## Routing in Multimodal Networks With Bicycles

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#### **Motivation**



- Bicycles An increasingly popular means of transport.
- Need to develop an algorithm that finds an A to B optimal path.
- Cyclists' paths preferences depend not only on distance but on a lot of other path features (slopes, traffic etc)

## **Personalized route planner for bicycles**

- Allows users to navigate road networks optimally.
- Based on individual driving styles as well as personal preferences.
- Takes as input
  - 1) A road network G = (V, E) and a set of cost functions  $c_1 , c_2 , ... c_r$  with  $c_i : E \rightarrow [0,\infty)$  for every metric i.
  - 2) A starting point
  - 3) A destination
  - 4) A set of weights (w<sub>1</sub>, w<sub>2</sub>, ..., w<sub>r</sub>) that determine which metrics the optimal path should be computed based upon.
  - 5) A set of parameters that determine some cyclist's individual riding features.
- Produces as output :
  - A set of vertices that corresponds to the path with minimal weighted cost

## State of the art - Graph Compression and Dijkstra Variations

- 1. Direct continuation of N. Vadakke-Palangatt and M. Zima PFE work
- Graph compression during preprocessing : A very popular approach (18.0 million vertices and 42.5 million edges : Memory, preprocessing time & Query time). Dijkstra 0.4 Gb, -, 2.2 s.
  - a. Hub-labeling 18.8 Gb, 0:37 h, 0.56 µs [D. Delling, A. Goldberg, R. Werneck, 2013]
  - b. Contraction Hierarchy 0.4 Gb, 0:05 h, 110 µs [Robert Geisberger, Peter Sanders, 2012]
  - c. Customizable Route Planning 0.9 Gb, 1:00 h, 1650 μs [D. Delling, A. Goldberg, T. Pajor, R. Werneck, 2014]
  - d. Pruned Landmark Labeling [T. Akiba, Y. Iwata, Y. Yoshida, 2013]
- 3. Dijkstra speedups : Bidirectional Dijkstra, Heuristic Based Dijkstra (A-star)
- 4. Tools used: a) OpenStreetMap b) QGIS c) Sagemath

## **Contribution during the internship**

#### Objectives:

- Find solution for optimal path search that takes into account features specific for cyclists;
- Create a real-world graph that contains these features;
- Make these features balanced in comparison to each other;
- Find different optimal paths for different users according to their preferences.

#### End result:

Working implementation of the algorithm on Nice's graph.

## The approach to Graph Compression : K-path Covers (Funke et al, 2014) & Cover Hierarchy (Akiba et al, 2016)

**k-Path Cover**: In Graph G(V, E), a set  $C \subseteq V$  such that  $C \cap P \neq \emptyset$  for any path P of length k.

- Minimum k-path Cover : A NP-Hard problem
- **k-all-path-cover hierarchy**: Based of vertex-covers (Akiba et al, 2016)
- Idea: Nth layers of vertex cover is the 2<sup>N</sup>-path cover of the original graph.



(a) Original graph  $G_0$ 



(b) Layers of 2-APCs



## Algorithm modification during internship



Densifying the graph

A vertex cover on densified graph

- Reason: to make the algorithm less dependent on graph topology;
- Before creating overlay layers leave as access nodes only those vertices which have more than 2 neighbours.
- Proved to be efficient.

## The approach to Graph Compression : Overlay Graphs

- Maintain all the routes possible among the compressed vertices.
- Route info: To relate edge to corresponding route in original graph.
- Cost info (dist, time etc): Single lookup retrieval of the route cost.



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## **Client - server socket system**

#### Server:

- operates in sage;
- contains precomputed overlay graph;
- receives queries from clients, processes them and sends responses back.
- responses contain lists of vertices that correspond to the optimal path and cost of paths

#### Client:

- operates in QGIS;
- contains full graph;
- sends queries which contain:
  - source node;
  - destination node;
  - user weights;
  - parameters.
- receives responses and presents them to users.

Client and server communicate with each other using socket system

## **Client part**



- User interface for simple source and destination point, user weights selection;
- The result is presented as a line with highlighted nodes.

## **Retrieving the graph**

- Data downloaded from OpenStreetMap.
- Presented as Shapefile (.shp)
- Afterwards converted to Sage object (.sobj)
- Steps to retrieve a working graph:
  - 1) Convert multilines (lines with intermediate points) to edges;
  - Make the graph directed according to 'oneway' tag;
  - Make the graph strongly connected;
  - 4) Add cliques to squares they are denoted with 'place=square' tag.

	osm_id	name	highway	waterway	aerialway	barrier	man_made	other_tags
1	164192418							"boundary"=>"political","political_division"=>"canton"
2	135463949		pedestrian					
3	243900477		footway					"bicycle"=>"dismount","wheelchair"=>"yes"
4	560091690		pedestrian					"tunnel"=>"building_passage"
5	4246694	Descente Crotti	residential					"surface"=>"asphalt"
6	560091691		service					"oneway"=>"yes"
7	4246695	Descente du Marché	steps					"reg_name"=>"calada dóu Mercat"
8	560091692		pedestrian					"tunnel"=>"building_passage"
9	105040811		primary_link					"oneway"=>"yes","surface"=>"asphalt"
10	243591053		footway					"bicycle"=>"dismount","wheelchair"=>"yes"
11	560091693		pedestrian					
12	243591054		footway					"bicycle"=>"dismount","wheelchair"=>"yes"
13	560091694		pedestrian					"tunnel"=>"building_passage"
14	560091695		pedestrian					
15	243591056		footway					"bicycle"=>"dismount","wheelchair"=>"yes"
16	560091696		pedestrian					
17	560091697		pedestrian					"tunnel"=>"building_passage"
18	560091698		pedestrian					"tunnel"=>"building_passage"
19	560091699		pedestrian					
20	243591060	Allée Albert Camus	footway					"bicycle"=>"dismount","lit"=>"yes","wheelchair"=>"yes"
21	137418706							"height"=>"25"
22	560091700		pedestrian					

## Example of OpenStreetMap tags in edges represented in QGIS

## **Retrieving the graph**



The initial graph

#### The processed graph

## **User's input during query**

The user has to choose values for these metrics:

- Travel time [0..1] how fast a user can reach destination;
- **Comfort** [0..1] how comfortable is user's ride;
- Flatness [0..1] how many slopes will the route contain.

and these parameters:

- Speed (m/s);
- Uphill penalty how much uphill ride slows down the user;
- Downhill speed multiplier maximum value of downhill speed;
- Critical downhill grade value when maximum downhill speed is achieved.

More detailed information in appendix.

## Path features which affect cyclist's choice

Feature	Affects	Description	Tags in OpenStreetMap
Distance	time, comfort, flatness	length of the edge in m	
Slope	time, flatness	relation of vertices height difference and distance	
Surface	time, comfort	type of surface which affects speed and comfort (asphalt, cobblestone, gravel etc)	smoothness, surface, tracktype
Highway	comfort	type of road (cycleway, primary, residential, pedestrian etc)	highway, bicycle, cycleway
Slowdown	time	obstacles that make the cyclist stop (crossings, traffic signals, steps etc)	crossing, highway

### **Elevation data**

- Elevation data absent from OpenStreetMap;
- Used SRTM 1 Arc-Second Global from EarthExplorer;
- Every edge given slope value by this formula:

 $slope(u,v) = \frac{height(v) - height(u)}{distance(u,v)}$ 

• Positive if uphill, negative if downhill.



## **Slopes metric problems**





Gentle slopes are easier for cyclists even if they are longer.

Solution: Flatness coefficient which polynomially depends on gradient value. (Crispin H.V. Cooper, 2016)

Less continuous slopes are easier for cyclists.

Solution: Currently an open question.

## Analysis of the algorithm's performance on the map of Nice

Overlay layer	k	Constr. time (s)	# vertices	# edges	D avg	D max	Dijkstra (ms)	Search (ms)	Speed up
Initial	-	-	100768	200155	3.97	56	-	-	-
0	1	12.9	17513	43276	4.94	56	873	421	2.07
1	2	16.8	9356	35529	7.59	146	873	322	2.71
2	4	20.4	5800	36819	12.7	304	873	276	3.16
3	8	24.2	3899	51293	26.31	1183	873	240	3.64
4	16	38.3	2830	137394	97.1	9081	873	280	3.12
5	32	287	2190	694848	634.56	49508	873	614	1.42

- Overlay layer #3 has the best performance for the graph of Nice with 240 ms of search;
- Compare it to 358 ms by the previous version of the algorithm.

Test Machine: Linux machine with 2.10 GHz Intel(R) Core(TM) i3-2310M CPU and 4GB of memory.

# Analysis of the algorithm's performance on the map of New York City

Overlay layer	k	Constr. time (s)	# vertices	# edges	D avg	D max	Dijkstra (ms)	Search (ms)	Speed up
Initial	-	-	240474	431371	3.59	12	-	-	-
0	1	27.9	66435	167805	5.05	12	2226	1464	1.52
1	2	38.4	38874	153786	7.91	24	2226	1206	1.85
2	4	50.3	26105	171377	13.13	72	2226	947	2.35
3	8	61	19110	231061	24.18	172	2226	994	2.24
4	16	81	15100	375574	49.74	1129	2226	1061	2.1

- About 2.2 times larger than graph of Nice;
- Overlay layer #2 has the best performance for the graph of New York with 947 ms of search;
- The speed-up is worse than the graph of Nice had.

#### Conclusions

#### Achievements

- A shortest path algorithm that takes into account cyclist's needs is designed and successfully tested;
- Its performance was made more independent on graph topology;
- A real-world graph with data important for cyclists was created;
- Implemented the diameter search DiFUB algorithm.

## Further development

- Find a good default ratio between metrics;
- Develop a continuous slopes metric;
- Consider individual user's preferences;
- Expand the graph to the whole PACA region;
- Implement less curves metric;
- Implement several optimal paths search, not only one.

# Thank you!

# Appendices

### **The Vertex Cover Problem and solution**

#### Main Steps in my compression implementation :

- 1. Create an overlay graph where crossroads are access points
- 2. Find Vertex cover of the previous layer
- 3. Create overlay graph for the vertex cover.

#### My Solution: Custom implement a vertex cover heuristic (LR-deg).

**LR-deg**: Initialize Vertex Cover VC to an empty Set. For each  $v \in V$  (v picked in increasing order of degree), add Neighbor(v) to VC if v not already in VC.

#### Real time Querying: Funke's algorithm

A fast bidirectional dijkstra using access points (H. Bast et al, 2007)

#### Formulas used in the algorithm

group	entity	key	value	r time	r slowdown	r surface	r traffic
crossing	node	crossing	island	1	20	-1	-
crossing	node	crossing	traffic signals	1	30	-1	-
crossing	node	crossing	uncontrolled	1	15	-1	-
crossing	node	crossing	unmarked	1	20	-1	-
crossing	node	crossing	ves	1	15	-1	
crossing	node	crossing	zebra	1	15	-1	
crossing	node	highway	crossing	1	15	-1	
crossing	node	highway	traffic signals	1	30	-1	
dismount	Waw	hicycle	dismount	0.4	0	2	
diemount	way	footway	croccing	0.4	0	2	
dismount	way	footway	sidowalk	0.4	0	-1	0.5
dismount	way	highwow	footwark	0.4	0	-1	0.5
dismount	way	highway	footway	0.4	0	-1	0.5
diamount	way	highway	nodestrian	0.4	0	1	0.5
uismount fas bisuslas	way	ilignway	peuestian	0.4	0	-1	0.5
lor_bicycles	relation	route	bicycle	1	0	-1	0.9
lor_bicycles	way	bicycle	designated	1	0	-1	0.2
for_bicycles	way	Dicycle	permissive	1	0	-1	-
tor_bicycles	way	bicycle	yes	1	0	-1	-
tor_bicycles	way	cycleway	lane	1	0	-1	0.6
for_bicycles	way	cycleway	share_busway	1	0	-1	0.7
for_bicycles	way	cycleway	shared_lane	1	0	-1	0.8
for_bicycles	way	cycleway	track	1	0	-1	0.4
for_bicycles	way	cycleway:left	lane	1	0	-1	0.6
for_bicycles	way	cycleway:left	share_busway	1	0	-1	0.7
for_bicycles	way	cycleway:left	shared_lane	1	0	-1	0.8
for_bicycles	way	cycleway:right	lane	1	0	-1	0.6
for bicycles	way	cycleway:right	share_busway	1	0	-1	0.7
for bicycles	way	cycleway:right	shared lane	1	0	-1	0.8
for bicycles	way	highway	cycleway	1	0	-1	0.2
motor roads	way	highway	living street	1	0	-1	0.5
motor roads	way	highway	primary	1	0	-1	1
motor roads	way	highway	primary link	1	0	-1	1
motor roads	way	highway	residential	1	0	-1	
motor roads	way	highway	secondary	1	0	-1	
motor roads	way	highway	secondary link	1	0	-1	
motor roads	way	highway	service	1	0	-1	
motor roads	way	highway	tertiary	1	0	-1	
motor roads	way	highway	tertiary link	1	0	-1	
obstacles	Way	highway	stone	0.1	0	10	
obstacles	node	highway	elevator	1	75	7	
obstacles	nodo	highway	ctone	1	25	10	
offraad	noue	highway	bridlowow	0.7	25	10	
offroad	way	ngnway	ogrigultural	0.0	0	2	-
offroad	way	access	forestry	0.0	0	2	
offreed	way	ducess	noth	0.0	0	2	0.5
oliroad	way	nignway	paur	0.7	0	2	0.5
oiroau	way	nignway	lack	0.0	0	2	0.5
sunace	way	smoothness	bad	0.7	0	3	-
surface	way	smoothness	excellent	1	0	0.5	-
surrace	way	smoothness	norrible	0.5	0	2	-
surface	way	smoothness	Intermediate	0.8	0	1	-
surface	way	smoothness	very_bad	0.6	0	4	-
surface	way	surface	cobblestone	0.7	0	5	-
surface	way	surface	compacted	0.9	0	1.5	-
surface	way	surface	dirt	0.7	0	3	-
surface	way	surface	grass	0.65	0	5	-
surface	way	surface	gravel	0.5	0	5	-
surface	way	surface	ground	0.6	0	4	l.
surface	way	surface	mud	0.4	0	5	-
surface	way	surface	paving stones	0.75	0	1.5	
surface	way	surface	sand	0.6	0	4	-
surface	way	surface	setts	0.8	0	2	-
surface	way	surface	unnaved	0.75	0	4	
curfaco	MON	curfaco	an part of a	0.05	0	-	

#### Travel time value:

$$c_1(u,v) = \begin{cases} \frac{l(u,v) + a_l \cdot a(u,v) \cdot l(u,v)}{r_1(u,v)} + q(u,v) * s & \text{if } a_l > 0, \\ \frac{l(u,v)}{s_d(u,v,s_{dmax}) \cdot r_1(u,v)} + q(u,v) * s & \text{otherwise}, \end{cases}$$

where l(u, v) - edge distance, a(u, v) - edge slope, r1 - r\_time coefficient, q(u, v) - r\_slowdown value, s - speed (m/s), sd - downhill speed multiplier

 $s_d(u, v, s_{dmax}) := \begin{cases} s_d max & \text{if } d'(u, v) > d'_c, \\ \frac{(s_{dmax} - 1)d'(u, v)}{d'_c} + 1 & \text{otherwise}, \end{cases}$ 

where sdmax - maximum downhill speed multiplier, d'(u, v) - edge slope, d'c - critical d' value when speed equals sdmax

#### Comfort value:

 $c_2(u, v) = l(u, v) \cdot \max\{(r_s(u, v), r_t(u, v))\}$ 

```
where rs - r_surface, rt - r_traffic
```

#### Flatness value:

 $c_3(u,v) = \begin{cases} l(u,v) \cdot (10128.074 \cdot a(u,v)^3 - 140.785 \cdot a(u,v)^2 + 6.693 * a(u,v) + 1) & \text{if } a(u,v) > 0, \\ l(u,v) & \text{otherwise.} \end{cases}$ 

#### Analysis of the initial algorithm's performance on the map of Nice

Overlay layer	k	Constru ction time (s)	# vertices	# edges	D avg	D max	Dijkstra (ms)	Search (ms)	Speed up
0	1	7.67	100768	200155	3.97	56	873	1898	0.46
1	2	23.1	55552	119675	4.31	146	873	975	0.9
2	4	32.6	30646	79887	5.21	322	873	561	1.56
3	8	38.7	17169	71627	8.34	1034	873	458	1.91
4	16	49.2	9923	114459	23.07	4263	873	358	2.44
5	32	138	6070	365236	120.34	28338	873	466	1.87

#### Example of optimal path



- The optimal path avoids:
  - hills;
  - major roads;

#### **Diameters of the graph**



The paths which are:

- the most time-consuming;
- the least comfortable;
- the most hilly.