

# WP4: Algorithmic Solutions and Technical Recommendations for Optical Networks I



**CRESCCO IST-2001-33135**

Critical Resource Sharing for Cooperation in Complex System



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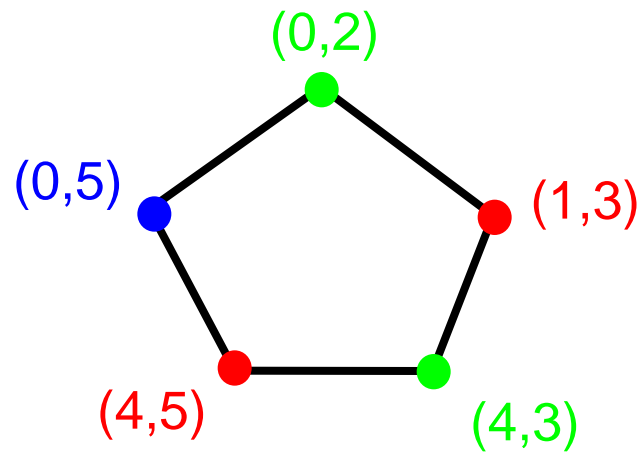
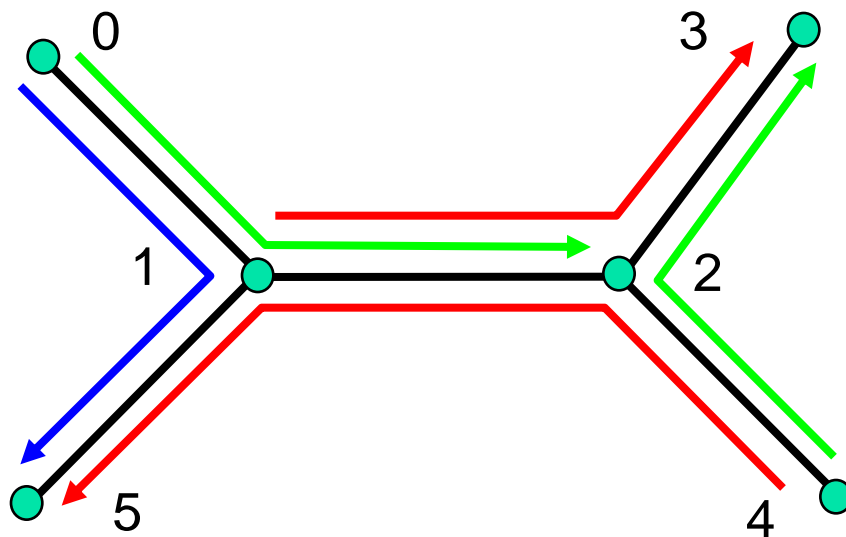
# All-Optical Network: Wavelength Assignment Problem

- Communications are routed along color-disjoint *lightpaths*
- Let  $\pi(G, I)$  be the minimum load for routing a given set of communication  $I$
- Let  $w(G, I)$  be the minimum number of wavelengths for a solution

**what is the relation between**

**$\pi(G, I)$  and  $w(G, I)$  ?**

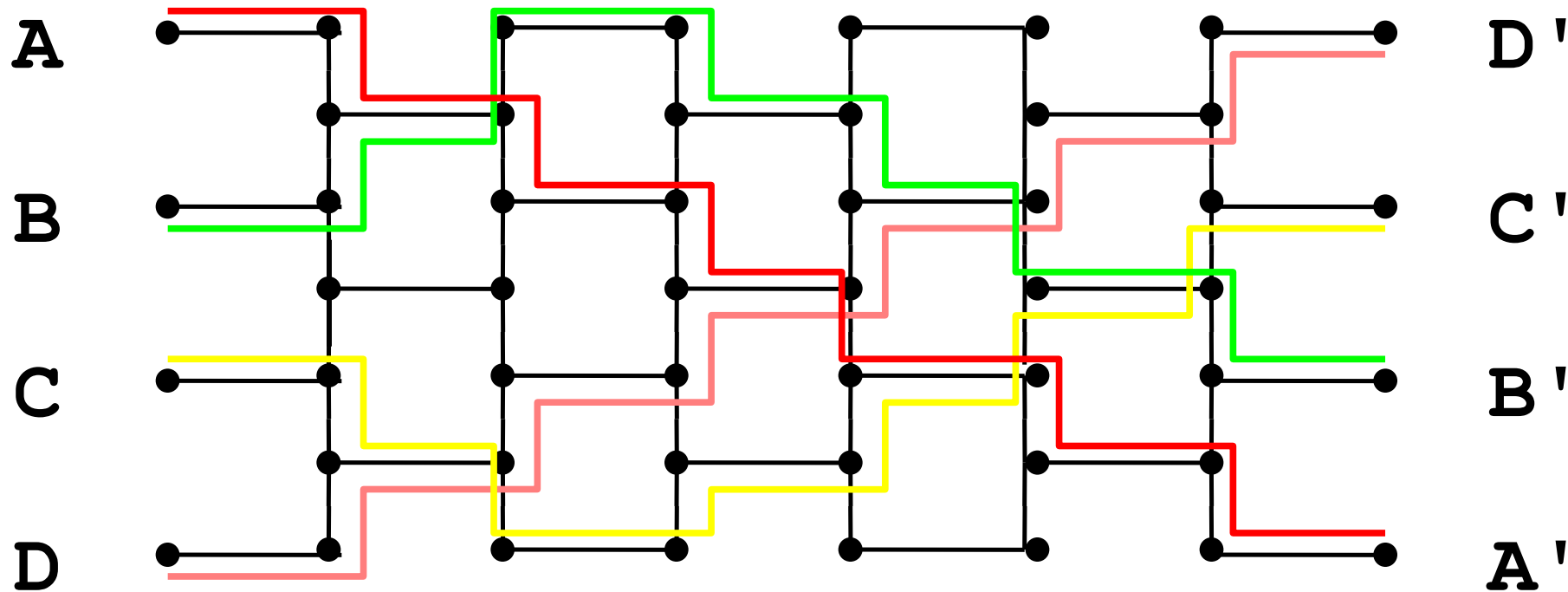
# Conflict Graph: coloring problem



Example where  $\vec{w}$  is not equal to  $\vec{\pi}$



$\pi = 2$  and  $w = |I|$





# New results for WAP

## Design and routing algorithms for multifiber networks

- conflict hypergraphs model
- coloring of hypergraphs (NP-complete)
  - randomized approximation algorithm
  - integer linear programming formulation
- combine routing and wavelength assignment
  - randomized rounding of the fractional multicommodity flow solution

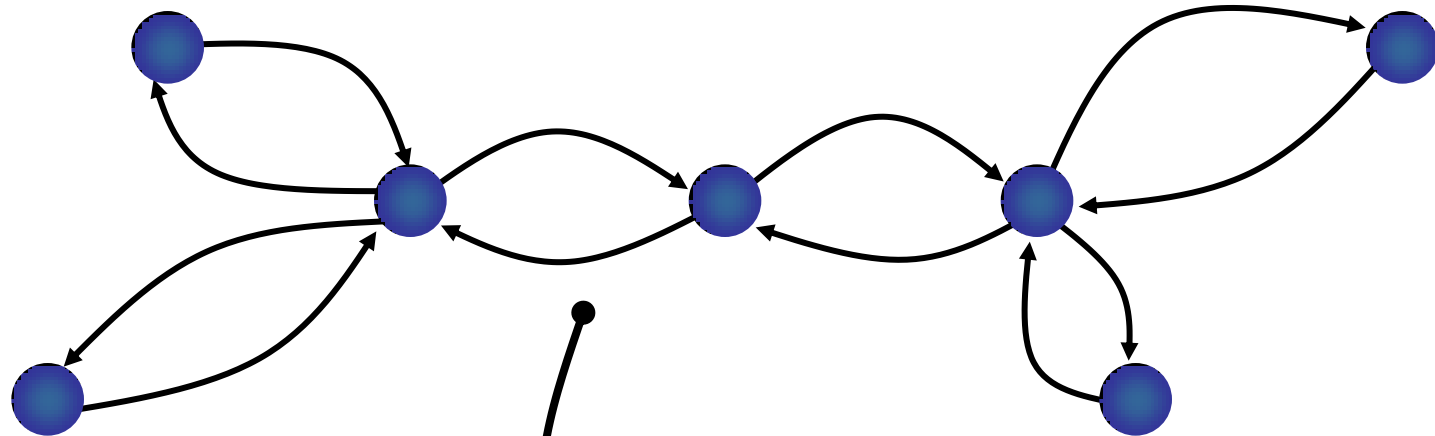


# New results for WAP

## Tree and ring topologies

- class of greedy algorithms using randomization
  - lower bounds and a  $7w/5 + o(w)$  algorithm on binary trees of depth  $o(L^{1/3})$  with high probability
- randomized WAP
  - randomized rounding on fractional path coloring
- polynomial time algorithms that compute almost optimal fractional path colorings in bounded degree trees and rings (including multifiber rings)

# Traffic Grooming



Network  $G=(V,E)$

$e \in E$



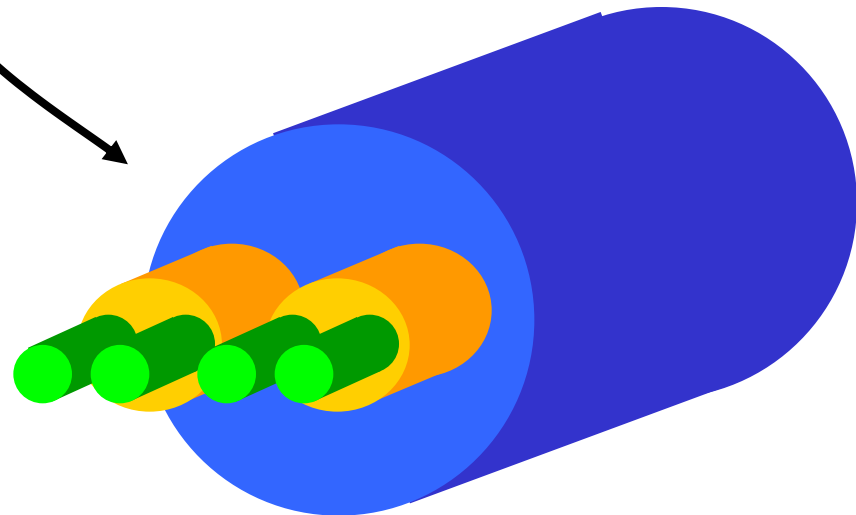
Fibre



Band

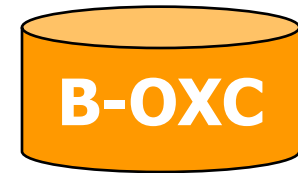
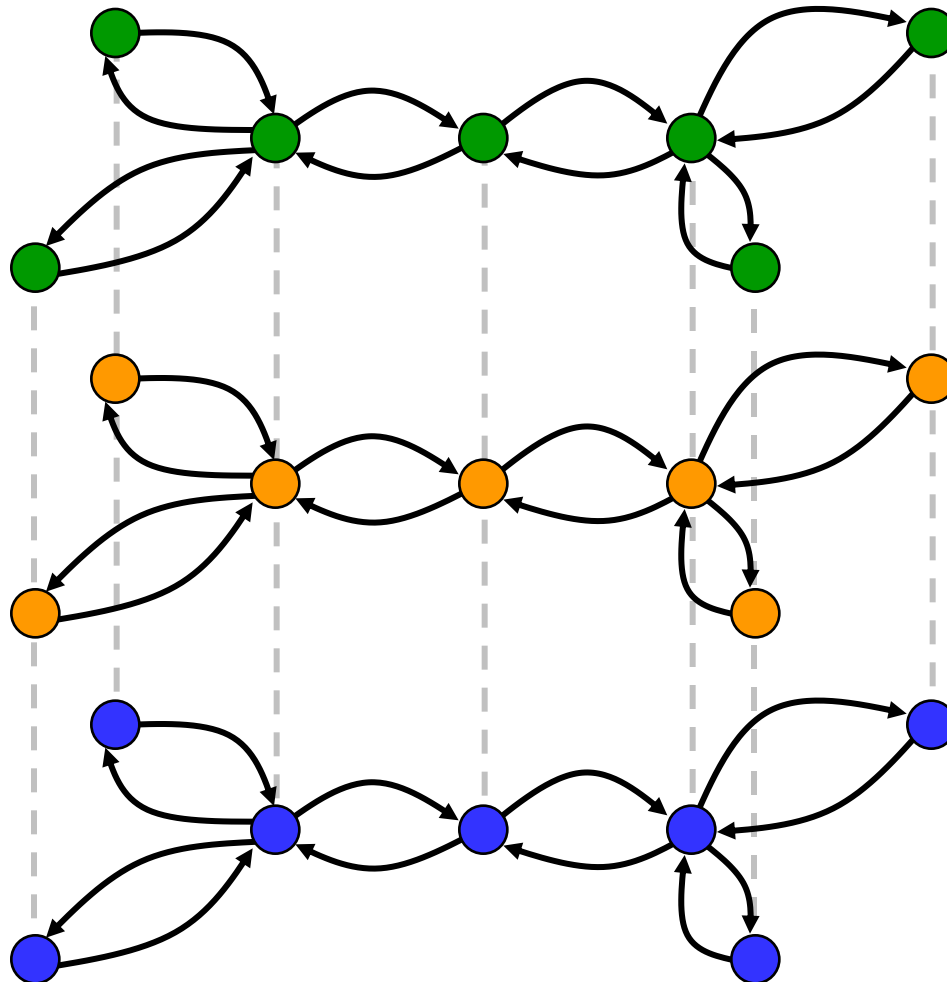


Wavelength



# Multi-Layer model

G: physical network



D: lightpaths



# Multi-Layer model

Cost of nodes:  
max degree  
of OXCs

Closed fiber

Opened fiber

G: physical network

F-OXC

B-OXC

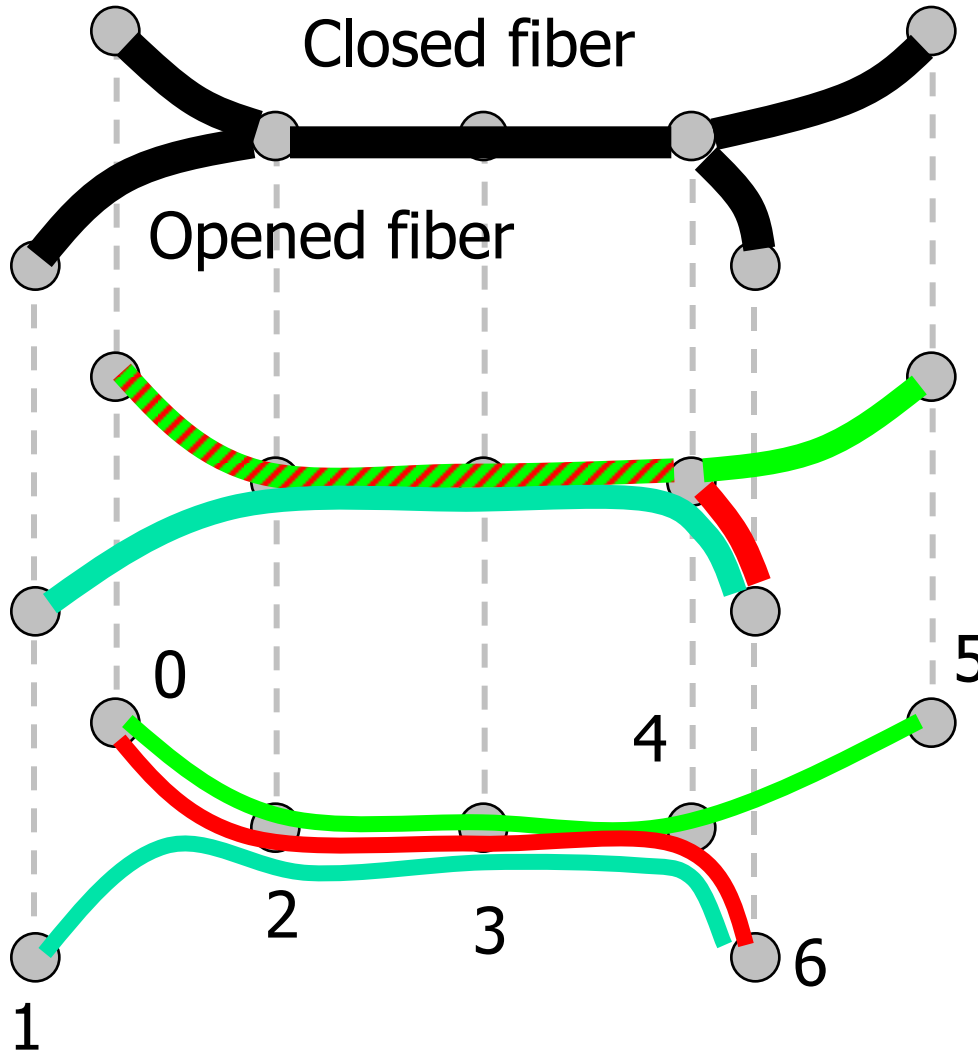
W-OXC

Capacities:

$W=1$

$B=2$

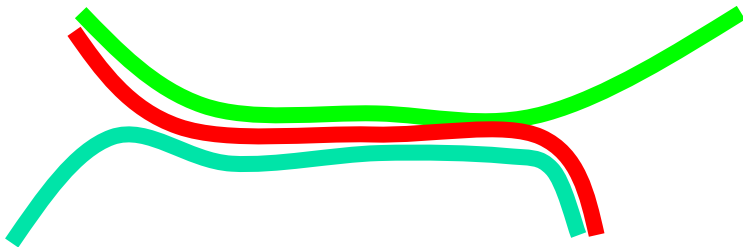
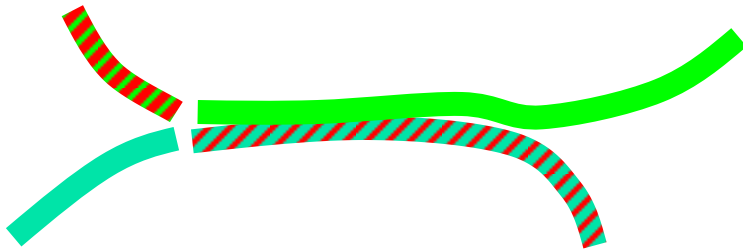
$F=2$



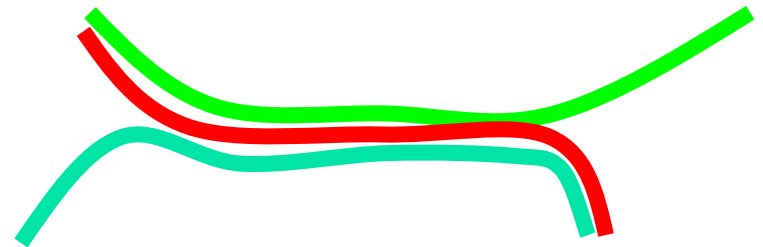
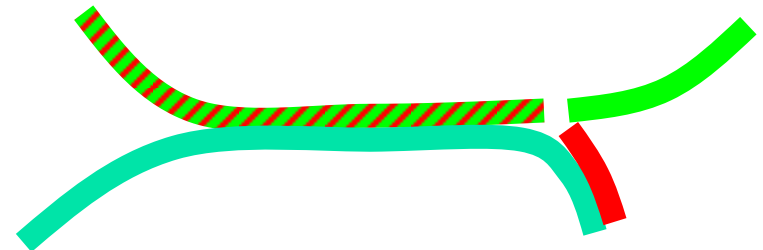
D: lightpaths

# Grooming example

Grooming (b)



Grooming (a)





# Results

- New model of pipes for multi-layer switches
  - ILP formulation
  - Efficient heuristics with a cost function on pipes
  - Experimental results
- Case for ATA traffic in SDH/SONET rings
  - Some optimal solutions by using results of design theory
    - Example: for  $C=3$  the problem consists in covering the ring with triangles



# Improved resources allocation algorithms

## ■ Packing problem

- approximation algorithm
- application to edge congestion problem in multicasting

## ■ Station Placement Problem

- dynamic programming algorithm (trees) and approximation algorithm : multicast, wavelength converters?

## ■ List coloring problem

- bounds for bi-partite graph



# Conclusion

- Identification of important techniques for allocation problem
  - Graph theory
  - LP-relaxation + randomization
  - Approximation algorithms
- Publications :
  - Theoretical computer science and Networks conferences
- Emergence of a new model: virtual topologies



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# Future work

- Survivability
  - Protection and restoration
- Dynamic traffic
- Experimentation / validation of algorithms

Changing order: demand  
order.  
Changing objective fun  
Routing 110 requests.  
Solving Cplex linear



Step: + 0 -

N\_0=>N\_5 :11

Set Path

Main

Protection



>N\_0 :11

Set Path

Main

Protection



>N\_6 :5

Set Path

Main

Protection



>N\_0 :5

Set Path

Main

Protection

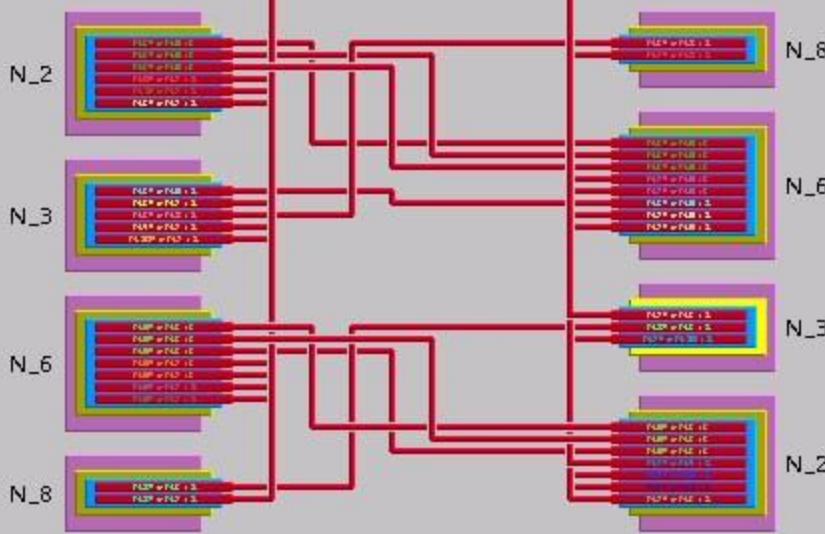


NodeN 7

Show Free    Oxc details    Node cost

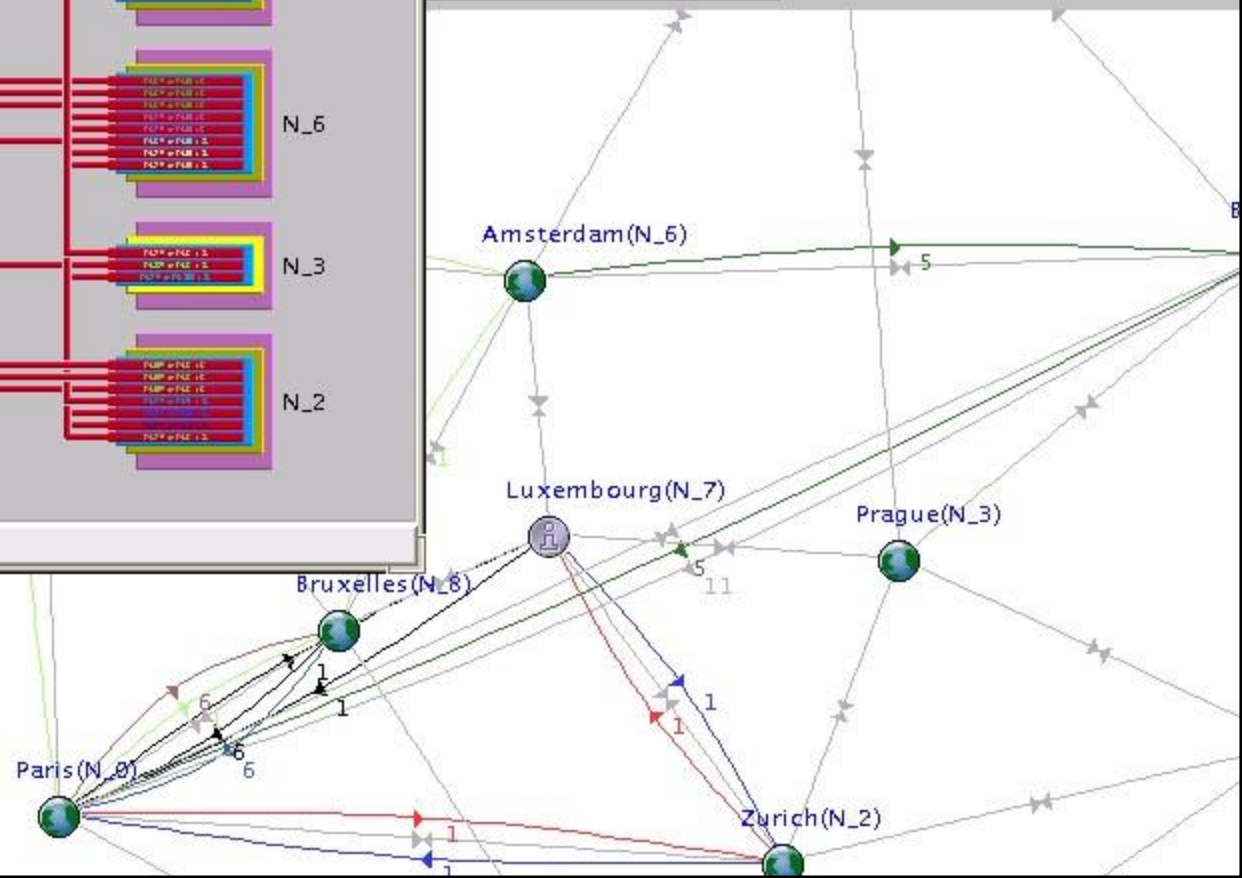
Node Luxembourg (N\_7)

In    Drop    Add    Out



FIBER\_7\_3\_0

COST (View #1)



Routing

Max Cplex

Optimality

Rate rc

Order

Object

omping by size



closing factor

0