

Analysis of the phenomenon of several slow start phases in TCP

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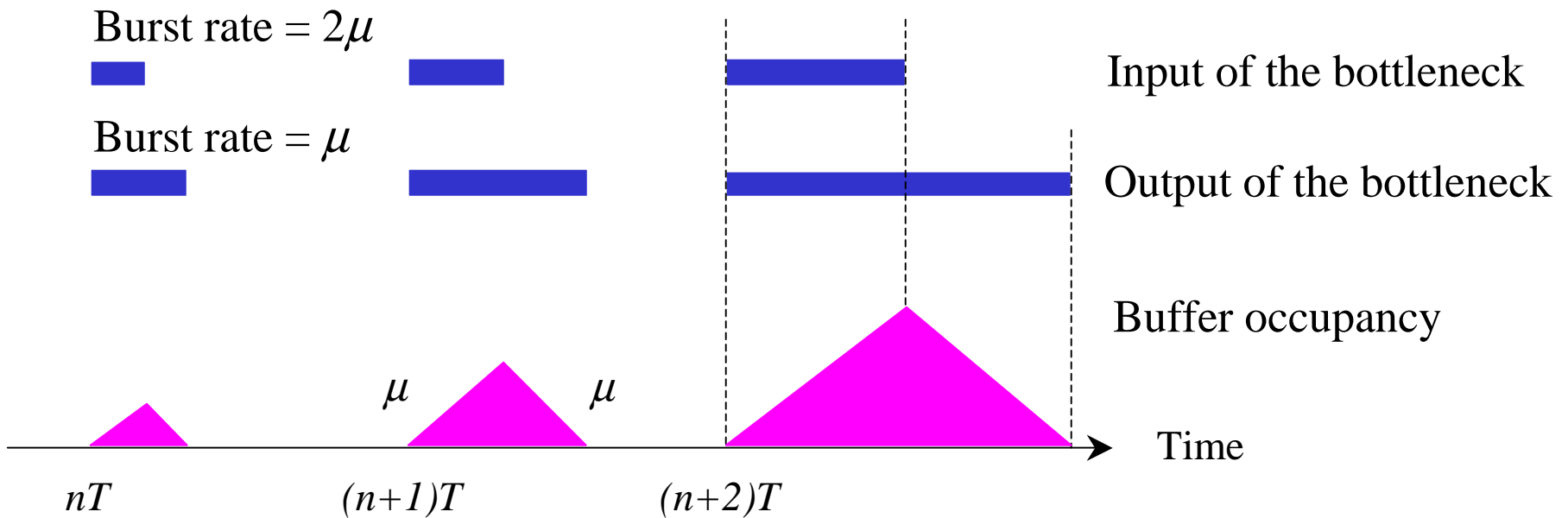
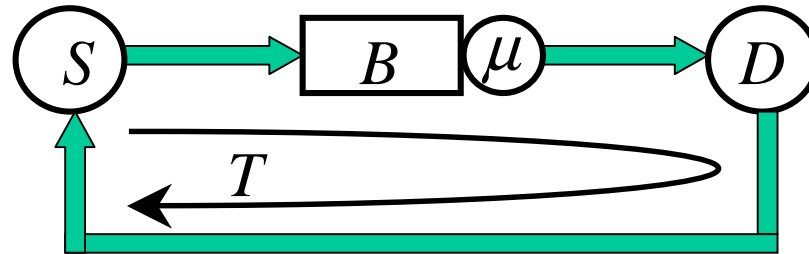
Burstiness of TCP during slow start

- Causes:
 - The fast window increase (one packet per ACK)
 - The ACK-clock based transmission

 The source transmits at twice the available bandwidth (case of no Delay ACK)


- Result:
 - In case of small buffers in routers, an early buffer overflow and an underestimation of the available bandwidth

The model



The overflow window W_B

Due to these bursts, B overflows at a window
[Lakshman et al. 1997, Barakat et al. 1999]

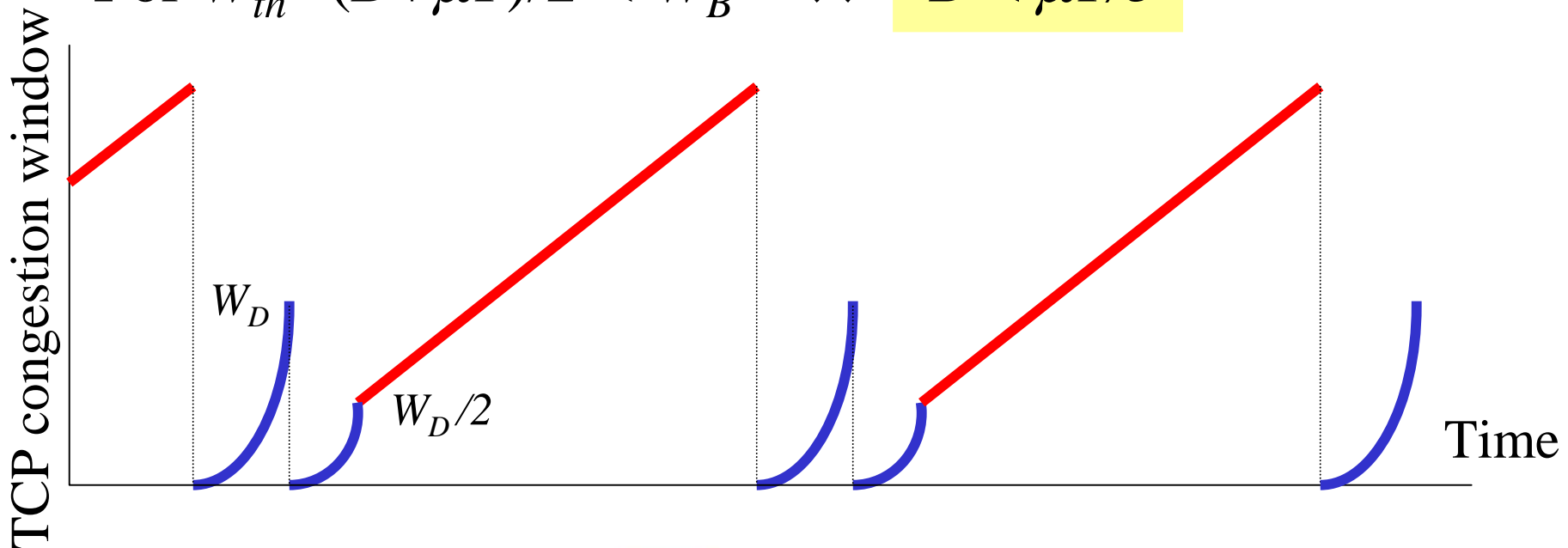

$$W_B = 2^{n_B-1} + B$$
$$2^{n_B-2} < B \leq 2^{n_B-1}$$

And losses are detected at a window


$$W_D = \min(W_{th}, 2W_B)$$

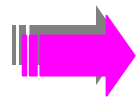
Double slow start phenomenon

- Observed in [Altman et al. 1995, Lakshman et al. 1997] for the Tahoe version of TCP.
- For $W_{th} = (B + \mu T) / 2 < W_B \iff B < \mu T / 3$

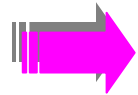


Triple slow start phenomenon

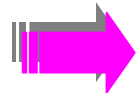
Although $W_D/2 < W_B$, three consecutive slow start phases are seen in some simulations



Cause: Packets stored in the receiver buffer at the beginning of the second slow start phase



Explanation: Different window evolution during the second slow start phase and thus different W_B



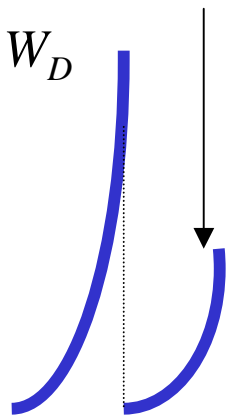
Consequence: Further deterioration in performance

Analysis of triple slow start

Methodology

A detailed analysis of the window evolution and the buffer occupancy during the second slow start phase

Second slow start



We distinguish four consecutive phases

- 1- The window is multiplied by 1.5 every T
- 2- A burst of packets of size $2^{n_B-1} - B + 2$
- 3- The window is multiplied by 1.5 every T
- 4 - Normal slow start: The window is multiplied by two every T

Results

- Losses cannot appear before phase 4
- In phase 4, losses cannot appear if $B > \mu T/7$
- For $B < \mu T/7$, a sufficient condition for losses

$$\left(2^{n_B-1} - B + 2\right) \left[2(1.5)^n - 1\right] > 2B - 1$$

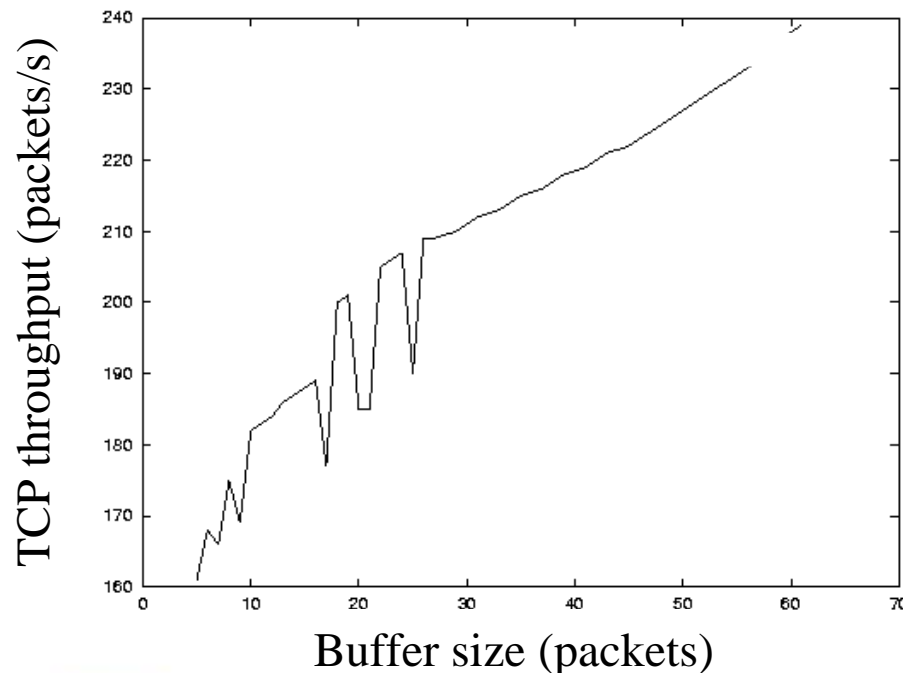
where n is the smallest integer number satisfying

$$\left(2^{n_B-1} - B + 2\right) \left[(1.5)^{n+1} - 1\right] > B + 1$$

Periodicity of the phenomenon

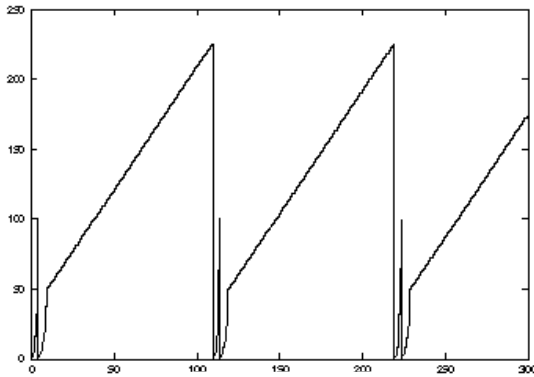
While increasing B from 0 to $\mu T/7$, the losses during the second slow start phase appear and disappear in a periodic manner causing an oscillation in the throughput

ns simulator
 $\mu=1.5$ Mbps
 $T=560$ ms
MSS = 512 bytes

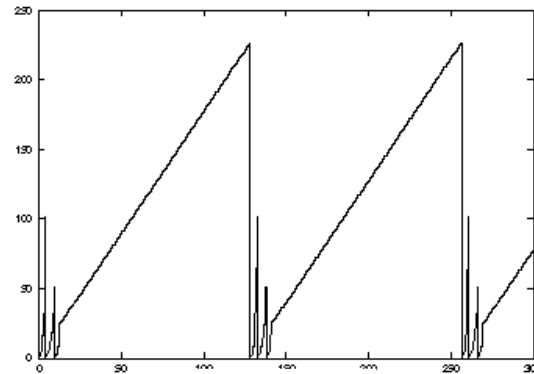


Results

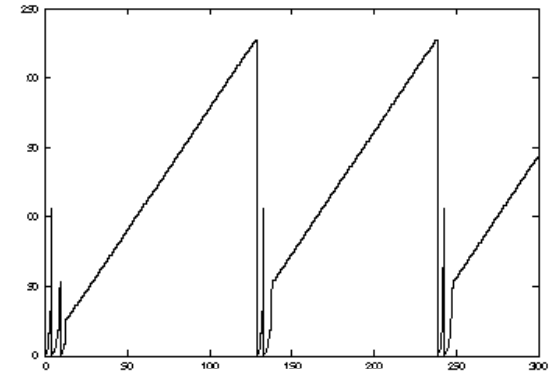
TCP congestion window vs. Time



$B=19$ packets



$B=20$ packets



$B=22$ packets

Conclusions:

- The problem of TCP slow start burstiness (C. Barakat and E. Altman, “Performance of short TCP transfers”, Networking 2000, May 2000)
- Slow start has different behavior when some packets are stored in the receiver buffer