Lecture 9.

Direct Datagram Forwarding:

Address Resolution Protocol (ARP)

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Problem statement

➜ **Routing decision for packet X has two possible outcomes:**

⇨ You are arrived to the final network: go to host X

⇨ You are not arrived to the final network: go through router interface Y

➜ **In both cases we have an IP address on THIS network. How can we send data to the interface?**

➜ **Need to use physical network facilities!**
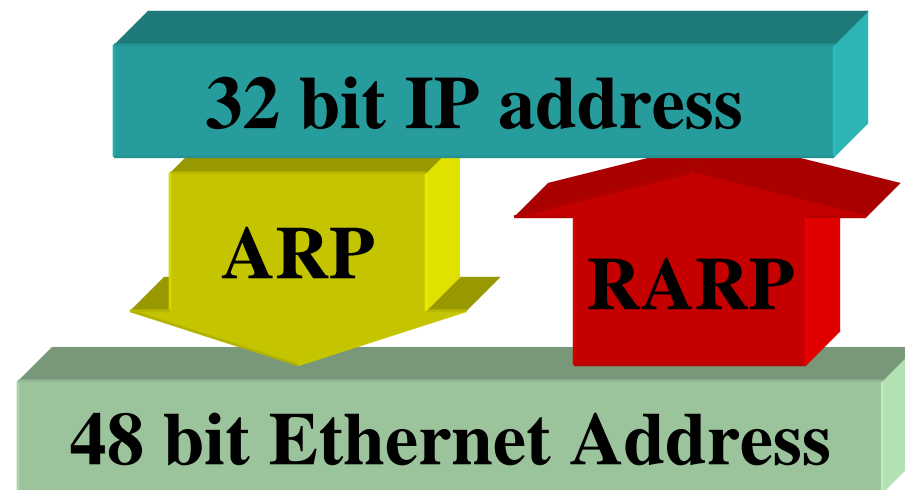
Reaching a physical host

➔ **IP addresses only make sense to TCPIP protocol suite**

➔ **physical networks have their own hardware address**

⇨ e.g. 48 bits Ethernet address, 16 or 48 bits Token Ring, 16 or 48 bit FDDI, ...

⇨ datalink layers may provide the basis for several network layers, not only IP!

Address Resolution Protocol
RFC 826

Here described for Ethernet, but more general: designed for any datalink with broadcast capabilities

32 bit IP address

ARP

RARP

48 bit Ethernet Address

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Manual mapping

→ **A possibility, indeed!!**

⇨ Nothing contrary, in principle

→ actually done in X.25, ISDN (do not support broadcast)

⇨ Simply keep in every host a mapping between IP address and hardware address for every IP device connected to the considered network

→ **drawbacks**

⇨ tedious

⇨ error prone

⇨ requires manual updating

→ e.g. when attaching a new PC, must touch all others...

ARP

➔ **Dynamic mapping**

⇨ not a concern for application & user

⇨ not a concern for system administrator!
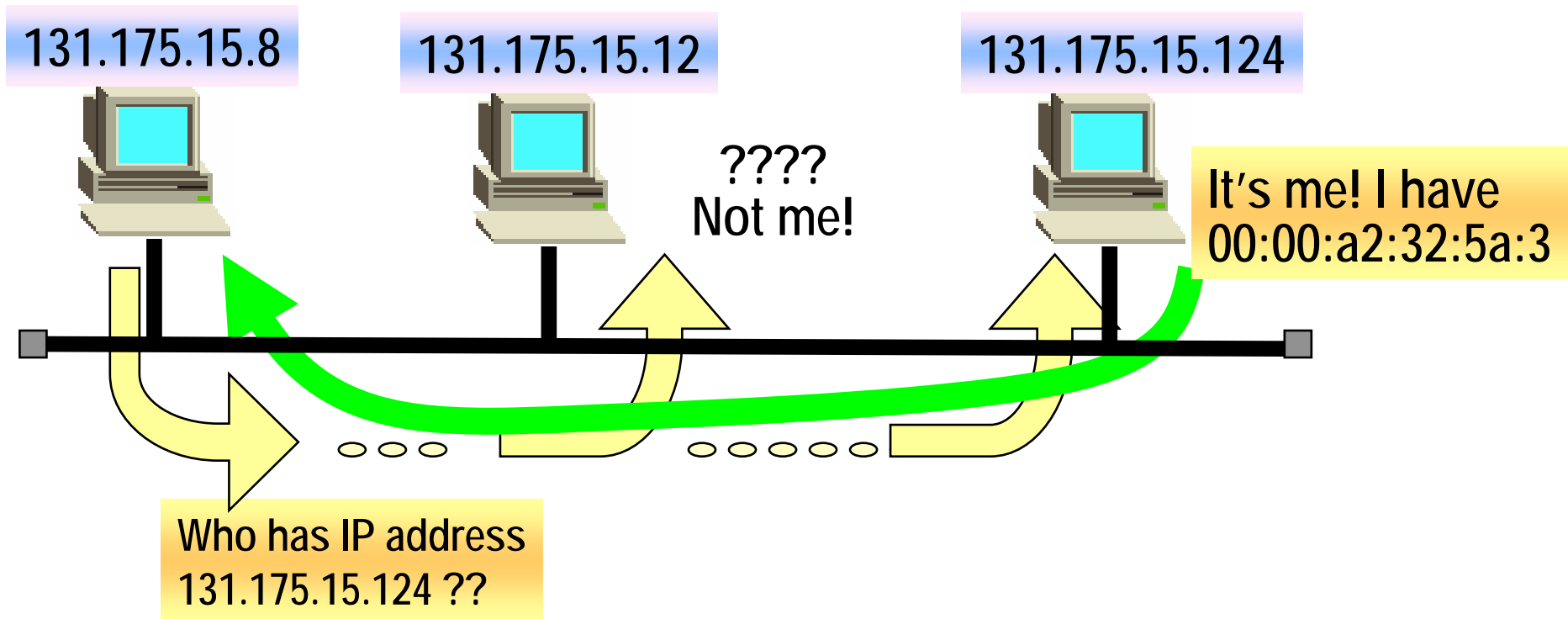
➔ **Any network layer protocol**

⇨ not IP-specific

➔ **supported protocol in datalink layer**

⇨ not a datalink layer protocol !!!!

➔ **Need datalink with broadcasting capability**

⇨ e.g. ethernet shared bus

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ARP idea

131.175.15.8

131.175.15.12

131.175.15.124

???? Not me!

It's me! I have 00:00:a2:32:5a:3

Who has IP address 131.175.15.124 ??

→ **Send broadcast request**

→ **receive unicast response**

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ARP cache

➔ **Avoids arp request for every IP datagram!**

⇨ Entry lifetime defaults to 20min

➔ deleted if not used in this time

➔ 3 minutes for "incomplete" cache entries (i.e. arp requests to non existent host)

➔ it may be changed in some implementations

» in particularly stable (or dynamic) environments

⇨ *arp -a* to display all cache entries (arp –d to delete)

try a traceroute or ping to check ARP caching!

➔ First packet generally delays more

➔ includes an ARP request/reply!

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ARP request/reply
Incapsulation in Ethernet Frame

6 bytes	6 bytes	2 bytes	28 bytes (for IP)	4 bytes
Ethernet destination address	Ethernet source address	frame type	ARP Request / Reply	CRC

➔ **Ethernet Destination Address**
⇨ ff:ff:ff:ff:ff:ff (broadcast) for ARP request

➔ **Ethernet Source Address**
⇨ of ARP requester

➔ **Frame Type**
⇨ ARP request/reply: 0x0806
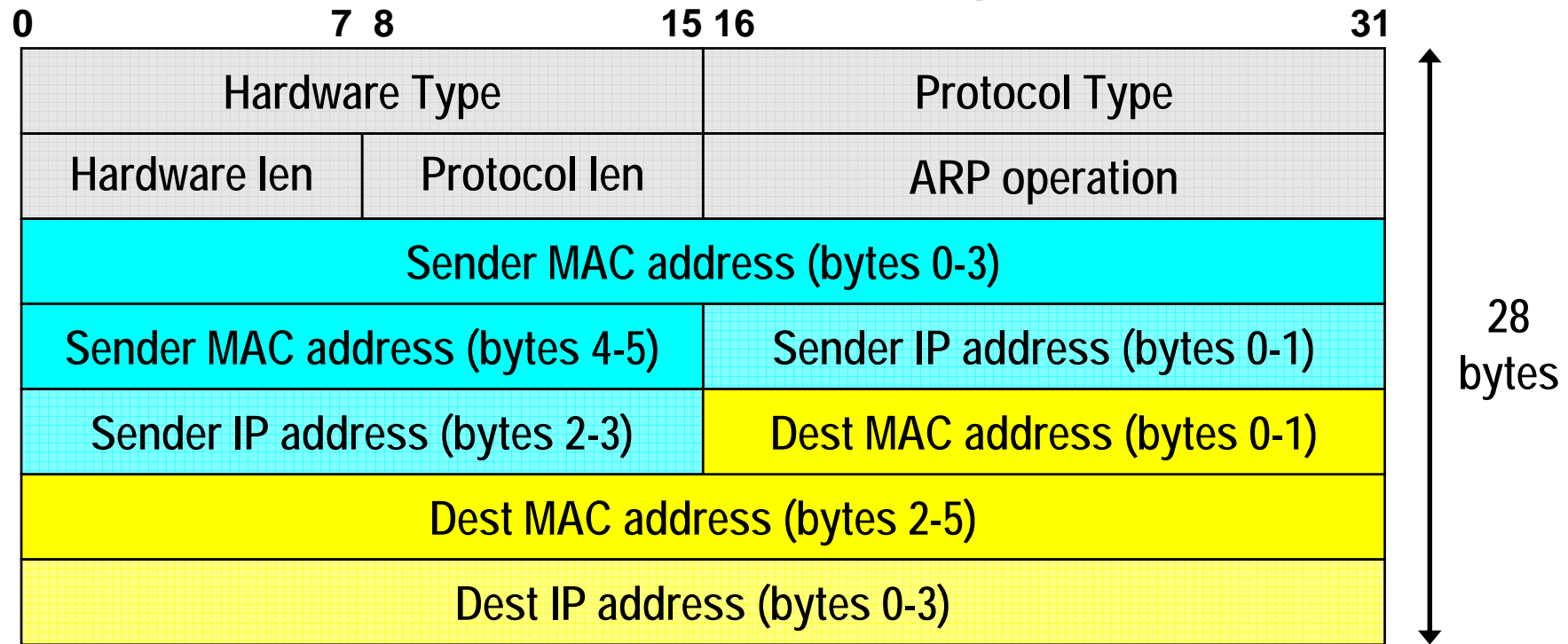⇨ RARP request/reply: 0x8035
⇨ IP datagram: 0x0800

⎫
⎬ Protocol demultiplexing codes!
⎭

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ARP request/reply format

0 ... 7	8 ... 15	16 ... 31
Hardware Type		Protocol Type
Hardware len	Protocol len	ARP operation
Sender MAC address (bytes 0-3)		
Sender MAC address (bytes 4-5)		Sender IP address (bytes 0-1)
Sender IP address (bytes 2-3)		Dest MAC address (bytes 0-1)
Dest MAC address (bytes 2-5)		
Dest IP address (bytes 0-3)		

28 bytes

Hardware type: 1 for ethernet
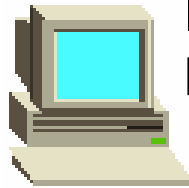
Protocol type: 0x0800 for IP (0000.1000.0000.0000)

⇨ *the same of Ethernet header field carrying IP datagram!*

Hardware len = length in bytes of hardware addresses (6 bytes for ethernet)

Protocol len = length in bytes of logical addresses (4 bytes for IP)

ARP operation: 1=request; 2=reply; 3/4=RARP req/reply

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Sample ARP request/reply

IP: 131.175.15.8
MAC: 0:0:8c:3d:54:1

IP: 131.175.15.24
MAC: 0:4f:33:3:ee:67

Ethernet Packet: ARP REQUEST ➡️

⬅️ Ethernet Packet: ARP reply

ARP REQUEST		dest MAC	ARP reply	
FF:FF:FF:FF:FF:FF		dest MAC	00:00:8c:3d:54:01	
00:00:8c:3d:54:01		src MAC	00:4f:33:03:ee:67	
0x0806		ARP frame type	0x0806	
0x0001	0x0800	Ethernet / IP	0x0001	0x0800
0x06	0x04	0x0001	MAC=6 / IP=4 / rq=1,rpl=2	0x06	0x04	0x0002
00:00:8c:3d:54:01		src MAC	00:4f:33:03:ee:67	
131.175.15.8		src IP	131.175.15.24	
00:00:00:00:00:00		dest MAC	00:00:8c:3d:54:01	
131.175.15.24		dest IP	131.175.15.8	
checksum		Ethernet checksum	checksum	

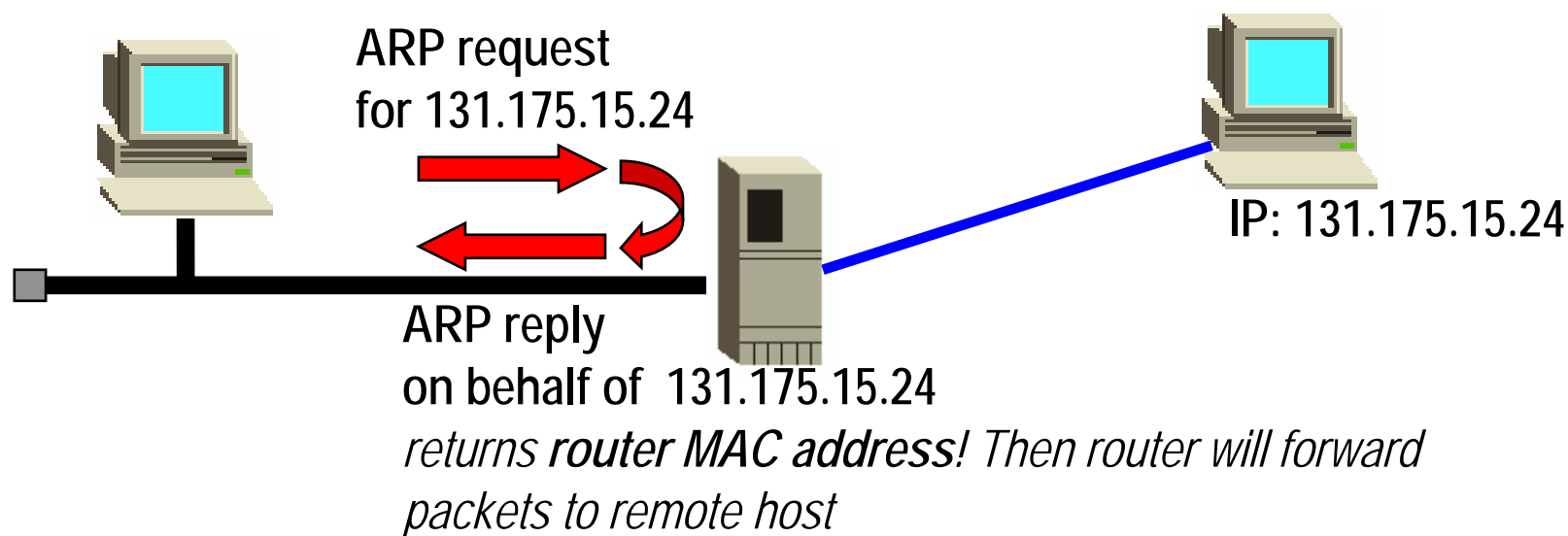G.Bianchi, G.Neglia, V.Mancuso

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ARP cache updating

➔**ARP requests carry requestor IP/MAC pair**

➔**ARP requests are broadcast**

⇨thus, they MUST be read by everyone

➔**Therefore, it comes for free, for every computer, to update its cache with requestor pair**

➔*Cannot do this with ARP reply, as it is unicast!*

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Proxy ARP

➔ **Device that responds to an ARP request on behalf of some other machine**

⇨ allows having ONE logical (IP) network composed of more physical networks

⇨ especially important when different techologies used (e.g. 100 PC ethernet + 2 PC dialup SLIP)

ARP request
for 131.175.15.24

IP: 131.175.15.24

ARP reply
on behalf of 131.175.15.24
*returns **router MAC address**! Then router will forward packets to remote host*

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Gratuitous ARP

➡ **ARP request issued by an IP address and addressed to *the same* IP address!!**

⇨ Clearly nobody else than ME can answer!

⇨ WHY asking the network which MAC address do I have???

➡ **Two main reasons:**

⇨ determine if another host is configured with the same IP address

→ in this case respond occurs, and MAC address of duplicated IP address is known.

⇨ Use gratuitous ARP when just changed hardware address

→ all other hosts update their cache entries!

→ A problem is that, despite specified in RFC, not all ARP cache implementations operate as described….

ARP: not only *this* mechanism!

➔ **Described mechanism for broadcast networks (e.g. based on shared media)**

➔ **Non applicable for non broadcast networks**

⇨ in this case OTHER ARP protocols are used

➔ e.g. distributed ARP servers

➔ e.g. algorithms to map IP address in network address

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Getting an IP address:

Reverse Address Resolution Protocol (RARP)

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The problem

➔ **Bootstrapping a diskless terminal**

⇨ this was the original problem in the 70s and 80s

➔ **Reverse ARP [RFC903]**

⇨ a way to obtain an IP address starting from MAC address

➔ **Today problem: dynamic IP address assignment**

⇨ limited pool of addresses assigned only when needed

➔ **RARP not sufficiently general for modern usage**

⇨ BOOTP (Bootstrap Protocol - RFC 951): significant changes to RARP (a different approach)

⇨ DHCP (Dynamic Host Configuration Protocol - RFC 1541): extends and replaces BOOTP

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RARP packet format
almost identical to ARP. Differences:

6 bytes	6 bytes	2B	28 bytes (for IP)	4 bytes
Dest addr	Src addr	ftyp: 0x 8035	RARP Request / Reply	CRC

0	7 8	15 16	31
Hardware Type		Protocol Type	
Hardware len	Protocol len	oper: 3 (RARP req) or 4 (RARP reply)	
Sender MAC address (bytes 0-3)			
Sender MAC address (bytes 4-5)		Sender IP address (bytes 0-1)	
Sender IP address (bytes 2-3)		Dest MAC address (bytes 0-1)	
Dest MAC address (bytes 2-5)			
Dest IP address (bytes 0-3)			

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RARP problems

➔ **Network traffic**

⇨ for reliability, multiple RARP servers need to be configured on the same Ethernet

➔ to allow bootstrap of terminals even when one server is down

⇨ But this implies that ALL servers simultaneously respond to RARP request

➔ contention on the Ethernet occurs

➔ **RARP requests not forwarded by routers**

⇨ being hardware level broadcasts...

RARP fundamental limit

➔ **Allows only to retrieve the IP address information**

⇨ and what about all the remaining full set of TCPIP configuration parameters???

➔ Netmask?

➔ name of servers, proxies, etc?

➔ other proprietary/vendor/ISP-specific info?

➔ **This is the main reason that has driven to engineer and use BOOTP and DHCP**

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BOOTP/DHCP approach

➔ **Requests/replies encapsulated in UDP datagrams**
  ⇨ may cross routers
  ⇨ no more dependent on physical medium

➔ **request addressing:**
  ⇨ destination IP = 255.255.255.255
  ⇨ source IP = 0.0.0.0
  ⇨ destination port (BOOTP): 67
  ⇨ source port (BOOTP): 68

➔ **router crossing:**
  ⇨ router configured as BOOTP relay agent
  ⇨ forwards broadcast UDP requests with destination port 67

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BOOTP parameters exchange

➔ **Many more parameters**

⇨ client IP address (when static IP is assigned)

⇨ your IP address (when dynamic server assignment)

⇨ gateway IP address (bootp relay agent - router - IP)

⇨ server hostname

⇨ boot filename

➔ **Fundamental: vendor-specific information field (64 bytes)**

⇨ seems a lot of space: not true!

⇨ DHCP uses a 312 vendor-specific field!

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Vendor specific information
format allows general information exchange

Tag 1 byte	Len 1 byte	Parameter exchanged

➔ **E.g.: subnet mask:**

    ⇨ tag=1, len=4, parameter=32 bit subnet mask

➔ **e.g.: time offset:**

    ⇨ tag=2, len=4, parameter=time
    (seconds after midnight, jan 1 1900 UTC)

➔ **e.g. gateway (variable item)**

    ⇨ tag=3, len=N, list of gateway IPaddr (first preferred)

➔ **e.g. DNS server (tag 6)**

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