

Defining a Specification Language for Distributed Components

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The Problem

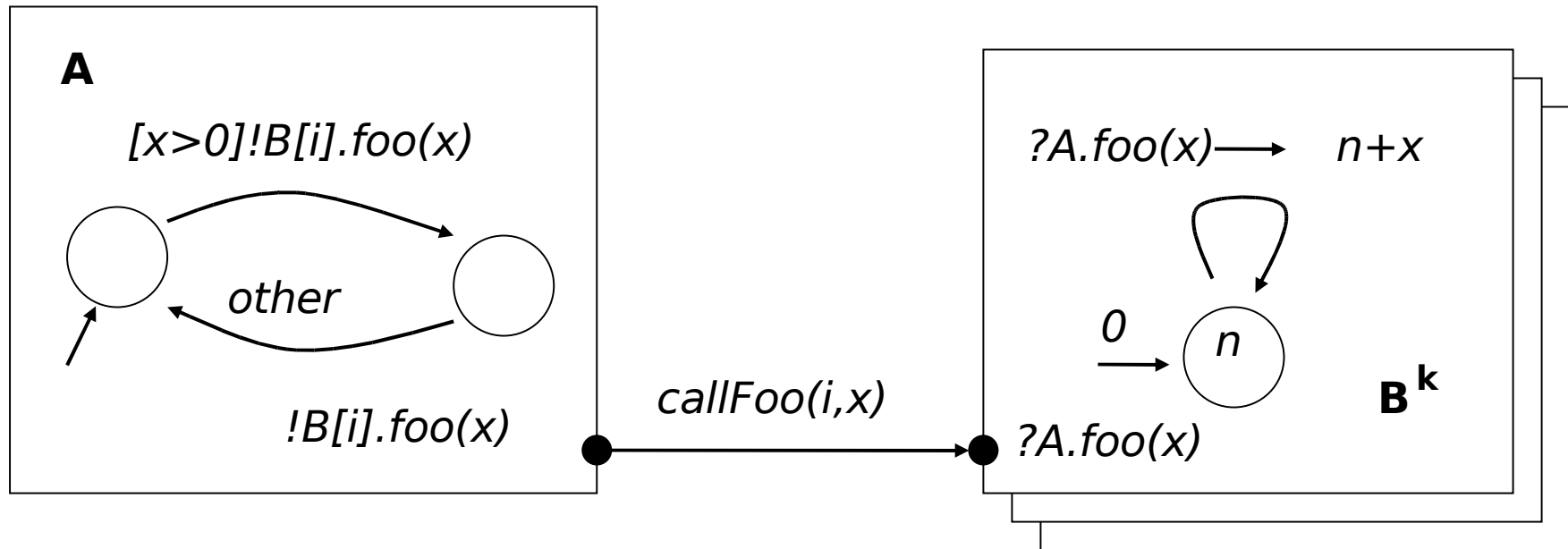
- What to specify?
 - Distributed Components – GRIDs, GCM
 - Asynchronous calls
 - Multiple components
 - Collective Interfaces
 - Dynamic structures
- Who are the users?
 - Software Engineers
 - UML, Java, XML

Basic Ideas

- Specify the assembly of hierarchical components
 - Components, required and provided interfaces
 - Bindings
- Specify behaviour
 - Potential messages, synchronisation points
 - Collective interfaces (Multicast, Gathercast)

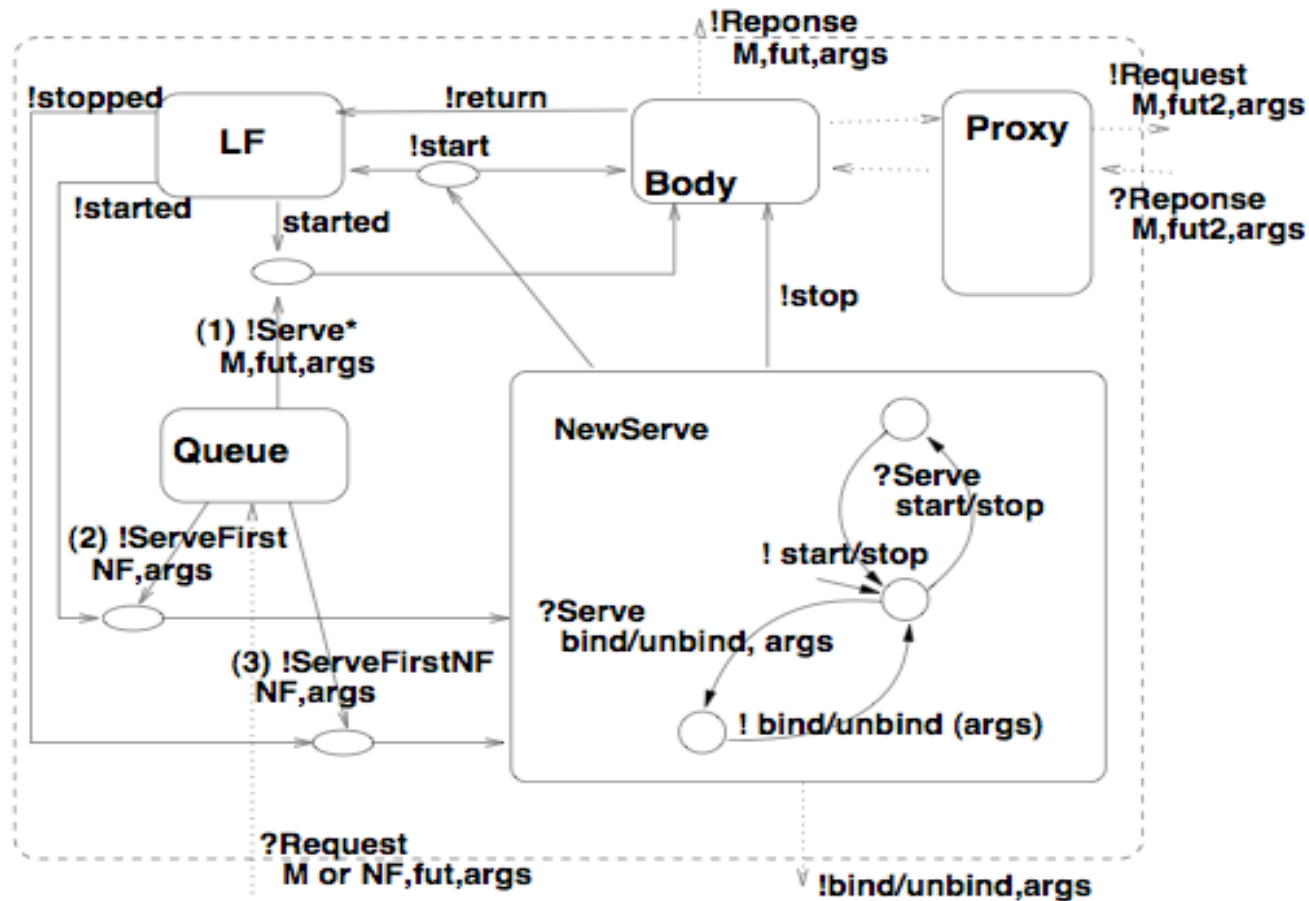
Basics

- Semantic model: pNets



Primitive Component

- Previously defined models [FACS'05/06]



What's wrong?

- Current approach (VERCORS)
 - Fractal ADL + behaviour (LOTOS)
 - Only static architectures
 - Different syntaxes and semantics
 - Up to the user to write “component-friendly” LOTOS code
 - Not clear at all what can and cannot be done
 - We have models, but we don't have a high-level language to express the behaviour!

What's wrong?

- Missing info in ADL
 - Useful only for component description and static deployment
 - Inadequate when dealing with multiple components
 - How to specify a binding between 2 specific components 'i' and '(i+1)%n'?
 - Difficult to soundly attach the component behaviour

CoCoME

- Provide a common example for Component-based software development (CBSD)
- Evaluate and compare the practical appliance of existing component models

Point-of-Sale (POS)

- M cashdesks, N stores
- N local databases, 1 main database
- Mandatory Features
 - Multiple components, collective interfaces
 - Asynchronous communications
- Optional Features
 - Dynamic reconfiguration
 - Fault tolerance

Language Goals

- Friendly syntax to Software Engineers
- Generate Java code (ProActive control flow)
 - At least for simulation
- Generate behavioural model (pNets)
 - Only a subset of specifications will be verifiable

Fundamentals

- Component Software
- ProActive Middleware
 - Single control thread (Primitives)
 - No explicit synchronisation primitives
 - Transparent to the user
 - RDV, futures, wait-by-necessity
 - i.e., Data-flow synchronisation

Key!

Scope

- Considered

- Architecture (composites, primitives)
- Control flow, Data flow
- Synchronisation points
- Dynamic structures

Design

- Not considered

- Data computation
- Physical infrastructure
- Non-functional exceptions (network problems)

Implementation

Middleware

How?

- Control flow, Data flow
 - Sequence, branching
 - Method calls
 - Return values
 - Data usage
 - Synchronisation points

Choices...

- Language semantics

- UML

- Component diagrams
 - Statechart diagrams

CTTool

- Java-like

- Architecture definition
 - Code

JDCSpec

Choices...

- 'Futures'

- Explicit

- !call(method,args);
 - ?use(value);

pNets

- Implicit

- value = interface.method(args);
 - print(value);

JDCSpec

Why implicit synchronisations?

- Allow switching from async to sync calls
 - Reuse specification for async and sync implementations
 - Check the effect of exceptions
 - Exceptions -- synchronous / barriers

User need

- Transparent first order continuations

```
client = databaseltf.getClient(id);  
managerltf.check(client);
```

ProActive

Data types

- First order types
 - Integers
 - Intervals
 - Enumerated types
 - Records (without recursion)
 - Arrays

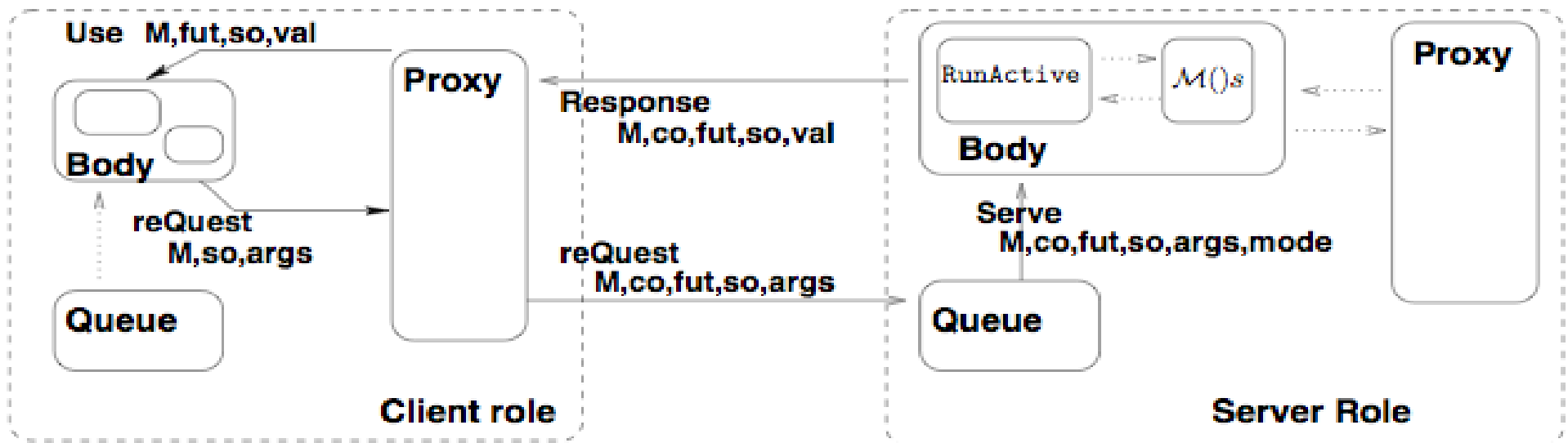
Data domains

- Definition of data domains inside the specification
 - Required for the verification and the code generation
- Finite model left as a second abstraction
 - Specific to the formula to prove

Proposition

- Start from a simplified Java
 - Simple data types
- Augment with component primitives
 - “component”, “interface”
 - “bind”, “unbind”, “start”, “stop”, ...
- Ideally, the user should only care about the business code
 - Implicit synchronisations
 - Transparent futures

A ProActive view



What's inside a primitive?

- Implicit
 - Infinite queue, proxies
- Explicit (given by the user)
 - Activity (Serving Policy)
 - Methods (defined by the server interfaces)
 - Local variables
 - Access to client interfaces
 - Persistent variables
 - Visible to all methods and to the activity

Restrictions

- A method may only call client interfaces or local methods (no service method calls)
- A method cannot manipulate the queue

What's inside a composite?

- Implicit
 - Infinite queue
- Explicit (given by the user)
 - Activity (Serving Policy)
 - Direct Subcomponents
 - External interfaces
 - Deploy policy
 - Bindings
 - Life-cycle management

Breaking the rules...

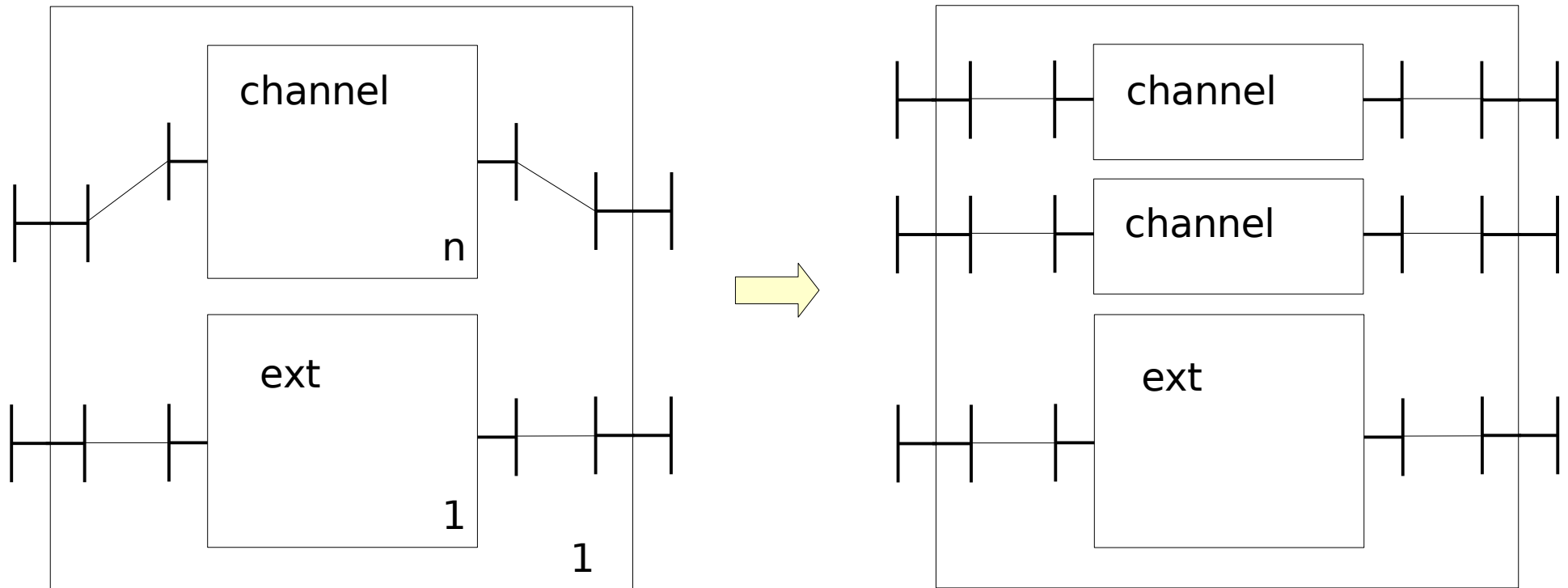
- Sometimes we don't want “transparency”
 - Wait for a value
 - `future.waitForValue();`
 - Check if value has arrived
 - `future.isAwaited();`
 - Collective interfaces
 - `listOfValues.waitForAny();`
 - `listOfValues.waitForN();`

Communication policies

- All calls are asynchronous unless stated
 - Explicit in method signature

```
interface AnInterface {  
    sync void foo(); // mandatory synchronous call  
    A bar();  
}  
  
// explicitly force synchronous call  
A a = interface.bar();  
a.waitForValue();
```

Multiple Components



Collective Communication

- Collective interfaces

```
interface printerCtrl[numberOfClients];  
bind(channel[i].printerCtrlI, this.printerCtrl[i]);
```

Multiple component

Collective interface

- Multicast interfaces

```
interface serversItf[numberOfServers];  
serversItf.flush();
```

Exceptions

- try/catch annotators
 - Exception may not leave block
 - No automatic continuation for futures carrying exceptions
 - Implicit barrier before leaving block

```
try {
    a = interface1.f();      // interface1.f() throws 'Exception'
    b = interface1.g();
    c = interface2.h(a);    // waits for 'a' & throws 'Exception'
    interface3.i(b);        // asynchronous
} catch (Exception e) {    // barrier for exceptions
    ...
}
// do something else AFTER the barrier is released
```

Indeterministic choice

- Hide computation

```
bool valid = clientItf.m();
```

```
Domain is enum {"A", "B"};
```

```
__INTERNAL(valid);
```

```
condition = __ANY(Domain);
```

```
switch (condition) {
```

```
  case A:
```

```
    // do something
```

```
    break;
```

```
  case B:
```

```
    // do something
```

```
    break;
```

```
}
```

Block

nism

What else to specify?

- Specialised Membrane behaviour
 - User defined controllers
 - Non-Functional components

Example

```
behavior component CashDeskApplication is Primitive  
{
```

```
    // primitive attributes
```

```
    CashState cashState;
```

```
    ...
```

```
    // set default policy as FIFO
```

```
    ServingPolicy = FIFO;
```

```
    // initialize the component
```

```
    init() {
```

```
        cashState = IDLE;
```

```
    }
```

Example

```
server interface cashDeskControlIf {
    void productBarcodeScannedEvent(int barcode) {
        switch (cashState) {
            case IDLE:
                // ignore signal
                break;
            case STARTED:
                Product product =
                    cashDeskConn.getProductWithStockItem(barcode);
                __INTERNAL(product);
                cashDeskConn.runningTotalChanged(product);
                break;
        }
    }
}

void paymentModeEvent(PaymentMode mode) {
    paymentMode = mode;
}
```

Summary

- Specification
 - Single language for both structural and behavioural aspects
 - Sound with component problematic
 - Can be verified (most cases)
 - Effects of sync/async calls can be checked

Summary

- Compared with pNets:
 - Higher-level, without ambiguity in what a label/sentence stands for
 - Specific to component-based software
 - (Hopefully) quicker learning curve

Summary

- Generated code:
 - Control code is correct-by-construction
 - Data-flow synchronisation!
 - Avoids writing cumbersome code
 - In particular when dealing with exceptions
 - Suitable for simulation
 - Executable

Additional thoughts

- Keep the specification close to the implementation
 - Attached to the ADL
 - Annotations / comments in Java code