Mobile Computing

General introduction

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BRICS

Mobile Computing – p. 1

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- Better refer to *process algebras* or *process calculi*
- Mobility = any formalism able to describe concurrent, distributed and dynamically-reconfigurable systems
- In fact, two notions:
 - Mobility of names, *labile* systems
 - Mobility of processes, *motile* systems

In the old times

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- Petri nets (Petri, early 60s): process synchronization, asynchronous events, concurrent operations, and conflicts or resource sharing
- CSP (Hoare, 1978), CCS (Milner, 1980): parallel and independent processes, synchronizing over *channels*, through *channel names*:

 $\bar{a}.P \mid a.Q \rightarrow P \mid Q$

Milner, Parrow, Walker, 1989

• Adds the possibility of *communication*:

 $\bar{a}b.P \mid a(x).Q \rightarrow P \mid Q\{x \mapsto b\}$

• the possibility of passing previously unknown names allows to *dynamically reconfigure* systems

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- Extensive formal developments (semantics, types, equivalences...)
- Still an active research area
- Implementations: Pict (Turner, 1995), Nomadic Pict, Blue, TyCo, JoCaml
- Used in industry: Microsoft's BizTalk

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→
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π₁ (Amadio, Prasad), Dpi (Hennessy, Riely), join-calculus, etc...

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 $@_{l}[\bar{a}b.P] | @_{l}[a(x).Q] | @_{p}[\bar{a}c.R]$

$@_l[P] | @_l[Q\{x \mapsto b\}] | @_p[\bar{a}c.R]$

- π₁ (Amadio, Prasad), Dpi (Hennessy, Riely), join-calculus, etc...
- Often *higher-order* communication (i.e. migration of processes or *agents*)

Mobile ambients

Cardelli, Gordon 1998

• Real "physical" mobility:



Mobile ambients

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• Real "physical" mobility:



- In fact, two notions:
 - Mobile computation: mobile code moving between different execution devices (agents...)
 - Mobile computing: computation in mobile devices (laptop, crossing of firewall...)

Mobile ambients

- Many variants: safe ambients, robust ambients, controlled ambients, boxed ambients, Seal...
- Related formal tools: type systems, equational theory, logics, expressiveness results...

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- Brane calculi (Cardelli 2003): interaction of biological membranes, inspired from ambients, providing new formal tools for biologists



- Static analysis of terms, in order to ensure some properties, w.r.t. some set of *typing rules*
- Examples:
 - Types of arguments for functions in C
 - Topic of conversation on a channel in π
 - Immobility of ambients

Equational theory

• When do we declare that two terms are *equivalent* ?

Equational theory

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- Why is it important ?

Equational theory

- When do we declare that two terms are *equivalent* ?
- Why is it important ?
- Examples: bisimilarity, strong/weak bisimilarity, trace equivalence, ...

That's enough. Show me something concrete now !