

Game Theory: introduction and applications to computer networks

Introduction

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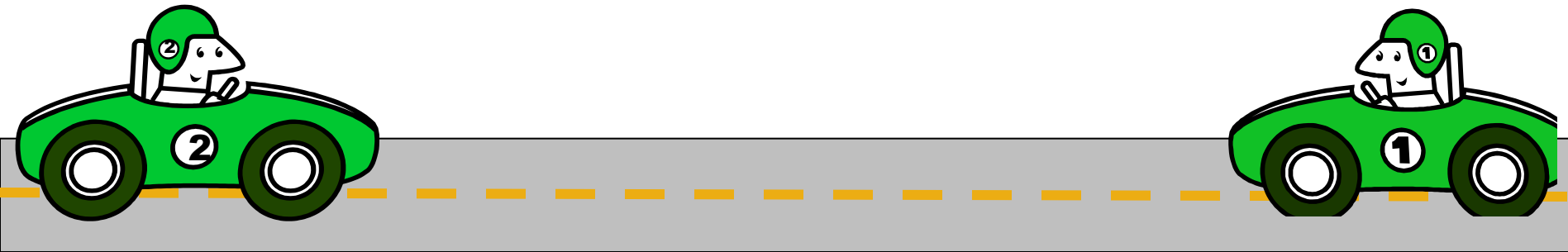
INRIA – EPI Maestro

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Part of the slides are based on a previous course
with D. Figueiredo (UFRJ) and H. Zhang (Suffolk University)

What is Game Theory About?

- ❑ Mathematical/Logical analysis of situations of conflict and cooperation



- ❑ Game of Chicken
 - driver who steers away loses
 - what should drivers do?
- ❑ Goal: to prescribe how rational players should act

What is a Game?

- ❑ A Game consists of
 - at least two players
 - a set of strategies for each player
 - a preference relation over possible outcomes
- ❑ Player is general entity
 - individual, company, nation, protocol, animal, etc
- ❑ Strategies
 - actions which a player chooses to follow
- ❑ Outcome
 - determined by mutual choice of strategies
- ❑ Preference relation
 - modeled as utility (payoff) over set of outcomes

Short history of GT

- Forerunners:
 - Waldegrave's first minimax mixed strategy solution to a 2-person game (1713), Cournot's duopoly (1838), Zermelo's theorem on chess (1913), Borel's minimax solution for 2-person games with 3 or 5 strategies (20s)
- 1928: von Neumann's theorem on two-person zero-sum games
- 1944: von Neumann and Morgenstern, *Theory of Games and Economic Behaviour*
- 1950-53: Nash's contributions (Nash equilibrium, bargaining theory)
- 1952-53: Shapley and Gillies' core (basic concept in cooperative GT)
- 60s: Aumann's extends cooperative GT to non-transferable utility games
- 1967-68: Harsanyi's theory of games of incomplete information
- 1972: Maynard Smith's concept of an Evolutionarily Stable Strategy
- Nobel prizes in economics
 - 1994 to Nash, Harsanyi and Selten for "their pioneering analysis of equilibria in the theory of non-cooperative games"
 - 2005 to Aumann and Schelling "for having enhanced our understanding of conflict and cooperation through game-theory analysis"
 - 2012 to Roth and Shapley "for the theory of stable allocations and the practice of market design"
- Movies:
 - 2001 "A beautiful mind" on John Nash's life
- See also:
 - www.econ.canterbury.ac.nz/personal_pages/paul_walker/gt/hist.htm

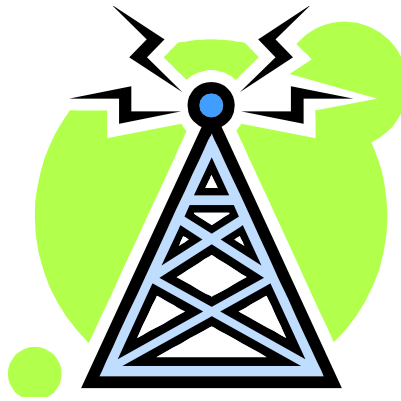
Applications of Game Theory

- Economy
- Politics (vote, coalitions)
- Biology (Darwin's principle, evolutionary GT)
- Anthropology
- War
- Management-labor arbitration
- Philosophy (morality and free will)
- National Football league draft

Applications of Game Theory

- “Recently” applied to computer networks
 - Nagle, RFC 970, 1985
 - “datagram networks as a multi-player game”
 - wider interest starting around 2000
- Which are the strategies available?
 - Network elements follow protocol!!!

Power games



$$SNIR_1 = \frac{H_{1,BS} P_1}{N + H_{2,1} P_2}$$



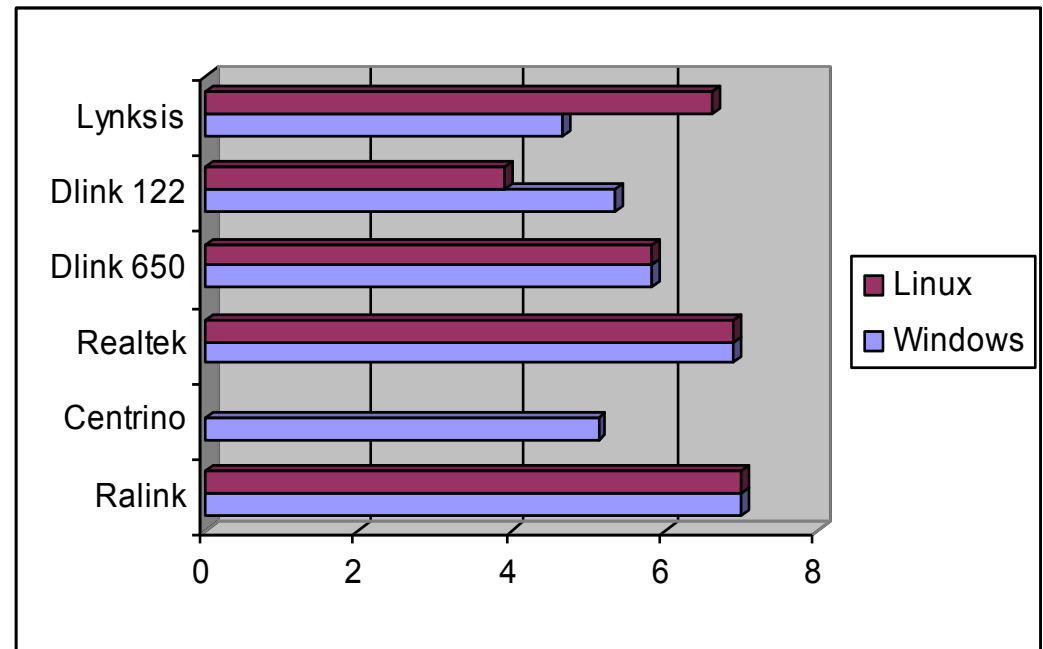
Medium Access Control Games



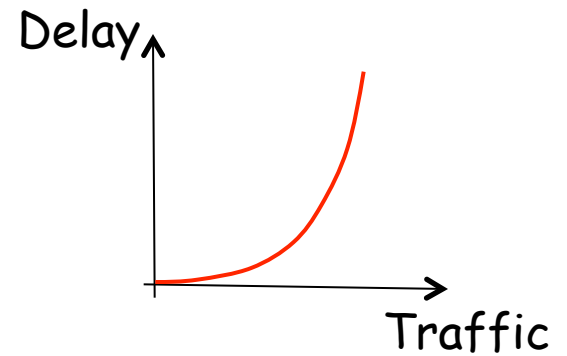
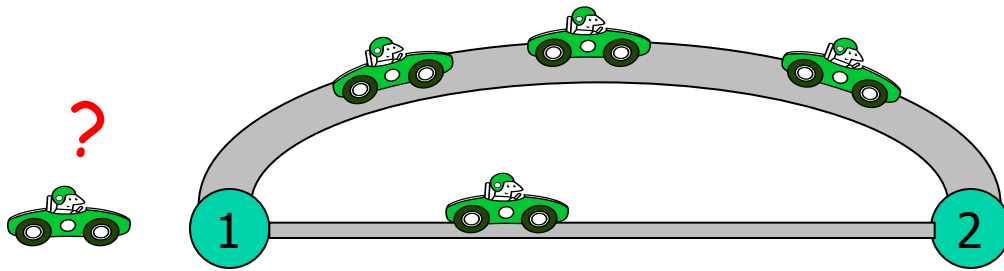
$$Thr_1 = \frac{p_1(1-p_2)P}{(1-p_1)(1-p_2)\sigma + [1-(1-p_1)(1-p_2)]T}$$

Medium Access Control Games

- Despite of the Wi-Fi certification, several cards exhibit very heterogeneous performance, due to arbitrary protocol implementations
 - “Experimental Assessment of the Backoff Behavior of Commercial IEEE 802.11b Network Cards,” G Bianchi et al, INFOCOM 2007

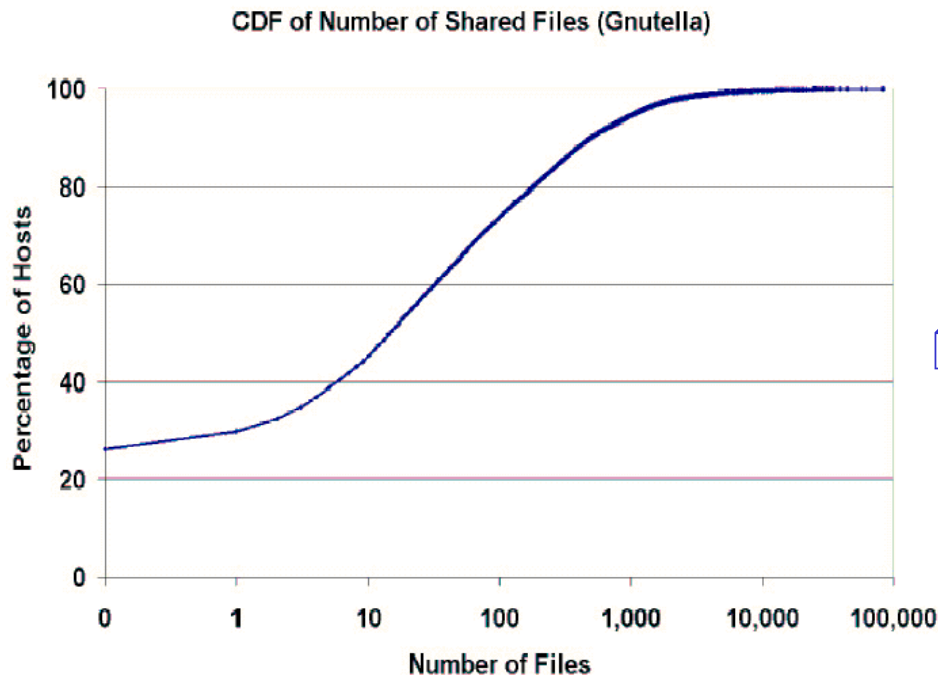


Routing games



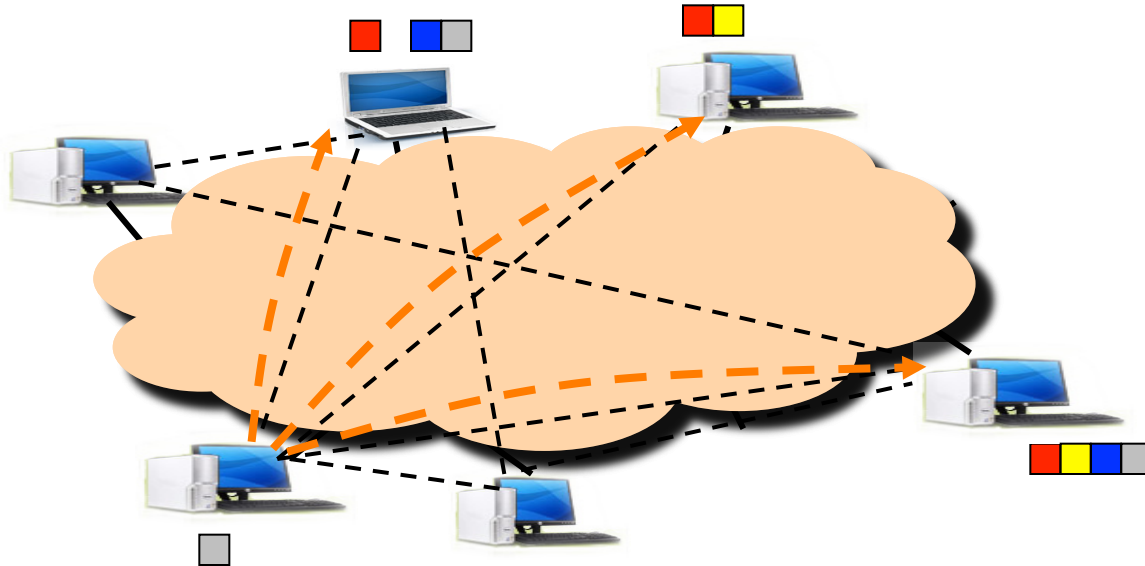
- Possible in the Internet (see later)

Free riders in P2P networks



- ❑ Individuals not willing to pay the cost of a public good, they hope that someone else will bear the cost instead
- ❑ Few servers become the hot spots:
Anonymous?,
Copyright?, Privacy?
Scalability?, Is it P2P?

Connection games in P2P



- Each peer may open multiple TCP connections to increase its downloading rate

Diffusion of BitTorrent variants

- Try to exploit BitTorrent clients weaknesses



BitThief



- Are they really dangerous?
 - Evolutionary game theory says that
Yes they can be

Space for GT in Networks

- ❑ User behaviors (to share or not to share)
 - Client variants
- ❑ Protocols do not specify everything...
 - power level to use
 - number of connections to open
- ❑ ...and/or are not easy to enforce
 - how control a P2P network
 - not-compliant WiFi implementation
- ❑ ...and software easy to modify

Limitations of Game Theory

- ❑ Real-world conflicts are complex
 - models can at best capture important aspects
 - ❑ Players are considered rational
 - determine what is best for them given that others are doing the same
 - Men are not, but computers are more
 - ❑ No unique prescription
 - not clear what players should do
- ❑ But it can provide intuitions, suggestions and partial prescriptions
 - the best mathematical tool we have

Syllabus

- References
 - [S] Straffin, *Game Theory and Strategy* (main one, chapters indicated)
 - [EK] Easley and Kleinberg, *Network Crowds and Markets*
 - [OR] Osborne and Rubinstein, *A course in game theory*, MIT Press
- Two-person zero-sum games
 - Matrix games
 - Pure strategy equilibria (dominance and saddle points), [S2]
 - Mixed strategy equilibria, [S3]
 - Game trees (?), [S7]
- Two-person non-zero-sum games
 - Nash equilibria...
 - ...And its limits (equivalence, interchangeability, Prisoner's dilemma), [S11-12]
 - Subgame Perfect Nash Equilibria (?)
 - Routing games [EK8]
- Auction theory [EK9]

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Two-person zero-sum games

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Matrix Game (Normal form)

Strategy set for Player 1

Player 2, Colin

Strategy set for Player 2

		Player 2, Colin		
		A	B	C
Player 1, Rose	A	(2, 2)	(0, 0)	(-2, -1)
	B	(-5, 1)	(3, 4)	(3, -1)

Payoff to Player 1

Payoff to Player 2

□ Simultaneous play

- players analyze the game and then write their strategy on a piece of paper

More Formal Game Definition

□ Normal form (strategic) game

- a finite set N of players
- a set strategies S_i for each player $i \in N$
- payoff function $u_i(s)$ for each player $i \in N$
 - where $s \in S = \times_{j \in N} S_j$ is an outcome
 - sometimes also $u_i(A, B, \dots)$ $A \in S_1, B \in S_2, \dots$
 - $u_i : S \rightarrow \mathfrak{R}$

Two-person Zero-sum Games

- One of the first games studied
 - most well understood type of game
- Players interest are strictly opposed
 - what one player gains the other loses
 - game matrix has single entry (gain to player 1)
- A “strong” solution concept

Let's play!


		Colin			
		A	B	C	D
Rose	A	12	-1	1	0
	B	5	1	7	-20
	C	3	2	4	3
	D	-16	0	0	16

- Divide in pairs, assign roles (Rose/Colin) and play 20 times
- Log how many times you have played each strategy and how much you have won

Analyzing the Game

		Colin			
		A	B	C	D
Rose	A	12	-1	1	0
	B	5	1	7	-20
	C	3	2	4	3
	D	-16	0	0	16

dominated
strategy
(dominated by B)



Dominance

- Strategy S (*weakly*) dominates a strategy T if every possible outcome when S is chosen is at least as good as corresponding outcome in T , and one is strictly better
 - S strictly dominates T if every possible outcome when S is chosen is strictly better than corresponding outcome in T
- Dominance Principle
 - rational players never choose dominated strategies
- Higher Order Dominance Principle
 - iteratively remove dominated strategies

Higher order dominance may be enough

		Colin			
		A	B	C	D
Rose	A	12	-1	1	0
	B	3	1	4	-18
	C	5	2	4	3
	D	-16	0	5	-1

Higher order dominance may be enough

GT prescribes:

Rose C - Colin B

		Colin			
		A	B	C	D
Rose	A	12	-1	1	0
	B	3	1	4	-18
	C	5	2	4	3
	D	-16	0	5	-1

(Weakly)
Dominated
by C

A priori
D is **not**
dominated
by C

Strictly
dominated
by B

... but not in the first game

		Colin			
		A	B	C	D
Rose	A	12	-1	1	0
	B	5	1	7	-20
	C	3	2	4	3
	D	-16	0	0	16

dominated
strategy
(dominated by B)

Analyzing the Reduced Game: Movement Diagram

		Colin		
		A	B	D
Rose	A	12	-1	0
	B	5	1	-20
	C	3	2	3
	D	-16	0	16

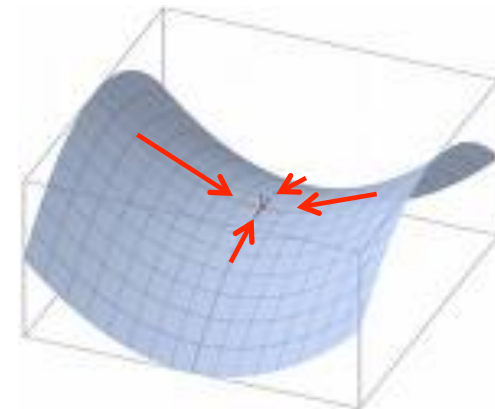
Outcome (C, B) is
“stable”

- *saddle point of game*
- *mutual best responses*

Saddle Points

- An outcome (x,y) is a *saddle point* if the corresponding entry $u(x,y)$ is both less than or equal to any value in its row and greater than or equal to any value in its column
 - $u(x,y) \leq u(x,w)$ for all w in $S_2 = S_{\text{Colin}}$
 - $u(x,y) \geq u(v,y)$ for all v in $S_1 = S_{\text{Rose}}$

	A	B	D
A	12	-1	0
B	5	1	-20
C	3	2	3
D	-16	0	16



Saddle Points Principle

- Players should choose outcomes that are saddle points of the game
 - Because it is an equilibrium...
 - ... but not only