Formalization of Rings with Explicit Divisibility in Type Theory

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Introduction

- SSReflect
- Rings with explicit divisibility
 - ► GCD domains
 - Bézout domains
 - Euclidean rings
- Smith normal form
 - Constructive PIDs

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SSReflect

- Extension to Coq
- George Gonthier: formalization of the four color theorem
- Cayley-Hamilton theorem, decidability of ACF...
- Feit-Thompson theorem (part of the classification of finite simple groups)

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SSReflect

- Small Scale Reflection
- New tactics and tacticals
- Library of mathematical theories

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DvdRing

A ring R has explicit divisibility if it has a divisibility test that give witnesses:

$$a \mid b \leftrightarrow \exists x. b = xa$$

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• \mathbb{Z} and k[x] where k is a field (e.g. \mathbb{Q})

DvdRing

Demo!

GCD domains

 GCD domain: Every pair of elements have a greatest common divisor

$$\forall a \, b. \exists g. \, g \mid a \land g \mid b \land \forall g'. \, g' \mid a \land g' \mid b \to g' \mid g$$

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Properties of GCD domains

Definition

The gcd of the coefficients of $p \in R[x]$ is called the *content* of p, written cont(p)

Definition

 $p \in R[x]$ is primitive if cont(p) = 1

Theorem

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Gauss lemma: cont(pq) = cont(p)cont(q)
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Theorem

Every polynomial $p \in R[x]$ can be written as p = cont(p)q with q primitive

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Properties of GCD domains

► Using this one can give an algorithm for computing the gcd of p, q ∈ R[x]

- Give a proof that if R is a GCD domain then R[x] also is
- Don't use field of fractions!
- ▶ Can compute gcd in $\mathbb{Z}[x_1, \ldots, x_n]$ and $k[x_1, \ldots, x_n]$

Bézout domains

- Non-Noetherian analogue of principal ideal domains
- PID: Every ideal is principal
 - Quantification over all ideals
- Bézout domain: Every finitely generated ideal is principal
- Equivalent definition:

$$\forall a b. \exists x y. ax + by = gcd(a, b)$$

Euclidean rings

- Euclidean norm, $f : R \to \mathbb{N}$
- ► Euclidean division: ∀ab.∃qr.a = bq + r and either f(r) < f(b) or r = 0</p>
- Examples: \mathbb{Z} with absolute value and k[x] with degree

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Bézout domains and Euclidean rings

Theorem Every Euclidean ring is a Bézout domain Theorem Every Bézout domain is a GCD domain Theorem \mathbb{Z} is a Euclidean ring Theorem k[x] is a Euclidean ring

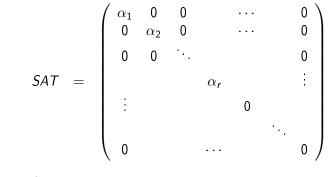
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Smith normal form

- Generalization of Gauss elimination algorithm
- Elements from a PID and not just a field
- Compute homology groups of simplicial complexes
 - "Homology is a rigorous mathematical method for detecting and categorizing holes in a shape." - Wikipedia

Smith normal form

► Let A be a nonzero m × n matrix over a PID. There exists invertible m × m and n × n matrices S,T such that



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and $lpha_i \mid lpha_{i+1}$, $1 \leq i < r$

Smith normal form

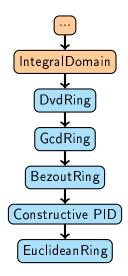
- Open question for Bézout domains
- Need constructive approximation of PIDs

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Constructive PIDs

- ▶ Mines, Richman, Ruitenburg: Bézout domains such that if we have a sequence u(n) with u(n + 1) | u(n) then there exists k such that u(k) | u(k + 1)
- In type theory this can be represented as that strict divisibility is well founded

Summary



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Conclusions and further work

Have: Divisibility theory in SSReflect

• Formalized algorithms for \mathbb{Z} and k[x]

- ► Todo: *R* GCD domain implies *R*[*x*] GCD domain
- Have: Smith normal form algorithm in Haskell
- Todo: Formalize Smith normal form algorithm

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Questions?

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