# Revisiting OCaml 

Lecture 2

Formal Languages and Compilers 2011

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## How to run OCaml

- Run the interpreter with
ocaml
- Save the file in "myfile.ml", let the interpreter run it from file ocaml
\#use "myfile.ml"
- Compilation of a single module ocamlc -c myfile.ml
Results in myfile.cmo
- Then use the compiled file in the interpeter:
ocaml
\#load "myfile.cmo";;
open Myfile;;
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# Value binding and pattern matching 

- let ( $x, y$ ) = ("hi", (1,2));;
- let (a, (b,c)) = (z, ( 3,4 ));;


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- let h::t = [4] $::[5 ; 6] ; ;$


## Value binding and pattern matching

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- let (a, (b,c)) = (z, (3,4));;
- let h:: $\dagger=[4 ; 5 ; 6] ;$;
- let h::t = [4] $::[5 ; 6] ; ;$
- let $x=1$ and $y=2$ in $x^{*} y ;$;
- let $\mathrm{a}=3$ and $\mathrm{b}=4$ in $\mathrm{c}=\mathrm{a}+\mathrm{b} ; ;$
- let $a=3$ and $b=4$ in $c=a+b$ in $c+2 ;$;


## Functions

- fun $x$-> $\left(x^{*} 2, x^{*} 4, x^{*} 8\right) ;$;
- let $\mathrm{f} x=\mathrm{x}^{*} 2 ;$;
- let $y=(f 2)$ in $y^{*} 2 ;$;


## Functions

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else 0;;


## Functions

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- let $\mathrm{f} x=\mathrm{x}^{*} 2$;
- let $y=(f 2)$ in $y^{*} 2 ;$;
- let $f x=$ if $x>0$ then $x$
else 0;;
- String.length;;
- String.contains;;

Lists

- List.rev;;
- List.hd;;
- List.t|;

Lists

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- List.hd;;
- List.t|;
- List.hd [1;2;3];;
- List.hd (List.tl [4;5;6]);;

Lists

- List.rev;;
- List.hd;;
- List.t|;
- List.hd [1;2;3];;
- List.hd (List.tl [4;5;6]);;
- List.append;;
- the same as list1@list2
- [1;2;3]@[4;5];;


## Recursive functions

$$
\begin{aligned}
& \text { let rec } f 1=\text { function } \\
& \quad \mid 0->0 \\
& \mid n->n+f 1(n-1)
\end{aligned}
$$

let rec $f 2 \mathrm{n}=$ match n with

$$
\begin{aligned}
& \mid 0->0 \\
& 1 \mathrm{n}->\mathrm{n}+\mathrm{f} 2 \mathrm{n}-1
\end{aligned}
$$

let rec f3 n m = match n with

$$
\text { | } 0 \text {-> m }
$$

$$
\text { |n }->\text { f3 }(n-1) m+n
$$

## Try an exercise!

- Given a list of string I, define a function find that builds a new list that contains elements from / such that the length of each element is less or equal than 3.
- The order of elements should be preserved.
- For example, if I = ["12"; "abcd"; "www"; "456"]
then result is [" 12 "; "www"; "456"]


## Compilers and Interpreters

Lection 2

## Running OCaml

- Run the interpreter with
- ocaml
- Exit the interpreter:
- \# quit;;

- Compilers:
- ocamlc compiles in bytecode
- Compilation of a single module
- ocamlc -c <fileName>.ml
- Produces <fileName>.cmo


## Compiler

| Source |
| :---: | :---: | :---: |
| code |\(\xrightarrow{Single translation} \xrightarrow[\begin{array}{c}Executable <br>

code\end{array}]{ }\)
in high-level
language
in machine language

- If an error is found, the source code is not converted


## Interpreter


in high-level
language
in some intermediate language

- If an error is found in a statement, the interpreter stops working and shows an error


## Compiler vs. Interpreter



## Compiler vs. Interpreter

- Compiler characteristics:
- spends a lot of time analyzing and processing the program
- the resulting executable is some form of machine- specific binary code
- the computer hardware interprets (executes) the resulting code
- program execution is fast


## Compiler vs. Interpreter

- Interpreter characteristics:
- relatively little time is spent analyzing and processing the program
- the resulting code is some sort of intermediate code
- the resulting code is interpreted by another program
- program execution is relatively slow


## Some real life examples

- C++ compiler
- Java with its Java Virtual Machine (JVM) is something in between, more similar to interpreter

- Java compiler transforms source program to Java bytecode
- JVM is an interpreter of the bytecode
- JIT (Just-In-Time) compiles parts of the bytecode to executable code


## Structure of a compiler



## Front-end structure



## Back-end structure

- is responsible for emitting the final (executable) version of the source program. Typical parts of the back end are responsible for:
- instruction selection
- register allocation
- memory management
- instruction scheduling


## Front-end and back-end



- Reuse the same front-end for different machines
- Reuse the same back-end for different source languages


## References

- CS544:
http://web.cs.wpi.edu/~gpollice/cs544-f05/ CourseNotes/maps/Class 1/ Compilervs.Interpreter.html

