

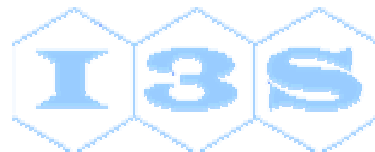


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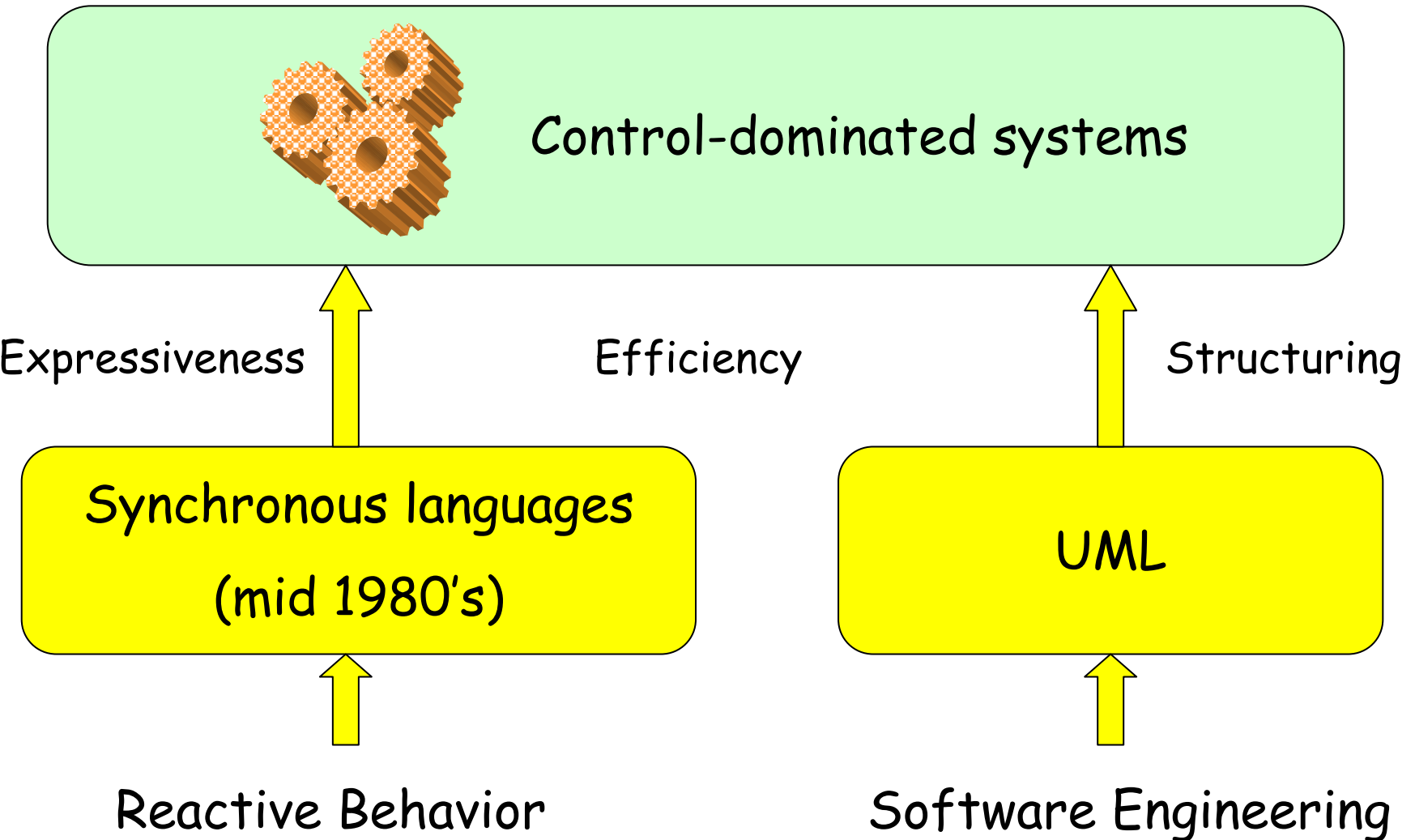


Synchronous Interface Behavior

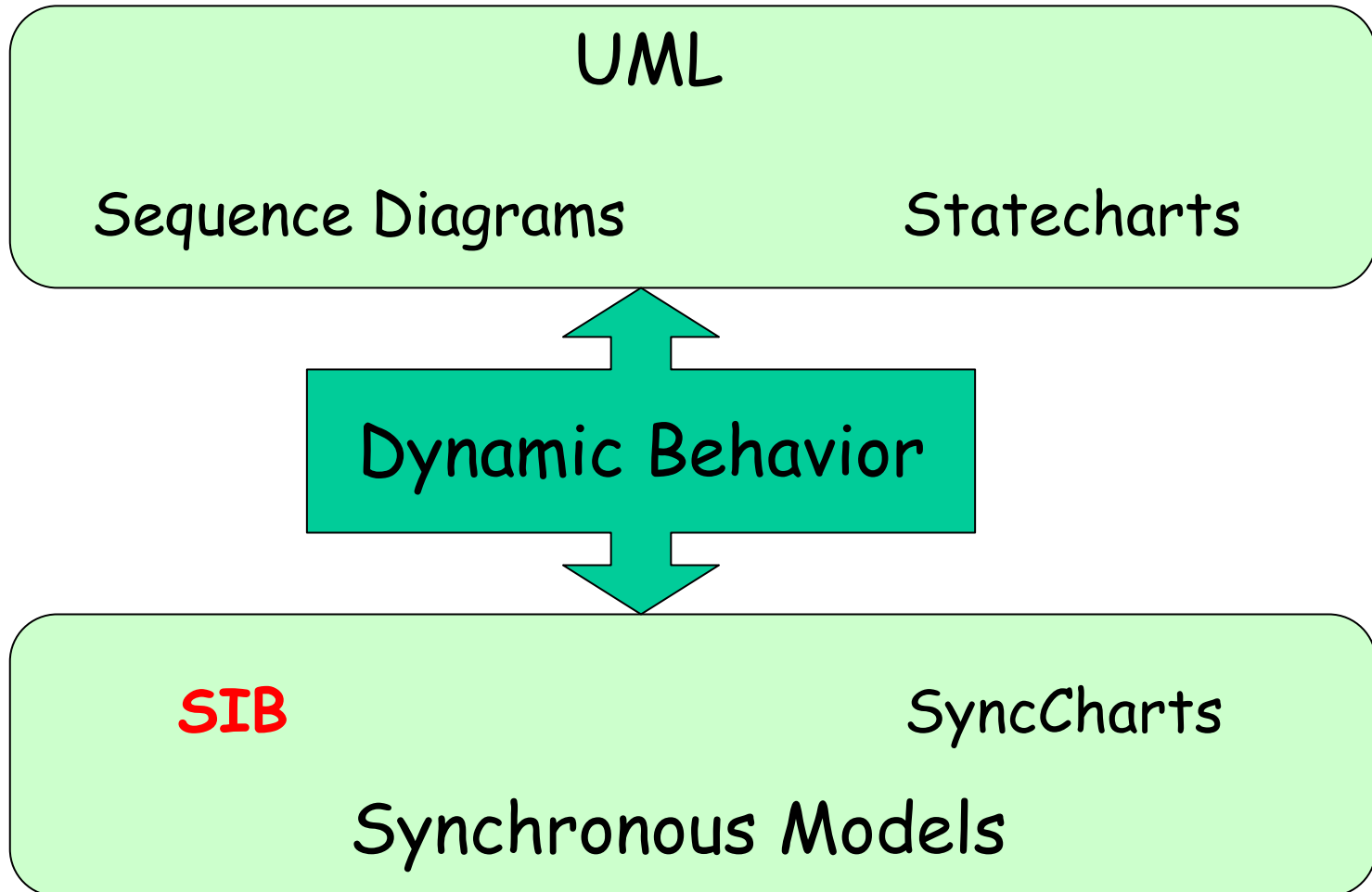
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Motivation



Motivation (cont'd)



Sequences

Sequence
Diagrams

simplicity
legibility

Message Sequence
Charts

structure
semantics

**Synchronous Interface
Behaviors**

Esterel programs

The Synchronous Approach

- **Abstract** and ideal view of reactive systems
- **Observation**: Discrete instants
- **Execution** model: 0-duration
- **Communication**: Instantaneous broadcast of signals
- **Semantics**: mathematically defined
- Support for validation

SIB: overview

- Syntax
 - Graphical language
 - Block-structured
 - A few constructs
- Semantics
 - Mathematical semantics
Synchronous process algebra

a SIB $\xrightarrow{\text{Semantically equivalent}}$ an Esterel program

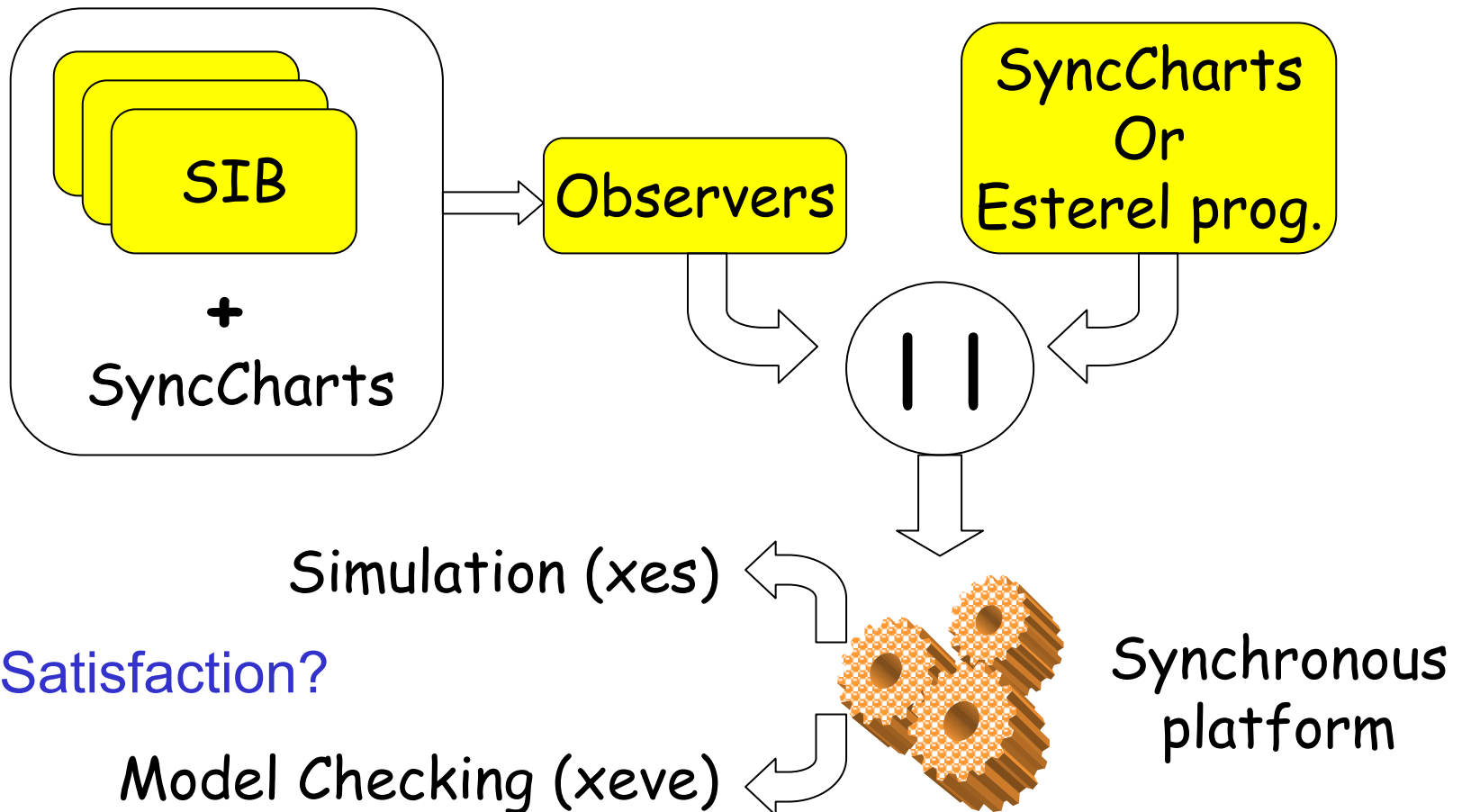
SIB: What for?

- Expressing expected behaviors
 - Discrete-event model
 - Partial order of events (signals)
 - Temporal constraints
- Interactive simulation
- Property checking

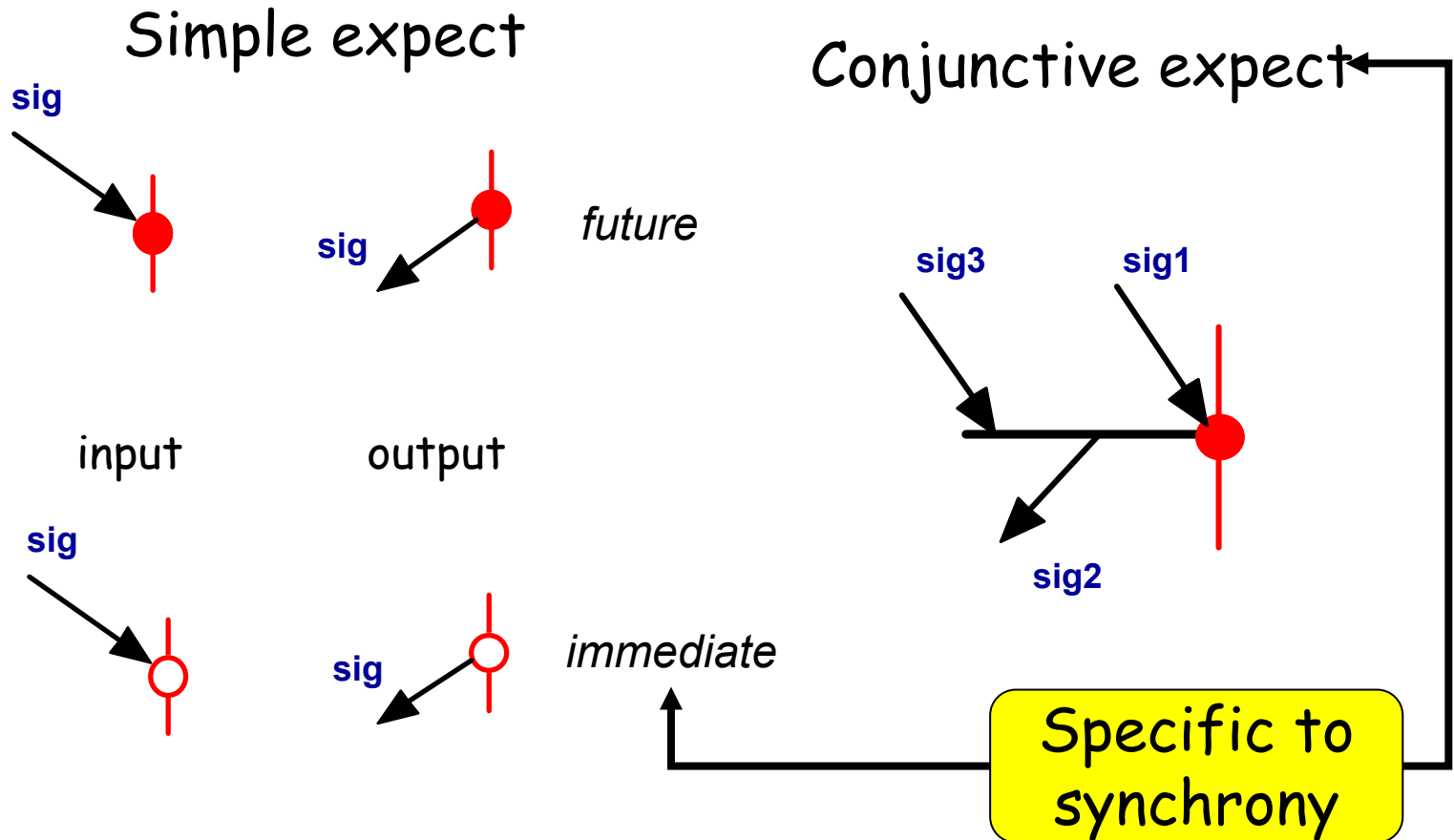
SIB: How?

Expected behavior

Controller (to check)



Syntax: Expect

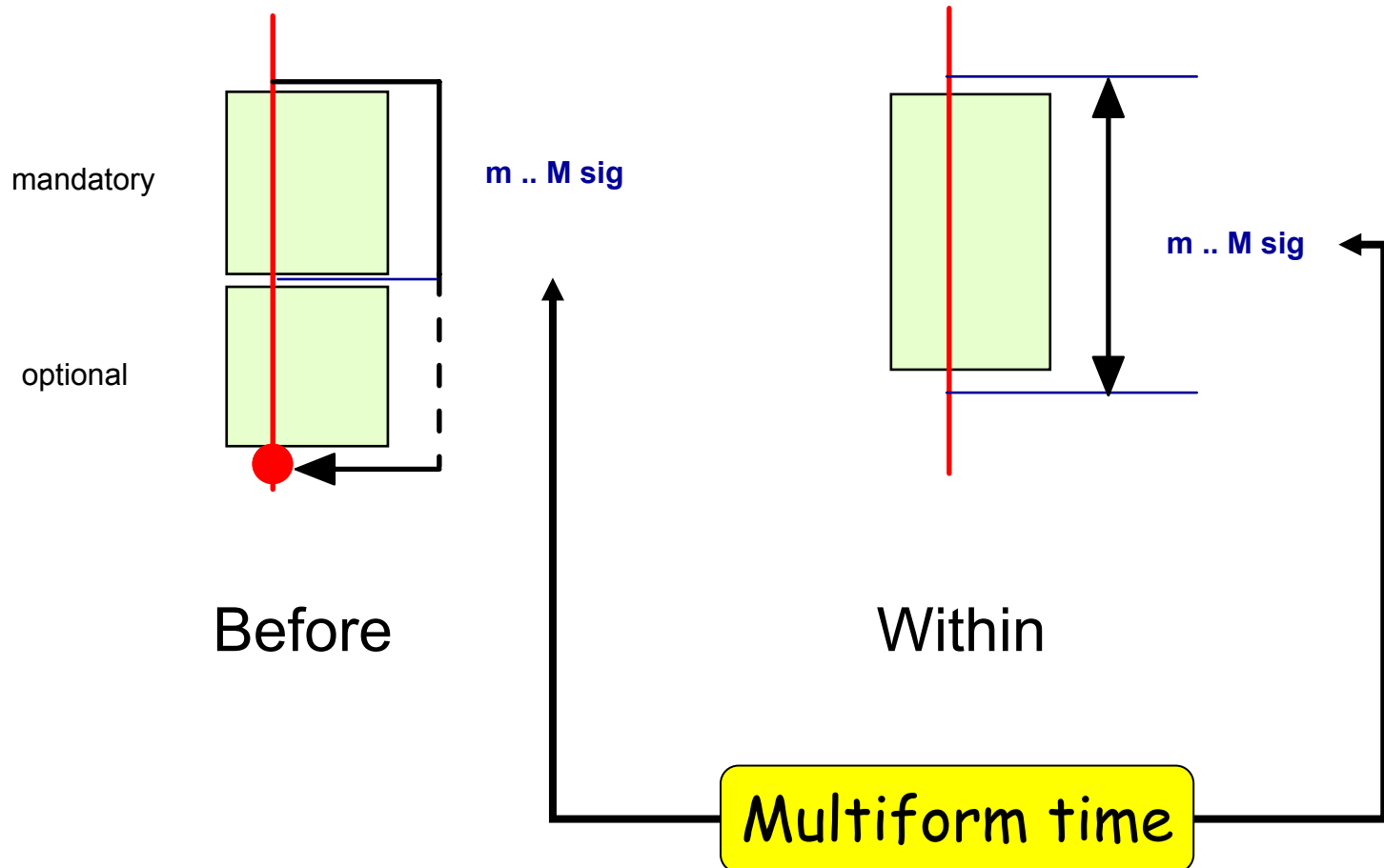


Other Constructs

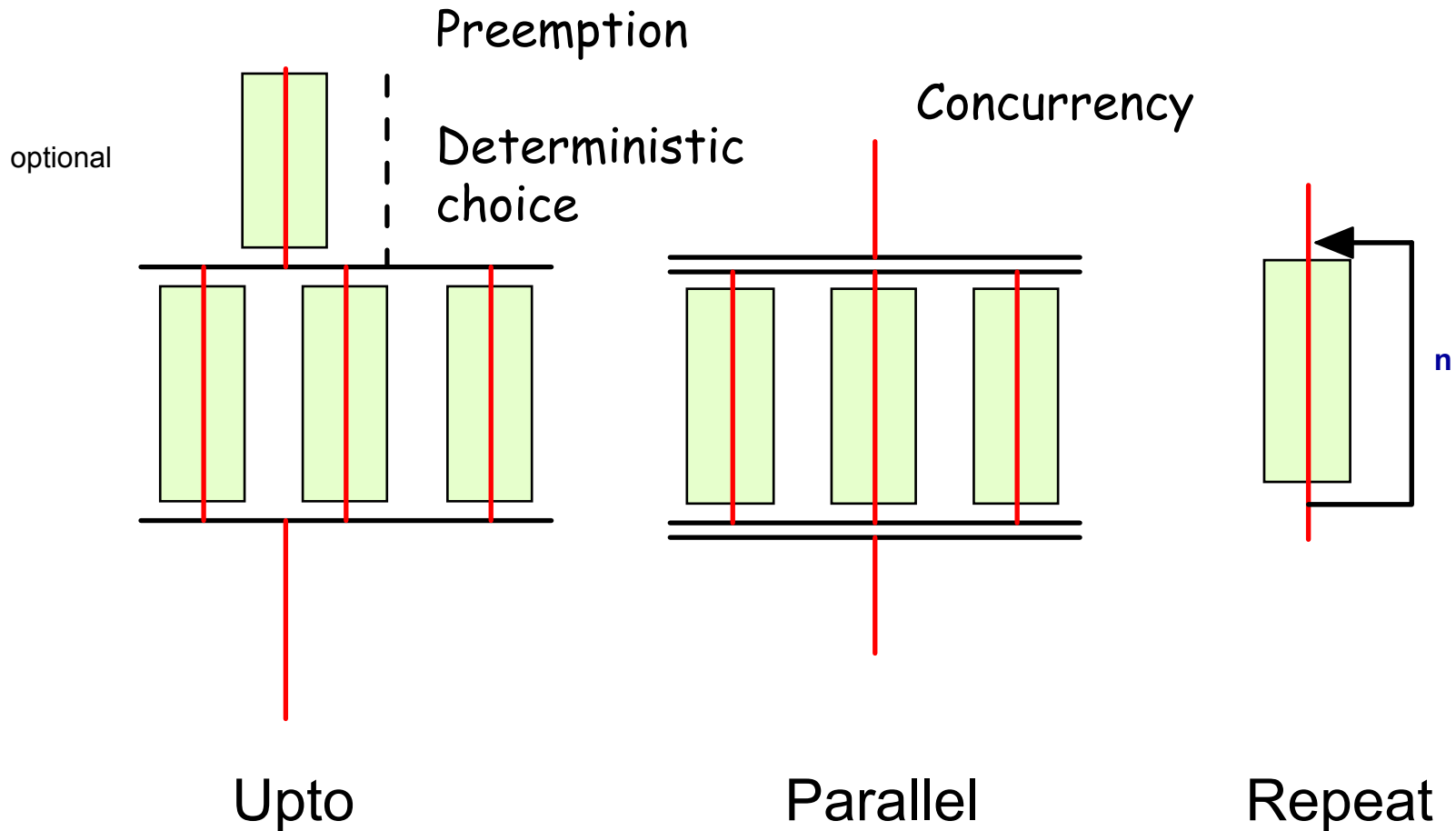
(block-based description)

- Temporal constraints
 - Before (with some degree of non determinism)
 - Within
 - Watchdog
- General control
 - Upto (pre-emption + deterministic choice)
 - Parallel
 - Repeat

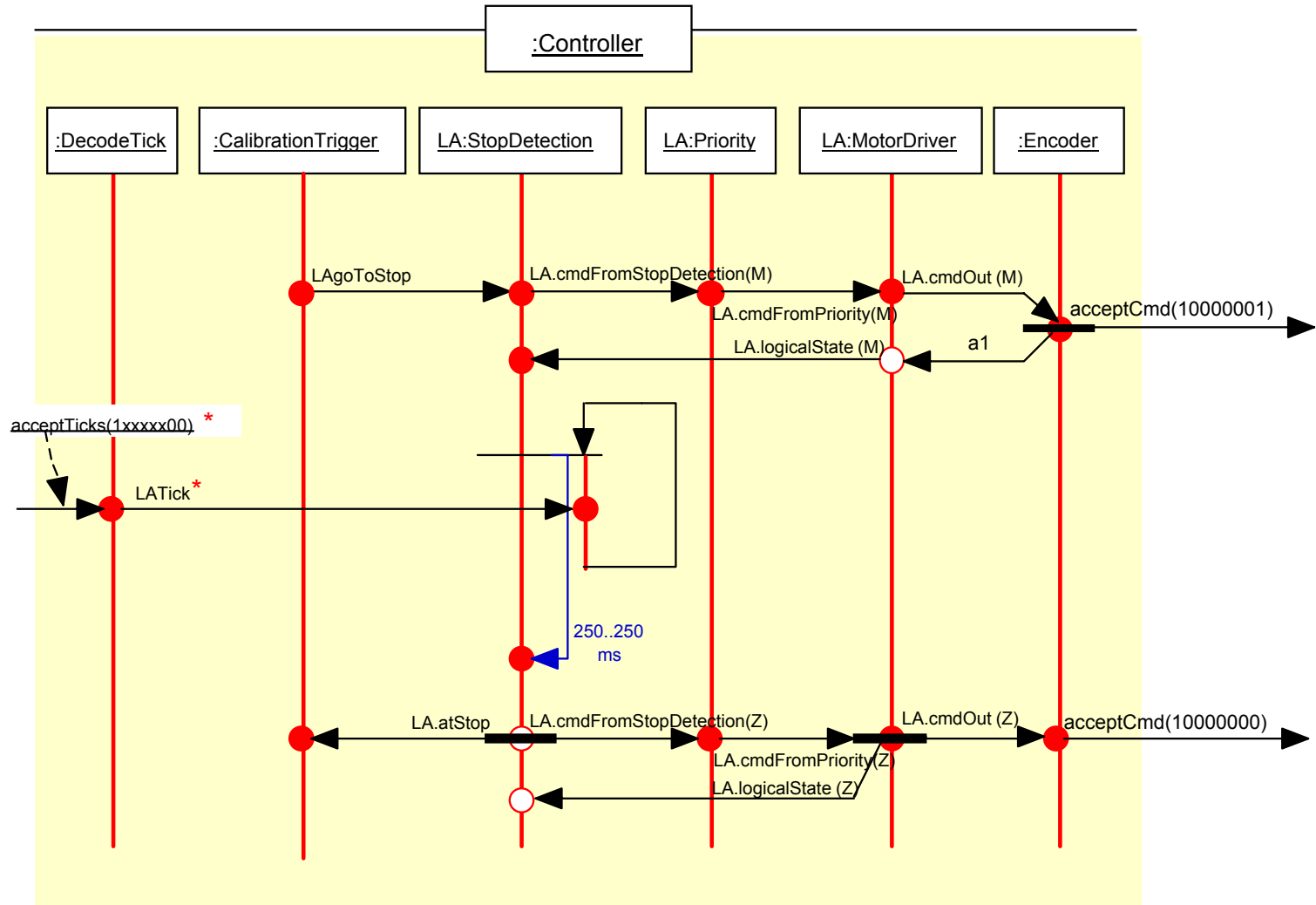
Syntax: temporal constraints



Syntax: Other constructs



Example



Semantics

A synchronous process algebra.

Let \mathcal{I} be the set of signals seen by the application.
Let **Observed** and **Timebase** two sets of signals s.t.

$$\text{Observed} \cup \text{Timebase} \subseteq \mathcal{I}$$

$$\mathcal{O} = \{ \text{Active}, \text{Accepted}, \text{Not_Applicable} \}$$

Let $I_1 \bullet I_2 \bullet \dots \bullet I_n \bullet \dots \in \mathcal{I}^*$ be an input sequence

The behavior of sib p is defined by a sequence of reactions:

$$p = p_0 \xrightarrow[I_1]{O_1} p_1 \xrightarrow[I_2]{O_2} p_2 \cdots p_{n-1} \xrightarrow[I_n]{O_n} p_n \cdots$$

Semantics (Cont'd)

$$p = p_0 \xrightarrow[I_1]{O_1} p_1 \xrightarrow[I_2]{O_2} p_2 \cdots p_{n-1} \xrightarrow[I_n]{O_n} p_n \cdots$$

For some $n \in \mathbb{N} \cup \{\omega\}$

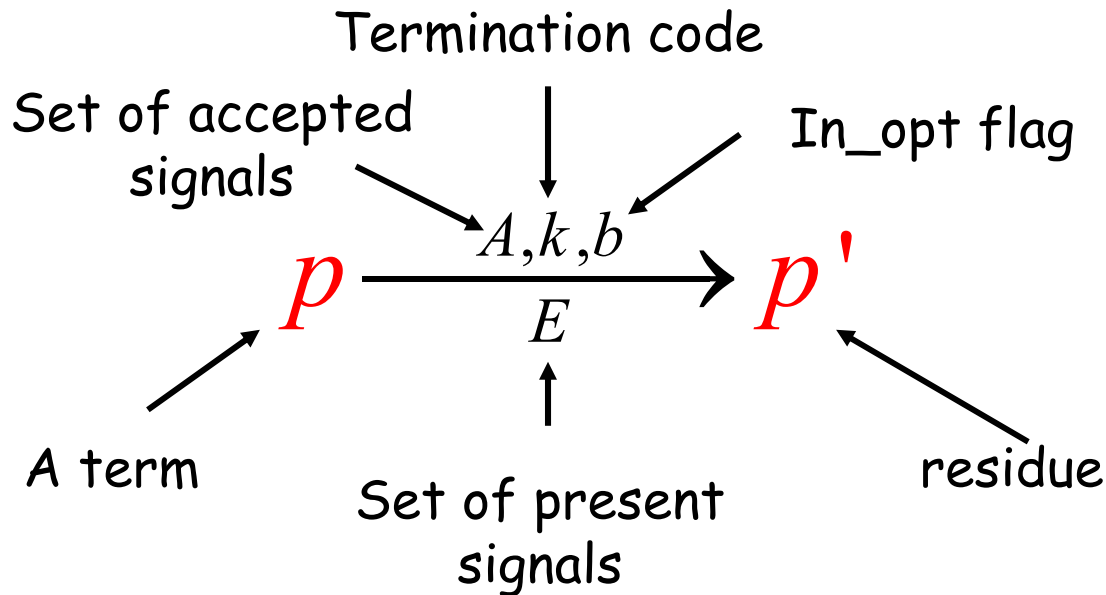
and for $O_j \in \{\{Accepted\}, \{Not_Applicable\}, \{Active\}\}$

If n is finite, then the execution of p terminates at instant n .

A reaction is computed by induction on the structure of the term, using rewriting rules.

$$p \xrightarrow[E]{A, k, b} p' \quad \text{Structural transition}$$

Semantics (Cont'd)



$$I'_j = I_j \cap \text{Observed}$$

if $(I'_j \setminus A \neq \emptyset) \vee (k = \omega)$ then $O_j = \{\text{Not_Applicable}\}$ and p terminates

elseif $(I'_j \setminus A = \emptyset) \vee (k = 0)$ then $O_j = \{\text{Accepted}\}$ and p terminates

else $O_j = \{\text{Active}\}$ and the execution pauses till the next instant.

Tools

- A Tcl-Tk based implementation
 - Graphical editor
 - Compiler to Esterel code
 - Automated instrumentation for
 - Interactive simulation
 - Model checking
 - Circuit generation (blif)
- Inherits from the Esterel S/W environment

Conclusion

- Control-dominated systems
 - Expression of the **expected behavior**,
 - and **validation**
- Sequence Diagrams: no clear semantics,
Synchronous Interface Behavior (SIB)
 - **Formal** foundation
 - **Synchronous** programming platform
- Still to assess
 - Scalability
 - User-friendliness