



SOFTWARE VERIFICATION AND COMPUTER PROOF (lesson 2)

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Lessons

1. Writing programs in Coq
2. Writing proofs in Coq
3. Proving your recursive programs correct
4. Proofs in arithmetic and more
5. More data types

Today: writing proofs in Coq

- Morning: proofs with simple connectives
- Afternoon: equality and quantifiers

Propositions

- Propositions do have a type too: **Prop**
- For this morning we assume P Q R are predicates
Variables P Q R : Prop.
- We focus on the logical connectives

Connectives

- Conjunction of P and Q is written $P \wedge Q$
- Disjunction of P and Q is written $P \vee Q$
- Negation of P is written $\sim P$
- P implies Q is written $P \rightarrow Q$

Connectives and related commands

connective	use	prove
$P \wedge Q$	destruct	split
$P \vee Q$	destruct	left, right
$P \rightarrow Q$	apply	intros
$\sim P$	case	

Proof commands

1. I'll present them with a demo
2. You will have a short reference for the exercises
3. You will find more examples on the course notes
4. The full documentation can be found in the user manual of Coq <https://coq.inria.fr/distrib/current/refman/>

Demo

Equality

- The most useful predicate is equality, E.g. $x = 3$
- Equality means:
“the two expressions *compute* the same value”.
- E.g.
 $3 + 7 = 10$ (* nothing to prove *)
 $x=y \rightarrow y=z \rightarrow x=z$ (* something to prove *)

Commands for equality

predicate	use	prove
$a = b$	rewrite \rightarrow ... rewrite \leftarrow ...	reflexivity

Quantifiers

- Universal quantification is written **forall** $(n : T), \dots$
- Existential quantification is written **exists** $(n : T), \dots$
- E.g. **forall** $(n : \text{nat}), \text{exists } (m : \text{nat}), \text{leb } n (n + m) = \text{true}$.

More on quantifiers

- Quantifiers can talk about almost everything, not just numbers, lists, etc. E.g.

$\text{forall } (P : \text{Prop}) (Q : \text{Prop}), ((P \rightarrow Q) \wedge P) \rightarrow Q$

- Moreover we can use some syntactic sugar. E.g.

$\text{forall } (P Q : \text{Prop}), (P \rightarrow Q) \wedge (P \rightarrow Q)$

$A \leftrightarrow B$ instead of $(A \rightarrow B) \wedge (B \rightarrow A)$

Commands for quantifiers

quantifier	use	prove
forall n, P	apply apply ... with ($n := v$)	intros
exists n, P	destruct	exists

Searching

- To complete the exercises *you must search* the library for some theorems about arithmetic
- They are too “hard” to be proved by you (for today)
- The command **SearchAbout pattern** finds all theorems whose statement matches pattern. E.g.

SearchAbout (_ * _).

SearchAbout (_ + _ = _).

SearchAbout (_ + _ * _).

Demo

That is all for today!