



## ProActive Parallel Suite



Denis Caromel Arnaud Contes

Univ. Nice

ActiveEon

# Agenda

- ▶ **ProActive and ProActive Parallel Suite**
- ▶ Programming and Composing
  - ProActive Core
  - High Level Programming models
  - ProActive Components
- ▶ Deployment Framework
- ▶ Development Tools



# Unification of Multi-Threading and Multi-Processing

## *Multi-Threading*

Multi-Core Programming

### ▶ **SMP**

- ❑ Symmetric Multi-Processing
- ❑ Shared-Memory Parallelism

▶ **Solutions : OpenMP, pThreads, Java Threads...**

## *Multi-Processing*

Distributed programming, Grid Computing

### ▶ **MPP**

- ❑ Massively Parallel Programming or
- ❑ Message Passing Parallelism

▶ **Solutions: PVM, MPI, RMI, sockets ,...**



**ProActive**  
Programming, Composing, Deploying on the Grid



# ProActive

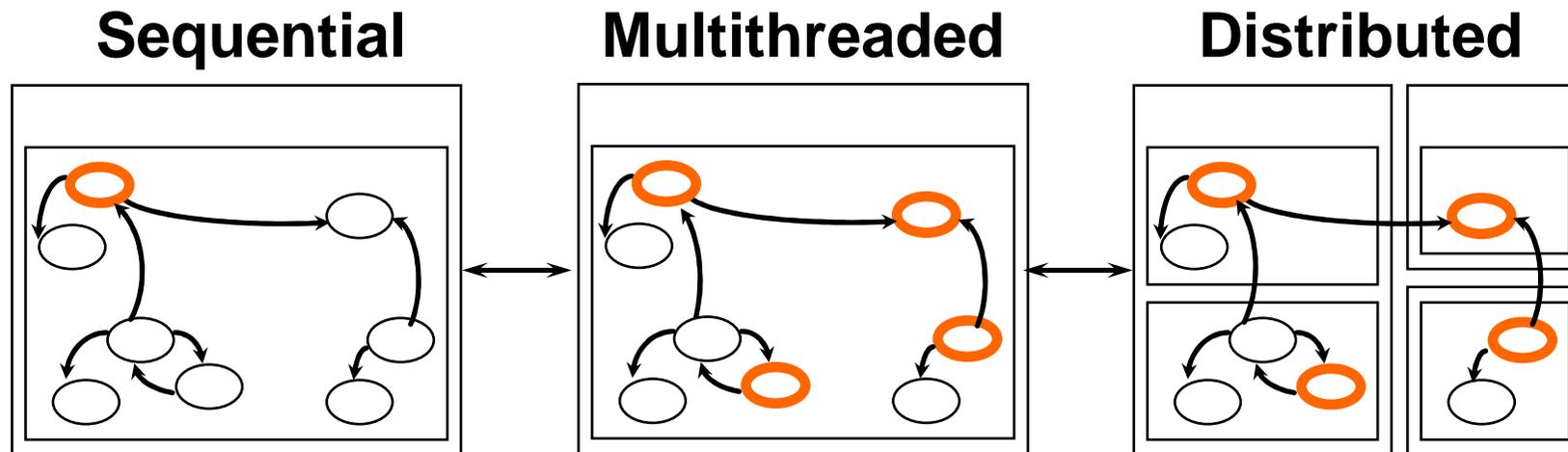
- ▶ ProActive is a JAVA **middleware** for **parallel, distributed** and **multi-threaded** computing.
- ▶ ProActive features:
  - ❑ A programming model
  - ❑ A comprehensive framework

*To simplify the programming and execution of parallel applications within multi-core processors, distributed on Local Area Network (LAN), on clusters and data centers, on intranet and Internet Grids.*



# Unification of Multi-threading and Multi-processing

## Seamless



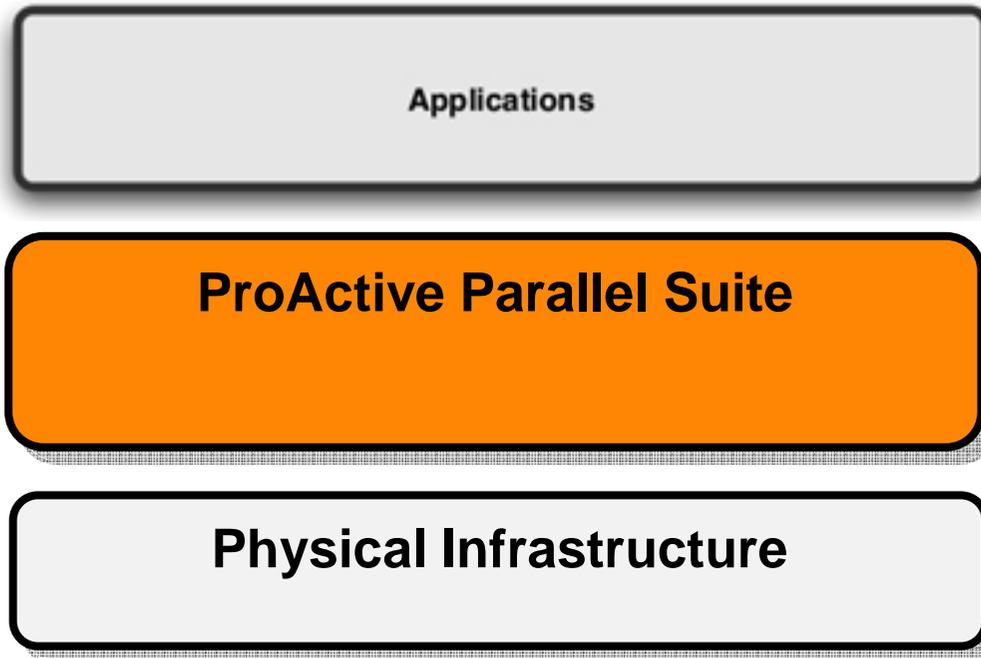
- ▶ Most of the time, activities and distribution are not known at the beginning, and change over time
- ▶ Seamless implies reuse, smooth and incremental transitions

# ProActive Parallel Suite

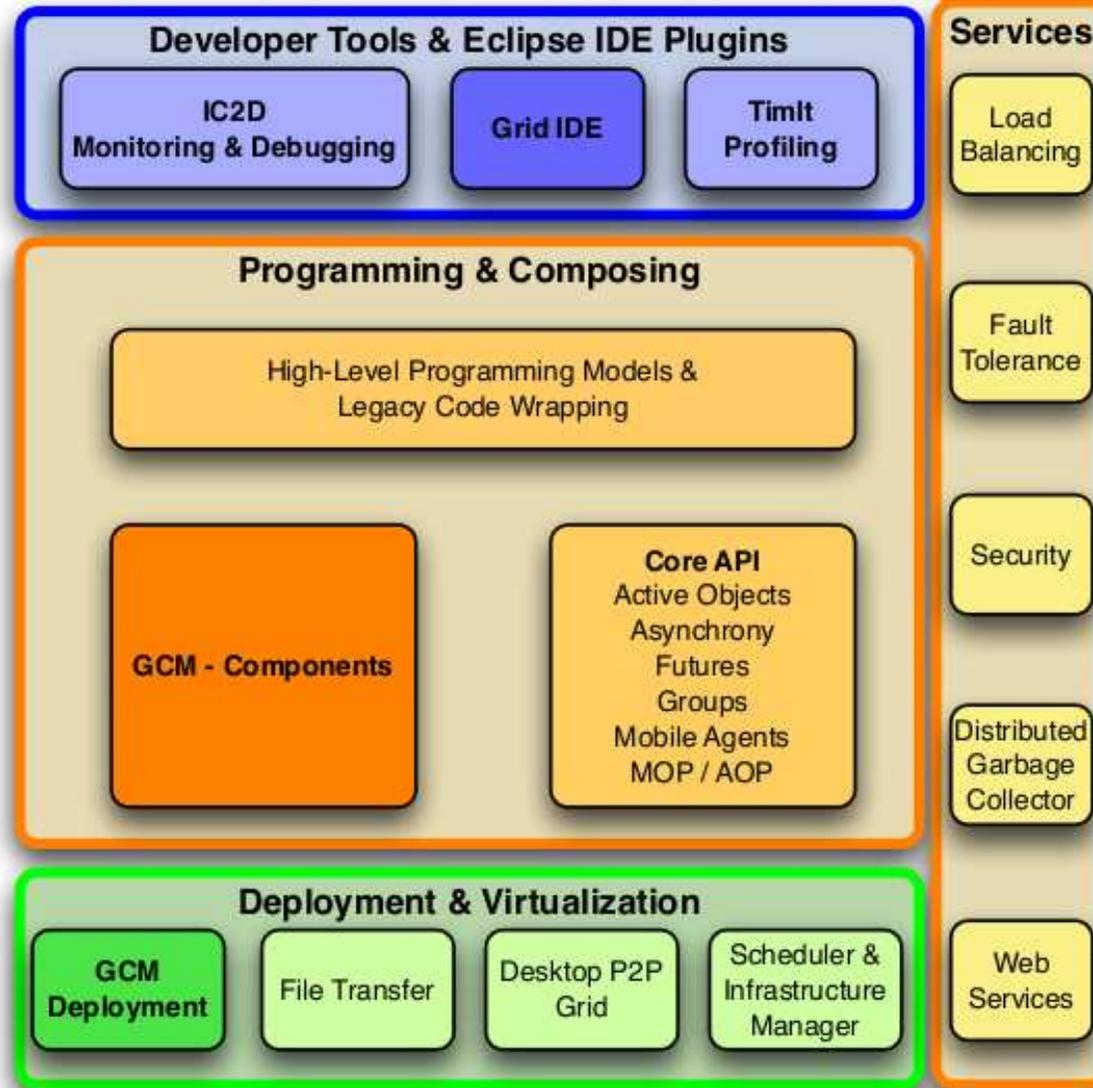
- ▶ ProActive Parallel Suite includes:
  - ❑ The ProActive middleware featuring services like:
    - Fault tolerance, Load balancing, Distributed GC, Security, WS
    - A set of parallel programming frameworks
    - A framework for deploying applications on distributed infrastructures
  - ❑ Software for scheduling applications and resource management
  - ❑ Software for monitoring and profiling of distributed applications
  - ❑ Online documentation
  - ❑ Full set of demos and examples



# ProActive Parallel Suite



# ProActive Parallel Suite



# Ways of using Proactive Parallel Suite?

- ▶ To **easily develop** parallel/distributed applications from scratch
- ▶ **Develop applications** using well-known **programming paradigms** thanks to our **high-level programming frameworks** (master-worker, Branch&Bound, SPMD, Skeletons)
- ▶ To **transform** your sequential mono-threaded application into a multi-threaded one (with minimum modification of code) and **distribute** it over the infrastructure.



# Ways of using Proactive Parallel Suite?

- ▶ To **wrap your native application** with ProActive in order to distribute it
- ▶ **Define jobs** containing your native-applications and use **ProActive to schedule** them on the infrastructure

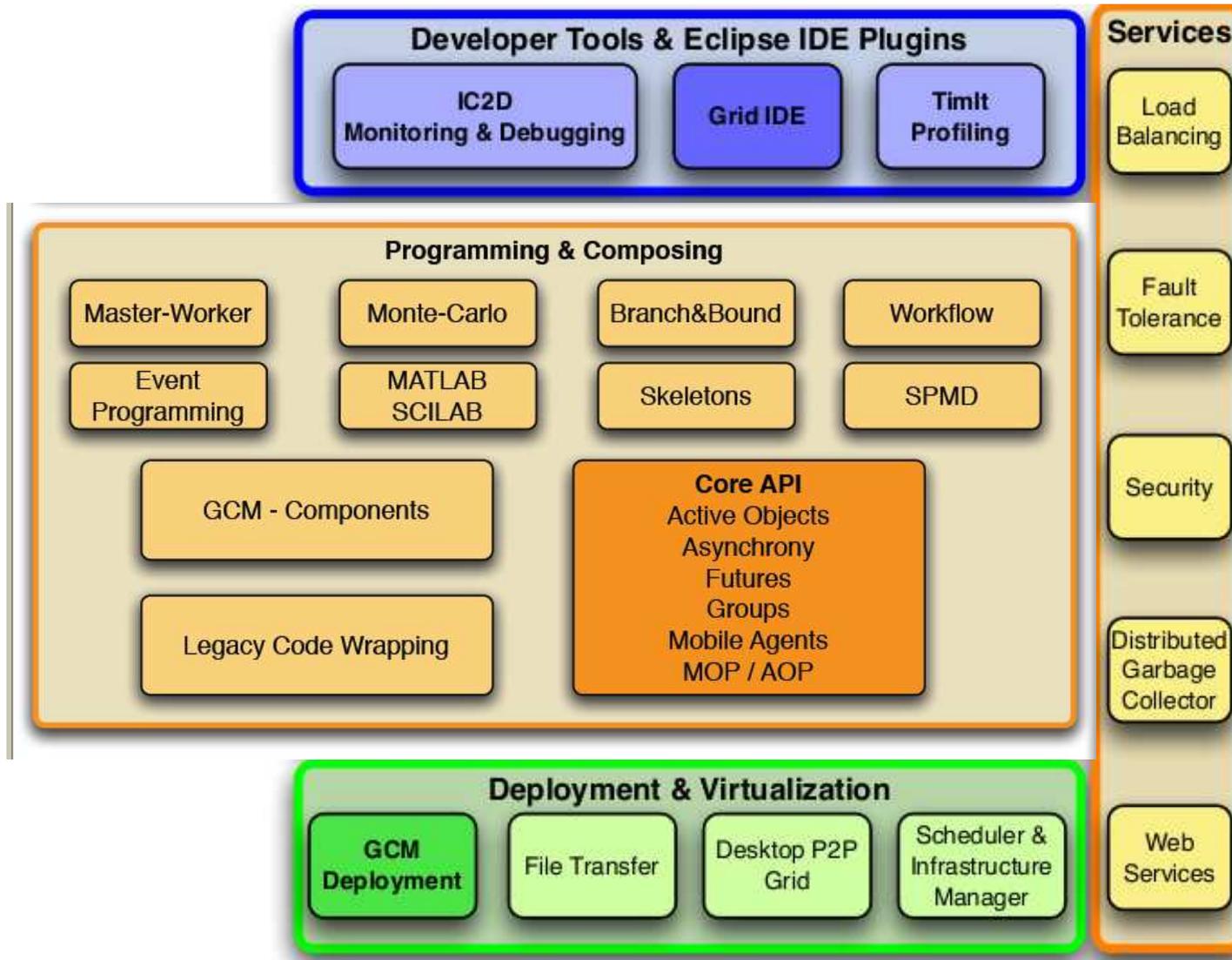


# Agenda

- ▶ ProActive and ProActive Parallel Suite
- ▶ Programming and Composing
  - ProActive Core
  - High Level Programming models
  - ProActive Components
- ▶ Deployment Framework
- ▶ Development Tools



# ProActive Parallel Suite



# ProActive Core

## ACTIVE OBJECTS



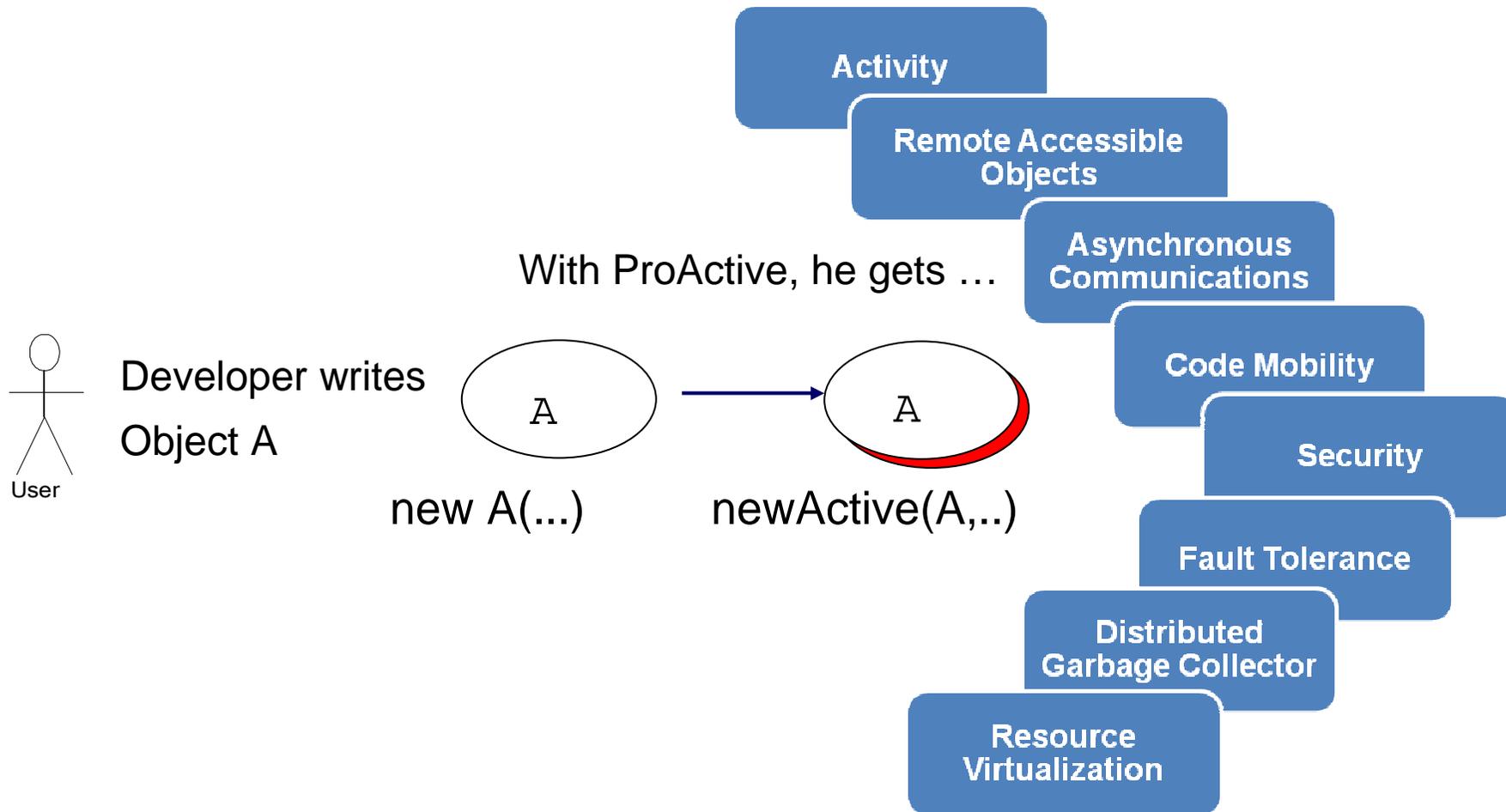
# ProActive

A 100% Java **API + Tools** for  
Parallel, Distributed Computing

- ▶ A programming model: Active Objects
  - Asynchronous Communications, Wait-By-Necessity, Groups, Mobility, Components, Security, Fault-Tolerance
- ▶ A formal model behind: Determinism (POPL'04)
  - Insensitive to application deployment
- ▶ A uniform Resource framework
  - Resource Virtualization to simplify the programming



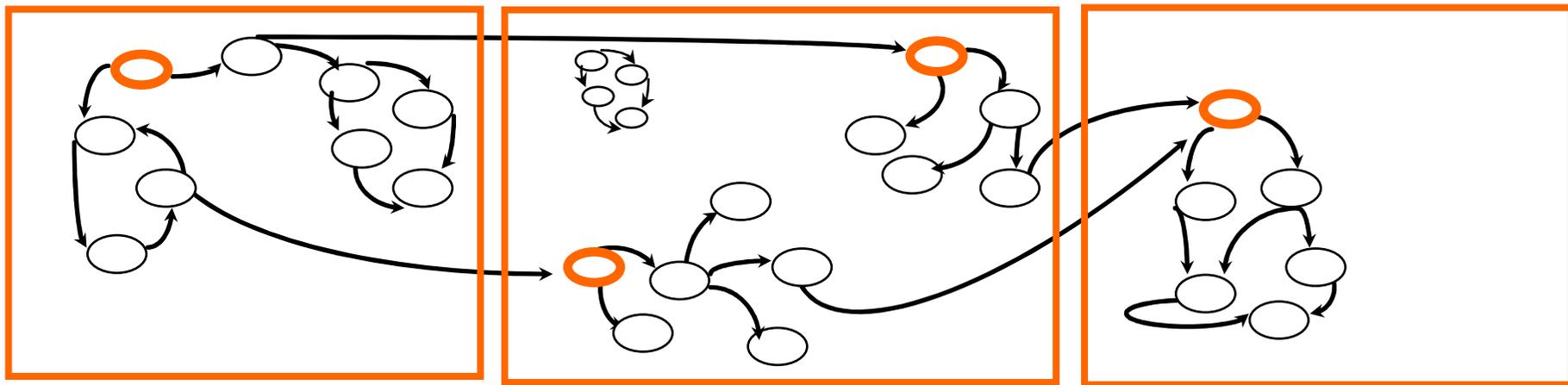
# Active Objects



# ProActive model : Basis

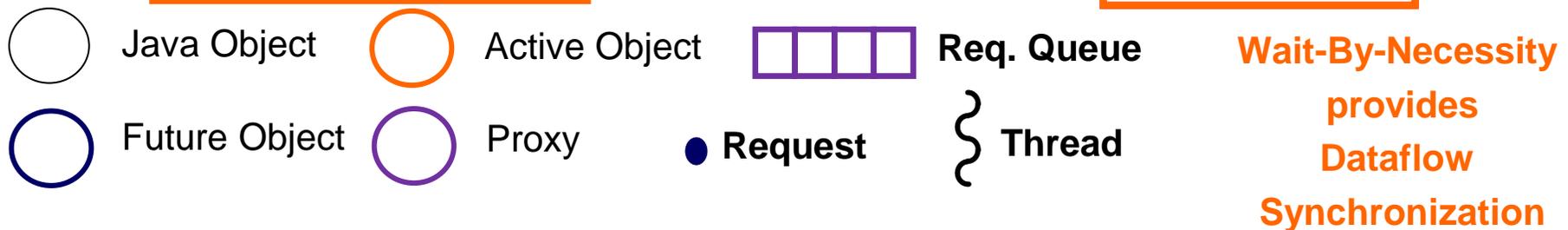
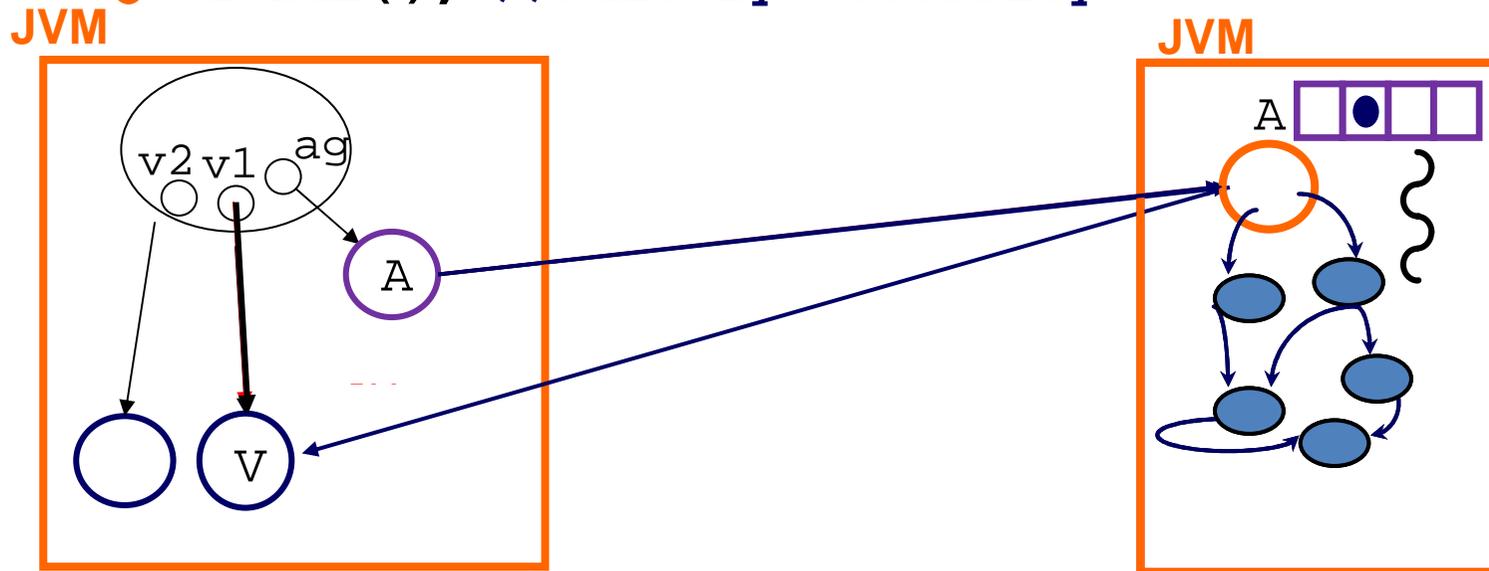
- ▶ Active objects
  - ❑ coarse-grained structuring entities (subsystems)
  - ❑ has exactly one thread.
  - ❑ owns many passive objects (Standard Java Objects, no thread)
  - ❑ No shared passive objects -- Parameters are deep-copy
- ▶ Remote Method Invocation
  - ❑ Asynchronous Communication between active objects
- ▶ Full control to serve incoming requests

JVM



# Active objects

- A ag = **newActive** ("A", [...], Node)
- V v1 = ag.foo (param);
- V v2 = ag.bar (param);
- ...
- v1.bar(); //Wait-By-Necessity



# ProActive : Reuse and seamless

- ▶ **Polymorphism** between standard and active objects
  - ❑ Type compatibility for classes (and not only interfaces)
  - ❑ Needed and done for the future objects also
- ▶ **Wait-by-necessity**: inter-object synchronization
  - ❑ Systematic, implicit and transparent futures
  - ❑ Ease the programming of synchronizations, and the reuse of routines



# Proofs in GREEK

$$\frac{(a, \sigma) \rightarrow_S (a', \sigma')}{\alpha[a; \sigma; \iota; F; R; f] \parallel P \longrightarrow \alpha[a'; \sigma'; \iota; F; R; f] \parallel P} \text{ (LOCAL)}$$

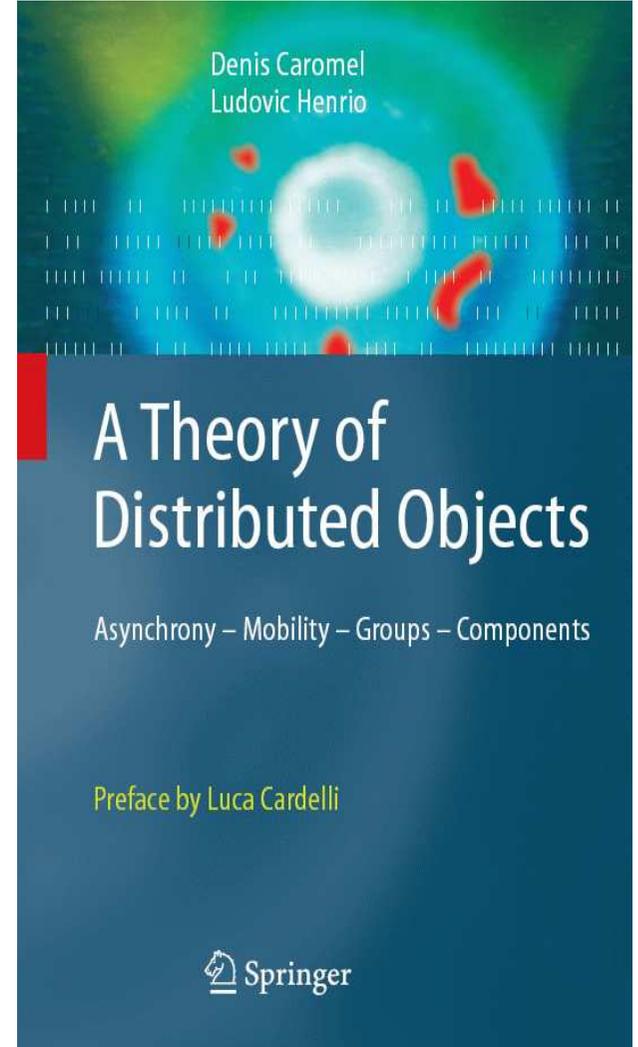
$$\frac{\begin{array}{l} \gamma \text{ fresh activity} \quad \iota' \notin \text{dom}(\sigma) \quad \sigma' = \{\iota' \mapsto AO(\gamma)\} :: \sigma \\ \sigma_\gamma = \text{copy}(\iota'', \sigma) \quad \text{Service} = (\text{if } m_j = \emptyset \text{ then } \text{FifoService} \text{ else } \iota''.m_j()) \end{array}}{\begin{array}{l} \alpha[\mathcal{R}[\text{Active}(\iota'', m_j)]; \sigma; \iota; F; R; f] \parallel P \\ \longrightarrow \alpha[\mathcal{R}[\iota']; \sigma'; \iota; F; R; f] \parallel \gamma[\text{Service}; \sigma_\gamma; \iota''; \emptyset; \emptyset; \emptyset] \parallel P \end{array}} \text{ (NEWACT)}$$

$$\frac{\begin{array}{l} \sigma_\alpha(\iota) = AO(\beta) \quad \iota'' \notin \text{dom}(\sigma_\beta) \quad f_i^{\alpha \rightarrow \beta} \text{ new future} \quad \iota_f \notin \text{dom}(\sigma_\alpha) \\ \sigma'_\beta = \text{Copy\&Merge}(\sigma_\alpha, \iota'; \sigma_\beta, \iota'') \quad \sigma'_\alpha = \{\iota_f \mapsto \text{fut}(f_i^{\alpha \rightarrow \beta})\} :: \sigma_\alpha \end{array}}{\begin{array}{l} \alpha[\mathcal{R}[\iota.m_j(\iota')]; \sigma_\alpha; \iota_\alpha; F_\alpha; R_\alpha; f_\alpha] \parallel \beta[a_\beta; \sigma_\beta; \iota_\beta; F_\beta; R_\beta; f_\beta] \parallel P \longrightarrow \\ \alpha[\mathcal{R}[\iota_f]; \sigma'_\alpha; \iota_\alpha; F_\alpha; R_\alpha; f_\alpha] \parallel \beta[a_\beta; \sigma'_\beta; \iota_\beta; F_\beta; R_\beta :: [m_j; \iota''; f_i^{\alpha \rightarrow \beta}]; f_\beta] \parallel P \end{array}} \text{ (REQUEST)}$$

$$\frac{R = R' :: [m_j; \iota_r; f'] :: R'' \quad m_j \in M \quad \forall m \in M, m \notin R'}{\alpha[\mathcal{R}[\text{Serve}(M)]; \sigma; \iota; F; R; f] \parallel P \longrightarrow \alpha[\iota.m_j(\iota_r) \uparrow f, \mathcal{R}[\Box]; \sigma; \iota; F; R' :: R''; f'] \parallel P} \text{ (SERVE)}$$

$$\frac{\iota' \notin \text{dom}(\sigma) \quad F' = F :: \{f \mapsto \iota'\} \quad \sigma' = \text{Copy\&Merge}(\sigma, \iota; \sigma, \iota')}{\alpha[\iota \uparrow (f', a); \sigma; \iota; F; R; f] \parallel P \longrightarrow \alpha[a; \sigma'; \iota; F'; R; f'] \parallel P} \text{ (ENDSERVICE)}$$

$$\frac{\sigma_\alpha(\iota) = \text{fut}(f_i^{\gamma \rightarrow \beta}) \quad F_\beta(f_i^{\gamma \rightarrow \beta}) = \iota_f \quad \sigma'_\alpha = \text{Copy\&Merge}(\sigma_\beta, \iota_f; \sigma_\alpha, \iota)}{\begin{array}{l} \alpha[a_\alpha; \sigma_\alpha; \iota_\alpha; F_\alpha; R_\alpha; f_\alpha] \parallel \beta[a_\beta; \sigma_\beta; \iota_\beta; F_\beta; R_\beta; f_\beta] \parallel P \longrightarrow \\ \alpha[a_\alpha; \sigma'_\alpha; \iota_\alpha; F_\alpha; R_\alpha; f_\alpha] \parallel \beta[a_\beta; \sigma_\beta; \iota_\beta; F_\beta; R_\beta; f_\beta] \parallel P \end{array}} \text{ (REPLY)}$$



ProActive Core

# MIGRATION: MOBILE AGENTS



# Mobile Agents: Migration

- ▶ The active object migrates with:
  - ❑ its state
  - ❑ all pending requests
  - ❑ all its passive objects
  - ❑ all its future objects
- ▶ Automatic update of references:
  - ❑ requests (remote references remain valid)
  - ❑ replies (its previous queries will be fulfilled)
- ▶ Migration is initiated by the active object itself
- ▶ Can be initiated from outside through any public method

# Migration Strategies

## ▶ Forwarders

- Migration creates a chain of forwarders
- A forwarder is left at the old location to forward requests to the new location
- Tensioning: shortcut the forwarder chains by notifying the sender of the new location of the target (transparently)

## ▶ Location Server

- A server (or a set of servers) keeps track of the location of all active objects
- Migration updates the location on the server

## ▶ Mixed ( Forwarders / Local Server )

- Limit the size of the chain up to a fixed size

# ProActive Core

## PROACTIVE GROUPS



# ProActive Groups

- Manipulate groups of Active Objects, in a simple and typed manner:
  - ➔ Typed and polymorphic Groups of local and remote objects
  - ➔ Dynamic generation of group of results
  - ➔ Language centric, Dot notation
- Be able to express high-level collective communications (like in MPI):
  - broadcast,
  - scatter, gather,
  - all to all

```
A ag=(A)ProActiveGroup.newGroup(«A», {{p1}, ...}, {Nodes, ..});  
V v = ag.foo(param);  
v.bar();
```

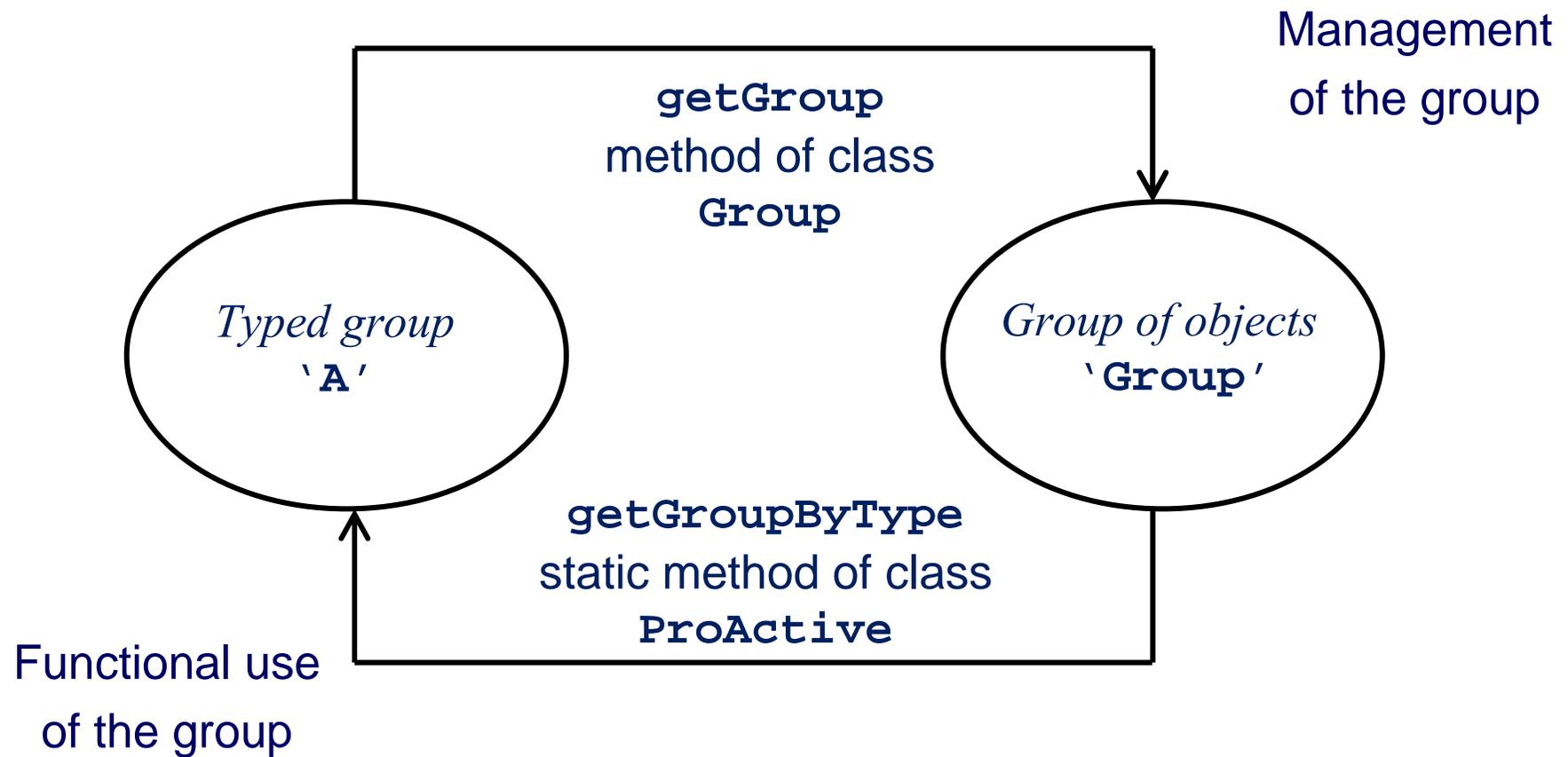


# ProActive Groups

- ▶ Group Members
  - ❑ Active Objects
  - ❑ POJO
  - ❑ Group Objects
- ▶ Hierarchical Groups
- ▶ Based on the ProActive communication mechanism
  - ❑ Replication of N 'single' communications
  - ❑ Parallel calls within a group (latency hiding)
- ▶ Polymorphism
  - ❑ Group typed with member's type



# Two Representations Scheme



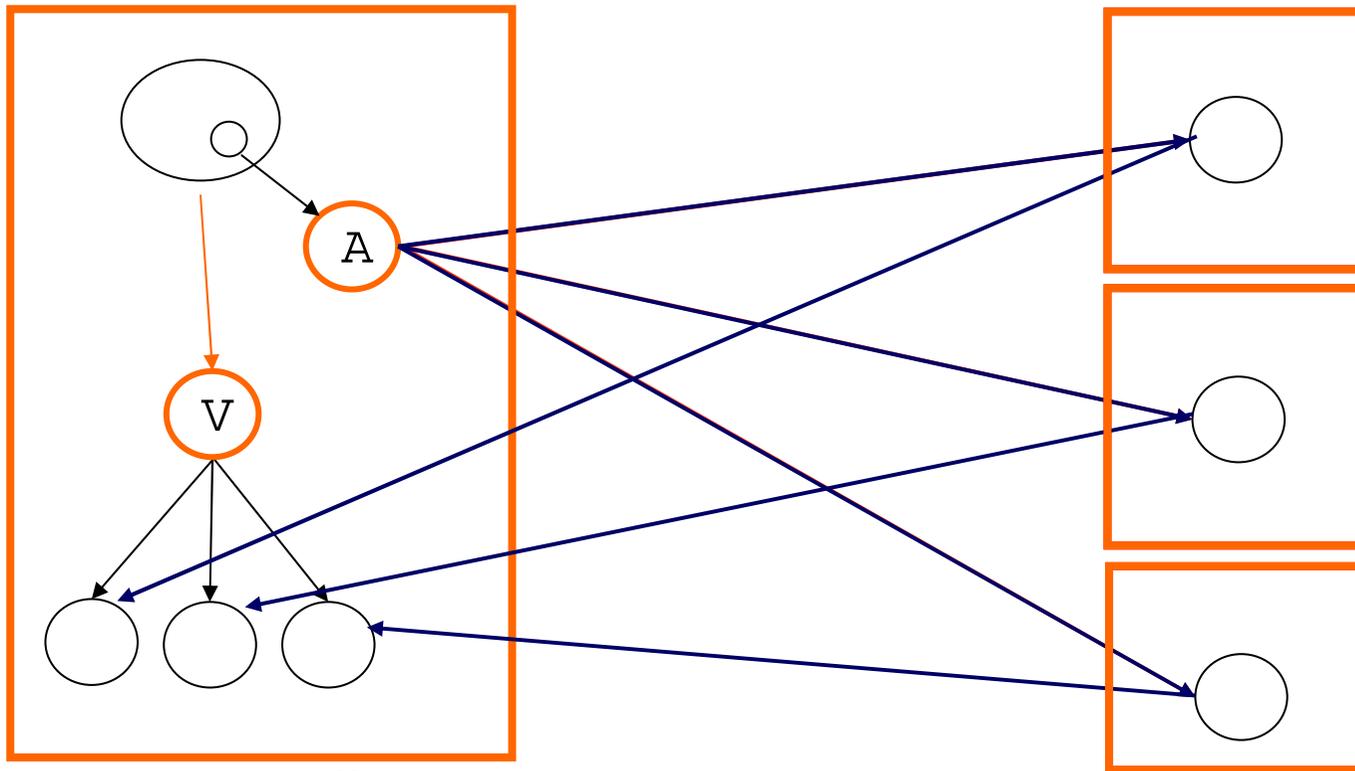
# Creating AO and Groups

● A ag = **newGroup** ("A", [...], Node[])

● V v = ag.foo(param);

● ● ● ●

JVM ● v.bar(); //Wait-by-necessity



○ Typed Group ○ Java or Active Object

# Typed Group as Result of Group Communication

- ▶ Ranking Property:
  - ❑ Dynamically built and updated
    - `B groupB = groupA.foo();`
  - ❑ Ranking property: order of result group members = order of called group members
- ▶ Explicit Group Synchronization Primitive:
  - ❑ Explicit wait
    - `ProActiveGroup.waitOne(groupB);`
    - `ProActiveGroup.waitAll(groupB);`
  - ❑ Predicates
    - `noneArrived`
    - `kArrived`
    - `allArrived, ...`



ProActive Core

# FAULT TOLERANCE SERVICE



# Fault-tolerance in ProActive

- ▶ Restart an application from latest valid checkpoint
  - ❑ Avoid cost of restarting from scratch
- ▶ Fault-tolerance is non intrusive
  - ❑ set in a deployment descriptor file
  - ❑ Fault-tolerance service attached to **resources**
  - ❑ **No source code alteration**
    - Protocol selection , Server(s) location, Checkpoint period



# Fault-tolerance in ProActive

- ▶ Rollback-Recovery fault-tolerance
  - ❑ After a failure, revert the system state back to some earlier and correct version
  - ❑ Based on periodical **checkpoints** of the active objects
  - ❑ Stored on a **stable** server
- ▶ Two protocols are implemented
  - ❑ Communication Induced Checkpointing (CIC)
    - + **Lower** failure free overhead
    - Slower recovery
  - ❑ Pessimistic Message Logging (PML)
    - Higher failure free overhead
    - + **Faster** recovery
- ▶ Transparent and non intrusive

# Built-in Fault-tolerance Server

- ▶ Fault-tolerance is based on a global server
- ▶ This server is provided by the library, with
  - ❑ Checkpoint storage
  - ❑ Failure detection
    - Detects fail-stop failures
  - ❑ Localization service
    - Returns the new location of a failed object
  - ❑ Resource management service
    - Manages a set of nodes on which restart failed objects



# ProActive Core

## SECURITY SERVICE



# ProActive Security Framework

## Issue

Access control, communication privacy and integrity

- ▶ Unique features
  - ❑ SPKI: Hierarchy of certificates
  - ❑ No security related code in the application source code
  - ❑ Declarative security language
  - ❑ Security at user- and administrator-level
  - ❑ Security context dynamic propagation
- ▶ Configured within deployment descriptors
  - ❑ Easy to adapt according the actual deployment



# ProActive Core

## WEB SERVICES



# Web Service Integration

## ▶ Aim

- Turn active objects and components interfaces into Web Services

→ interoperability with any foreign language or any foreign technology.

## ▶ API

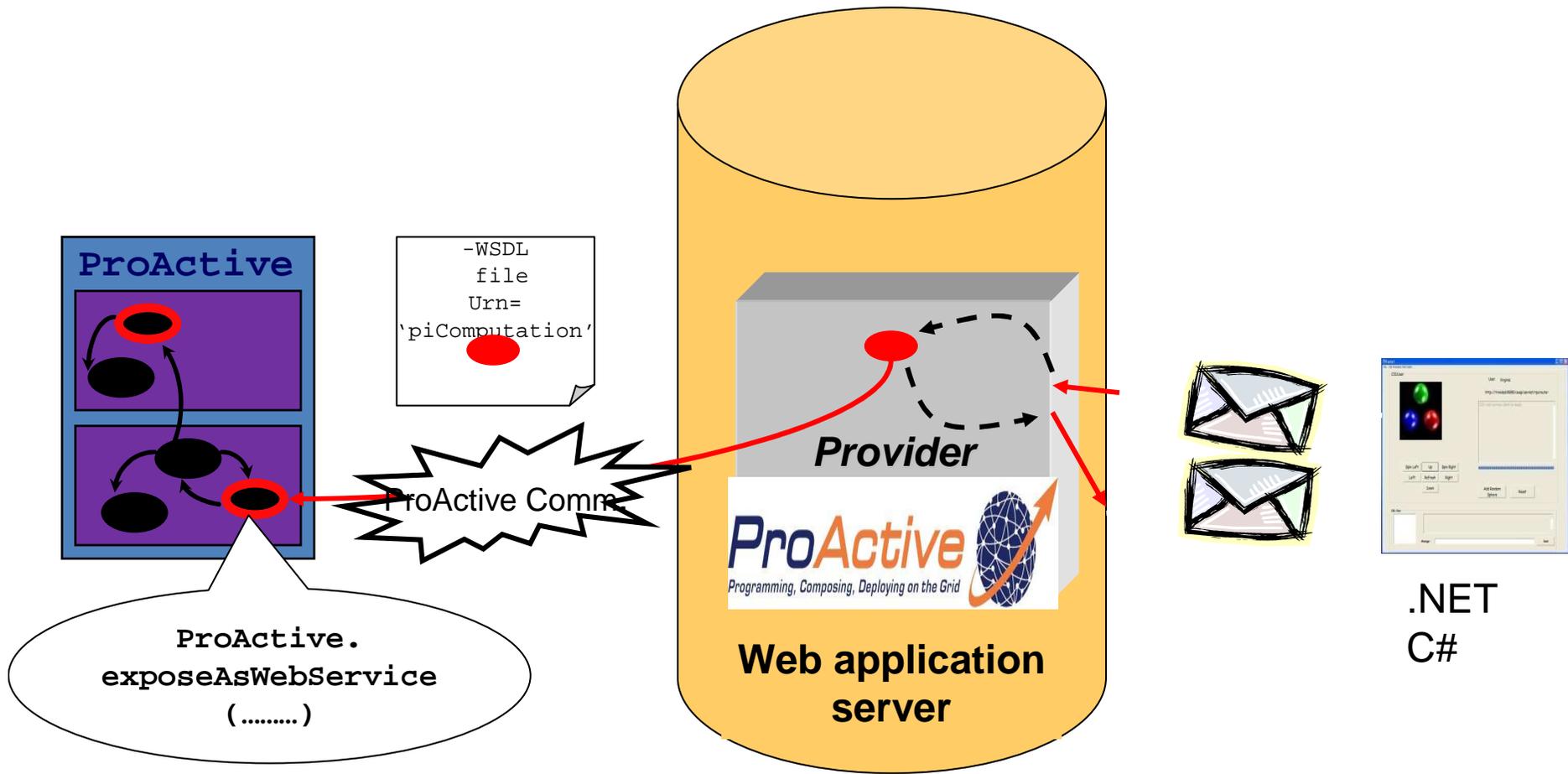
- Expose an active object as a web Service (the user can choose the methods he wants to expose)

- `exposeAsWebService(Object o, String url, String urn, String [] methods );`

- Expose component's interfaces as web services

- `exposeComponentAsWebService(Component component, String url, String componentName );`





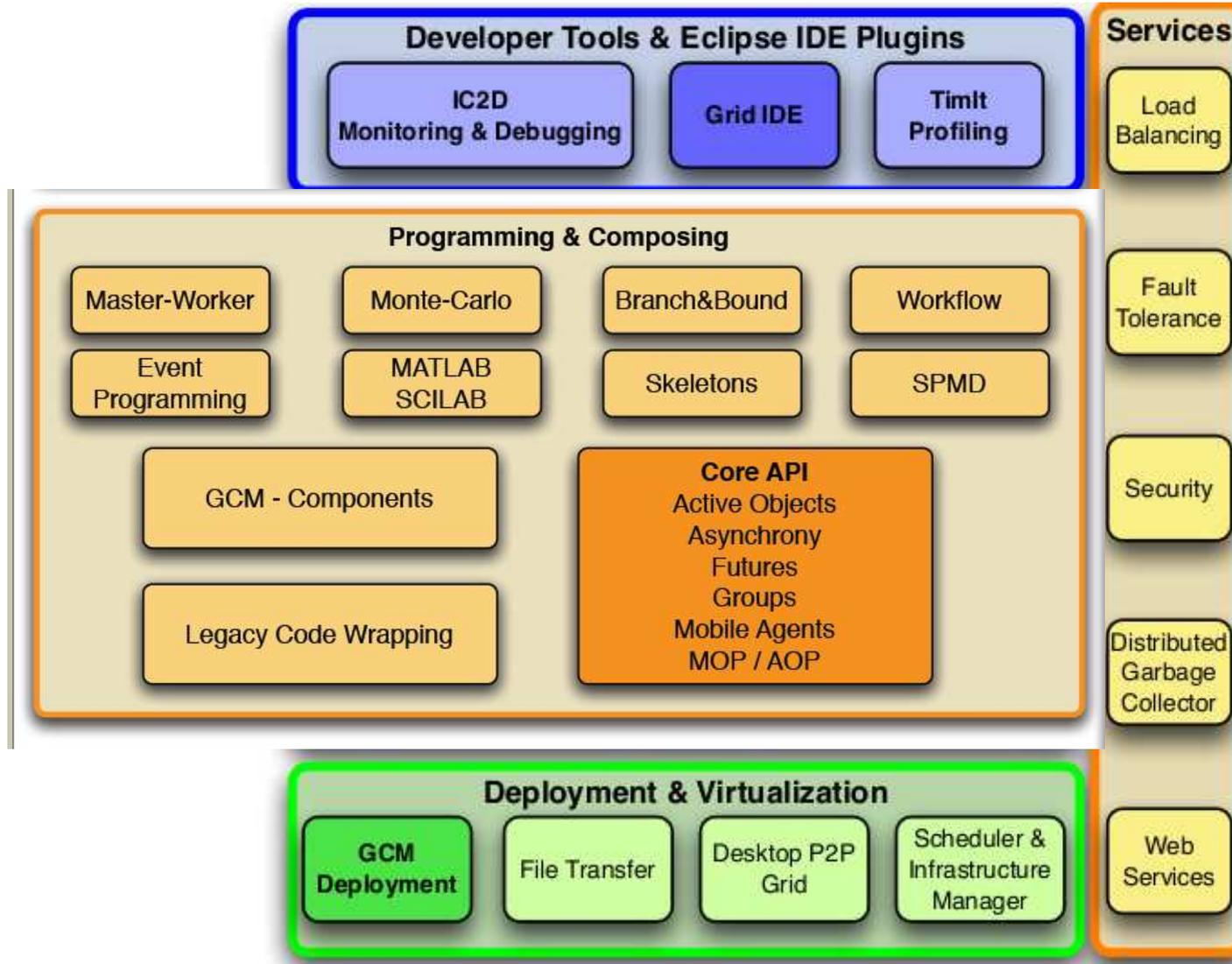
### 3. Client Call Introduce exposeAsWebService ()

# Agenda

- ▶ ProActive and ProActive Parallel Suite
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# ProActive Parallel Suite





# High Level Programming models

## Master-Worker Framework

# Motivations

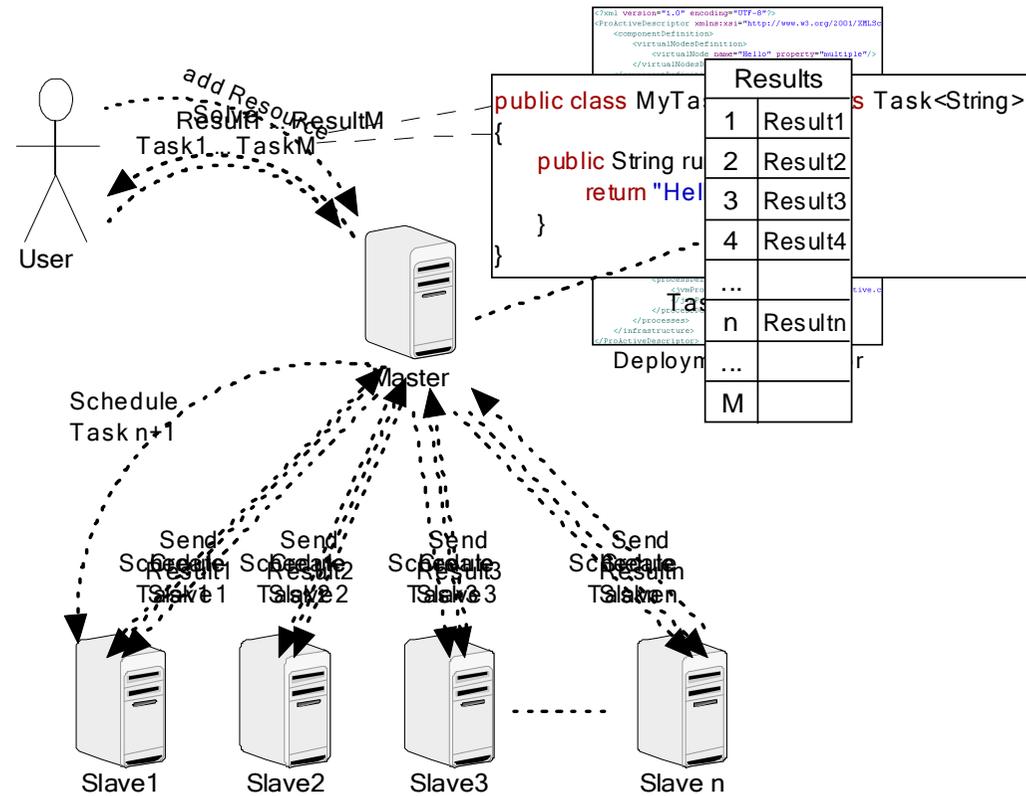
- ▶ Embarrassingly parallel problems : simple and frequent model
- ▶ Write embarrassingly parallel applications with ProActive :
  - ❑ May require a sensible amount of code (fault-tolerance, load-balancing, ...).
  - ❑ Requires understanding of ProActive concepts ( Futures, Stubs, Group Communication )



# Goals of the M/W API

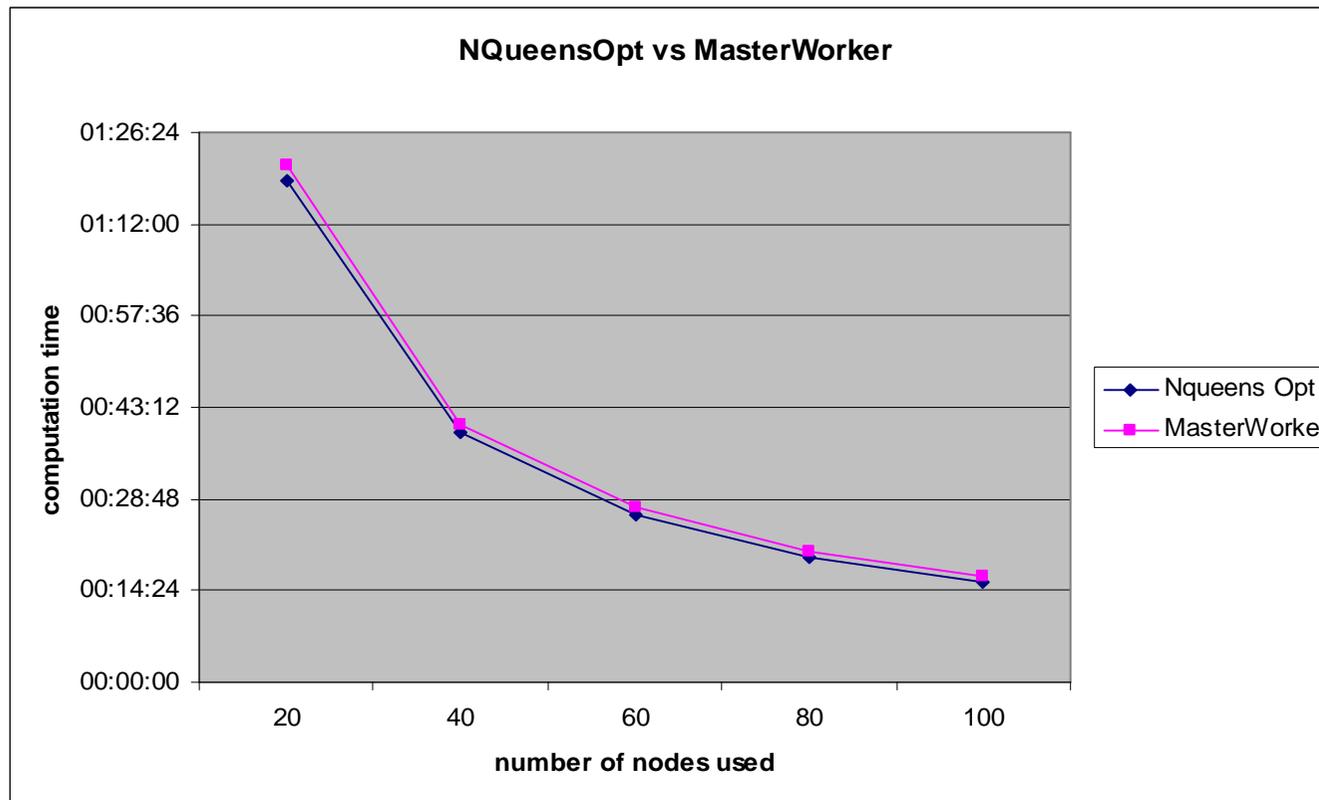
- ▶ Provide a easy-to use framework for solving embarrassingly parallel problems:
  - ❑ Simple Task definition
  - ❑ Simple API interface (few methods)
  - ❑ Simple & efficient solution gathering mechanism
- ▶ Provide automatic fault-tolerance and load-balancing mechanism
- ▶ Hide ProActive concepts from the user

# How does it work?



# Comparison between specific implementation and M/W

- ▶ Experiments with nQueens problem
- ▶ Runs up to 25 nodes





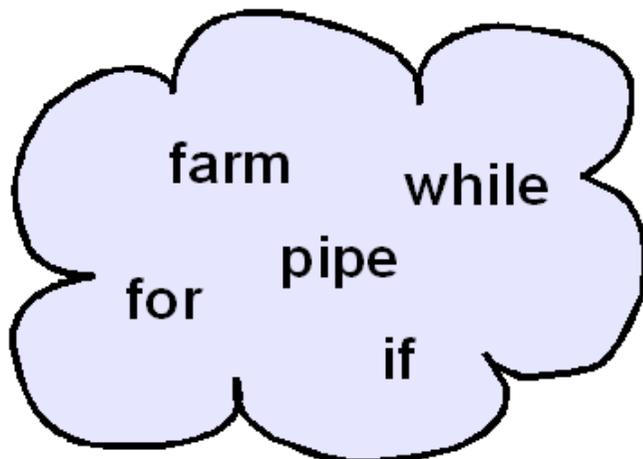
# High Level Programming models

## Skeletons Framework

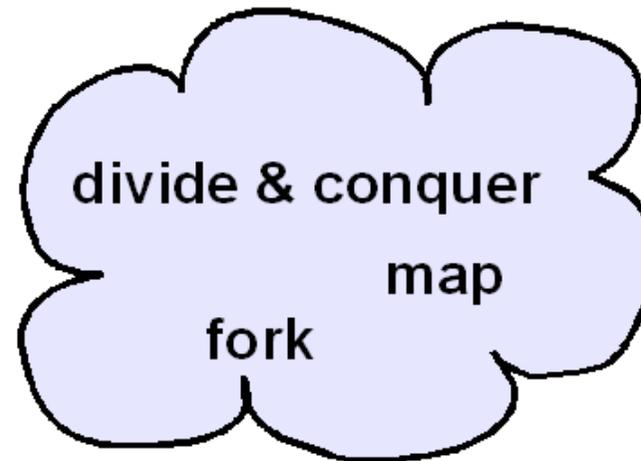
# Algorithmic Skeletons

- ▶ High Level Programming Model
- ▶ Hides the complexity of parallel/distributed programming.
- ▶ Exploits nestable parallelism patterns

Task Parallelism

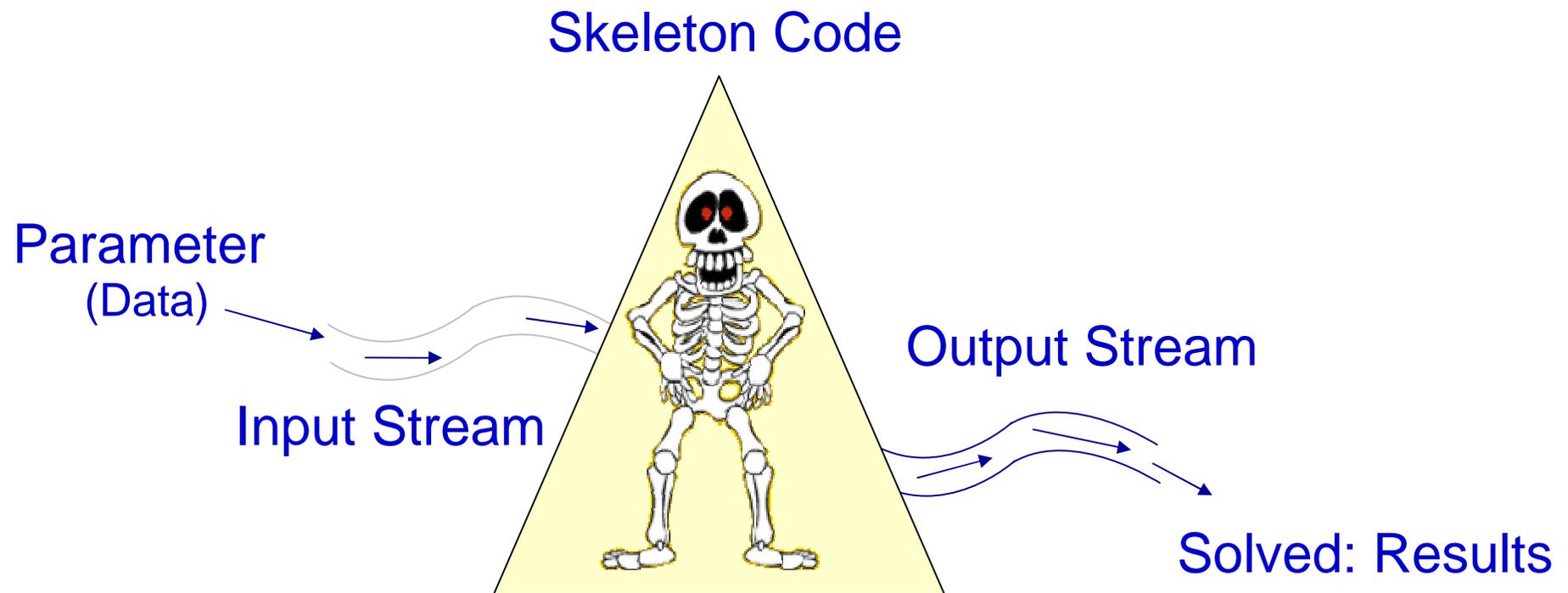


Data Parallelism



# Skeletons Big Picture

- ▶ Parameters/Results are passed through streams
- ▶ Streams are used to connect skeletons (CODE)



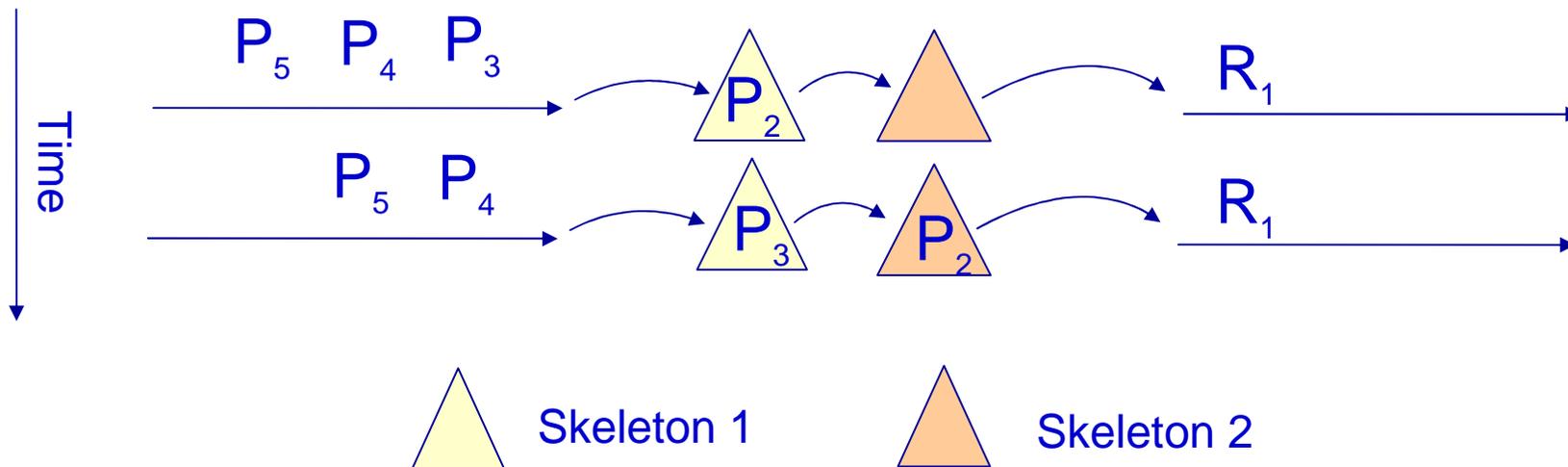
# Pipe Skeleton

- ▶ Represents computation by stages.
- ▶ Stages are computed in parallel for different parameters.

Input Stream

Execute Skeleton

Output Stream



# Simple use of Pipe skeleton

```
Skeleton<Eggs, Mix> stage1 =  
    new Seq<Eggs, Mix>(new Apprentice());
```

```
Skeleton<Mix, Omelette> stage2 =  
    new Seq<Mix, Omelette>(new Chef());
```

```
Skeleton<Eggs, Omelette> kitchen =  
    new Pipe<Eggs, Omelette>(stage1, stage2);
```



# High Level Programming models

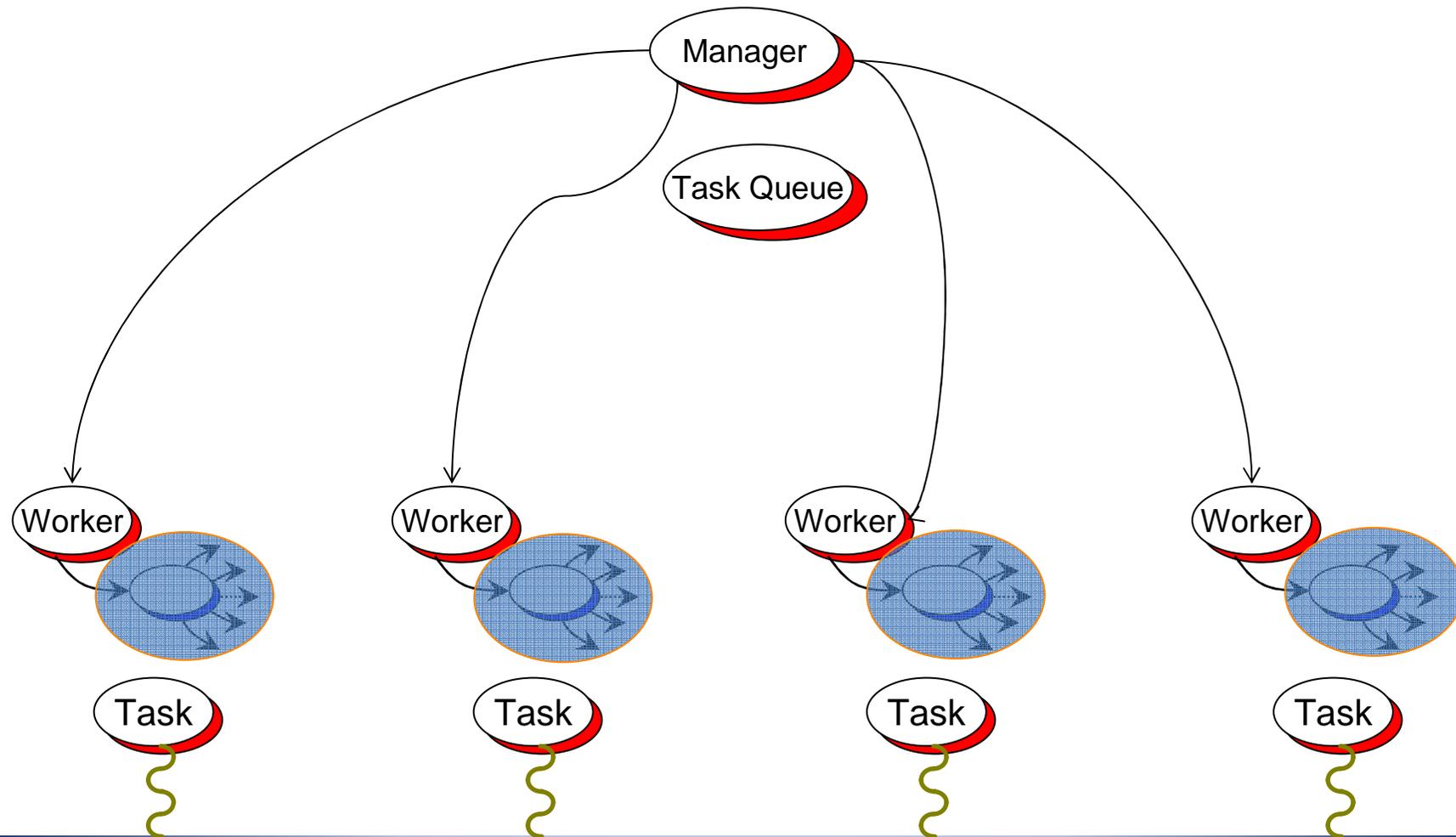
## Branch-and-Bound Framework

# Branch & Bound API (BnB)

- ▶ Provide a high level programming model for solving BnB problems:
  - ❑ manages task distribution and provides task communications
  
- ▶ Features:
  - ❑ Dynamic task split
  - ❑ Automatic result gather
  - ❑ Broadcasting best current result
  - ❑ Automatic backup (configurable)



# Global Architecture : M/W + Full connectivity





# High Level Programming models

OO-SPMD

# Object-Oriented Single Program Multiple Data

## ▶ Motivation

- Cluster / GRID computing
- SPMD programming for many numerical simulations
- Use enterprise technology (Java, Eclipse, etc.) for Parallel Computing

## ▶ Able to express most of MPI's

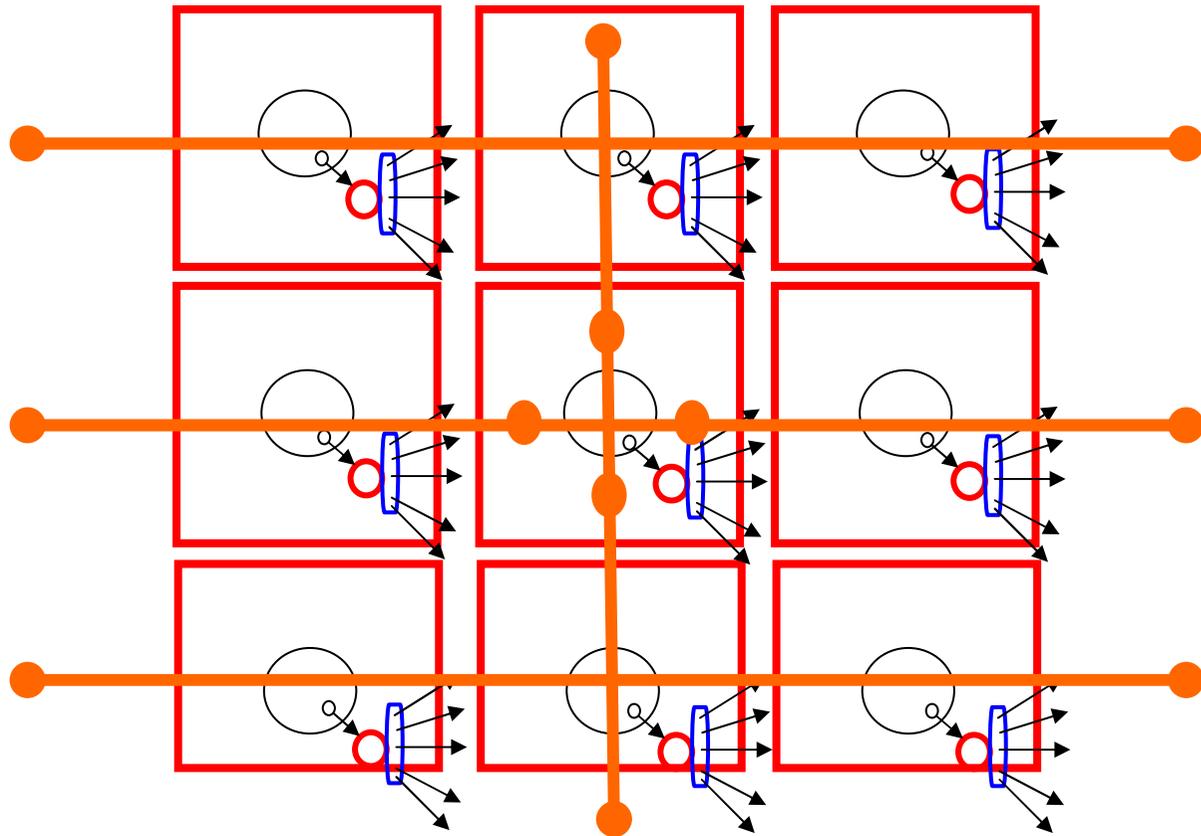
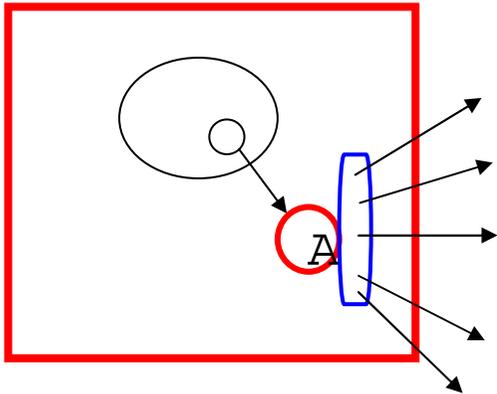
- Collective Communications (broadcast, gathercast, scattercast,..)
- Barriers
- Topologies

## ▶ With a small object-oriented API



# Execution example

- A ag = `newSPMDGroup ("A", [...], VirtualNode)`
  - // In each member
  - `myGroup.barrier ("2D"); // Global Barrier`
  - `myGroup.barrier ("vertical"); // Any Barrier`
  - `myGroup.barrier ("north","south","east","west");`

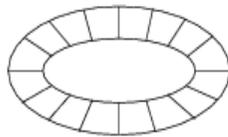


# Topologies

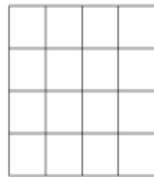
- ▶ Topologies are typed groups
- ▶ Customizable
- ▶ Define neighborhood



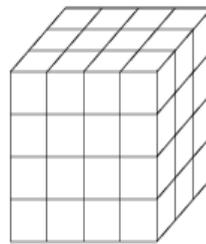
Line



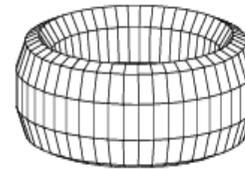
Ring



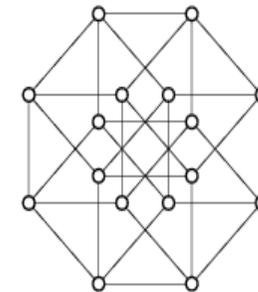
Plan



Cube



Torus



Hypercube

```
Plan plan = new Plan(groupA, Dimensions);  
Line line = plan.getLine(0);
```

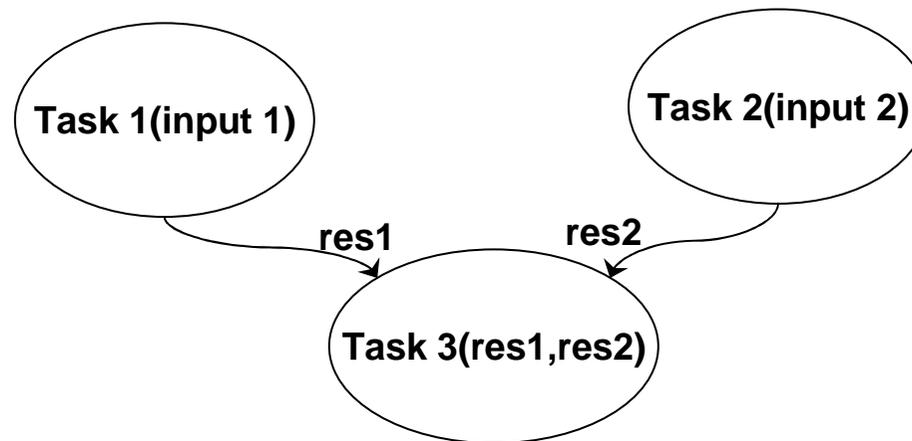


# High Level Programming models

## Scheduler

# Programming with flows of tasks

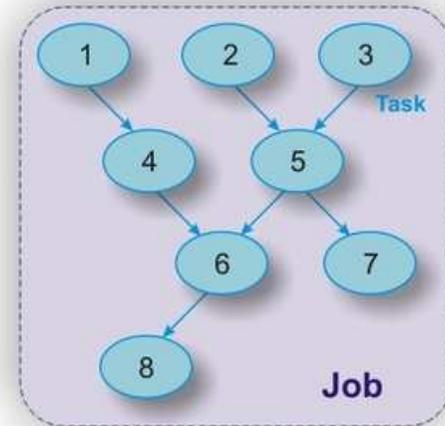
- ▶ Program an application as an ordered tasks set
  - ❑ **Logical flow** : Tasks execution are orchestrated
  - ❑ **Data flow** : Results are forwarded from ancestor tasks to their children as parameter



- ▶ The task is the smallest execution unit
- ▶ Two types of tasks:
  - ❑ Standard Java
  - ❑ Native, i.e. any third party application

# Defining and running jobs with ProActive

- ▶ A workflow application is a job
  - ❑ a set of tasks which can be executed according to a dependency tree
- ▶ Rely on ProActive Scheduler only
- ▶ Java or XML interface
  - ❑ Dynamic job creation in **Java**
  - ❑ Static description in **XML**
- ▶ Task failures are handled by the ProActive Scheduler
  - ❑ A task can be automatically re-started or not (with a user-defined bound)
  - ❑ Dependant tasks can be aborted or not
  - ❑ The finished job contains the cause exceptions as results if any

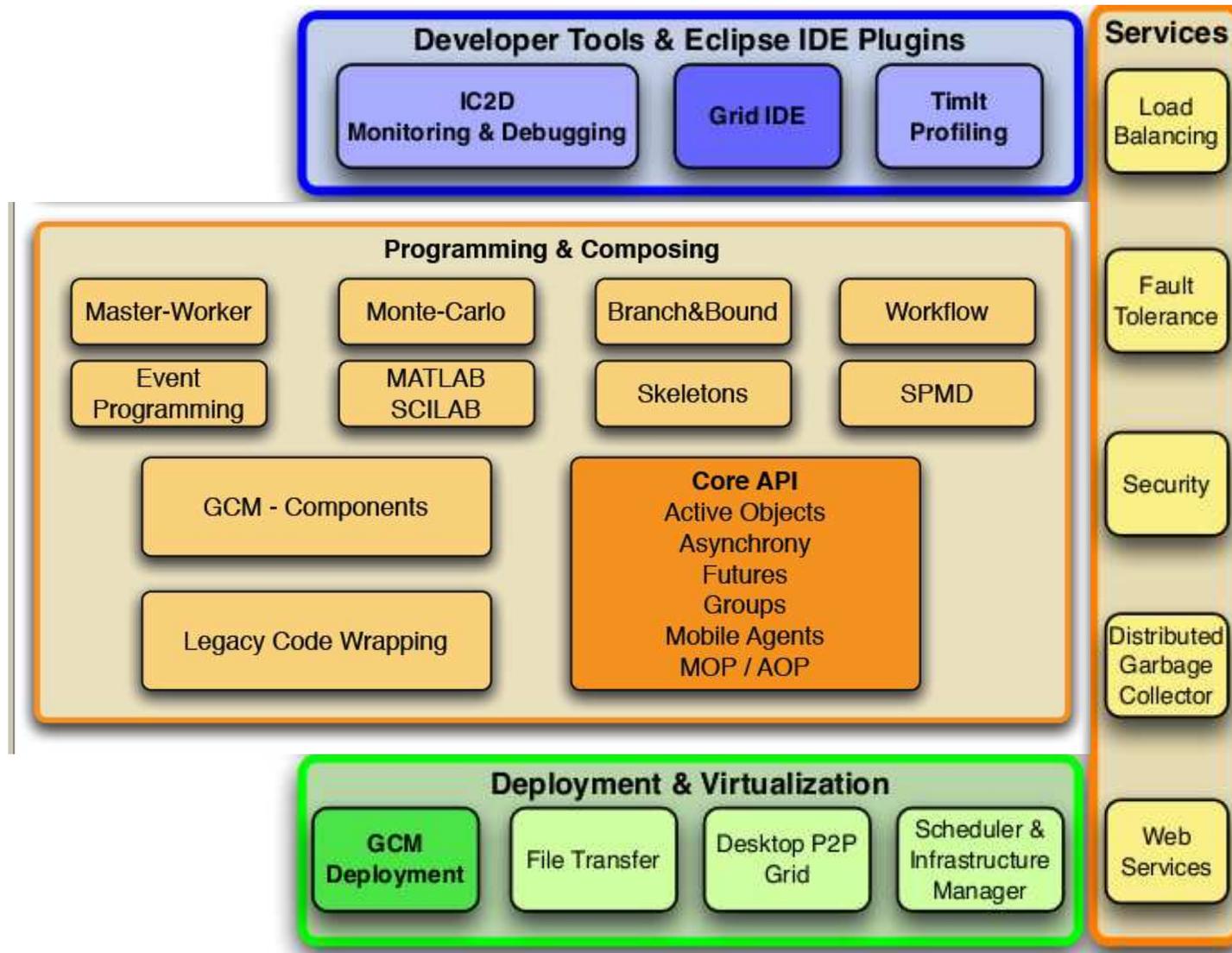


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# ProActive Parallel Suite



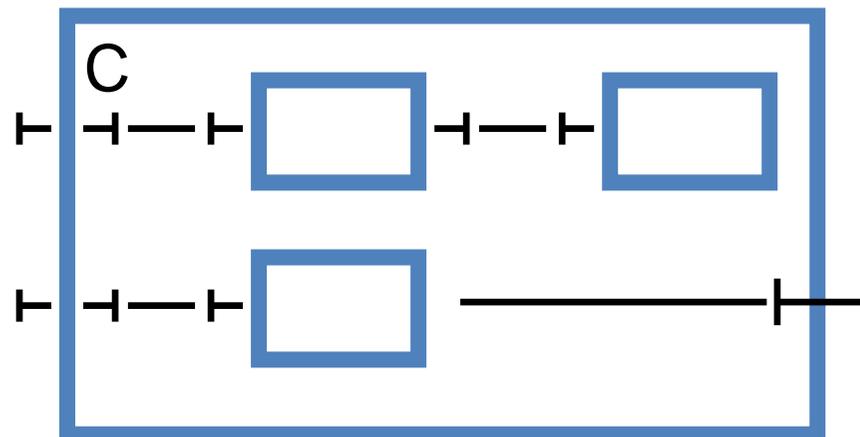
# A framework for Grid components

- ▶ Facilitating the design and implementation of complex distributed systems
- ▶ Leveraging the ProActive library  
ProActive components benefit from underlying features
- ▶ Allowing reuse of legacy components (e.g. MPI)
- ▶ Providing tools for defining, assembling and monitoring distributed components



# Component - What is it ?

- ▶ A component in a given infrastructure is:  
a **software module**,  
with a **standardized description** of what it **needs** and **provides**,  
to be manipulated by **tools** for **Composition** and **Deployment**



# ProActive Component Definition

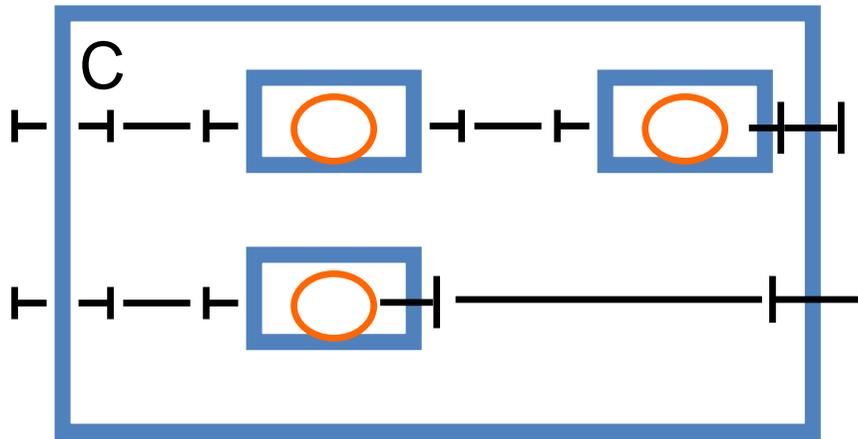
- ▶ A component is:
  - ❑ Formed from one (or several) Active Object
  - ❑ Executing on one (or several) JVM
  - ❑ Provides a set of server ports: Java Interfaces
  - ❑ Uses a set of client ports: Java Attributes
  - ❑ Point-to-point or Group communication between components
- ▶ Hierarchical:
  - ❑ Primitive component: define with Java code and a descriptor
  - ❑ Composite component: composition of primitive + composite
  - ❑ Parallel component: multicast of calls in composites
- ▶ Descriptor:
  - ❑ XML definition of primitive and composite (ADL)
  - ❑ Virtual nodes capture the deployment capacities and needs
- ▶ Virtual Node:
  - ❑ a very important abstraction for GRID components



# Components for the GRID

○ An activity, a process, ...  
potentially in its own JVM

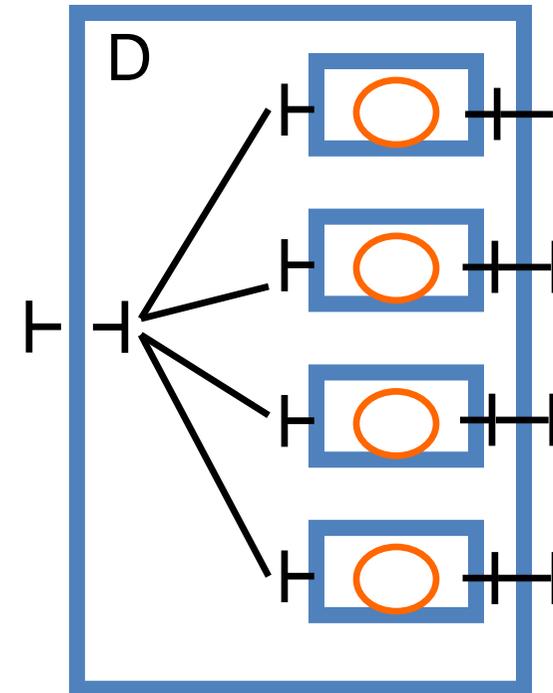
┌ ○ ─┘ 1. Primitive component



2. Composite component

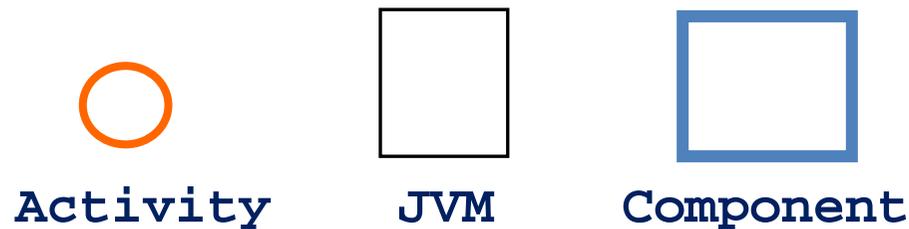
**Composite:** Hierarchical, and  
Distributed over machines

**Parallel:** Composite  
+ Broadcast (group)

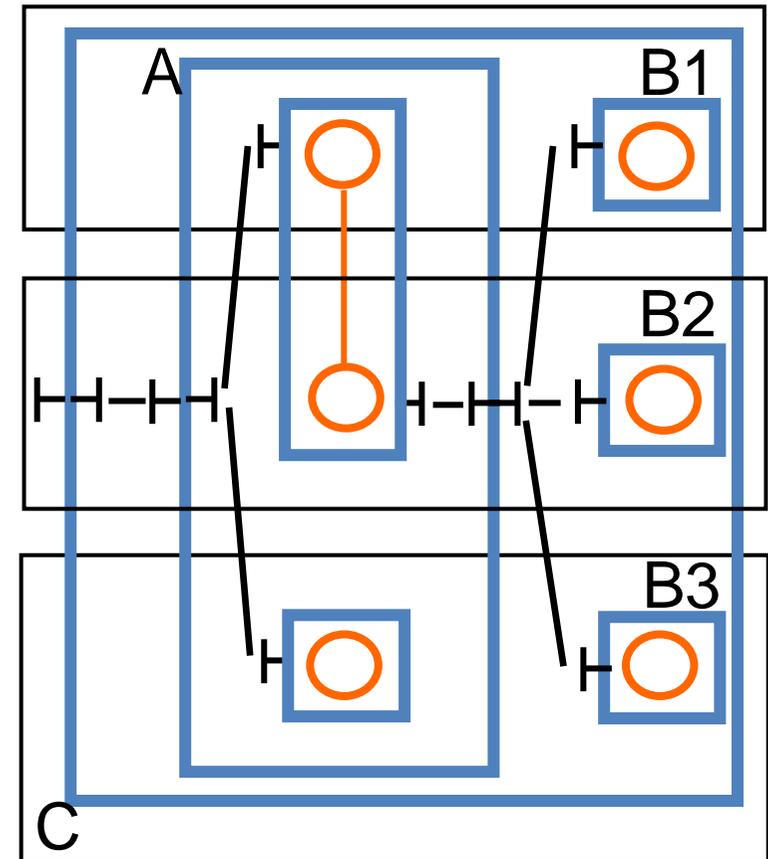


3. Parallel and composite  
component

# Components vs. Activity and JVMs



- ▶ Components are orthogonal to activities and JVMs
  - They contain activities, span across several JVMs
- ▶ Components are a way to globally manipulate distributed, and running activities

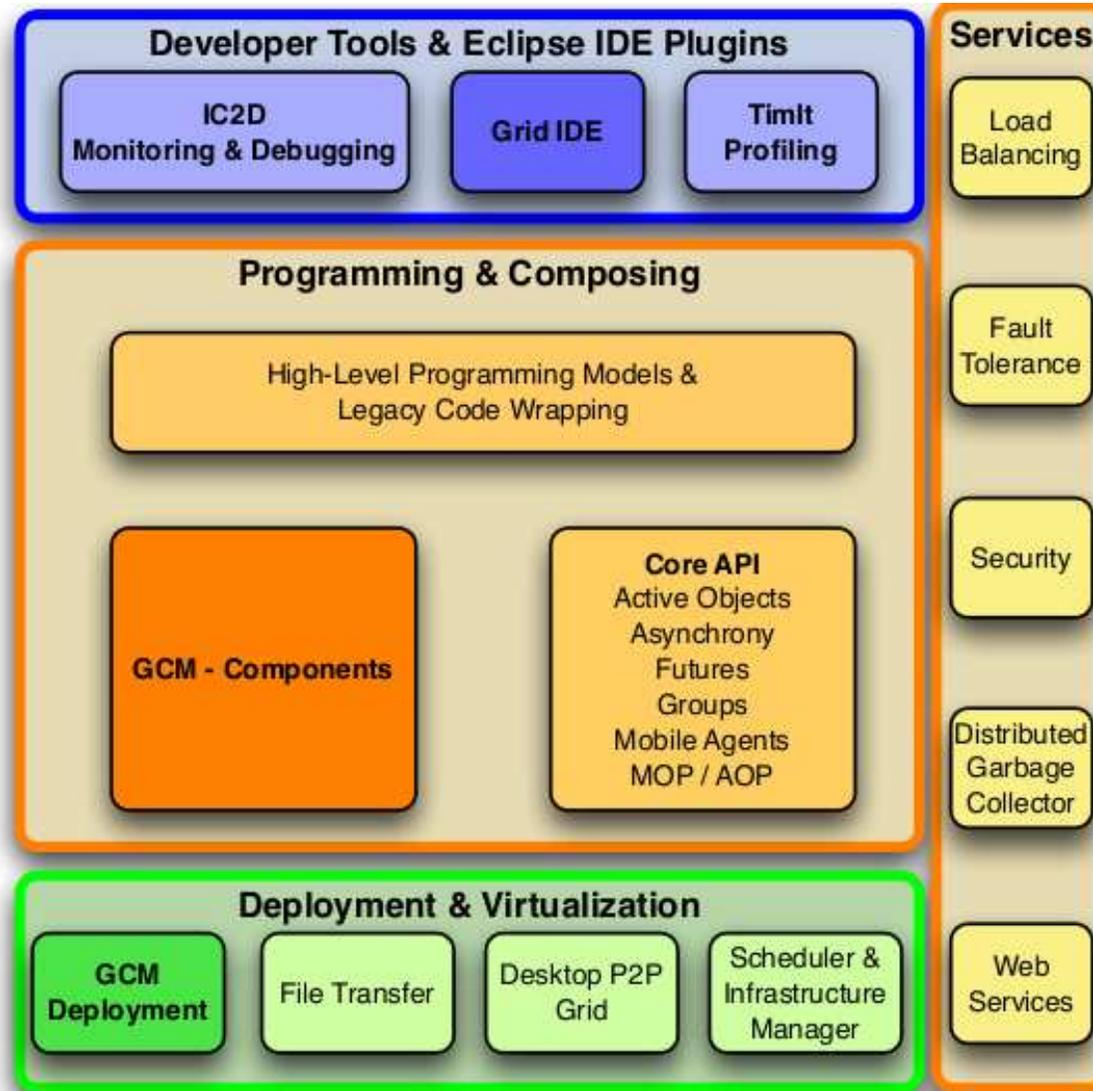


# Agenda

- ▶ ProActive and ProActive Parallel Suite
- ▶ Programming and Composing
  - ProActive Core
  - High Level Programming models
  - ProActive Components
- ▶ **Deployment Framework**
- ▶ Development Tools



# GCM Deployment



# Abstract Deployment Model

## Problem

Difficulties and lack of flexibility in deployment

Avoid scripting for configuration, getting nodes, connecting...

A key principle: Virtual Node (VN)

Abstract Away from source code:

- Machines names

- Creation/Connection Protocols

- Lookup and Registry Protocols

Interface with various protocols and infrastructures:

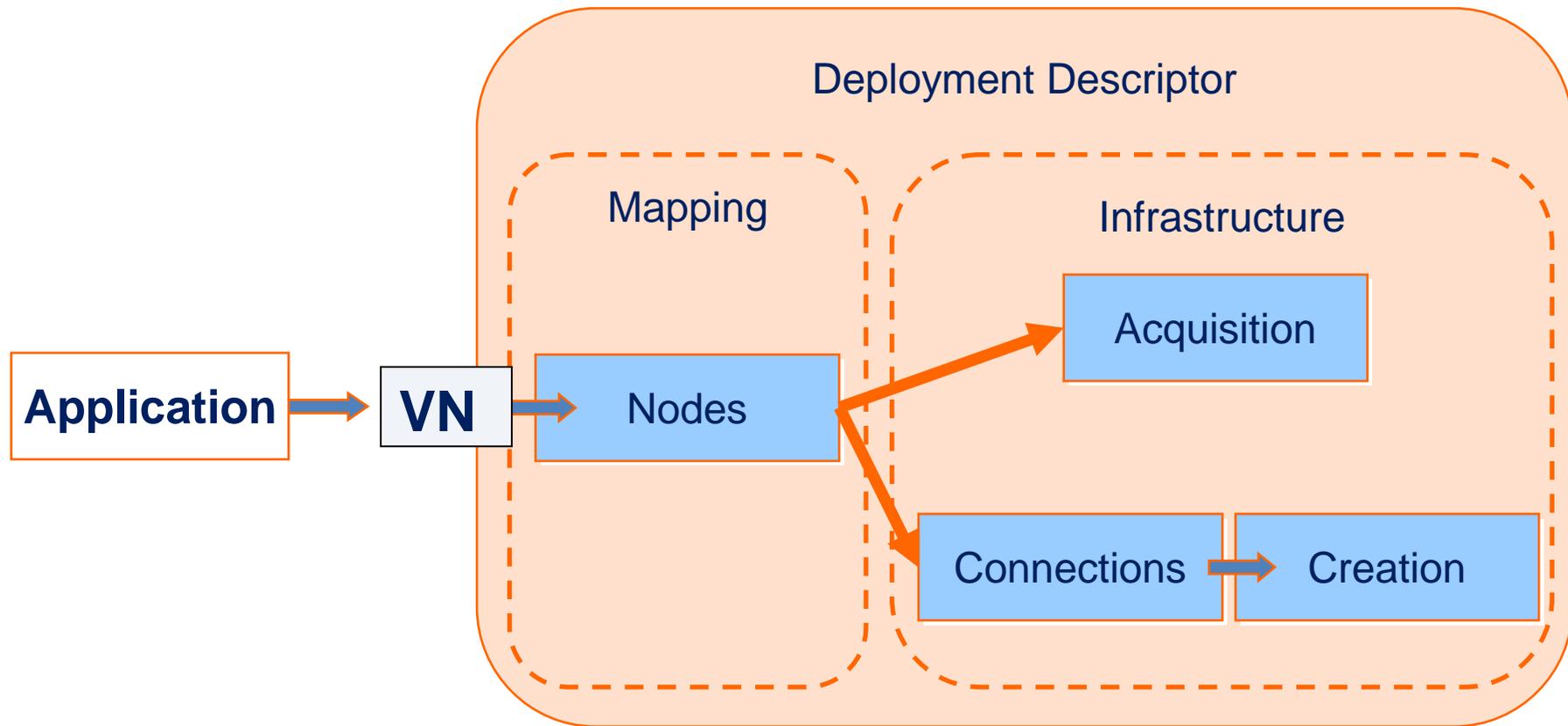
- Cluster: LSF, PBS, SGE , OAR and PRUN(custom protocols)

- Intranet P2P, LAN: intranet protocols: rsh, rlogin, ssh

- Grid: Globus, Web services, ssh, gsissh

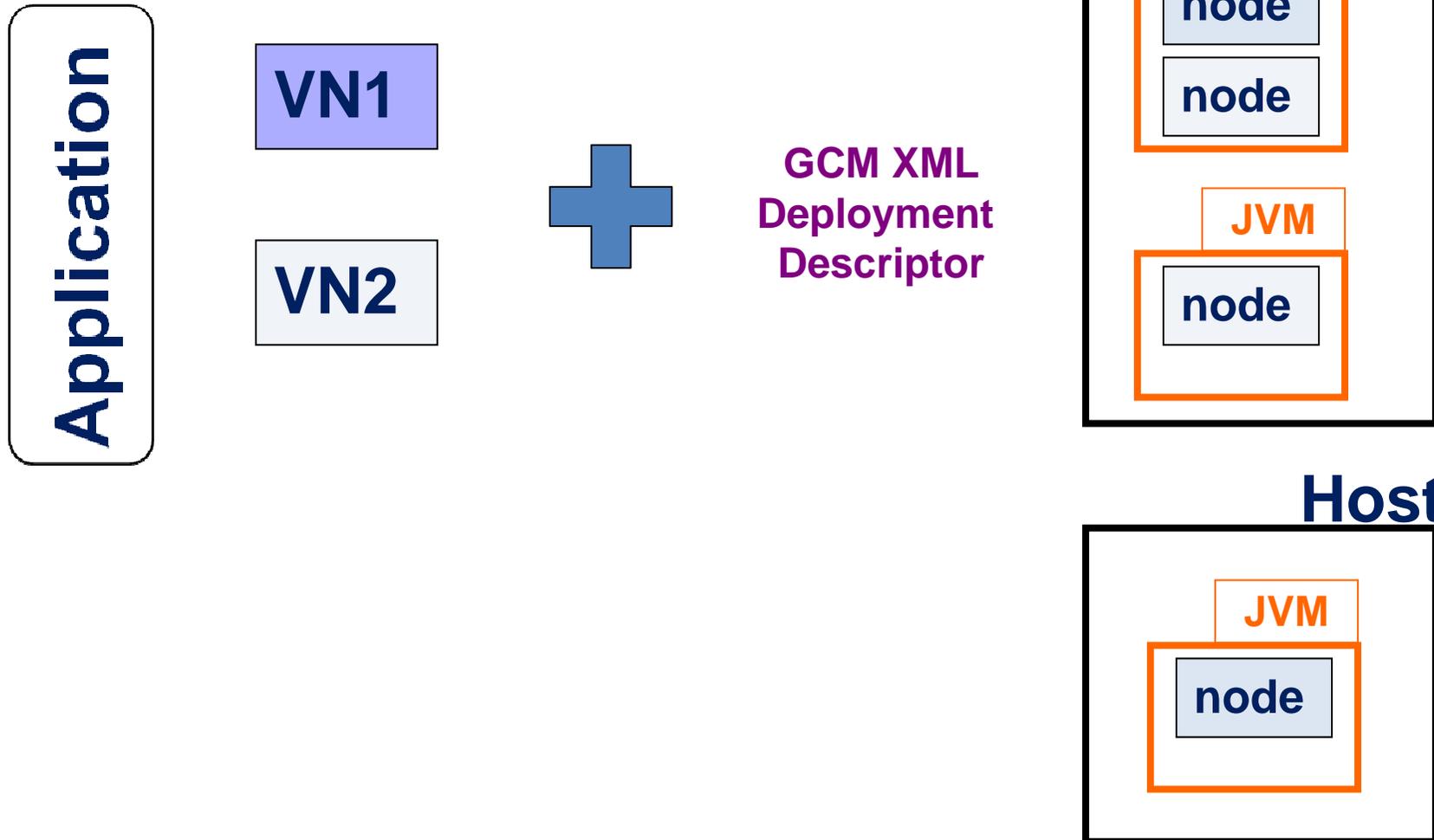


# Resource Virtualization

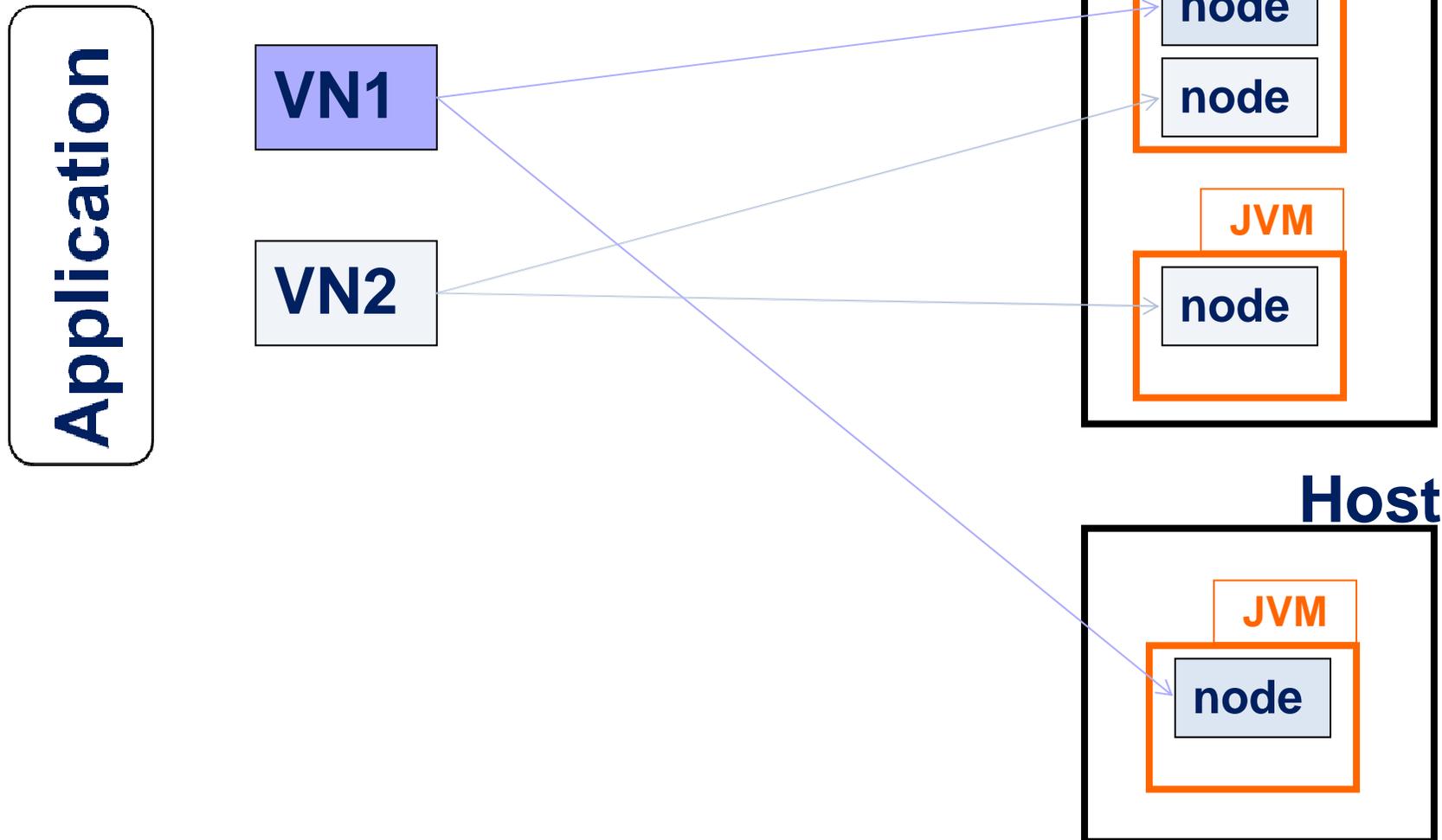


Runtime structured entities: 1 VN --> n Nodes in m JVMs on k Hosts

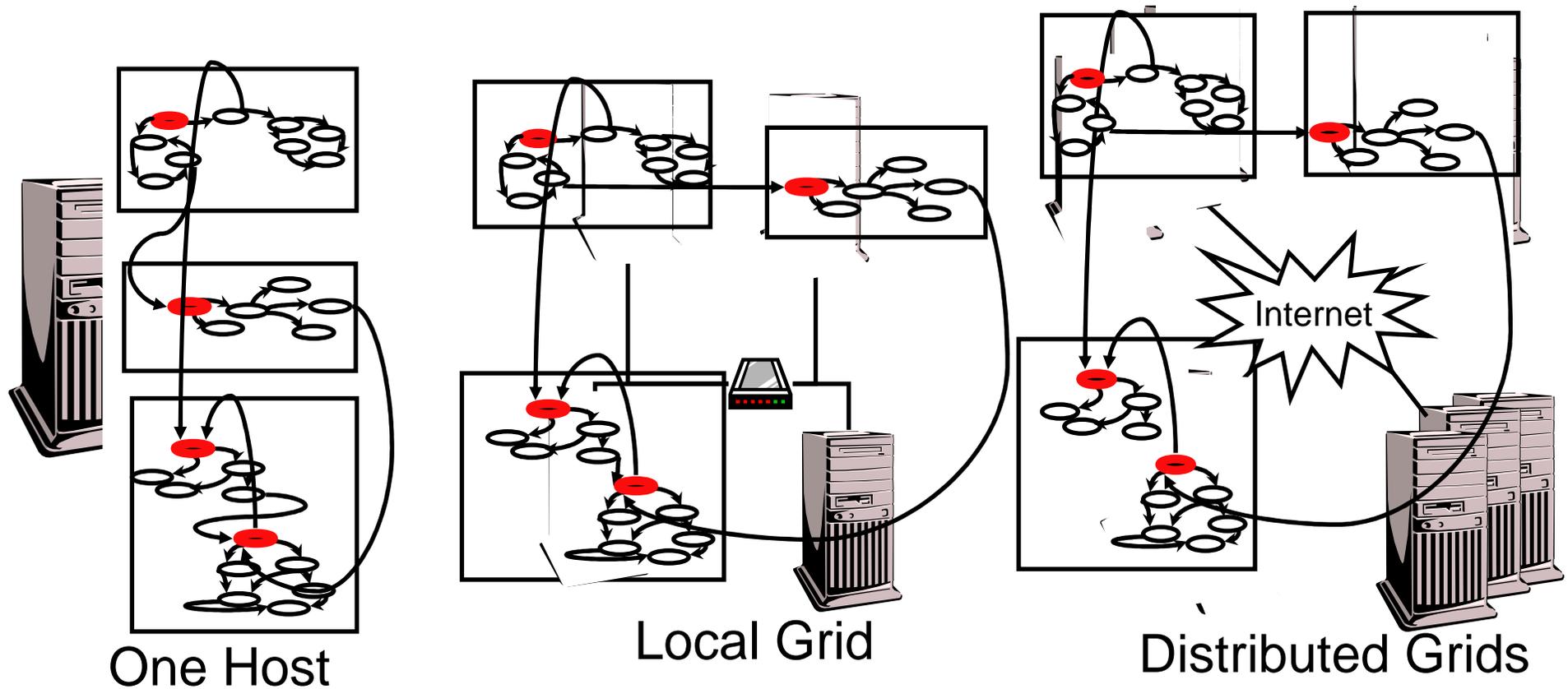
# Resource Virtualization Host



# Virtualization resources Host



# Multiple Deployments



# Rmissh : SSH Tunneling

## ▶ **A fact : overprotected clusters**

- Firewalls prevent incoming connections
- Use of private addresses
- NAT, IP Address filtering, ...

## ▶ **A consequence :**

- Multi clustering is a nightmare

## ▶ **Context :**

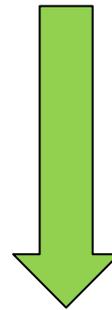
- SSH protocol : encrypt network traffic
- Administrators accept to open SSH port
- SSH provides encryption

## Rmissh : SSH Tunneling (2)

- ▶ Create a communication protocol within ProActive that allows firewall transversal
- ▶ Encapsulates rmi streams within ssh tunnels
- ▶ Avoid ssh tunneling costs when possible by first trying a direct rmi connection then fallbacking with rmissh



# The ProActive P2P

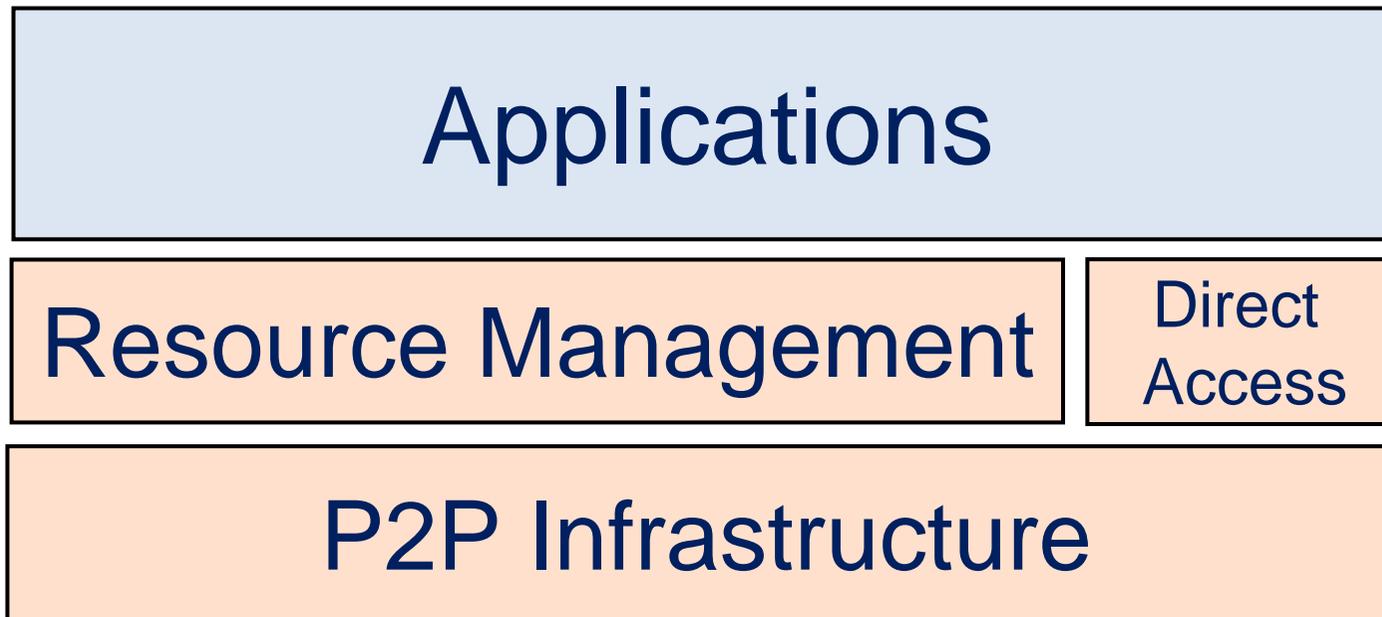


# The ProActive P2P

- ▶ Unstructured P2P
  - ❑ Easier to deploy/manage
  - ❑ Only 1 resource : CPU
- ▶ Java code
  - ❑ Each peer is written in Java and can run any Java application
- ▶ Direct communications
  - ❑ Peers are reachable using their name (URLs)
  - ❑ One peer can send/receive a reference on another peer



# The ProActive P2P (2)



# Infrastructure

- ▶ A peer is an Active Object in a JVM
- ▶ Each peer knows a limited number of other peers (bi-directional links)
  - ❑ Its *acquaintances*
  - ❑ The number is set by a variable (*NOA*)
- ▶ Goal of a peer
  - ❑ A peer will always try to maintain the number of its acquaintances equals to its *NOA*
- ▶ 2 basic operations
  - ❑ Adding an acquaintance
  - ❑ Removing an acquaintance

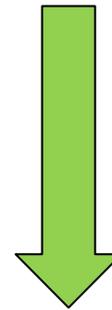


# Requesting Nodes

- ▶ To request a node
  - Contact only a Peer (URLs)
- ▶ The infrastructure will handle the reservation
- ▶ The application has to wait until the nodes are available
- ▶ Using the P2P network
  - Programmatically at runtime using the Java API
  - At Deployment time through the GCMDeployment



# Scheduler and Resource manager



# Scheduler / Resource Manager

## Overview

The screenshot displays three main windows from the ProActive Scheduler and Resource Manager:

- ProActive Scheduler (Left):** Shows a list of pending jobs. The 'Jobs' tab is active, displaying a table with columns: Id, State, User, Priority. The 'Console' tab shows a list of submitted tasks.
- ProActive Resource Manager (Middle):** Shows a 'Resource Explorer' tree view of nodes and their resources. A 'Statistics' window at the bottom shows:
 

name	value
# free nodes	95
# busy nodes	11
# down nodes	0
- Scheduler (Right):** Shows a list of finished jobs. The 'Scheduler' tab is active, displaying a table with columns: User, Priority, Name. A 'Result Preview' window at the bottom shows task details for ID 2008.

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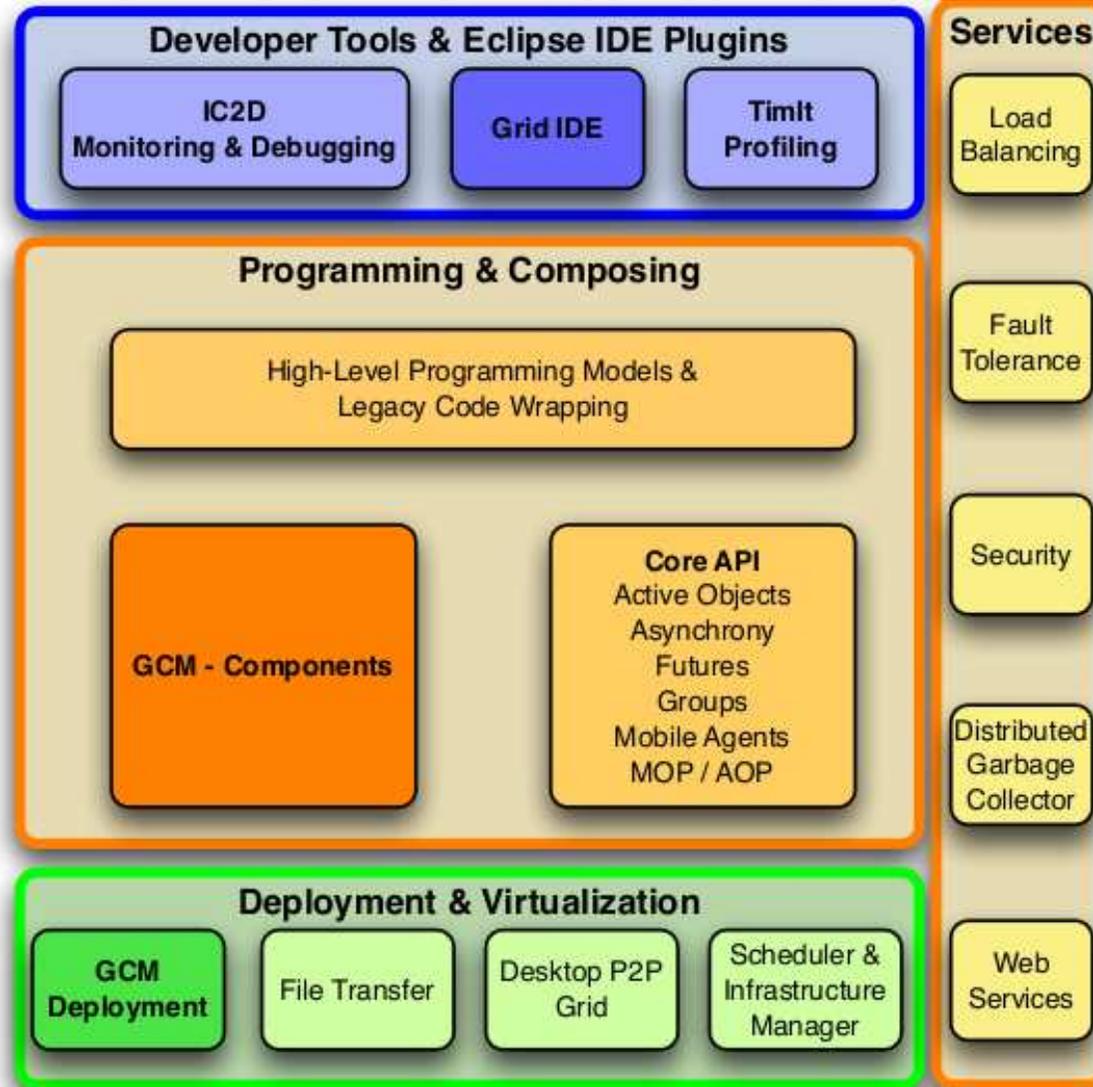


# Agenda

- ▶ ProActive and ProActive Parallel Suite
- ▶ Programming and Composing
  - ProActive Core
  - High Level Programming models
  - ProActive Components
  - Legacy code wrapping
- ▶ Deployment Framework
- ▶ **Development Tools**



# ProActive Parallel Suite



# IC2D

## Interactive Control & Debug for Distribution

- ▶ Basic Features:
  - Graphical visualization
  - Textual visualization
  - Monitoring and Control
- ▶ Extensible through RCP plug-ins
  - TimIt
  - ChartIt
  - P2P view
  - DGC view

# IC2D: Monitor your application

The screenshot displays the Eclipse IDE's Monitoring View and Job Monitoring View. The Monitoring View (left) shows a graph of virtual nodes and their interactions. The Job Monitoring View (right) shows a tree view of the application's job structure.

**Monitoring View:**

- Virtual nodes:  Renderer,  DefaultVN,  Dispatcher,  User.
- Graph components: PA\_JVM1357457629\_be..., Node Node60562498..., DinnerLayout#2, Table#3, Philosopher#4, Philosopher#5, Philosopher#6, Philosopher#7, Philosopher#8, PA\_JVM 436155261\_be..., Node Renderer127..., C3DRendering..., PA\_JVM 1672076495\_be..., Node Dispatcher 5..., C3DDispatche..., PA\_JVM 294719007\_be..., Node User16026446..., C3DUser#13, PA\_JVM 1631909824\_be..., Node Renderer1307..., C3DRendering..., duff.inria.fr:1099:OS und..., PA\_JVM1530781642\_du..., Node Renderer1174..., C3DRendering..., sidonie.inria.fr:1099:OS ..., PA\_JVM-772843461\_si..., Node Renderer151..., C3DRendering..., PA\_JOB-455186381\_si..., Node Node-455186381...
- Buttons:  Display topology,  Proportional,  Ratio,  Filaire,   Monitoring enable
- Console: Monitoring 15:09:15 => NodeObject id=Node 455186381 already monitored, check for new active objects

**Job Monitoring View:**

- Legend
- Job Monitoring View
- Tree structure: DefaultVN (JOB-1357457629), hebita.inria.fr:1099:OS un, PA\_JVM1357457629\_, Node Node60562498, DinnerLayout#2, Table#3 (JOB-13), Philosopher#4(J), Philosopher#5(J), Philosopher#6(J), Philosopher#7(J), Philosopher#8(J), sidonie.inria.fr:1099:OS u, Dispatcher (JOB--167207649), User (JOB--294719007), bebita.inria.fr:1099:OS un, PA\_JVM-294719007\_t, Node User1602644, C3DUser#13(JC), Renderer (JOB--1672076495), bebita.inria.fr:1099:OS un, PA\_JVM-1631909824\_

# TimIt: Automatic Timers in IC2D

The screenshot displays the IC2D monitoring interface. On the left, four horizontal bar charts show performance metrics for Worker#1, Worker#2, Worker#3, and Worker#4. An arrow points to the 'Application Level Timer' label above the Worker#1 chart. The charts include metrics like computePI\_rank, WaitForRequest, SendReply, AfterSerialization, Serialization, BeforeSerialization, SendRequest, Serve, and Total. The right side features a 'Monitoring#1' window with a 'Virtual nodes' diagram showing a hierarchy of nodes and workers. A legend on the far right defines object states and colors. The bottom console window shows a log of monitoring events.

**Application Level Timer**

**Worker#1 Performance (Snapshot: 18/06/07 15:27:18)**

computePI_rank1	25ms
WaitForRequest	4.43m
SendReply	28ms
AfterSerialization	29ms
Serialization	5ms
BeforeSerialization	15ms
SendRequest	49ms
Serve	1.99s
Total	4.46m

**Worker#2 Performance (Snapshot: 18/06/07 15:27:18)**

computePI_rank0	34ms
WaitForRequest	4.44m
SendReply	26ms
Serve	41ms
Total	4.47m

**Worker#3 Performance (Snapshot: 18/06/07 15:27:18)**

computePI_rank3	30ms
WaitForRequest	4.42m
SendReply	80ms
AfterSerialization	79ms
Serialization	5ms
BeforeSerialization	94ms
SendRequest	177ms
Serve	1.86s
Total	4.45m

**Worker#4 Performance (Snapshot: 18/06/07 15:27:18)**

computePI_rank2	40ms
WaitForRequest	4.43m
SendReply	37ms
AfterSerialization	152ms
Serialization	9ms
BeforeSerialization	25ms
SendRequest	188ms
Serve	1.95s
Total	4.46m

**Virtual nodes Diagram:**

- amda.inria.fr:1099:Linux
  - PA\_JVM1530790716\_am... (Worker#4)
  - PA\_JVM704475267\_amd... (Worker#2)
  - PA\_JVM1714913173\_am... (Worker#1)
  - PA\_JVM1122637657\_am... (Worker#1)
  - PA\_JVM2010164699\_am... (Worker#3)

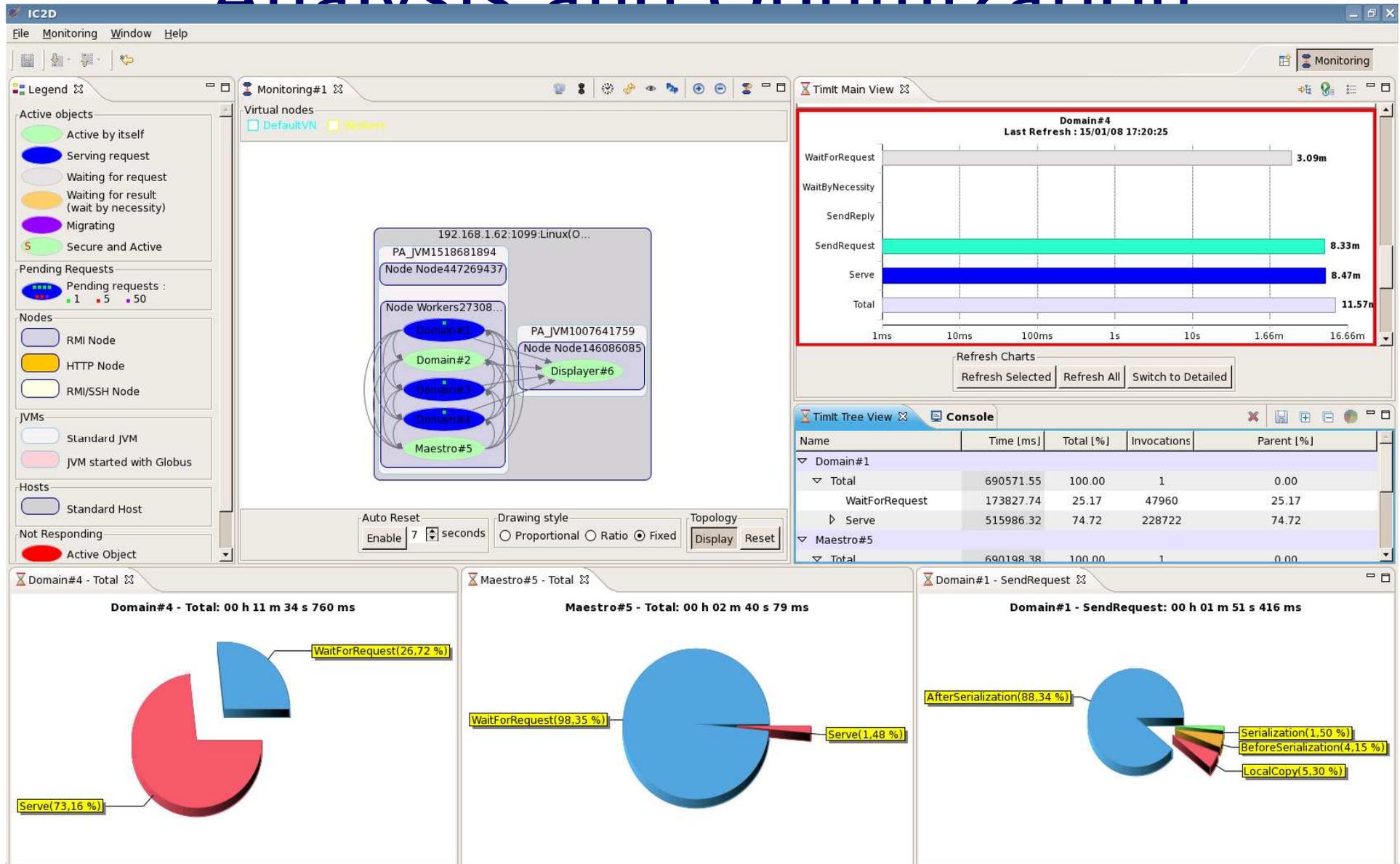
**Console Log:**

```

Monitoring
15:27:15 => VMObject id=PA_JVM2010164699_amda.inria.fr already monitored, check
15:27:15 => VMObject id=PA_JVM1530790716_amda.inria.fr already monitored, check
15:27:45 => Exploring Host amda.inria.fr with RMI on port 1099
15:27:45 => VMObject id=PA_JVM1714913173_amda.inria.fr already monitored, check
15:27:45 => VMObject id=PA_JVM1122637657_amda.inria.fr already monitored, check
15:27:45 => VMObject id=PA_JVM704475267_amda.inria.fr already monitored, check f
15:27:45 => VMObject id=PA_JVM2010164699_amda.inria.fr already monitored, check
15:27:45 => VMObject id=PA_JVM1530790716_amda.inria.fr already monitored, check
    
```

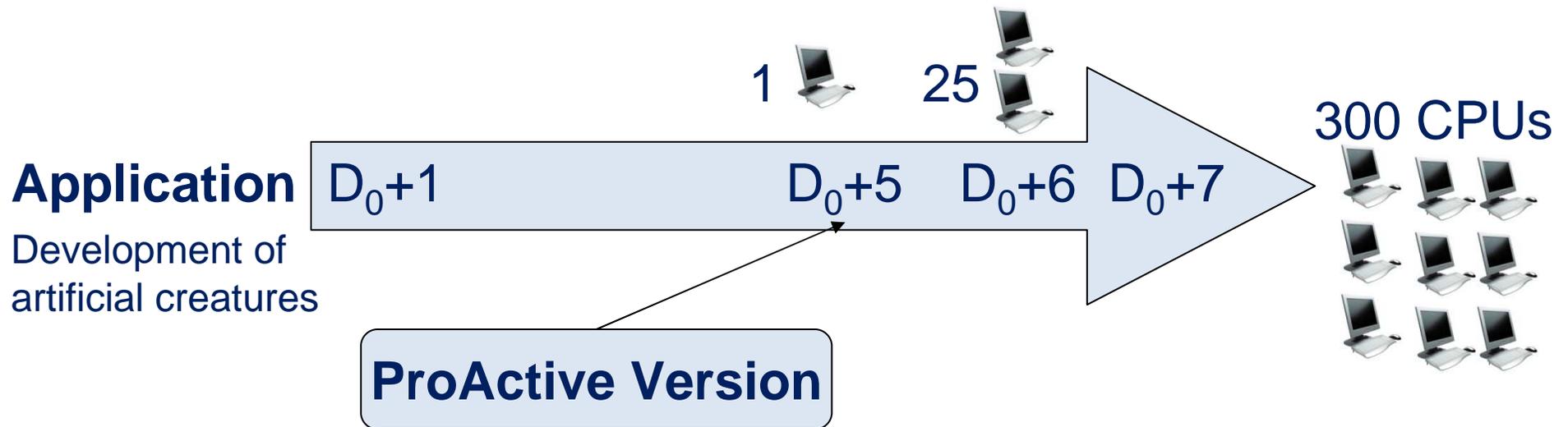


# Analysis and Optimization



# M/W Success Story: Artificial Life Generation

*Sylvain Cussat-Blanc, Yves Duthen – IRIT TOULOUSE*



<b>Initial Application (C++)</b>	<b>1 PC</b>	<b>56h52 =&gt; Crashed</b>
<b>ProActive Version</b>	<b>300 CPUs</b>	<b>19 minutes</b>

# Price-It workload distribution with ProActive

- ▶ Low level parallelism : shared memory
- ▶ Written in c++
- ▶ Originally written for microsoft compiler
- ▶ JNI, Com interface
- ▶ No thread safe
  
- ▶ Upgrading the code base to thread safe code might be costly
- ▶ Is there any easier and cheaper alternative to extract parallelism from Price-it Library ?

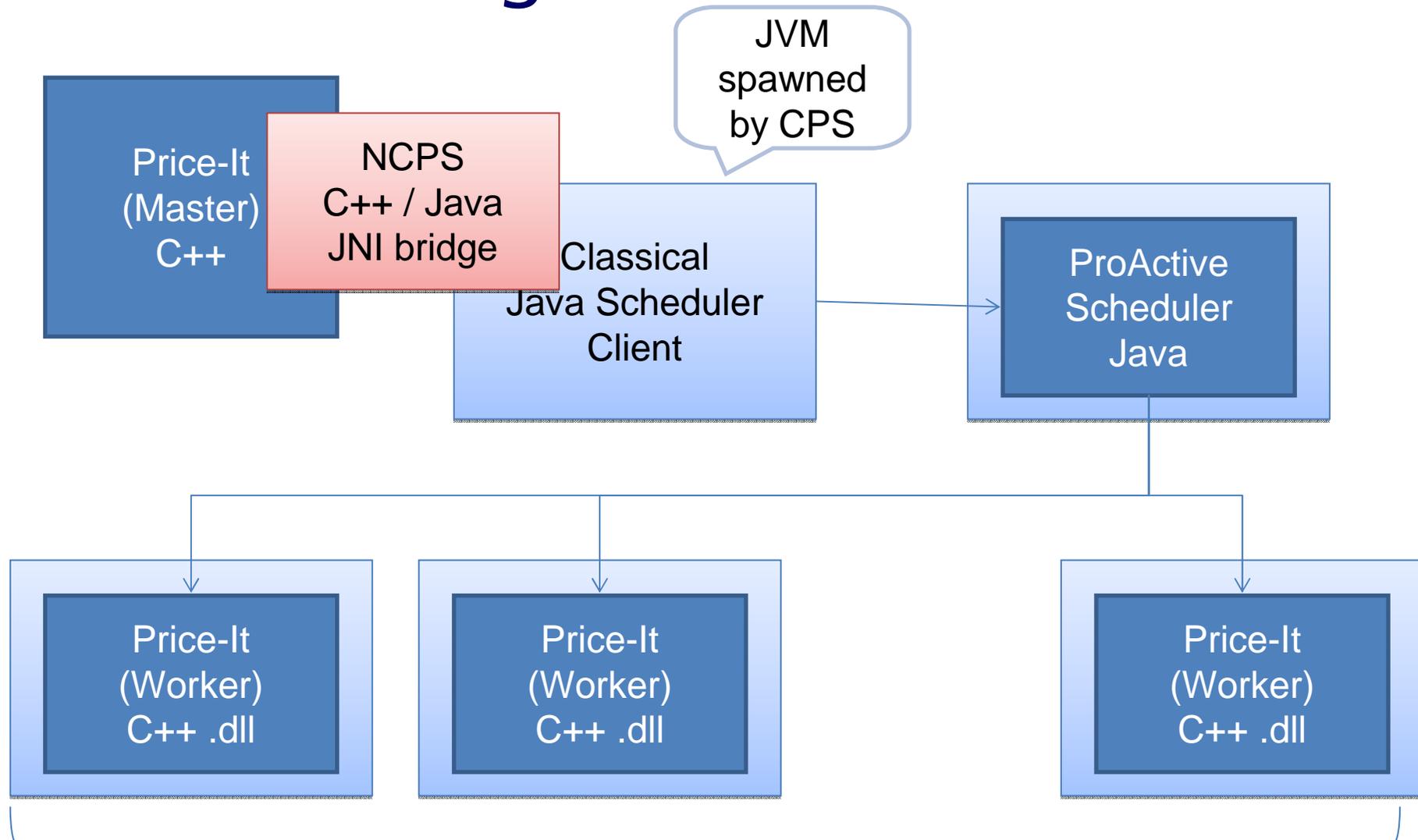


# CPS : C++ API Client for ProActive Scheduler

- ▶ CPS : Client for ProActive Scheduler
- ▶ Shipped as .so/.dll
- ▶ A set of C++ methods to submit jobs to the Scheduler
  - ❑ SchedulerClient::init() and dispose()
  - ❑ SchedulerClient::submitJob(Job\* jobPtr)
  - ❑ SchedulerClient::getJobResult(int jobId)
- ▶ Internally uses JNI



# Using CPS in Price-It



Workers are shipped as .dll then loaded by JVMs and executed through JNI

# Conclusion

- ▶ Simple and Uniform programming model
- ▶ Rich High Level API
- ▶ Write once, deploy everywhere (GCMD)
- ▶ Let the developer concentrate on his code, not on the code distribution
- ▶ Easy to install, test, validate on any network



# Now, let's play with ProActive...

- **Start** and **monitor** with IC2D the ProActive examples, and have a look at the **source code**

`org.objectweb.proactive.examples.*`

Features	Applications
Basics, Synchronization	Doctors problem ( <code>doctors.bat</code> ), Reader/Writer problem ( <code>readers.bat</code> ),...
Futures, Automatic Continuation	Binary Search Tree ( <code>bintree.bat</code> )
Migration	Migrating Agent ( <code>/migration/penguin.bat</code> )
Group	Chat ( <code>/group/chat.bat</code> )
Fault-Tolerance	N-body problem ( <code>/FT/nbodyFT.bat</code> )
All	Distributed 3D renderer ( <code>c3d*.bat</code> )

