithm $\mathcal{A}$.

- ▶ **An algorithm** $\mathcal{A}(G, T^\lambda, D)$ computes a D-family-matching $\mathcal{S}^\lambda$ constrained by T^λ.

▷ Generic algorithm for the D-family-matching problem:

Require: An intersection graph $G = (V, E, w)$, an integer $D \geq 1$, a property Π, a spanning tree generator $\mathcal{R}$, and an algorithm $\mathcal{A}$.

1: $\mathcal{M} := \emptyset$, $\lambda := 0$
2: **while** $\neg \, \Pi(\mathcal{M})$ **do**
3:    $\lambda := \lambda + 1$; Compute the spanning tree $T^\lambda := \mathcal{R}(G, \lambda)$
4:    Compute $\mathcal{S}^\lambda$ by using Algorithm $\mathcal{A}(G, T^\lambda, D)$; $\mathcal{M} := \mathcal{M} \cup \mathcal{S}^\lambda$
5: **return** $\mathcal{S} \in \mathcal{M}$ of maximum score

Results on spanning trees

Lemma 10. Let $D \in \mathbb{N}^+$. Let G be any intersection graph. Then, there exists a rooted spanning tree T of G such that $\Phi_D(G) = \Phi_D(G, T)$.

Proof.
See black board. $\qquad\qquad\qquad\qquad\qquad\qquad\qquad\qquad\qquad\qquad\qquad$ $\square$

Lemma 11 (Computation of $\Phi_D(G, T)$). Let $D \in \mathbb{N}^+$. Let $G = (V, E, w)$ be any intersection graph and T be any spanning tree of G. Then, there exists a $O(2^{D\Delta \log_2(\Delta)} n)$-time algorithm for the D-family-matching problem for G constrained by T.

Proof.
See paper. $\qquad\qquad\qquad\qquad\qquad\qquad\qquad\qquad\qquad\qquad\qquad\qquad\qquad$ $\square$

Statistical analysis: complementary topics

Various strategies

Three strategies:

- ▶ (Stable plateaus) compute a set of non-overlapping plateaus optimizing a functional favoring *long* and *thin* plateaus.

- ▶ (Prescribed num. plateaus) specify the number of plateaus to be obtained.

- ▶ (Hierarchical plateaus) perform a hierarchical decomposition into plateaus, which is of interest if there are several *vertical scales*.

▷ Focus on:

- ▶ local maxima of Φ

- ▶ plateaus

Stable plateaus: long plateaus of small height

▷ Quality measures for a plateau:

- ► $\tau_w(y)$: positive increasing function for the plateau width
- ► $\tau_h(y)$: positive increasing function for the plateau height

▷ Given:

- ► Let D_G be the diameter of the intersection graph; we assume $\{(D, \Phi_D)\}_{D=1,\dots,D_G}$
- ► Consider the set $\{\Phi_1(G), \dots, \Phi_{D_G}(G)\}$
- ► Let $|I_x|$ is the size of plateau I_x

Definition 12. Determine $\mu \in \{1, \dots, D_G\}$ plateaus (intervals) $I_1, \dots, I_\mu$ of $[1, D_G]$ with

- ► $I_1 \cup \dots \cup I_\mu = \{1, \dots, D_G\}$, $I_x \cap I_{x'} = \emptyset$ for every $1 \le x < x' \le \mu$,
- ► such that the following function is minimum:

$$-\sum_{x=1}^{\mu} \frac{\tau_w(|I_x|)}{\tau_h(\max_{D,D' \in I_x \cap \mathbb{N}} \Phi_{D'}(G) - \Phi_D(G))}$$

Stable plateaus: construction

Theorem 13. There is an $O(D_G^2)$-time complexity algorithm that computes an optimal solution for the Tradeoff-plateau problem.

▷ Algorithm: blackboard

Hierarchical plateaus

▷ Dendogram of plateaus:

► For two consecutive plateaus, each consisting of a set of values $\{(D, \Phi_D)\}$: *coherence measure* for the union of these two plateaus: the maximum difference between any two values Φ. on these plateaus.

► Merge two plateaus realizing the minimum value then yields a dendogram.

▷ Formally: build a rooted tree $T = (V, E)$ representing the hierarchical plateaus

► One leaf per possible value of D; $D_G - 1$ internal nodes (including the root). That is, let $(l_1, l_2, \ldots, l_d)$ be the $d = D_G$ initial plateaus each composed of 1 point.

► Perform the aforementioned binary merge.

Statistical analysis: complementary topics

Generic code and instantiation for experiments

▷ Implementation in the SBL:
http://sbl.inria.fr/doc/D_family_matching-user-manual.html.

▷ Implementation $STS(G, D)$ has the following ingredients:

- **(i)** the spanning tree generator $\mathcal{R}$ returns a *maximum spaning tree*, or a *random spanning tree*;

- **(ii)** the property $\Pi(\mathcal{M})$ returns true once we have computed a solution on the maximum spanning tree, as well as a solution on $n_i = (10,000)$ distinct random spanning trees (for a given n_i);

- **(iii)** $\mathcal{A}$: algorithm as in Theorem 6 with an additional step: edges for which both extremities belong to the same meta-cluster are added to the said meta-cluster. (In general, the intersection graph is indeed not a tree, so that such edges were unaccounted for.)

- The solution returned for a given graph G and a diameter D is the best yielded by the aforementioned $1 + n_i$ spanning trees.

Randomly edited clusterings: setup

▷ **Initial random clusterings:**

- $(t = 1\ 000, r = 20)$ and $(t = 3\ 000, r = 50)$.
- Generated with the Boltzmann sampler from Flajolet - Duchon et al
- Due to the randomness, the process is repeated $N_r = 10$ times for each pair (t, r).

▷ **Edited clusterings:** a copy F' of a clustering F is edited in two steps

- Union operations: e unions reduce the number of clusters to $r - e$
- Jittering: for each cluster, a fraction τ of its items are distributed amongst the remaining $k - 1$ clusters uniformly at random.

▷ **Values:** 9 scenarios for edits and jitters

- $e \in \{0, \lfloor r/4 \rfloor, \lfloor r/2 \rfloor\}$ and $\tau \in \{0.05, 0.1, 0.2\}$. (NB: for $e = 0$, F' is a jittered version of F (i.e. the numbers of clusters are identical.)
- yields $N_r \times \#(t, r) \times \#e \times \#\tau = 180$ comparisons, which are ascribed to 9 scenarii (3 values for $e \times 3$ values for τ) denoted $EeJy$, where $y = 100\tau$.

▷ **Comparson against VI:** comparison of normalized scores $\in [0, 1]$:

$$s_\Phi = 1 - \Phi_D(\cdot)/t \text{ versus } s_{VI} = VI/\log t.$$

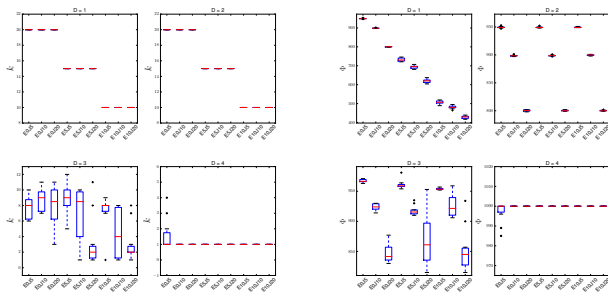# Randomly edited clusterings: results for ($t = 1000, r = 20$)

Figure: Algorithm $STS(G, D)$ for clusterings with ($t = 1\,000, r = 20$).
(Left) Best value for k as a function of the 9 scenarii. **(Right)** Scores s_Φ as a function of the 9 scenarii.

- $D \leq 2$: algo. finds the right number of clusters $\forall e$ (resp: 20, 15, 10)
- For $D = 2$: score $\Phi_D(\cdot)$ is almost perfect (≥ 800, wrt $t = 1000$)
- Across scenarii: scores hardly depend on the jitter level
- For $D = 3$: scores $\Phi_D(\cdot)$ varies significantly–but medians ok
- For $D = 4$: the algorithms output a full graph

Comparison with the Variation of Information: results

▷ Method: scatter plot of $s_\Phi = 1 - \Phi_D(\cdot)/t$ versus $s_{VI} = VI/\log t$

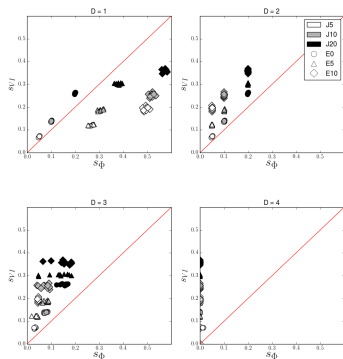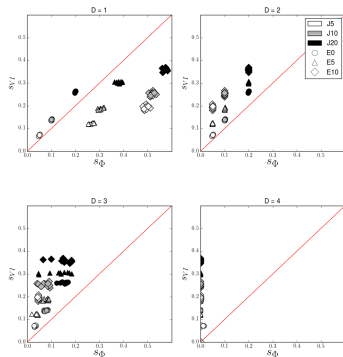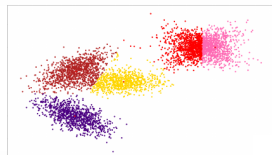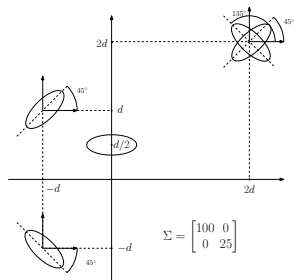NB: 1 symbol per scenario, one number of a symbol, number of repeats.

Figure: **Normalized score s_{VI} versus normalized score s_Φ of algorithm $STS(G, D)$.** Each marker is a different union scenario and each color represents a different jitter scenario following the legend on the upper right. We plot the $y = x$ function for reference.

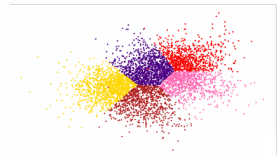# Comparison with the Variation of Information: results

- $D = 2$: s_Φ corresponds to a matching.
- $D = 2$, two key differences with VI: $s_\Phi \leq s_{VI}$; s_Φ constant against union operations. Both s_{VI} and s_Φ are affected by jittering.
- $D = 3$: higher variability in s_Φ; dependence on jittering and $\#$ union operations.
- For $D = 4$: $s_\Phi = 0$ ie the full intersection graph reported.

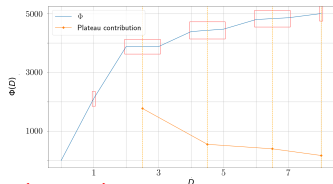# On the separability of clusters and D: setup

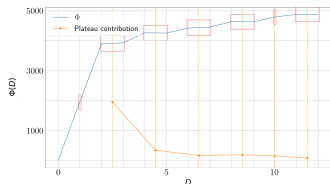Figure: **Parameterized dataset: mixture of 5 Gaussian blobs**. **(A)** Relative position of the five Gaussian blobs: function of d **(B, C, D)** $t = 5,000$, $d = 50, 20, 5$. Samples clustered with k-means++ ($k = 5$).
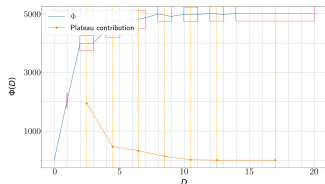
On the separability of clusters and D: plateaus

▷ **(A, d=50)**
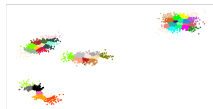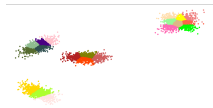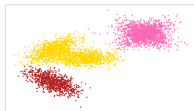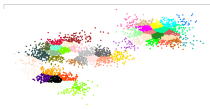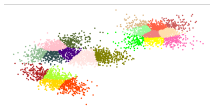

▷ **(C, d=5)**

▷ **(B, d=20)**

▶ **(A)** $d = 50$, $k = 4$ meta-clusters suggested for $D = 8$.

▶ **(B)** $d = 20$, $k = 3$ meta-clusters suggested for $D = 8$.

▶ **(C)** $d = 5$ No obvious choice for the number of meta-clusters.

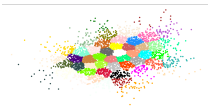# On the separability of clusters and D: illustrations

$d = 50$: $D = 8$, 4 m.c.

$d = 20$: $D = 8$, 3 m.c.

$d = 5$: no real hint

Final words on the choice of D

Fom Strehl et al (JMLR 2002): "*In fact, the* `right` *number of clusters in a dataset often depends on the* `scale` *at which the dataset is inspected*".

- Parameter D acts as a scale parameter providing information of the structure of the intersection graph.

- When this graph is dense or has a specific topology (star-shaped), trivial values of Φ are obtained for small values of D, and a unit change of D may trigger an abrupt change of Φ. However, in more complex situations, large values of D may be required.

- As a general strategy to choose D, we suggest identifying drops in Φ when decreasing D. Indeed, for any range of D corresponding to a *plateau* for Φ, the most significant value for D is the smallest one.

Outlook

- Interesting complexity issues: open
- Useful tool, available from
  https://sbl.inria.fr/doc/D_family_matching-user-manual.html
- Interesting connexions with model clustering in deep learning – amongst others

▷Ref: Cazals et al, ACM J. of Experimental Algorithms, 2019
▷Ref: Interactive Naming for Explaining Deep Neural Networks: A
Formative Study M Hamidi-Haines, Z Qi, A Fern, F Li arXiv preprint
arXiv, 2018.

Statistical analysis: complementary topics
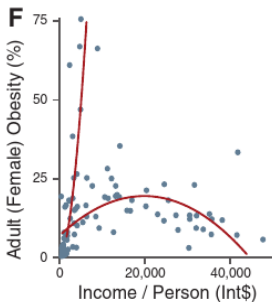
Detecting Novel Associations in Large Data Sets

The Maximal Information Coefficient (MIC)

Science 334, December 2011

9 authors, led by Michael Mitzenmacher and Pardis Sabeti, Harvard

Relationship Type	MIC	Pearson	Spearman	Mutual Information (KDE)	(Kraskov)	CorGC (Principal Curve-Based)	Maximal Correlation
Random	0.18	-0.02	-0.02	0.01	0.03	0.19	0.01
Linear	1.00	1.00	1.00	5.03	3.89	1.00	1.00
Cubic	1.00	0.61	0.69	3.09	3.12	0.98	1.00
Exponential	1.00	0.70	1.00	2.09	3.62	0.94	1.00
Sinusoidal (Fourier frequency)	1.00	-0.09	-0.09	0.01	-0.11	0.36	0.64
Categorical	1.00	0.53	0.49	2.22	1.65	1.00	1.00
Periodic/Linear	1.00	0.33	0.31	0.69	0.45	0.49	0.91
Parabolic	1.00	-0.01	-0.01	3.33	3.15	1.00	1.00
Sinusoidal (non-Fourier frequency)	1.00	0.00	0.00	0.01	0.20	0.40	0.80
Sinusoidal (varying frequency)	1.00	-0.11	-0.11	0.02	0.06	0.38	0.76

Correlations in 2D: the Pearson correlation coefficient

Does the knowledge of X provide information on Y?

▷ The Pearson coeff.: $\rho = cov(X,Y)/(\sigma_x\sigma_Y)$

▷ Anscombe's quartet: $\rho = 0.816$

▷ Properties of ρ:
– Coupled to linear regression $f_i = \alpha x_i + \beta$
  – $\alpha = \rho\sigma_Y/\sigma_X$
  – coeff. of determination: $R^2 = \rho^2$
– Not invariant to rotations
– Spearman's coeff: Pearson on ranks:
  for monotonic correlations

▷ Coeff of determination
R^2: variance explained by the model
$1 - R^2$: unexplained var. / noise level

Information Theory: Key Quantities

▷ Entropy of the r.v. X

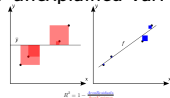$$H(X) = -\sum_{x \in \mathcal{X}} p(x) \log p(x)$$

$$H(X) \leq \log |\mathcal{X}|$$

▷ Joint and Conditional entropies

$$H(X, Y) = -\sum p(x, y) \log p(x, y)$$

$$H(Y \mid X) = \sum_{x \in \mathcal{X}} p(x) H(Y \mid X = x)$$

▷ Relative entropy: Kullback-Leibler divergence of two distributions on the same proba. space:

Def:
$$D(P, Q) = \sum_{x \in \mathcal{X}} p(x) \log(p(x)/q(x))$$

Prop.: $D(P, Q) \geq 0$

▷ Mutual information

$$I(X, Y) = \sum p(x, y) \log \frac{p(x, y)}{p(x)p(y)}$$

$$\begin{cases} I(X, Y) & = D(p(x, y), p(x)p(y)) \\ & = H(X) - H(X \mid Y) \\ & = H(X) + H(Y) - H(X, Y) \end{cases}$$
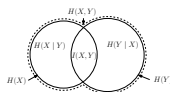

Notes: (i) Poincare formula for I
(ii) I as correlation: common entropy

http://en.wikipedia.org/wiki/Mutual_information
http://en.wikipedia.org/wiki/Kullback-Leibler_divergence

Maximal Information Coefficient (MIC): Definition

▷ **Input:** a 2D point cloud $D = \{(x_i, y_i)\}_{i=1,\ldots,n}$ and its bounding box
▷ **Grids:** $G_{x,y}$: grids of size $x \times y$ not necessarily regular
▷ **Joint proba/marginal of $D_{|G}$** : fraction of samples, out of n, in a cell/row/column

▷ **Def of MIC:**

$$I^*(D, x, y) = \max_{G \in G_{x,y}} I(D_{|G}) \qquad (8)$$

$$M_{xy} = \frac{I^*(D, x, y)}{\log \min(x, y)} \qquad (9)$$

$$MIC = \max_{xy < B(n) = n^{1-\epsilon}} M_{xy} \qquad (10)$$

▷ **Elementary properties:**

    – $M_{xy} \in [0, 1]$
    – $MIC(X, Y) = MIC(Y, X)$
    – MIC invariant to order preserv. transf.
       grids determined by abscissa / ordinates
    – MIC not invariant to rotations
       cf $y = x$ vs $y = c$

▷ **Note:** exploring all grids $\sim$ enclosing the data in a *tube*

Note: For the normalization of Eq. (9):
log min(x, y) rather than n: # cells
sub-linear, see Eq. (10).

MIC, illustrations (I): the functional noiseless case

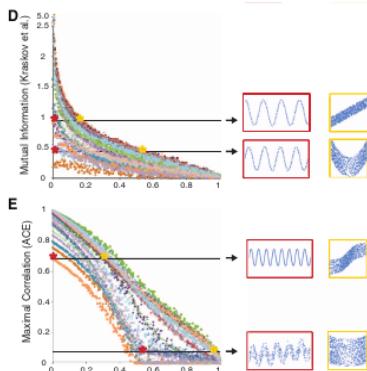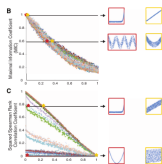▷ Ideal scores: (almost) one

A

Relationship Type	MIC	Pearson	Spearman	Mutual Information (KDE)	Mutual Information (Kraskov)	CorGC (Principal Curve-Based)	Maximal Correlation
Random	0.18	-0.02	-0.02	0.01	0.03	0.19	0.01
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Parabolic	1.00	-0.01	-0.01	3.33	3.15	1.00	1.00
Sinusoidal (non-Fourier frequency)	1.00	0.00	0.00	0.01	0.20	0.40	0.80
Sinusoidal (varying frequency)	1.00	-0.11	-0.11	0.02	0.06	0.38	0.76

MIC, illustrations (II): the functional noisy case
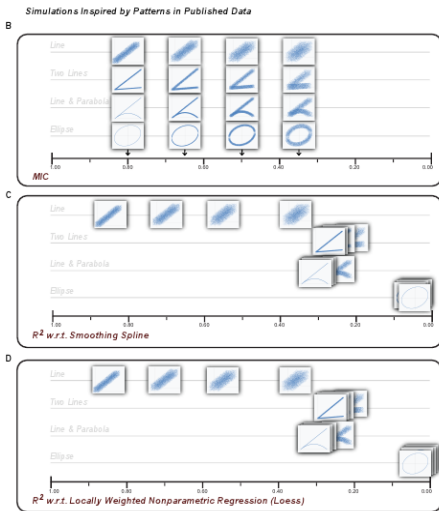
▷ Testing 27 functions with uniform vertical noise: *MIC = function of* $(1 - R^2)$
with R^2 the determination coeff of the data relative to the noiseless function

Bottomline is $\rightarrow MIC \sim R^2$: easy comparison of $\neq$ functions

MIC, illustrations (III): the non functional noisy case

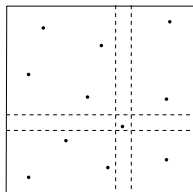▷ MIC also degrades *smoothly* as a function of the noise level

Details of the Definition

▷ Grid resolution B(n)
  – too low: searching for simple patterns
  – too high: high scores even for random data
    samples isolated in cell/column
▷ Normalization
  – grids with $\neq$ dimensions have $\neq$ mutual information
  – normalizing by $\log \min(x, y)$:
    Comparing grids of $\neq$ dimensions
    Ensures that almost all
      noiseless functions get MIC of one
      (finite union of differentiable curves)

MIC: Theorems

(Ten pages of proofs in the Supplemental)

▷ **Thm 1.** If X and Y independent R.V.: ApproxMIC converges to 0 in probability when $n \to \infty$
If X and Y are not independent R.V.: MIC bounded away from 0 almost surely.

▷ **Thm 2.** For any joint distribution (X, Y), MIC computed with a number of cells $B(n) = n^{1+\varepsilon}$ would yield $MIC \to 1$ almost surely.

▷ **Thm 3.** Let D consist of n samples drawn according to a distribution $(X, f(X))$, with f nowhere constant on $[0, 1]$. Then $MIC \to 1$ almost surely.

▷ **Thm 4.** If the support of $(x(t), y(t))$ is a finite union of smooth curves, nowhere flat (critical points of measure 0), then $MIC > 1 - \varepsilon$ for large n.

▷ **Thm 5.** MIC of a noisy functional $(X, f(X) + E_h)$, with E_h uniform noise in $[-h, h]$, is lower bounded by a (complex) functional of the R^2 between $f(X)$ and $f(X) + E_h$.

More Ingredients

- Computing MIC: algorithm ApproxMIC uses 2D dynamic programming

- p-value calculation for H0: *X and Y are statistically independent*
  Create surrogate datasets created with random permutations
  (eg of *X* for *Y* fixed)

- *MIC* $- \rho^2$ as a natural measure of linear dependence:
  Since MIC behaves as R^2 for functional relationships

- Symmetry of the matrix M_{xy}: hints at monotony
  Maximum Asymmetry Score $\mid m_{xy} - m_{yx} \mid$
  Hints at periodic relationships with non constant period

- Software MINE: `http://www.exploredata.net/`