Bandwidth Allocation Policies for Unicast and Multicast Flows

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Overview

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- Model
- Analytical results
- Simulations
- Practical aspects
- Conclusion



Introduction

Multicast is a very attractive data delivery.



- Multicast is implemented in new routers.
- Multicast is still not deployed due to:
 - lack of congestion control.
 - no incentive to use multicast.
- Our new bandwidth allocation policy gives an incentive to use multicast.

Model: Allocation policies

- Receiver Independent (RI):
 - Does not make any changes in the current bandwidth allocation policies. The benchmark.
- Linear Receiver Dependent (LinRD):
 - Gives to multicast the bandwidth used by the equivalent of unicast connection.
- Logarithmic Receiver Dependent (LogRD):
 - Rewards multicast with its bandwidth saving on the network.



Model: Criteria

- Receiver satisfaction:
 - Mean bandwidth.
- Fairness:
 - Standard deviation of the bandwidth seen by the receivers.
- Optimality:
 - receiver satisfaction and fairness are inconsistent.
 - We evaluate the trade-off between receiver satisfaction and fairness.



Model: assumptions

- Knowledge in every network node of:
 - every flow on an outgoing link.
 - the number of receivers per flow reached via an outgoing link.
- All flows are CBR.
- No arriving or departing flows.
- Each node makes the bandwidth allocation independently.
- A receiver sees the minimum allocated bandwidth along its path.
- The sources can send via cumulative layered transmission.



Analytical study

Star topology



Analytical study

Mean bandwidth



LinRD considerably increases receiver satisfaction.

Standard deviation



But LinRD is not fair.

Simulations

- Study on a hierarchical topology:
 - ◆ 1 WAN, 20 MANs, and 180 LANs.
- Create a unicast environment, which aim is to study the deployment of multicast, with 2000 unicast flows.
- Two scenarios:
 - One increasing multicast group.
 - An increasing number of multicast groups.



Mean bandwidth



The two RD policies highly improve receiver satisfaction.



Mean bandwidth for MC receivers



 LogRD increases receiver satisfaction. LogRD does not starve unicast flows.

Mean bandwidth for UC receivers

Standard deviation



LogRD is more fair than LinRD.

Minimum bandwidth



LinRD starves the worst case receiver.

- RI policy:
 - Receiver satisfaction and fairness are not influenced by an increase in the multicast group size.
- The receiver dependent policies:
 - Significantly increase receiver satisfaction.
- LinRD policy:
 - Leads to high unfairness.
 - Starves unicast flows.
- LogRD policy performs best:
 - Keeps fairness close to the one of RI.
 - Does not starve unicast flows.

Simulations: varying the # of MC groups

Mean bandwidth for MC receivers

Mean bandwidth



satisfaction.

Simulations: varying the # of MC groups

Standard deviation



 The LinRD policy starves the worst case receiver.

Minimum bandwidth

Simulations: varying the # of MC groups

- LogRD achieves the best trade-off between receiver satisfaction and fairness:
 - LogRD highly improves the mean bandwidth for the multicast receivers.
 - LogRD does not significantly decrease the minimum bandwidth compared to the RI policy.
- Mean bandwidth, for all the receivers, slightly better for the LogRD policy than for the others policies.
- Same standard deviation for the three policies for all the receivers.



Practical aspects

- Knowing the number of receivers downstream:
 - business model for charging.
 - Useful for feedback implosion avoidance.
- Introducing the LogRD inside routers:
 - WFQ to realise the bandwidth allocation.
- Introducing LogRD in a real network is practically feasible.



Conclusions

- Our LogRD bandwidth allocation policy performs best:
 - highly improves the receiver satisfaction for the MC receivers.
 - does not significantly affect the fairness.
- LogRD gives an incentive to use multicast.

Thanks



The hierarchical topology



