



Content Dissemination in Mobile Wireless Networks

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Context / Objectives

- ❑ **An activity running for couple of years, involving several people**
 - Amir Krifa, Francisco Santos, Benjamin Ertl, Thierry Turletti
 - Thrasyvoulos Spyropoulos (Eurecom)

- ❑ **To exchange data between mobile devices**
 - Without infrastructure – disaster, censorship, etc
 - Multi-hop mode: I give you, you give the others, and so on
 - Often called Delay Tolerant Networks, Pocket switched Networks, etc

- ❑ **Communication can be of different types**
 - Point-to-point: I send a content/message to someone
 - Broadcast
 - Publish-subscribe: One publishes, others search and request

Difficulties

- ❑ **Current Internet architecture does not work in this context**
 - No end-to-end path, no TCP/IP
 - Network topology changes frequently
 - Opportunistic encounters

- ❑ **Forwarding, replication and content-aware routing**
 - Either give the content to encountered devices and get rid of it
 - Or simply give them a copy
 - How to know?
 - Several solutions in the literature: spray-and-wait, age-based routing, community-based routing, routing by social networks, etc

- ❑ **Routing reduces the load, but does not provide explicit solutions in case of resource starvation**
 - TCP and buffer management vs. IP routing

Our framework

❑ Transform the problem into a resource management problem

- Set a global objective for the network
 - Maximize throughput, Minimize delay, being fair, etc
- Devices take local decisions that push the network to its desirable global behavior (and keep it there)
- Two main decisions:

When there is a need to drop some data, drop the least useful first
When there is a need to forward, forward the most useful first

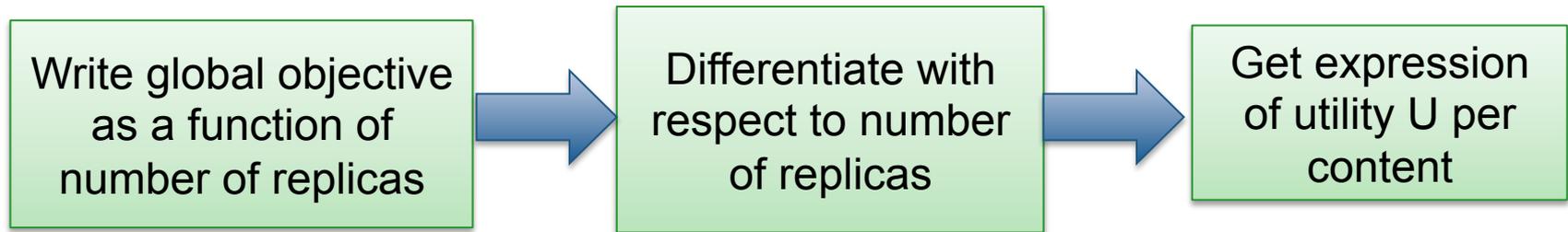
❑ Control variable: The number of replicas ...

❑ Control function: A utility per content

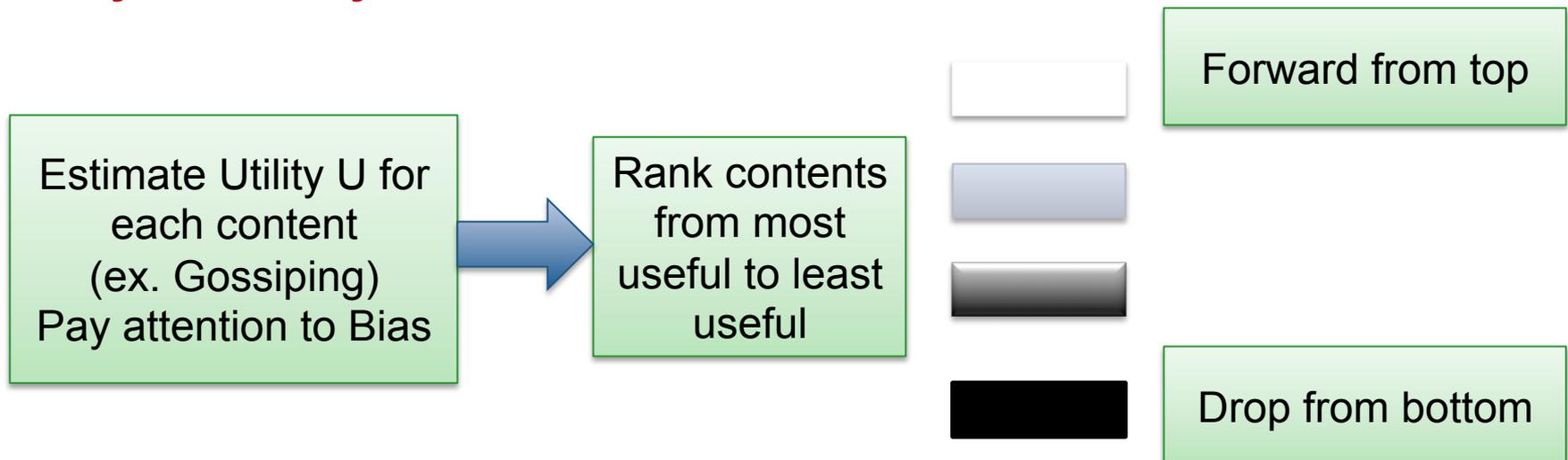
- The marginal gain/loss upon drop/replication

Algorithm in a nutshell

Beforehand



Locally, on the fly



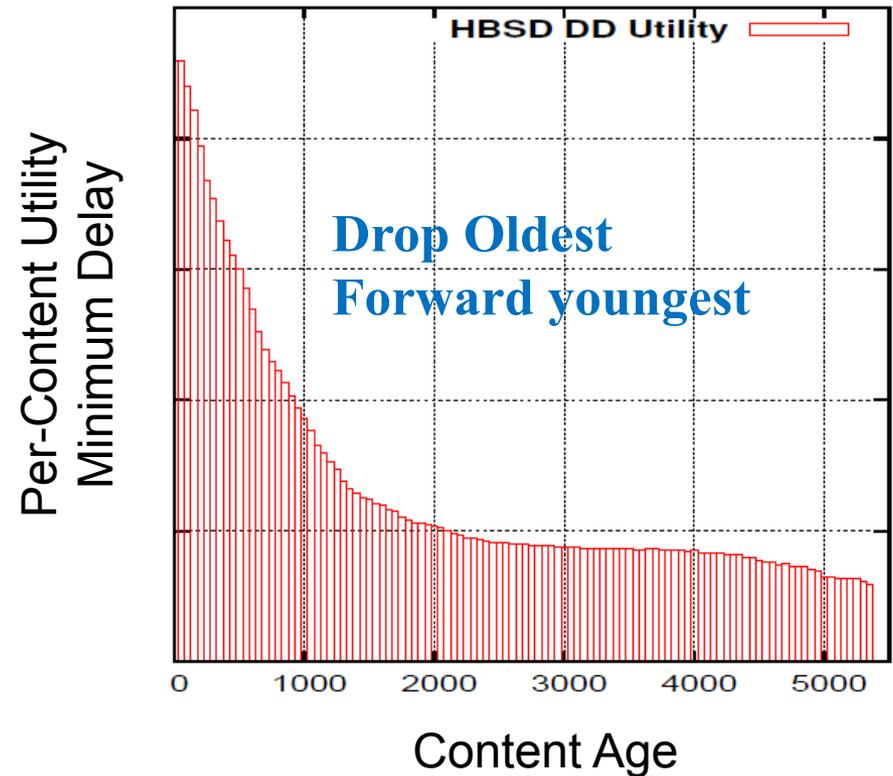
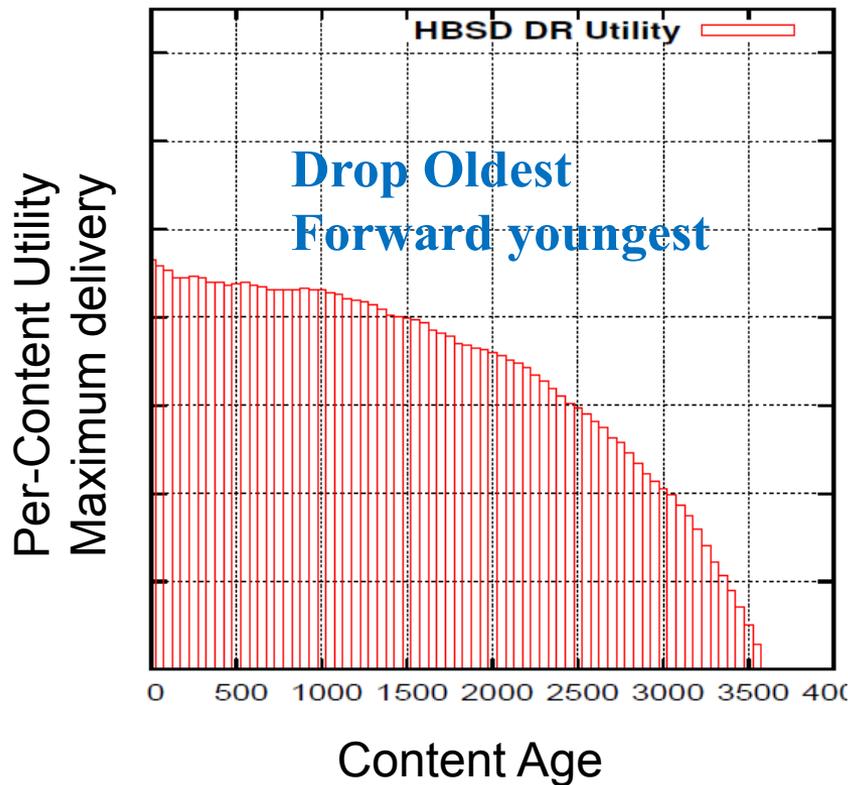
HBSD: The point-to-point case

- ❑ **History-Based Scheduling and Drop**
 - http://planete.inria.fr/HBSD_DTN2/
- ❑ **Contents/Messages appear at a device, try to find their way to some other destination device, then disappear**
- ❑ **Utility: A function of content age**
 - By gossiping, HBSD finds this function by itself
 - Same function for all contents
 - Age of content indicates its utility
- ❑ **Two variants: **Maximum delivery** and **Minimum delay****

Further details: Amir Krifa, Chadi Barakat, Thrasyvoulos Spyropoulos, “[Message Drop and Scheduling in DTNs: Theory and Practice](#)”, in IEEE Transactions on Mobile Computing, vol. 11, no. 9, September 2012.

Samples of utility functions

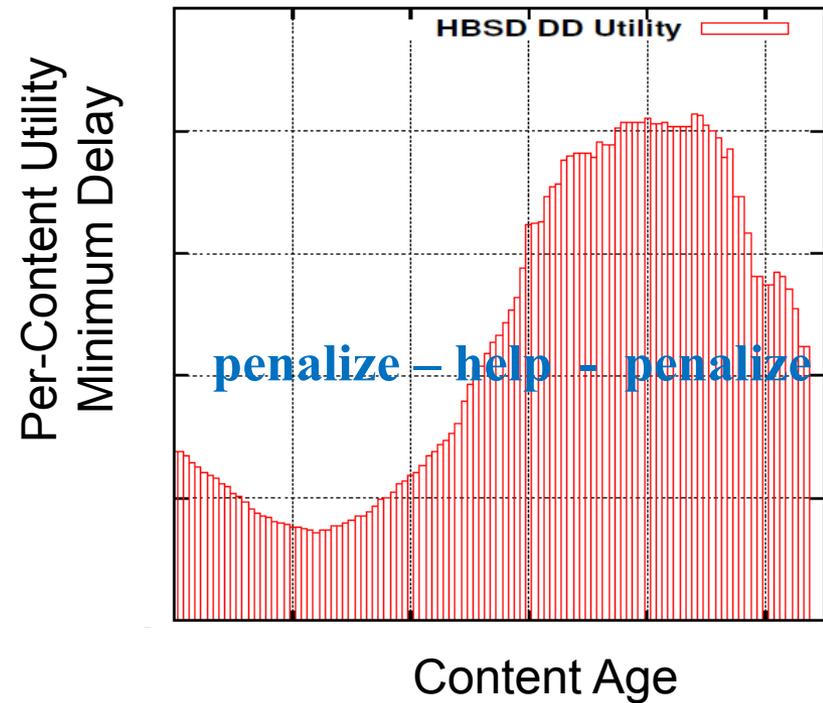
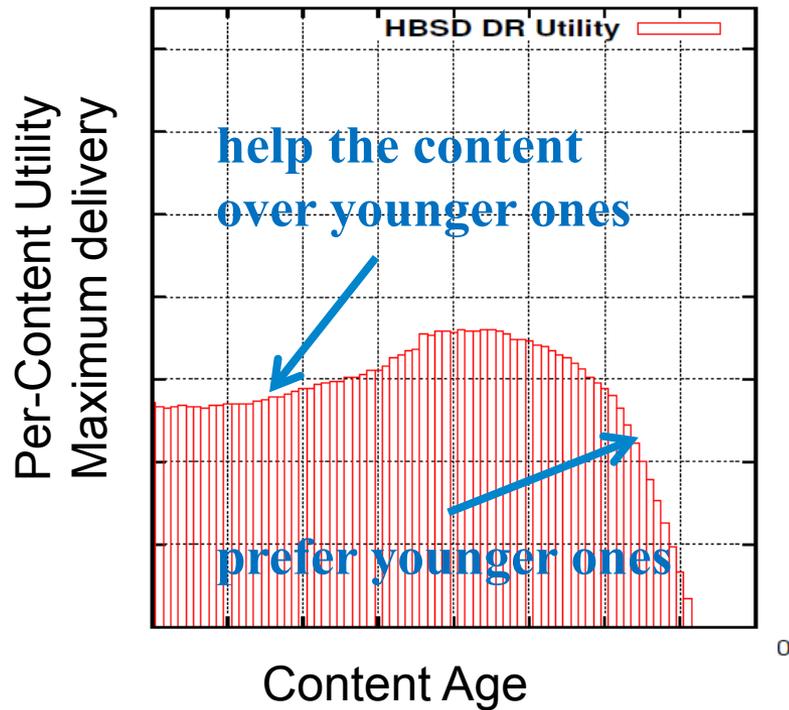
- For a lightly loaded network, utility decreases with age
 - Best is to schedule youngest first and to drop oldest first



Samples of utility functions

□ When the load increases, the shape of utilities changes

- Simple policies are no longer optimal
- HBSD can capture the optimal behavior whatever is the load



CODA: The publish-subscribe case (point-to-multipoint)

❑ Content Optimal Delivery Algorithm

- <http://hal.inria.fr/hal-00742734/>
- <http://planete.inria.fr/Software/CODA/>

❑ Developed within the CCN/NDN framework

- Contents have universal names. Ex: \inria\team\diana\codas
- Have different popularities (different request rates)
- Users issue requests for contents, contents flow back, intermediate devices cache

❑ Again, control the number of replicas per content

- Function of its popularity
- Collaborative network-wide caching, taking mobility into account (vs. LRU, LFU)

❑ First solution that maximizes the number of satisfied requests (throughput)

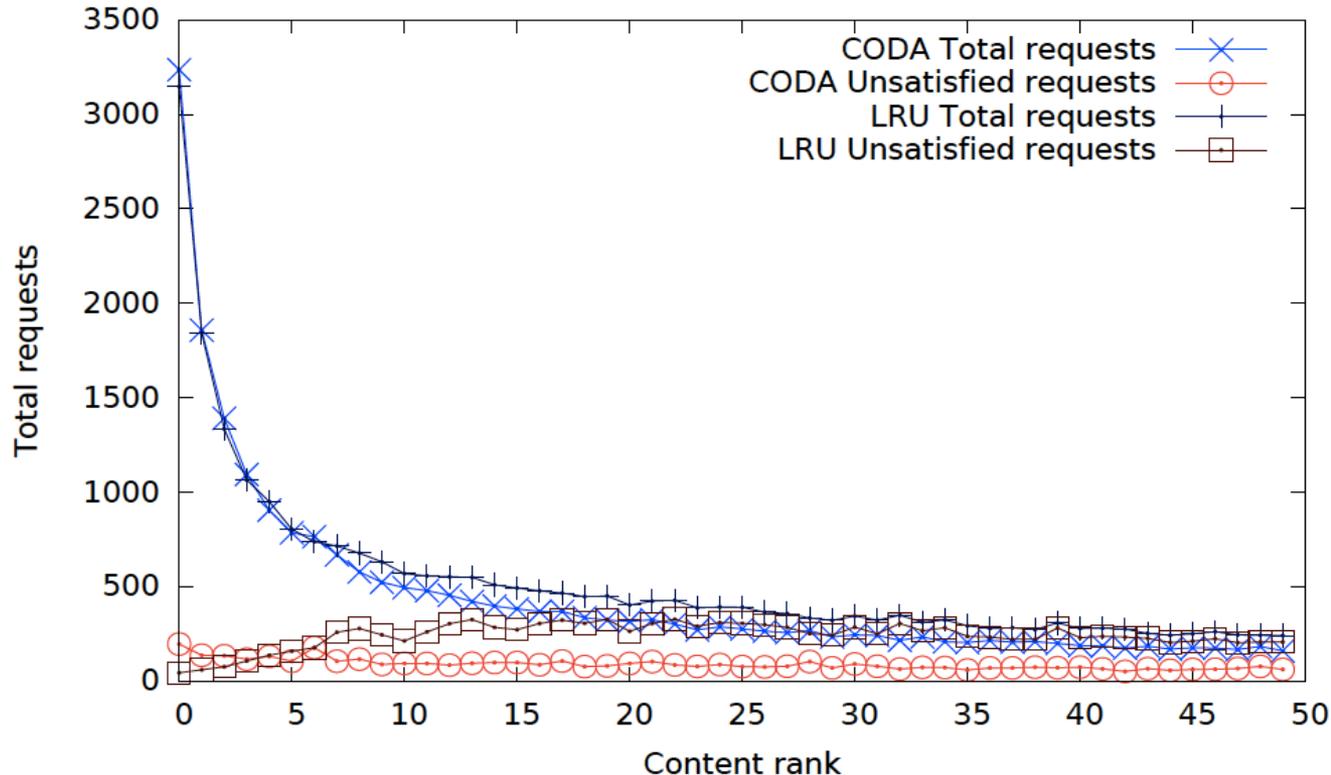
- Under some assumptions:

$$\text{Utility} \approx \text{miss rate} = \text{popularity} - \text{delivery}$$

$$\text{Maximize Throughput} = \text{Equalize miss rate}$$

CODA: The publish-subscribe case (point-to-multipoint)

- ❑ Observed 50% more throughput than LRU
- ❑ And better protection of non popular contents
- ❑ Global network performance can be easily tuned



Open issues

Future activity will build upon CODA and transform it into a general solution for storage management and scheduling in mobile content-centric networks

- Scalability of signaling
- Fairness vs. Throughput
- Convergence vs. Reactivity
- Collaboration enforcement
- Heterogeneity of devices (battery, storage, etc)
- Large scale experiments

merci

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