

Short exam : 1 hour

No electronic/manuscript document is allowed. You can answer in french or english. Most of the questions may be answered in at most one sentence.

1 Graph colouring

Question 1 Give the definition of the degree of a vertex $v \in V$ in a graph $G = (V, E)$.

A graph is **cubic** if all its vertices have degree exactly 3.

Question 2 Let $k \in \mathbb{N}$. Give the definition of a proper k -colouring of a graph $G = (V, E)$.

The **chromatic number** $\chi(G)$ of a graph G is the smallest integer k such that G admits a proper k -colouring.

Question 3 — Draw a cubic graph G with $\chi(G) = 4$.

- Draw a cubic graph G different from K_4 (the complete graph with 4 vertices).
- Prove that, for every cubic graph G , $\chi(G) \leq 4$.

2 Planar graphs

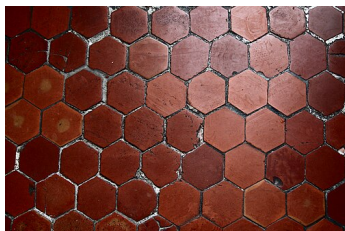
Question 4 — Give the definition of a connected graph.

- Give the definition of an outerplanar graph.
- Draw $K_{3,3}$ on the torus (i.e., on a donuts) without crossing edges.

Let $G = (V, E)$ be a connected planar graph with $|F|$ faces.

Question 5 Give (without proof) the Euler's Formula linking $|F|$, $|V|$ and $|E|$ in connected planar graphs.

The Euler's Formula is actually also true for graphs embedded (drawn) without crossing on a sphere (not only on a plane).



(a) Tiling of the plane with hexagones.



(b) Soccer ball.

FIGURE 1 – Examples.

The goal of the following questions is to explain how a soccer ball is built. Did you notice that a soccer ball consists of hexagons (generally in white) and of pentagons (generally in black)? Why not only with hexagons? How many pentagons?

Recall that a **vertex cover** of a graph $G = (V, E)$ is a set $Q \subseteq V$ of vertices that “touches” all edges. Formally, $Q \subseteq V$ is a vertex cover of G if, for every edge $e \in E$ (seen as a set of 2 vertices), $e \cap Q \neq \emptyset$.

In what follows, we consider a weight function $w : V \rightarrow \mathbb{R}^+$ over the vertices of G . The weight of a subset $Q \subseteq V$ of vertices is $w(Q) = \sum_{v \in Q} w(v)$, i.e., the sum of the weights of the vertices of Q .

The goal of what follows is to design an algorithm that computes a vertex cover of minimum weight in a weighted tree. From now on, let $T = (V, E)$ be a tree rooted in a vertex $r \in V$ and with a weight function $w : V \rightarrow \mathbb{R}^+$.

Question 9 Let $v \in V$ with children v_1, \dots, v_d and let Q be a vertex cover of T that does not contain v (i.e., $v \notin Q$). Explain (in one sentence) why $\{v_1, \dots, v_d\} \subseteq Q$.

Given a vertex v of the rooted tree (T, r) with weight w , let $vc(T_v)$ denote the smallest weight of a vertex cover of T_v . Let $vc^+(T_v)$ denote the smallest weight of a vertex cover of T_v that contains v (i.e., consider all vertex covers of T_v that contains v and consider one such vertex cover with smallest weight), and let $vc^-(T_v)$ denote the smallest weight of a vertex cover of T_v that does not contain v (i.e., consider all vertex covers of T_v that do not contain v and consider one such vertex cover with smallest weight).

Question 10 For every vertex $v \in V$, explain (in one sentence) why $vc(T_v) = \min\{vc^+(T_v), vc^-(T_v)\}$.

Question 11 Let v be a leaf of T (so, T_v is the tree that contains only the vertex v). Give the values of $vc(T_v)$, of $vc^+(T_v)$ and of $vc^-(T_v)$ as a function of $w(v)$.

Question 12 Let $v \in V$ with children v_1, \dots, v_d . Using question 9, explain (in one sentence) why $vc^-(T_v) = \sum_{1 \leq i \leq d} vc^+(T_{v_i})$.

Question 13 Let $v \in V$ with children v_1, \dots, v_d . Explain (in one/two sentences) why $vc^+(T_v) = w(v) + \sum_{1 \leq i \leq d} vc(T_{v_i})$.

Question 14 From previous questions, design an algorithm (write the pseudo code) that takes a weighted tree as input and computes the minimum weight of a vertex cover of T .

Question 15 What is the time-complexity of your algorithm (as a function of the number n of vertices of the input tree)?