Hierarchical Components for the GRID

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0. GRIDs
1. *ProActive:* Asynchronous Distributed Objects
2. Groups
3. Components

ProActive:
A Java API + Tools for Parallel, Distributed Computing

- A uniform framework: An Active Object pattern
- A formal model behind: Prop. Determinism, insensitivity to deploy.

Programming Model:

- **Remote Objects (Classes, not only Interfaces, Dynamic)**
- **Asynchronous Communications, Futures, Wait-By-Necessity**
- **Groups, Mobility, Components, Security, Fault Tolerance: Checkpoints**

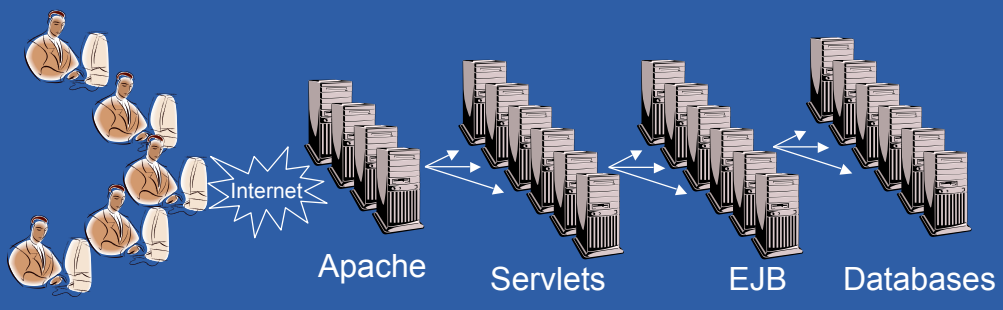
Environment:

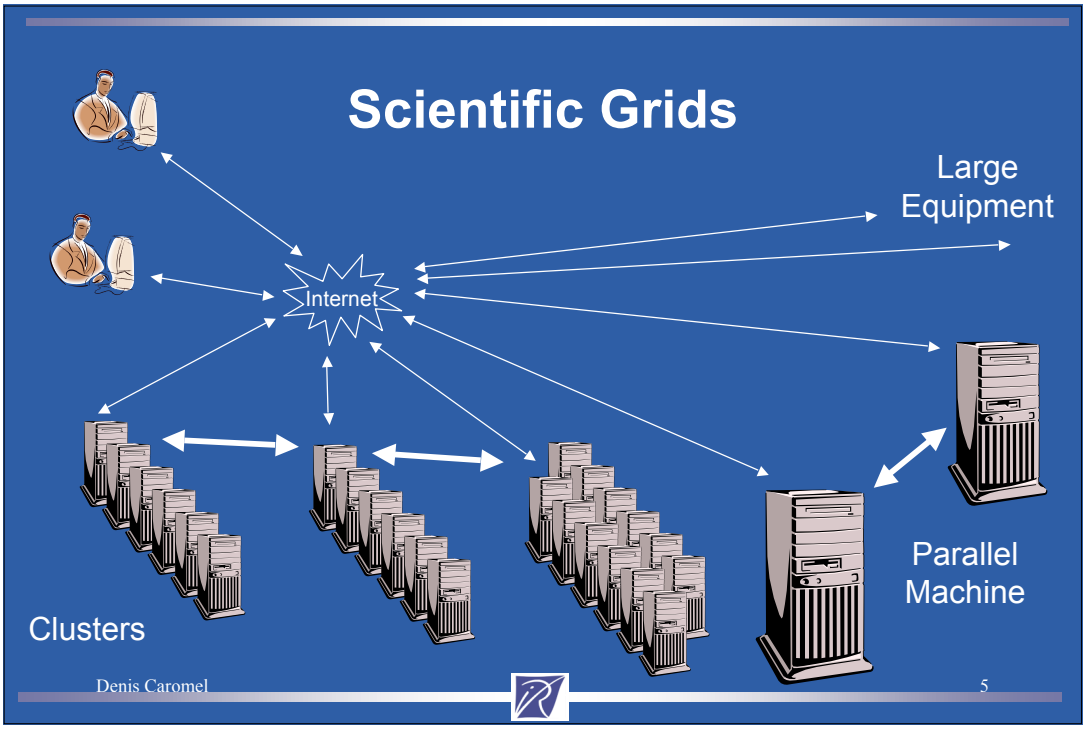
- **XML Deployment Descriptors,** `Web Service Export., HTTP, ssh Tunneling`
- **Various protocols:** `rsh,ssh,LSF,PBS,Globus,Sun Grid Engine, sshGSI`
- **Visualization and monitoring:** IC2D

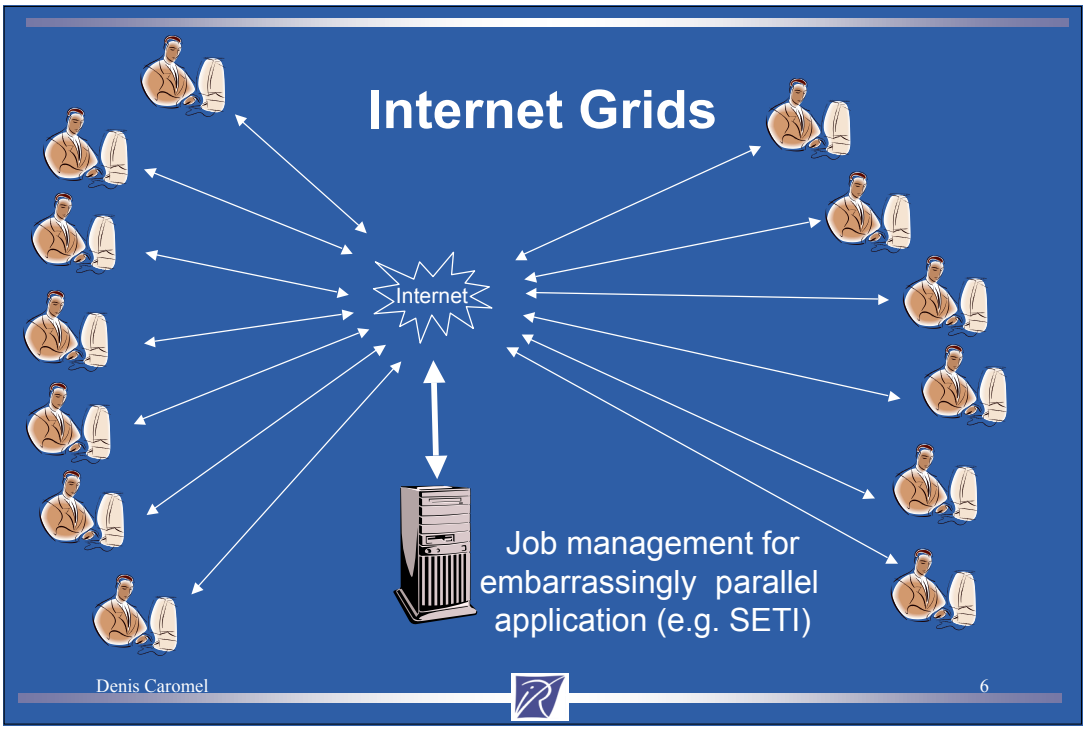
In the www. ObjectWeb .org Consortium (Open Source middleware)
since April 2002 (LGPL license)

GRIDs

Enterprise Grids




Scientific Grids
Large
Equipment
Internet
Clusters
Parallel
Machine
Denis Caromel
5


Internet Grids
Internet
Job management for
embarrassingly parallel
application (e.g. SETI)
Denis Caromel
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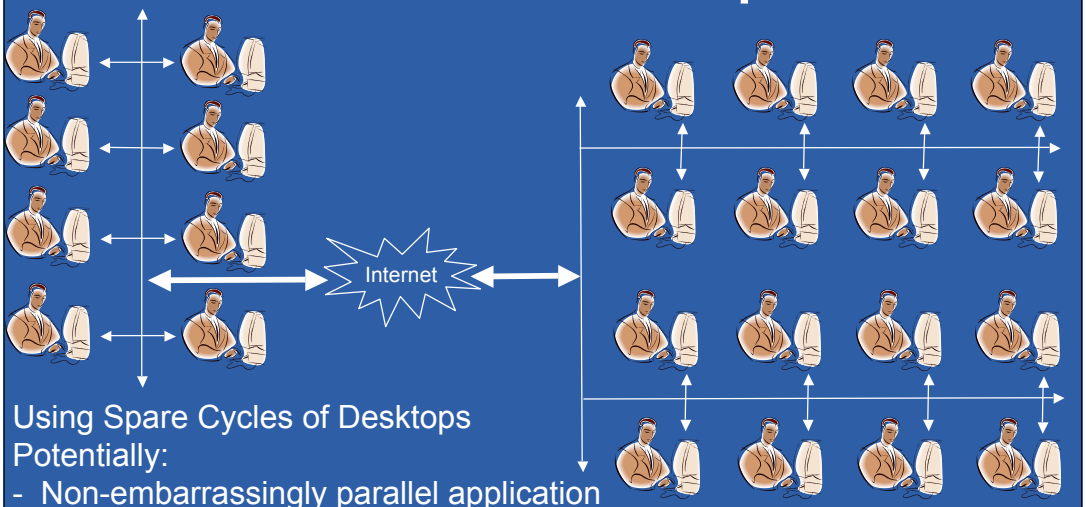
Intranet Grids - Desktop Grids



Using Spare Cycles of Desktops
Potentially:
- Non-embarrassingly parallel application
- Across several sites

The multiple GRIDs

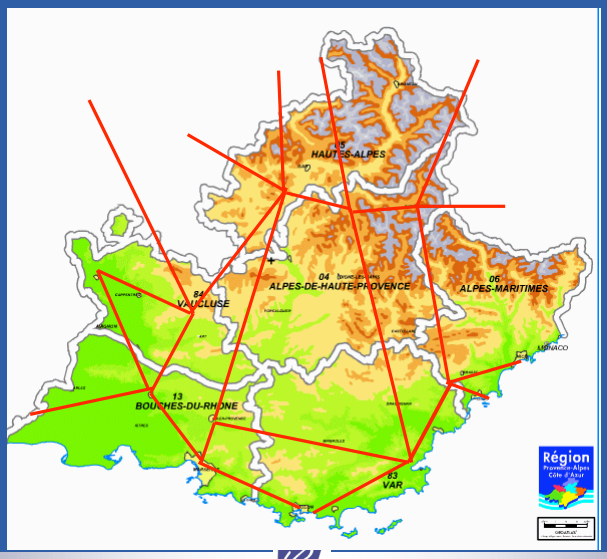
- Scientific Grids
- Enterprise Grids
- Internet Grids
- Intranet Desktop Grids

Strong convergence in process!

At least at the infrastructure level, i.e. WS

Grid: from enterprise ... to regional

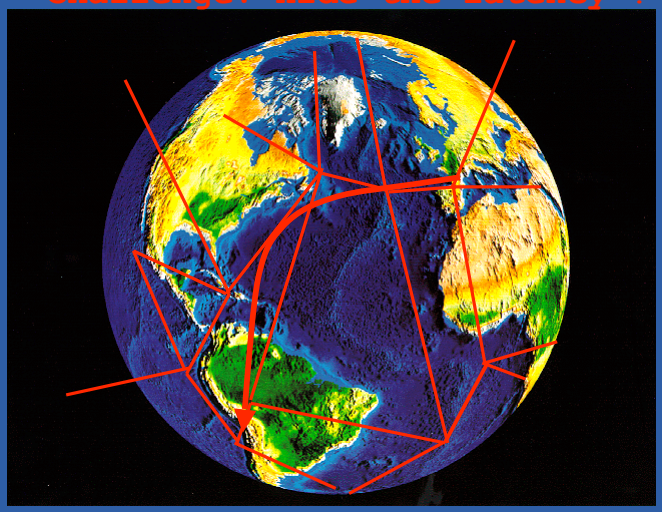
Very hard deployment problems … right from the beginning



Grid: from regional ... to worldwide

Communication Nice - Santiago: 70 ms Light Speed

Challenge: Hide the latency !



Define adequate programming model

Programming

W r a p p i n g

Composing

Deploying

Figures: Web Page Hits: ~ 3 000 / month,
Downloads: 150-300 / month, Users: ?? .us, .mx, .br, .cl, .ch, .it, .fr, ...

Distributed Objects

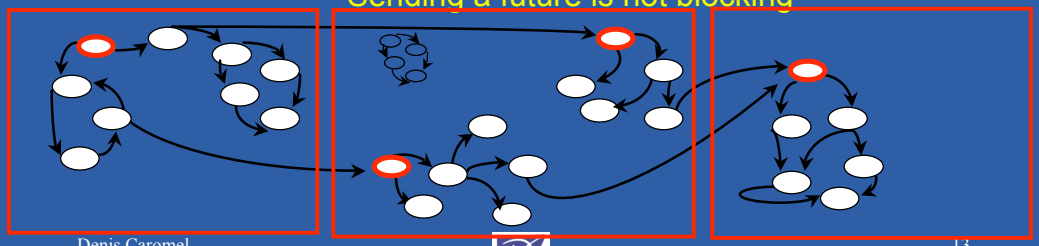
ProActive Programming

ProActive model

Java RMI (Remote Method Invocation = Object RPC = o.foo(p))

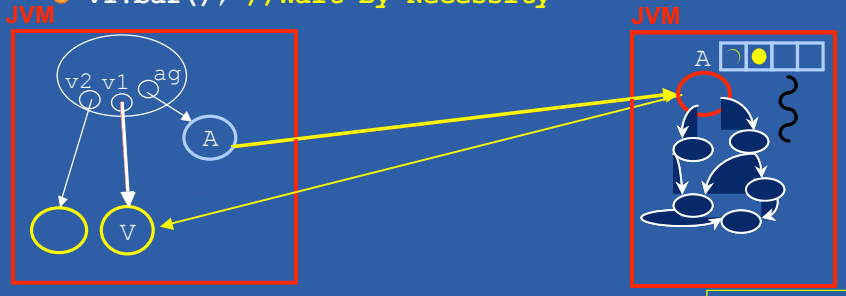
plus a few important features:

- Asynchronous Method calls towards Active Objects:
  Implicit Futures as RMI results
- Wait-By-Necessity:
  - Automatic wait upon the use of an implicit future
  - First-Class Futures:
    - Futures passed to other activities
    - Sending a future is not blocking



ProActive : Active objects

```
A ag = newActive ("A", […], VirtualNode)
V v1 = ag.foo (param);
V v2 = ag.bar (param);
...
v1.bar(); //Wait-By-Necessity
```

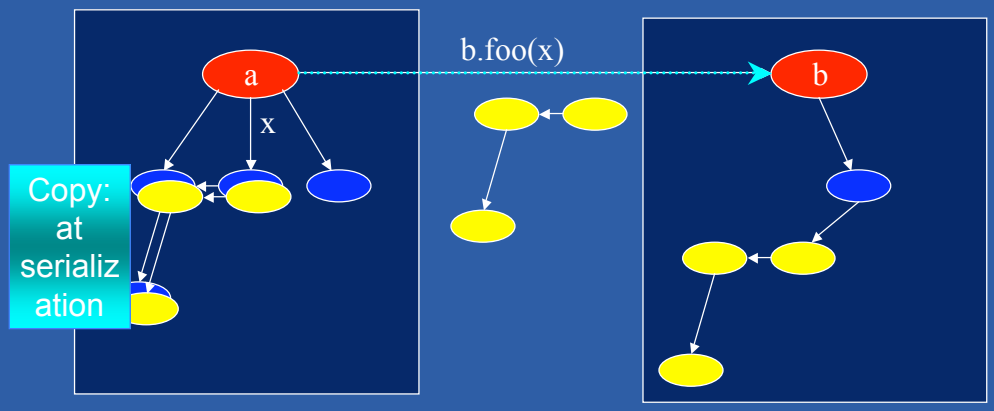


Java Object Active Object Req. Queue

Future Object Proxy Request Thread

Wait-By-Necessity is a Dataflow Synchronization

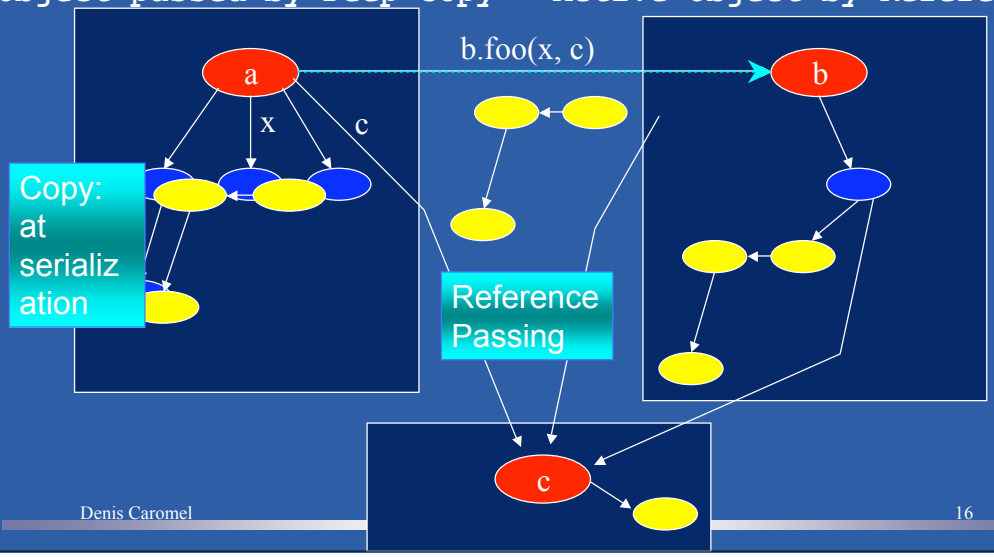
Call between Objects: Parameter passing: Copy of Java Objects



(Deep) Copies evolve independently -- No consistency

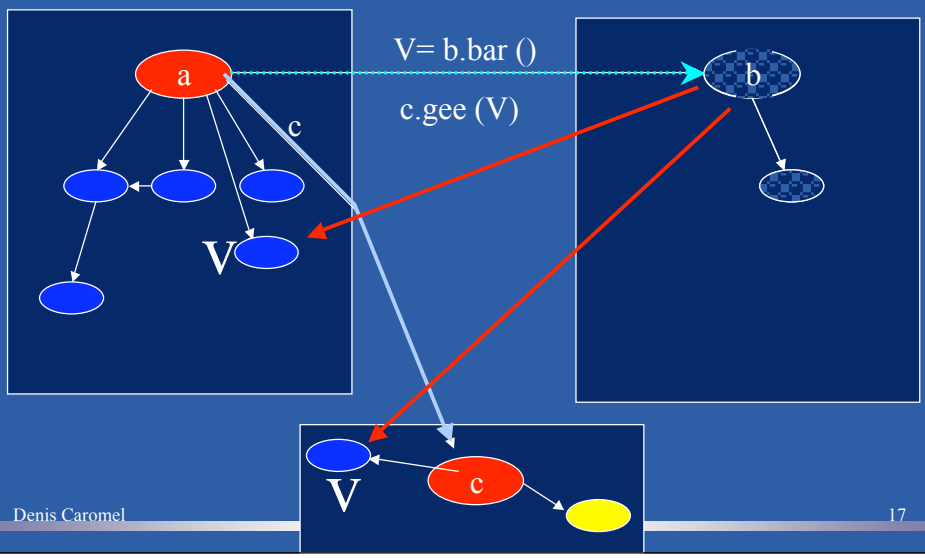
Call between Objects: Parameter Passing: Active Objects

Object passed by Deep Copy - Active Object by Reference



Wait-By-Necessity: First Class Futures

Futures are Global Single-Assignment Variables



ProActive : Explicit Synchronizations

```
A ag = newActive ("A", […], VirtualNode)
V v = ag.foo(param);
...
v.bar(); // Wait-by-necessity
```

Single Future Synchronization:

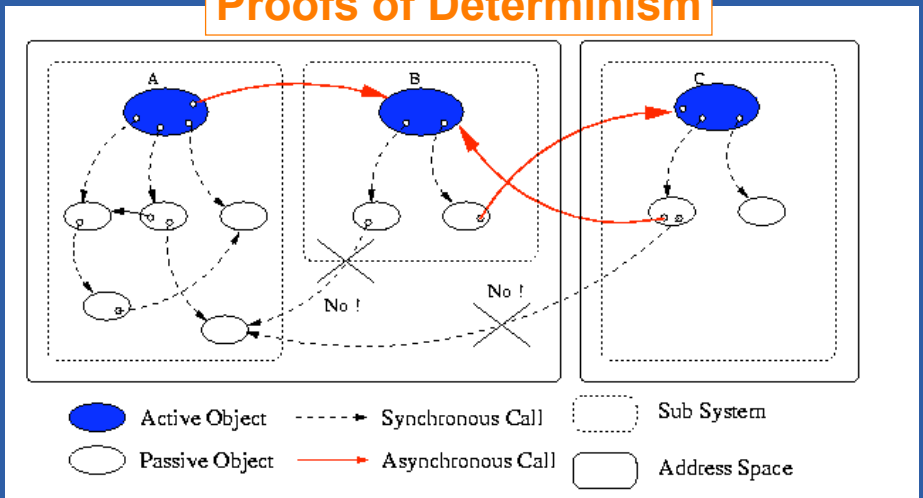
- `ProActive.isAwaited (v); // Test if available`
- `.waitFor (v); // Wait if not available`

Vectors of Futures:

- `.waitForAll (Vector); // Wait all of them`
- `.waitForAny (Vector); // Get One`

Standard system at Runtime: Asynchrony, WbN, ... but no sharing

Proofs of Determinism



Groups

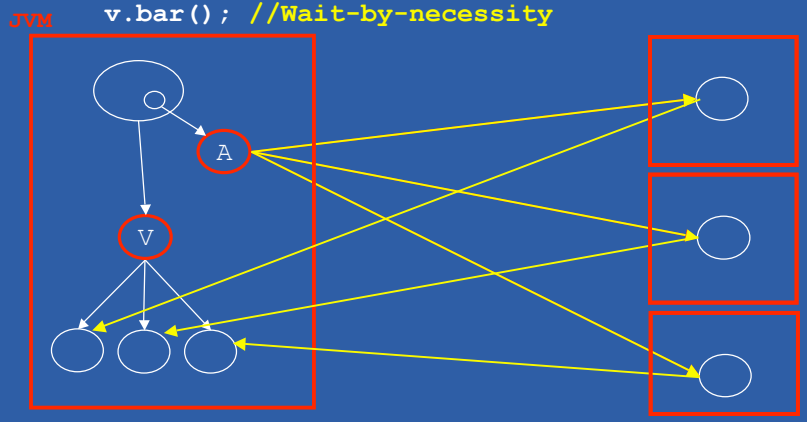
Collective Communications: Groups

- Manipulate groups of Active Objects, in a simple and typed manner:
  - ⇒ Typed and polymorphic Groups of active and remote objects
  - ⇒ Dynamic generation of group of results
  - ⇒ Language centric, Dot notation
- Be able to express high-level collective communications (like in MPI):
  - broadcast,
  - scatter, gather,
  - all to all

```
A ag=(A)ProActiveGroup.newActiveGroup(«A»,{{p1},...},{Nodes,..});
V v = ag.foo(param);
v.bar();
```


Creating AO and Groups

```
A ag = newActiveGroup ("A", […], VirtualNode)
V v = ag.foo(param);
...
v.bar(); //Wait-by-necessity
```



Object-Oriented
Typed Group Communications

Group, Type, and Asynchrony
are crucial for Cpt. and GRID

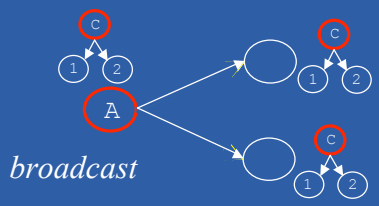
Broadcast or Scatter

Broadcast is the default behavior

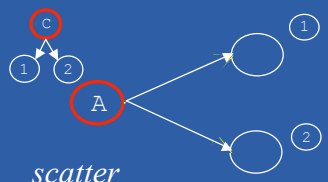
Scatter is also possible

- use a group as parameter
- Scattered depends on rankings

```
gA.bar(gC);    // broadcast gC
ProActive.setScatterGroup(gC);
ga.bar(gC);    // scatter gC
```



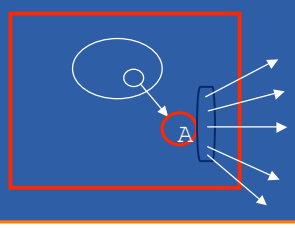
broadcast

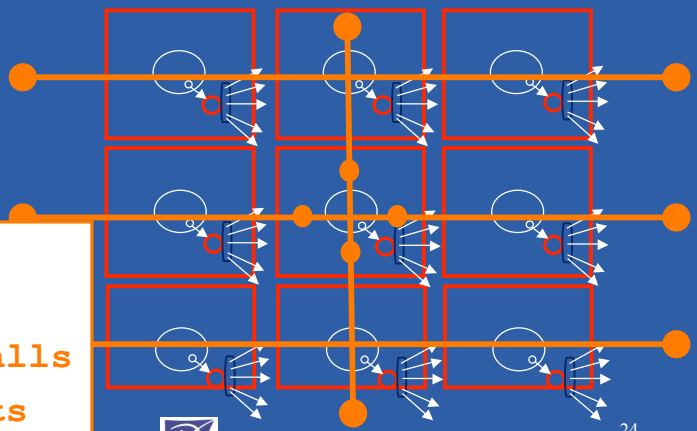


scatter

OO SPMD

```
A ag = newSPMDGroup ("A", […], VirtualNode)
        // In each member
        myGroup.barrier ("2D"); // Global Barrier
        myGroup.barrier ("vertical"); // Any Barrier
        myGroup.barrier ("north","south","east","west");
```





Still,

not based on raw messages, but

Typed Method Calls

==> Components

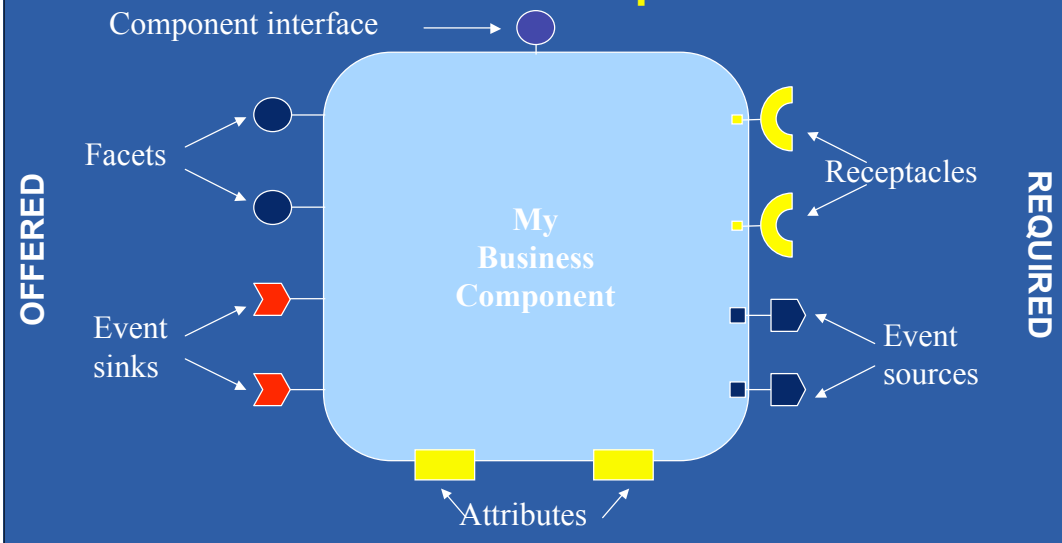
Parallel, Distributed, Hierarchical

Components

for the Grid

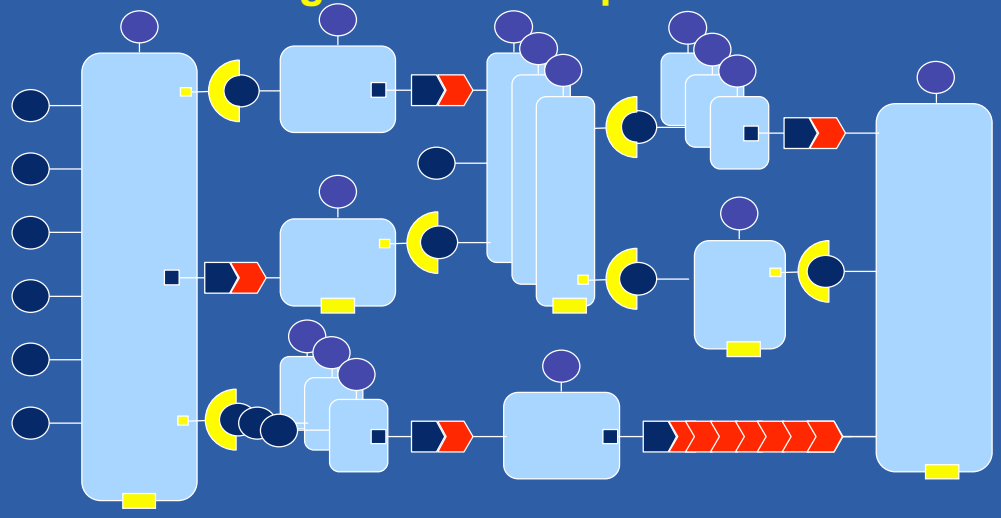
Composing

A CORBA Component



Courtesy of Philippe Merle, Lille, OpenCCM platform

Building CCM Applications = Assembling CORBA Component Instances



Provide + Use, but flat assembly

Component Orientedness

- Level 1: Instantiate - Deploy - Configure
  - Simple Pattern
  - Meta-information (file, XML, etc.) — JavaBeans, EJB
- Level 2: Assembly (flat)
  - Server and client interfaces — CCM
- Level 3: Hierarchic
  - Composite — Fractal, ProActive, ...
- Level 4: Distributed + Reconfiguration
  - Binding, Inclusion, Location — ProActive + On going work

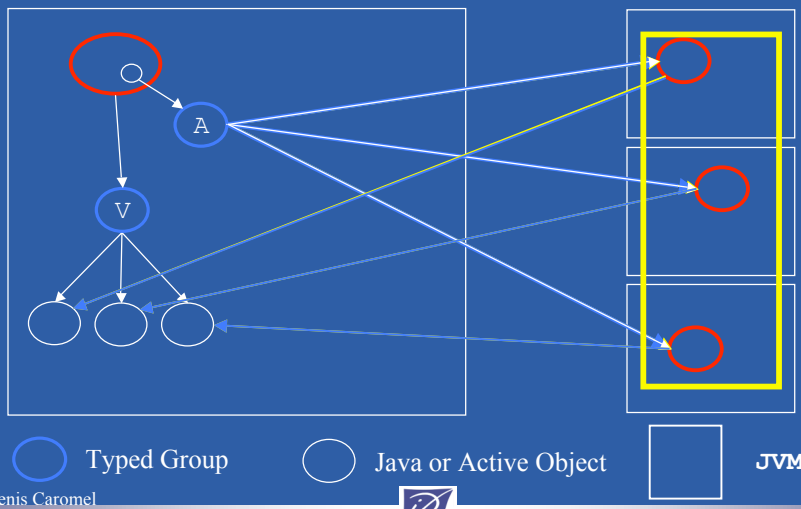
Interactions / Communications: `ProActive`

Functional Calls: service, event, stream

Non-Functional: instantiate, deploy, start/stop, inner/outer, re-bind

Distributed Components (1)

```
ComponentIdentity Cpt = newActiveComponent (params);
A a = Cpt … .getFcInterface ("interfaceName");
V v = a.foo(param);
```

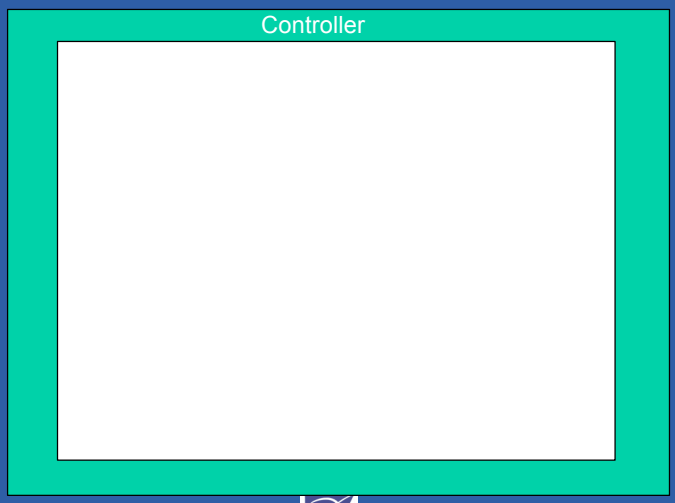


Example of component instance

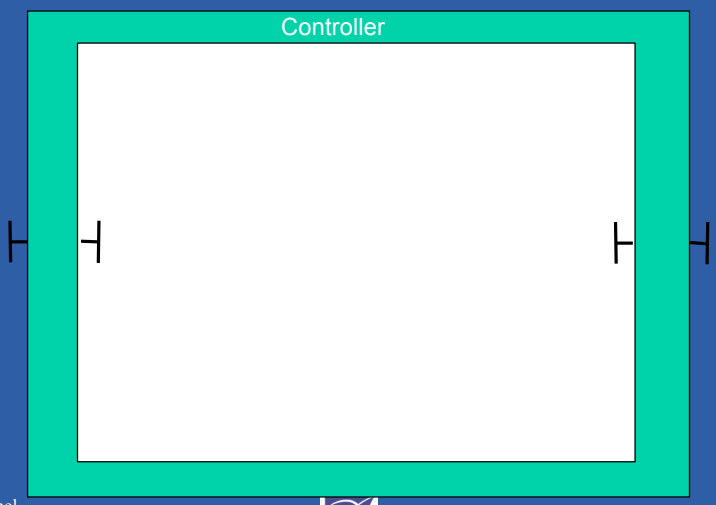
Typed Group Java or Active Object JVM

The Fractal model: Hierarchical Component

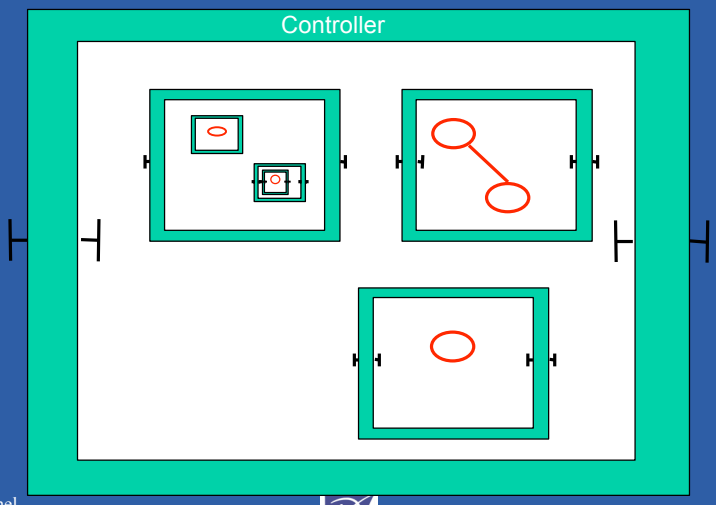
Defined by E. Bruneton, T. Coupaye, J.B. Stefani, FT, et al.



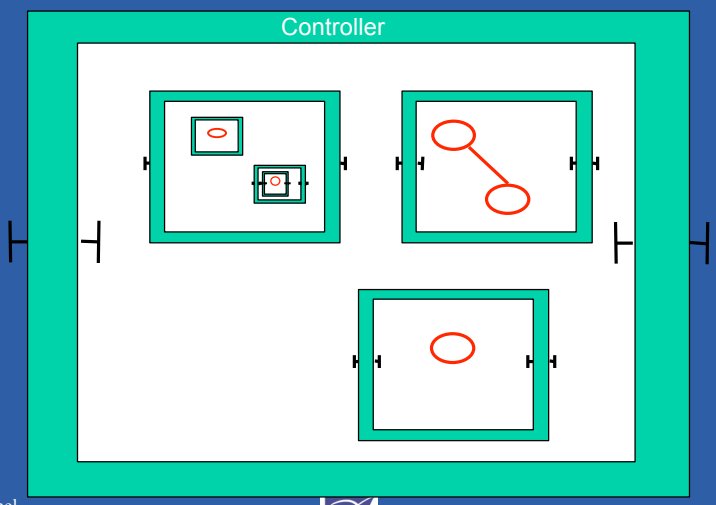
Interface = access point



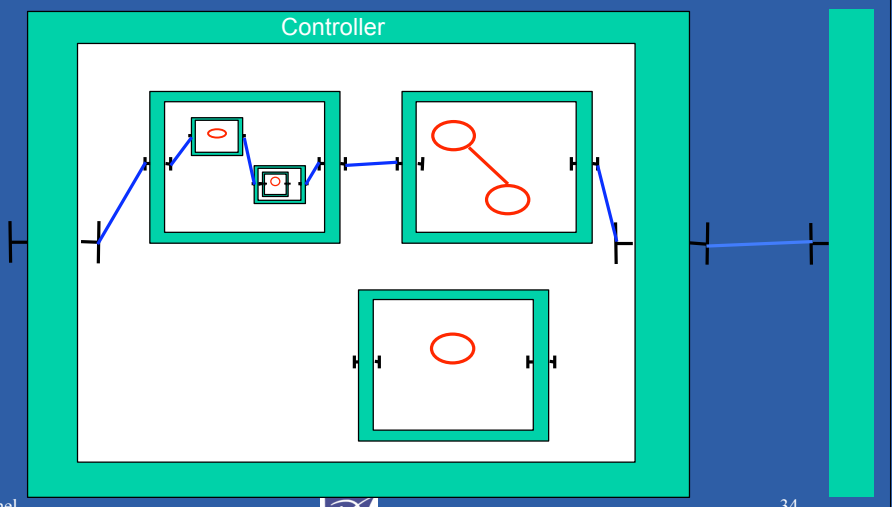
Hierarchical model : composites encapsulate primitives encapsulate Java code



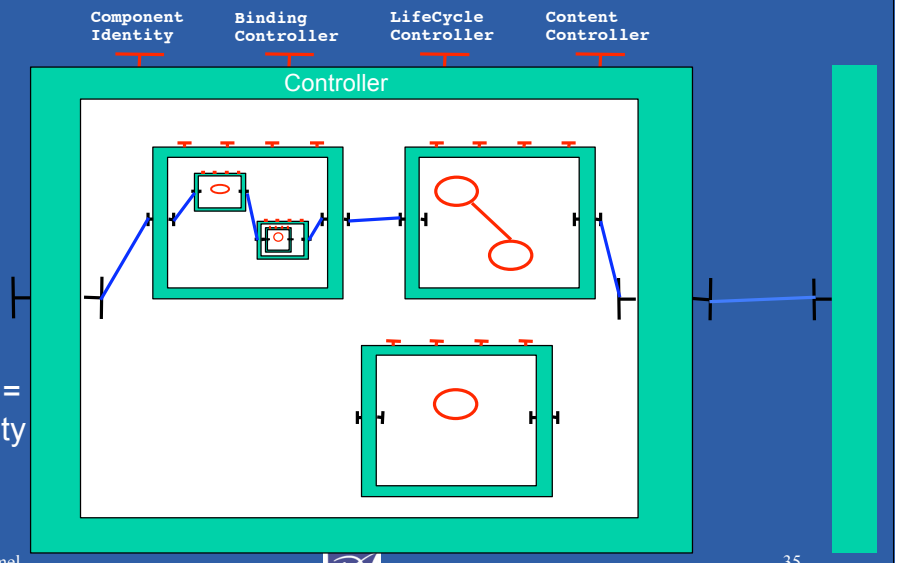
Binding = interaction



Binding = interaction



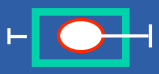
Controllers : non-functional properties



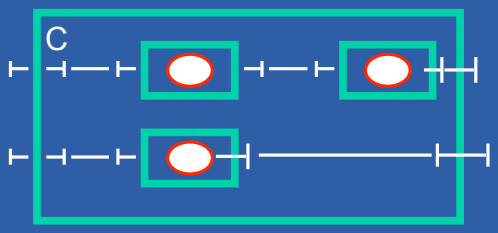
Component = runtime entity

ProActive Components for the GRID

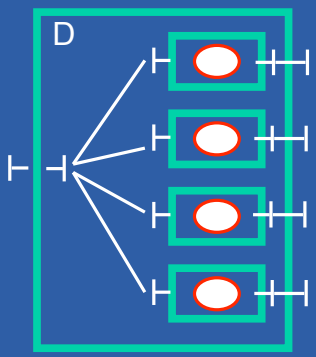
An activity, a process, … potentially in its own JVM



1. Primitive component



2. Composite component



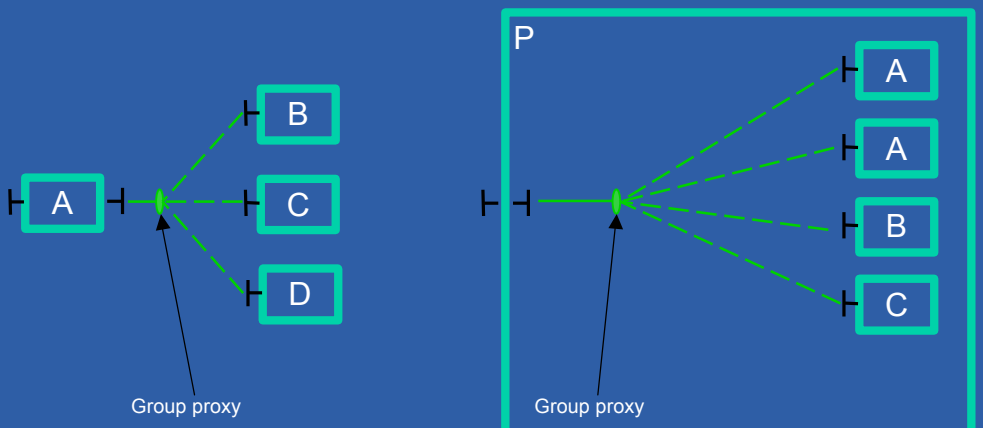
3. Parallel and composite component

Composite: Hierarchical, and Distributed over machines

Parallel: Composite + Broadcast (group)

Groups in Components (1)

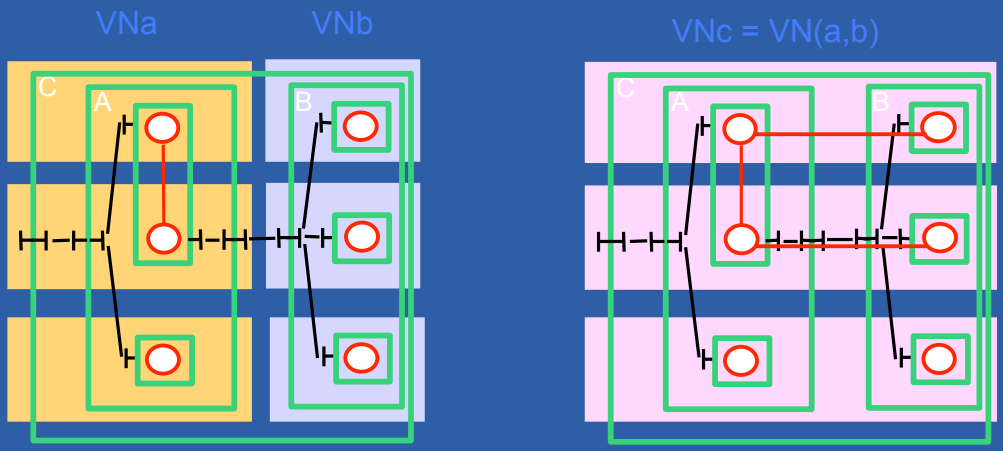
A parallel component!



Broadcast at binding, on client interface

At composition, on composite inner server interface

XML Deployment (Not in source)



Separate or Co-allocation

ProActive Component Definition

A component is:

- Formed from one (or several) Active Object
- Executing on one (or several) JVM
- Provides a set of server ports: Java Interfaces
- Uses a set of client ports: Java Attributes
- Point-to-point or Group communication between components

XML Example

Hierarchical:

- Primitive component: define with Java code and a descriptor
- Composite component: composition of primitive + composite
- Parallel component: multicast of calls in composites

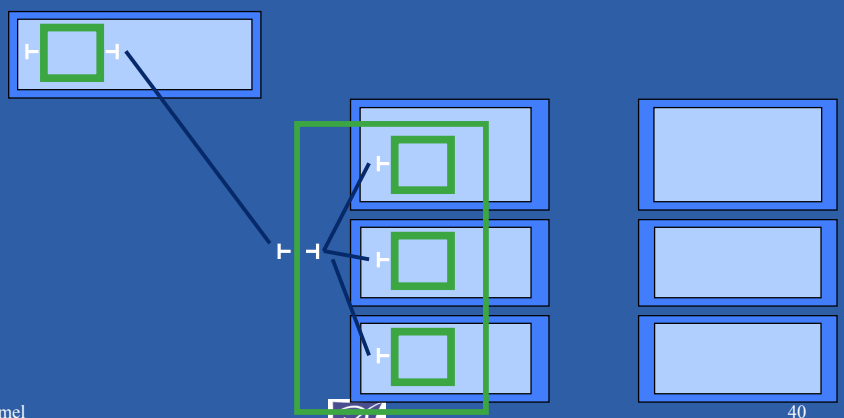
Descriptor:

- XML definition of primitive and composite (ADL)
- Virtual nodes capture the deployment capacities and needs

Virtual Node is a very important abstraction for GRID components

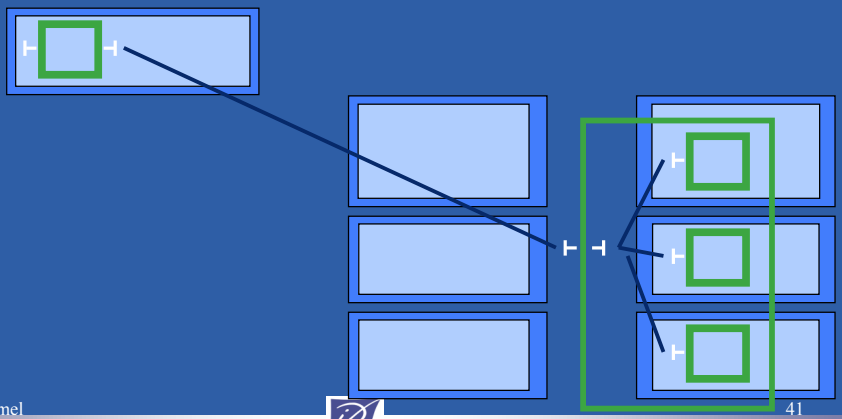
Migration Capability of composites

Migrate sets of components, including composites



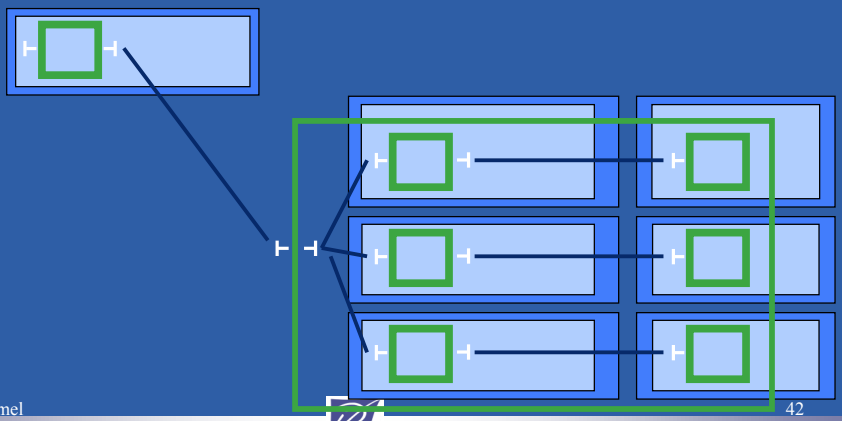
Migration Capability of composites

Migrate sets of components, including composites



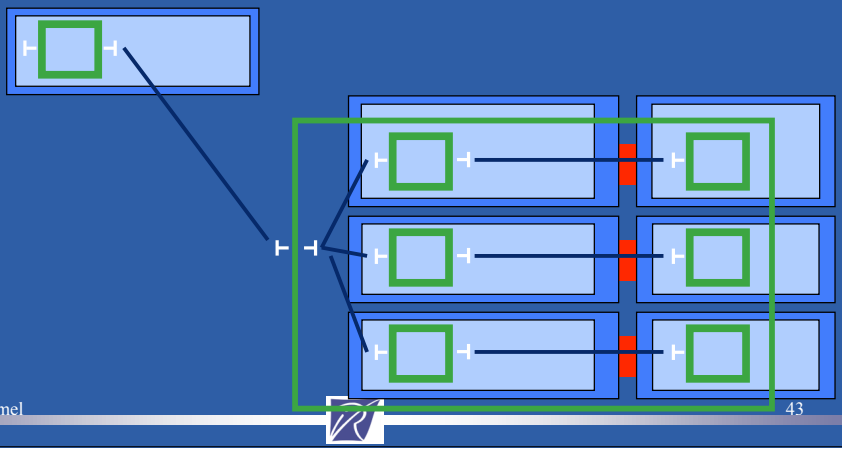
Co-allocation, Re-distribution

e.g. upon communication intensive phase



Co-allocation, Re-distribution

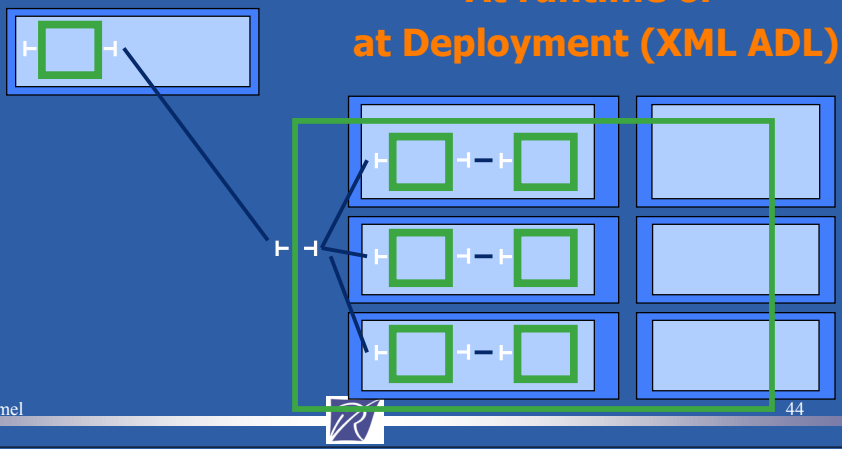
e.g. upon communication intensive phase



Co-allocation, Re-distribution

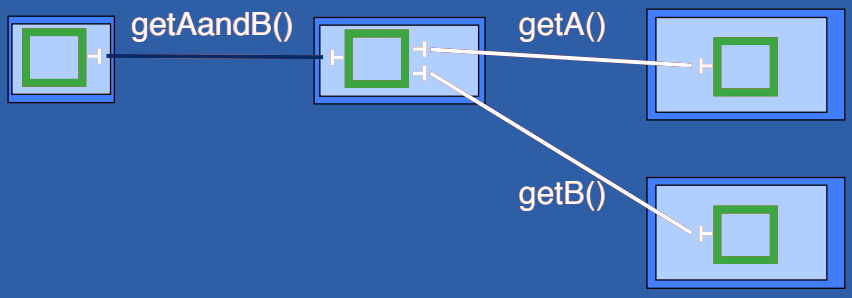
e.g. upon communication intensive phase

At runtime or
at Deployment (XML ADL)



Functionalities : Without First Class Futures

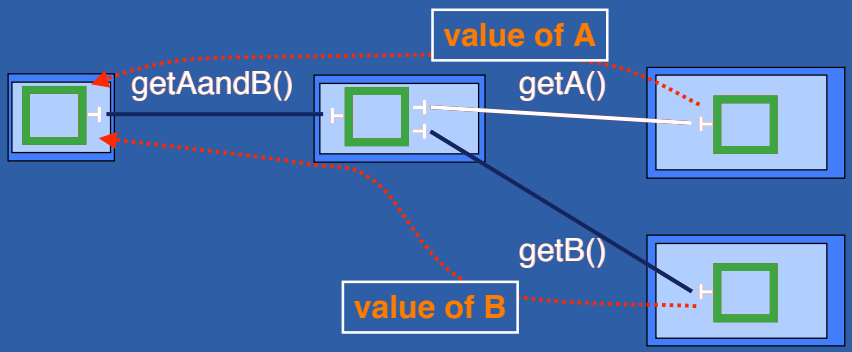
Or in the case of **Synchronous** method calls



Functionalities : With First Class Futures

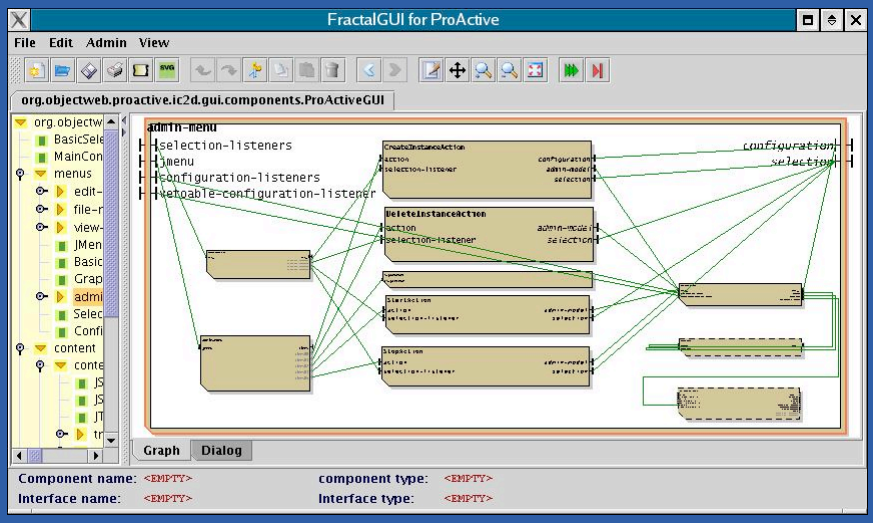
Non-blocking method calls

Example 2 : **Asynchronous** method calls with full-fledge **Wait-By-Necessity**

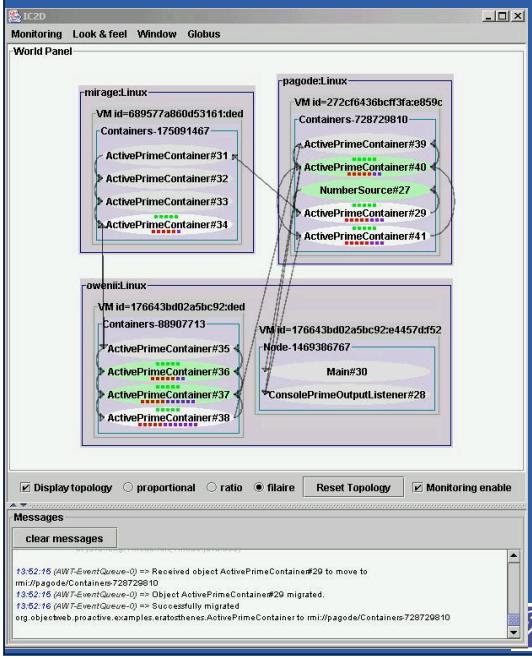


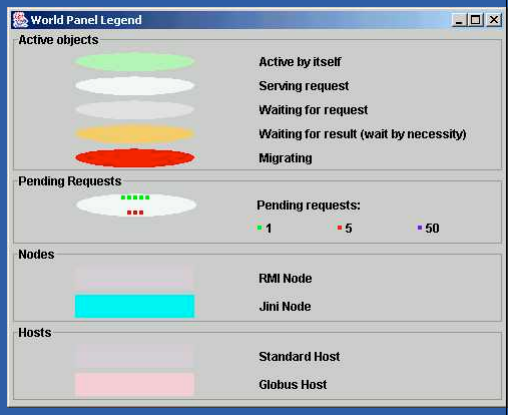
Assemblage are not blocked with Asynchrony + WbN

On-going work : GUI



IC2D: Interactive Control and Debugging of Distribution





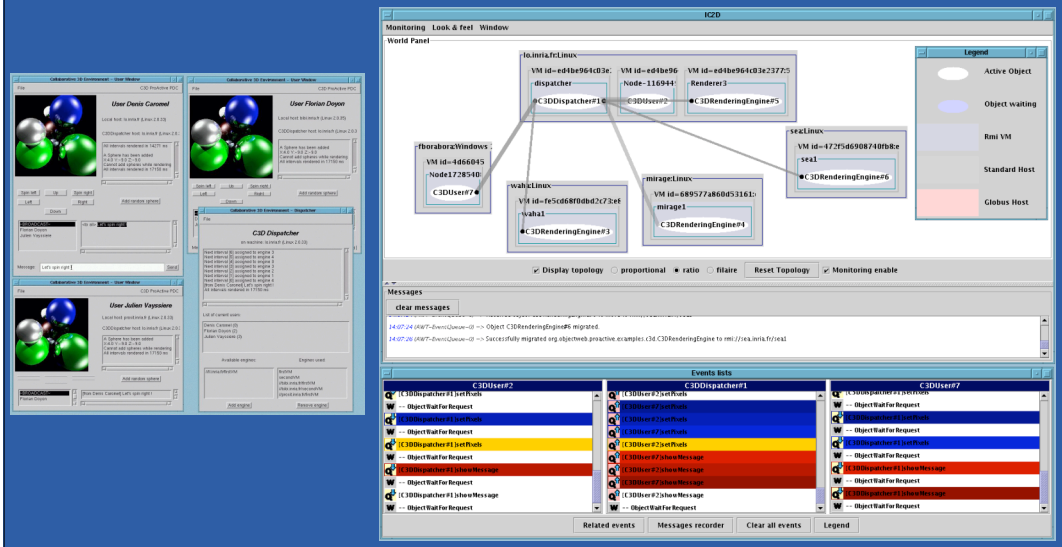
With any ProActive application

Features:

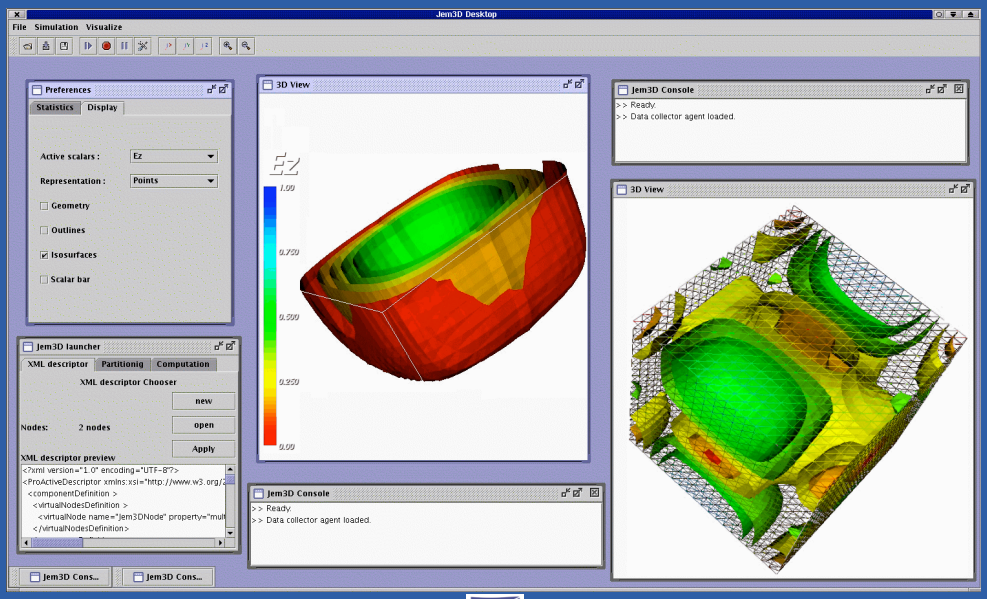
Graphical and Textual visualization

Monitoring and Control

C3D Monitoring: graphical and textual com.



Jem3D



JEM 3D : Java 3D Electromagnetism

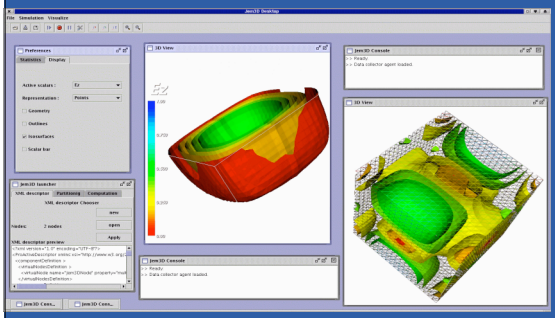
together with Said El Kasmi, Stéphane Lanteri (caiman)

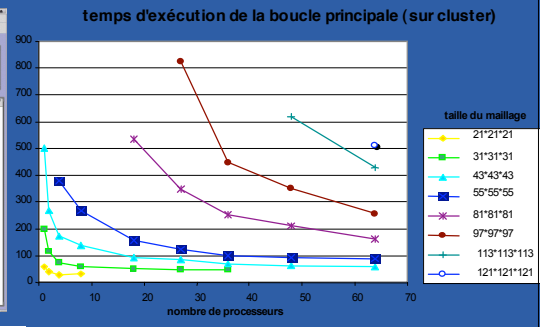
Maxwell 3D equation solver, Finite Volume Method (FVM)

Pre-existing Fortran MPI version: EM3D (CAIMAN team @ INRIA)

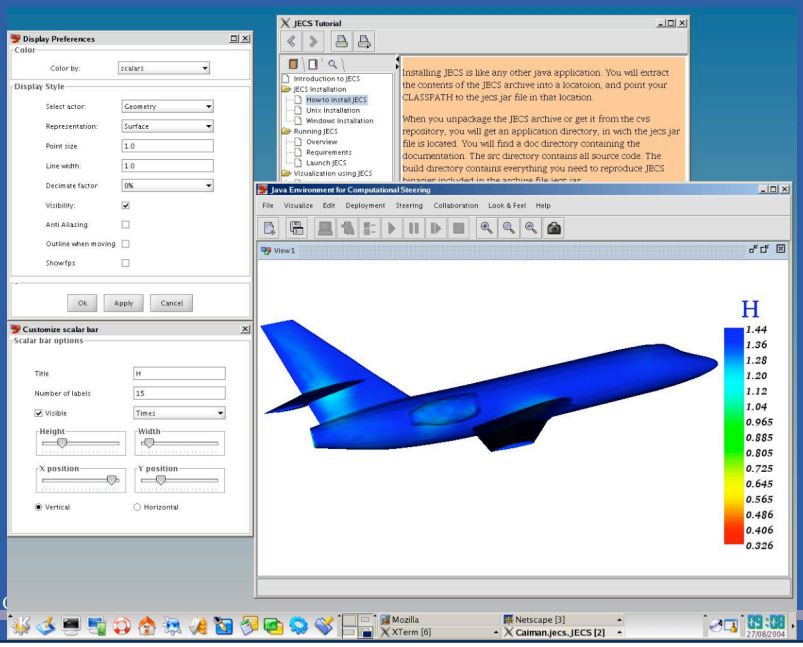
Up to 294 machines at the same time (Intranet and cluster)

Large data sets: 150x150x150 (100 million facets)





JECS : A Generic Version of Jem3D



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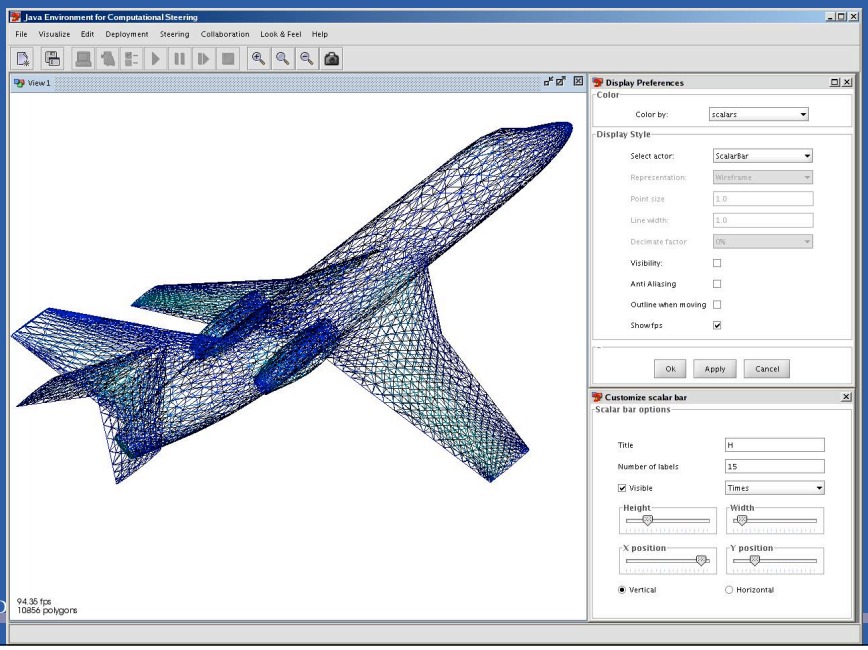
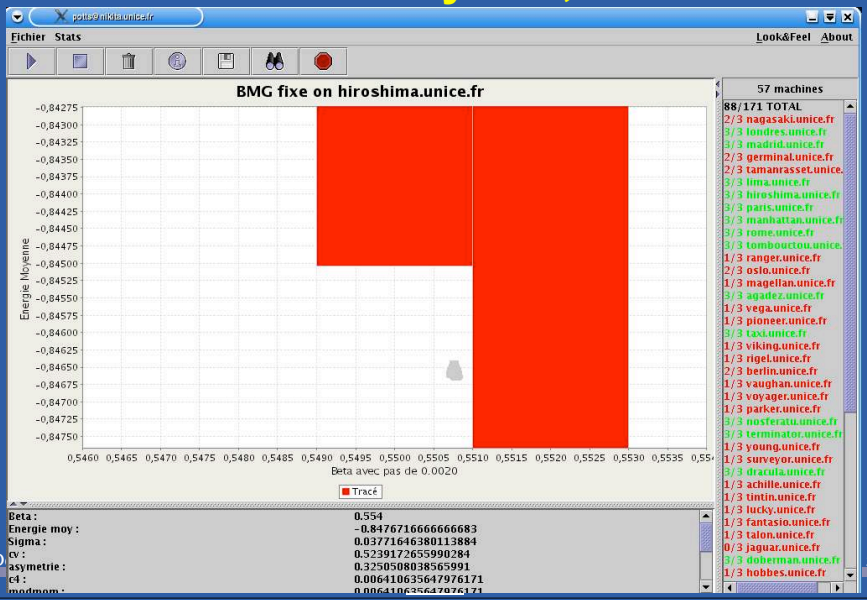


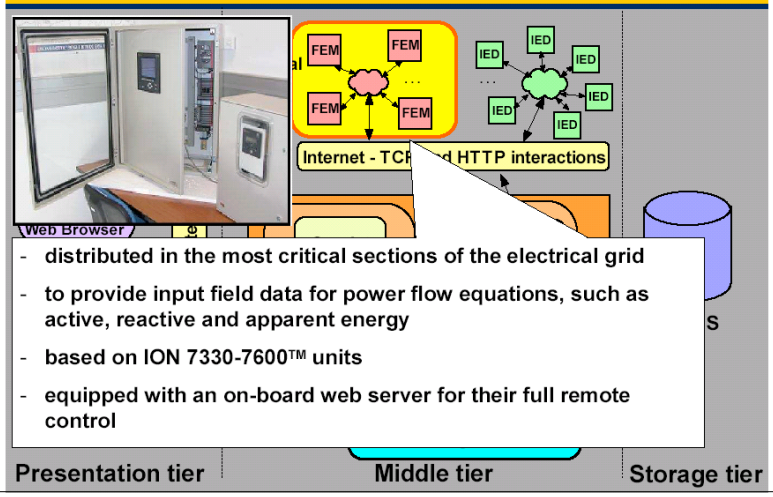
Monte Carlo Simulations, Non-Linear Physics, INLN



Electric Network Planning, E. Zimeo et al., Benevento (Naples), Italy

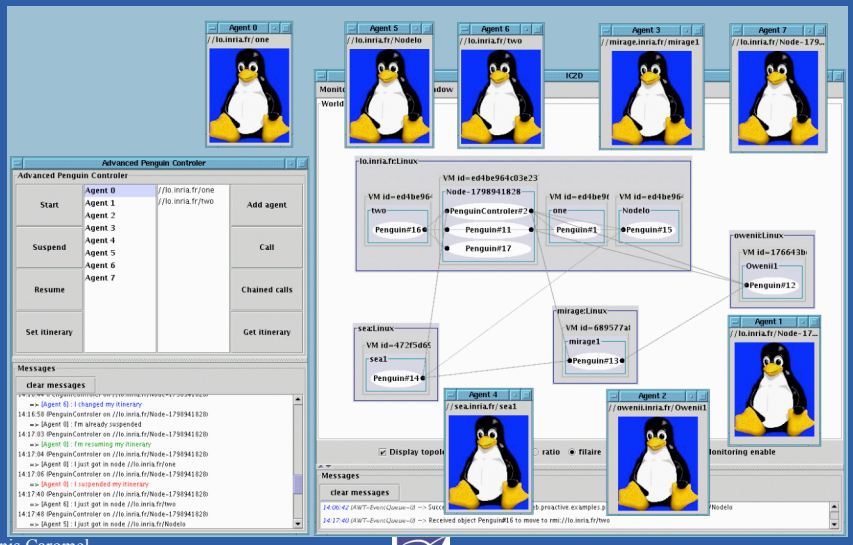
On-line Power Systems Security Analysis (OPSSA)

A network of field power meters (FEMs)

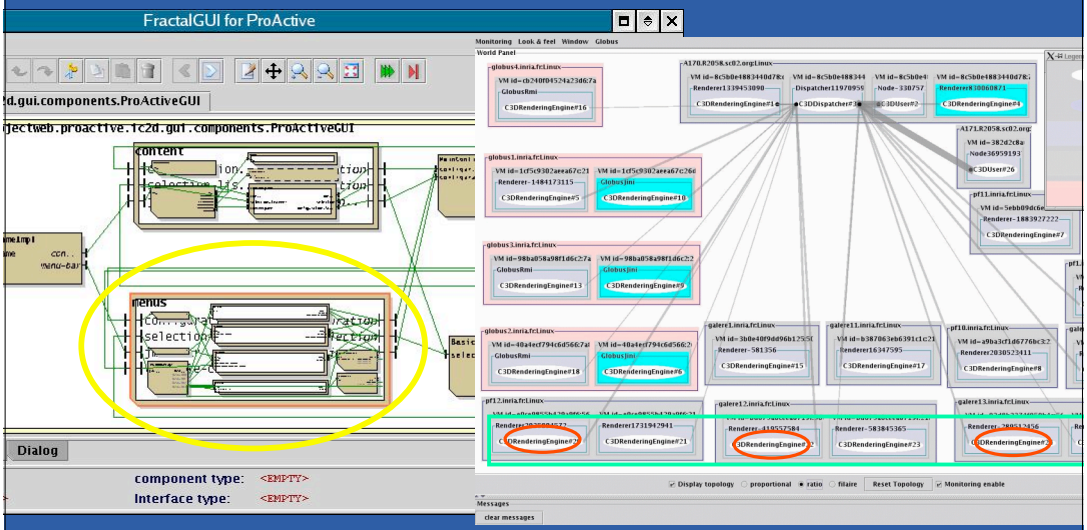


- distributed in the most critical sections of the electrical grid
- to provide input field data for power flow equations, such as active, reactive and apparent energy
- based on ION 7330-7600™ units
- equipped with an on-board web server for their full remote control

Mobile Application executing on 7 JVMs



Perspective for Components - PSE
Graphical Composition, Monitoring, Migration



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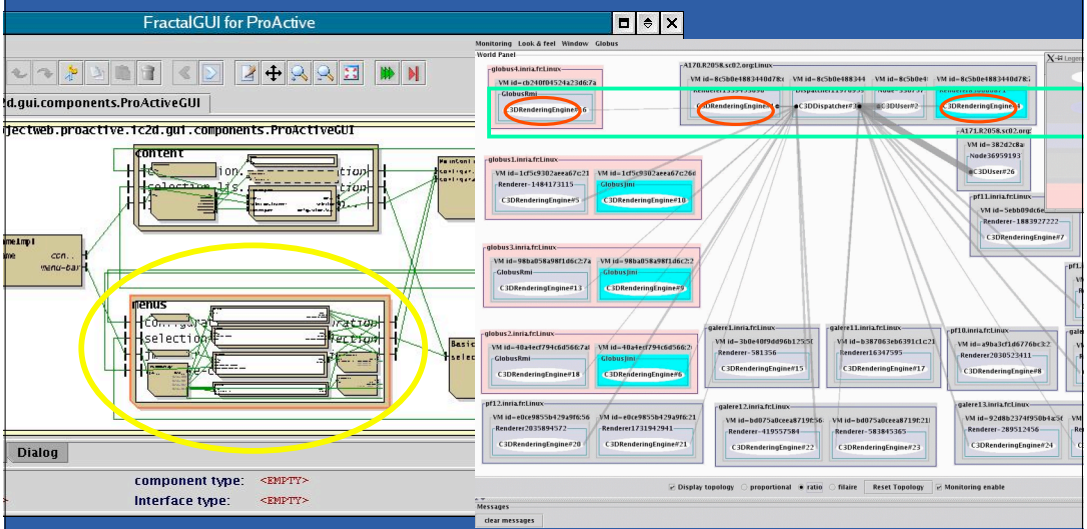

Conclusions and A Few Directions

ProActive
Programming, Composing, Deploying on the Grid

A Strong Programming Model + Components

FACTS AND FIGURES

5 years of computation in 17 days in Desktop P2P

Deployed at once on 600 CPUs (Plugtests on ssh, Globus, LSF, ...)

Conclusions and A Few Directions

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FACTS AND FIGURES

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Deployed at once on 600 CPUs (Plugtests on ssh, Globus, LSF, ...)

(Close to) Beating Fortran on an Electromagnetic Application

PERSPECTIVES FOR COMPONENTS

Safe Reconfiguration

How to specify for components: QoS, Ranking, etc. ?

A great alchemy for the Grid:

Asynchrony + Wait By Necessity + Groups + Components

Conclusion - Beating Fortran ?

Current status:

- Sequential Java vs. Fortran code: 2 times slower
- Large data sets in Java ProActive: 150x150x150 (100 million facets)
- Large number of machines: up to 294 machines in Desktop P2P
- Speed up on `16 machines:`
  - Fortran: 13.8
  - ProActive/Ibis: 12
  - ProActive/RMI: 8.8

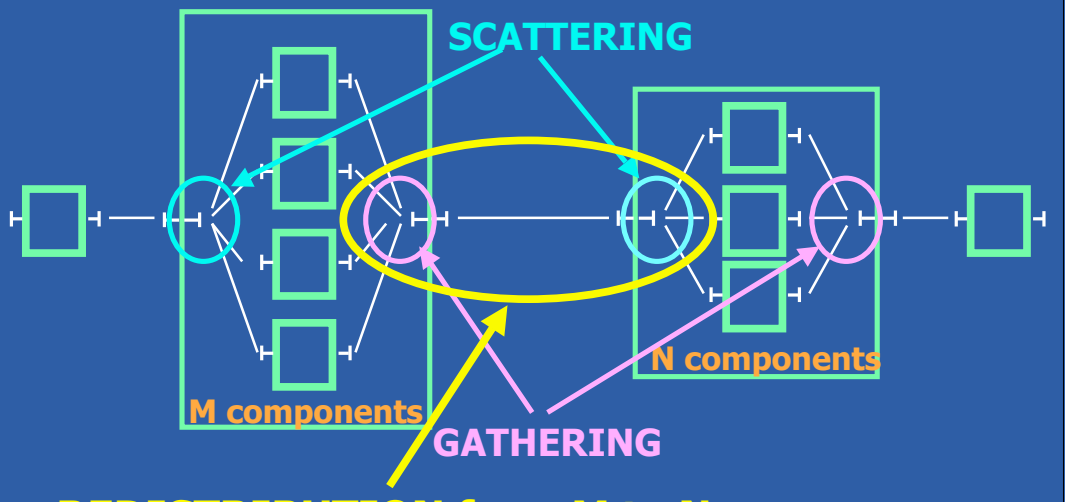
Grid on 5 clusters (DAS 2): Speed up of 100 on 150 machines

Fortran: no more than 40 proc. …

Beating Fortran MPI with Java ProActive? X/40 (14/16) = 2X/ n (100/150)

Yes, starting at 105 machines !

On Going : M x N Communications + Redistribution



Adaptive Feature: Multi-transports layer RMI, RMI-ssh, …, Ibis, HTTP XML, ...

Adaptive choice of transport layer between:

- RMI
- ssh/RMI

Also available with static configuration:

- Ibis (TCP, Myrinet, etc.)
- HTTP
- … ssh/HTTP

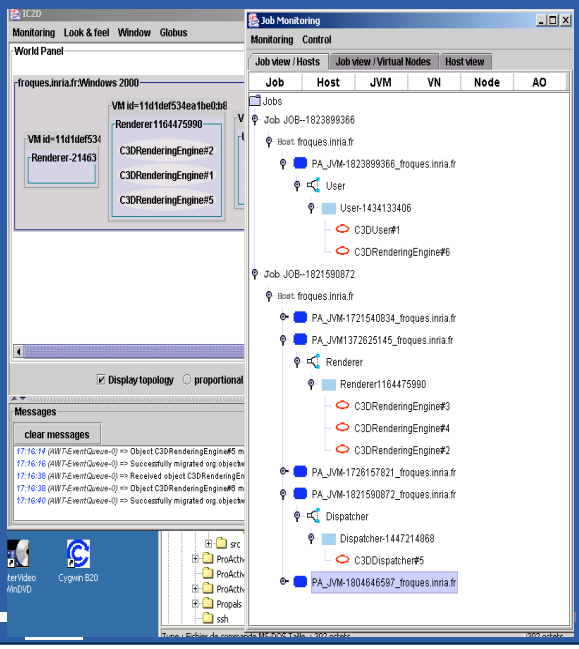
Short Term Perspective:

Fully Adaptive Choice between all transports

IC2D: Basic features cont.

Job Management:

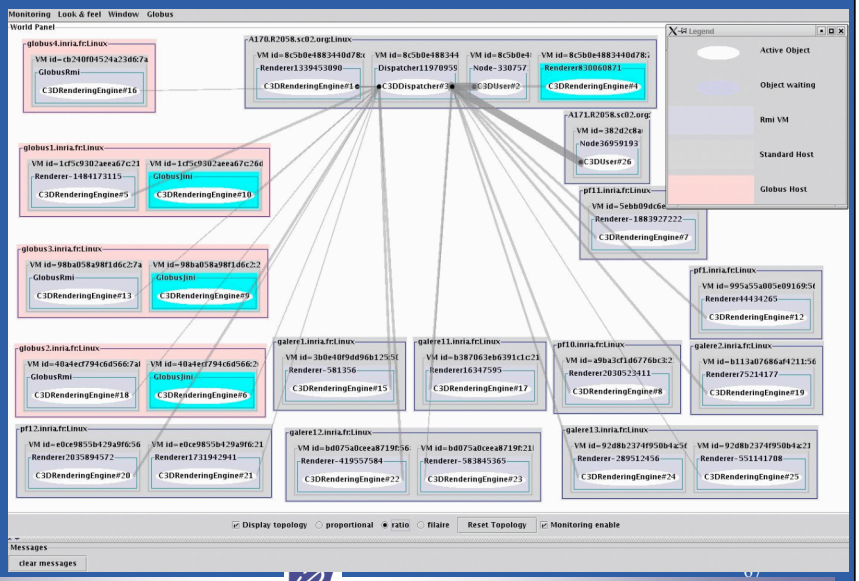
- JVM, AO per Job ID
- Textual visualisation,
- control (kill all, etc.)



Monitoring of RMI, Globus, Jini, LSF cluster Nice -- Baltimore

ProActive IC2D:

Width of links proportional to the number of com-munications



ProActive:

A Java API + Tools for Parallel, Distributed Computing

- A uniform framework: An Active Object pattern
- A formal model behind: Determinism, Insensitivity to deployment

Programming Model:

- **Remote Objects (Classes, not only Interfaces, Dynamic)**
- **Asynchronous Communications, Automatic dataflow synchro: Futures**
- **Groups, Mobility, Components, Security**

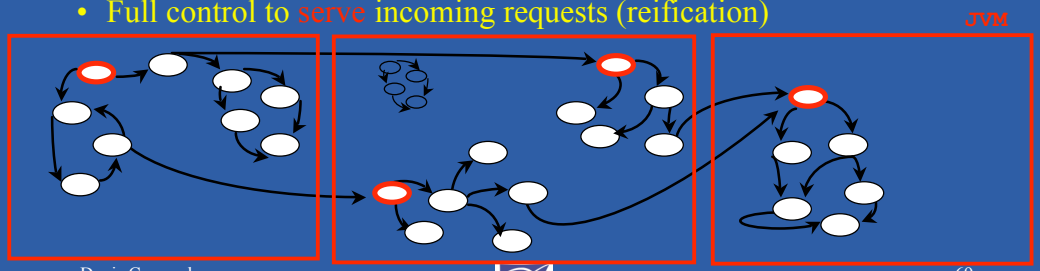
Environment:

- **XML Deployment Descriptors**
- **Interfaced with various protocols:** `rsh,ssh,LSF,Globus,Jini,RMIregistry`
- **Visualization and monitoring: IC2D**

In the www. ObjectWeb .org Consortium (Open Source middleware) since April 2002 (LGPL license)

ProActive : model

- Active objects : coarse-grained structuring entities (subsystems)
- Each active object: - possibly owns many passive objects
  - has exactly one thread.
- No shared passive objects -- Parameters are passed by deep-copy
- Asynchronous Communication between active objects
- Future objects and wait-by-necessity.
- Full control to serve incoming requests (reification)



ProActive model (2)

Java RMI (Remote Method Invocation = Object RPC = o.foo(p))

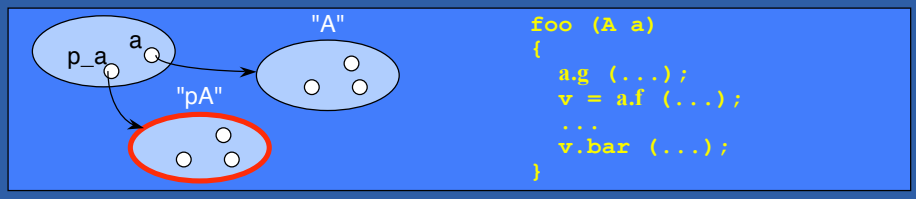
plus a few important features:

- Sequential Object: a single thread with FIFO service
- Asynchronous Method calls towards Active Objects:
  Implicit Futures as method results
- Wait-By-Necessity:
  - Automatic wait upon a strict operation on an unknown future
  - First-Class Futures:
    - Futures can be passed to other activities
    - Sending a future to another machines is not blocking

ProActive : Reuse and seamless

Two key features:

- Polymorphism between standard and active objects
  - Type compatibility for classes (and not only interfaces)
  - Needed and done for the future objects also
  - Dynamic mechanism (dynamically achieved if needed)



- Wait-by-necessity: inter-object synchronization
  - Systematic, implicit and transparent futures

    Ease the programming of synchronizations, and the reuse of routines

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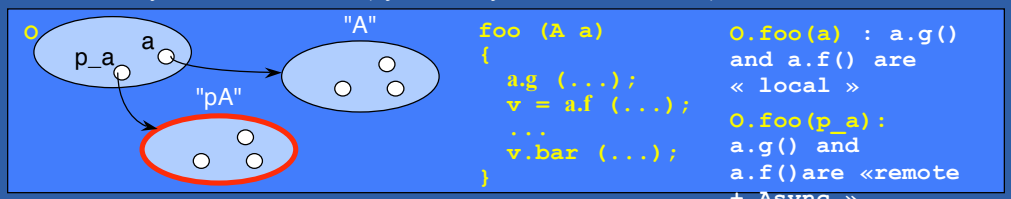


- Wait-by-necessity: inter-object synchronization
  - Systematic, implicit and transparent futures (“value to come”)

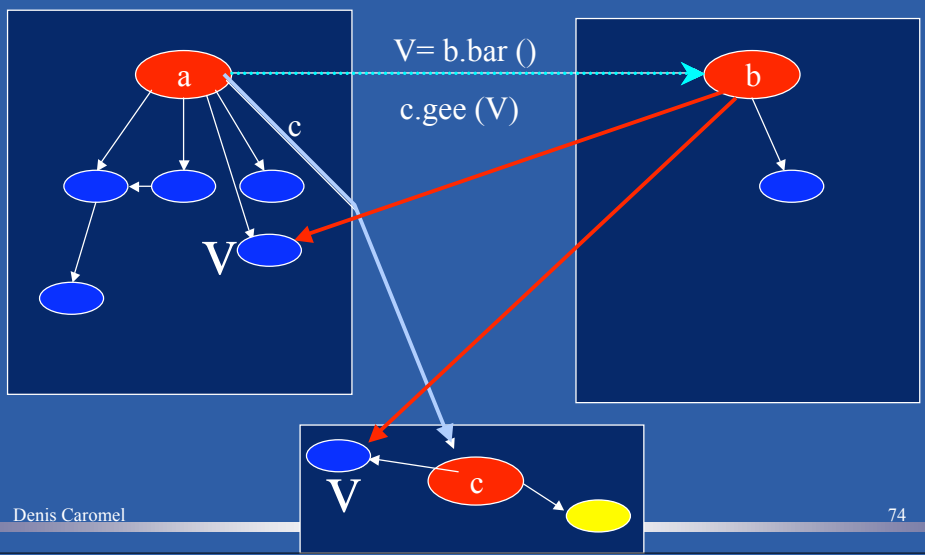
    Ease the programming of synchronizations, and the reuse of routines

First-Class Futures

Update

Wait-By-Necessity: First Class Futures

Futures are Global Single-Assignment Variables



Future update strategies

No partial replies and requests:

- No passing of futures between activities, more deadlocks

Eager strategies: as soon as a future is computed

- Forward-based:
  - Each activity is responsible for updating the values of futures it has forwarded
- Message-based:
  - Each forwarding of future generates a message sent to the computing activity
  - The computing activity is responsible for sending the value to all

Mixed strategy:

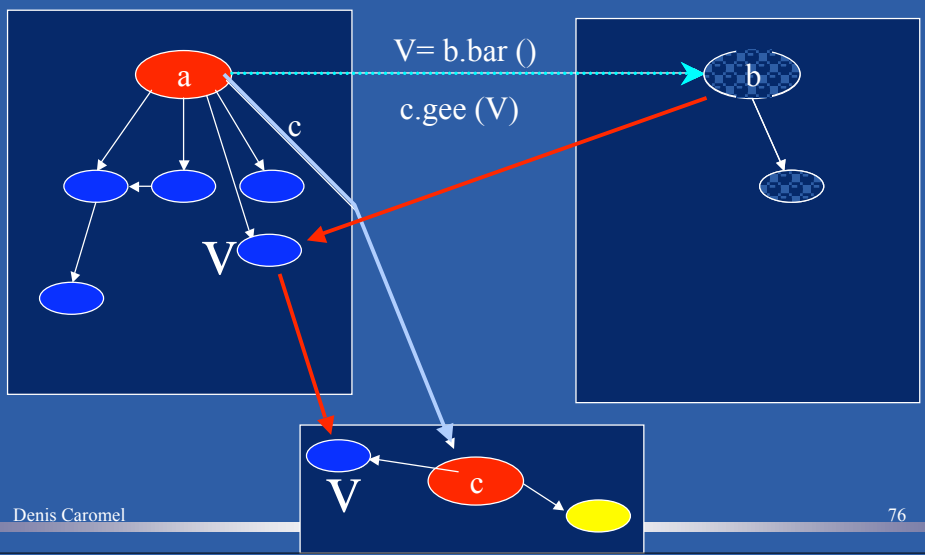
- Futures update any time between future computation and WbN

Lazy strategy:

- On demand, only when the value of the future is needed (WbN on it)

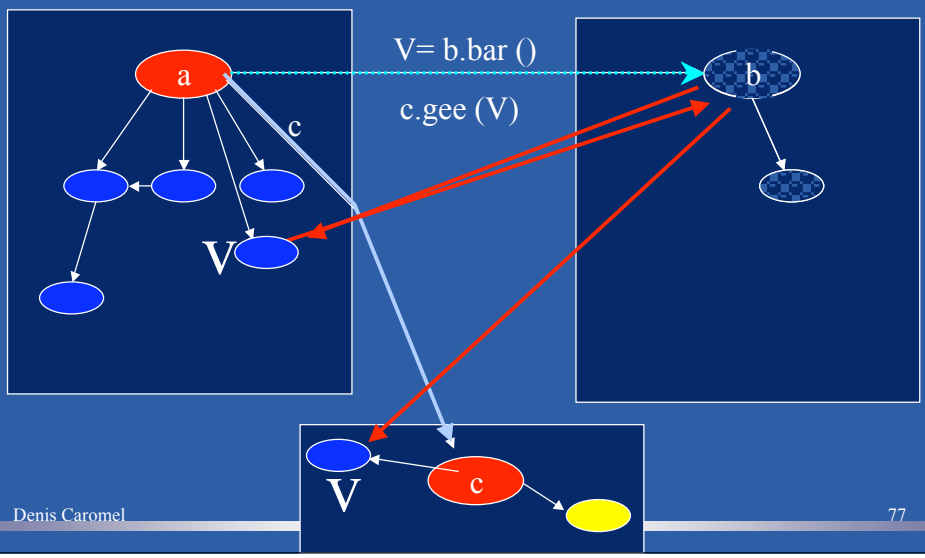
Wait-By-Necessity: Eager Forward Based

AO forwarding a future: will have to forward its value



Wait-By-Necessity: Eager Message Based

AO forwarding a future: send a message



Wait-By-Necessity: Lazy Strategy

An Active Object requests a Future Value when needed