

# Interval Analysis in Coq

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Tool used to handle inaccuracies in computations.

$$-\pi * \sqrt{2} \approx -3.14 * 1.41 = -4.4274$$

$$[-3.15, -3.14] * [1.41, 1.42] = [-4.473, -4.4274]$$

If we know the bounds on the input data we can compute the bounds on the result.

# Interval arithmetic, more formally

## Definition

**interval** := closed, bounded, connected, nonempty subset of  $\mathbb{R}$

$$x := [\underline{x}, \bar{x}] = \{\tilde{x} \in \mathbb{R} \mid \underline{x} \leq \tilde{x} \leq \bar{x}\}, \quad \text{where } \underline{x}, \bar{x} \in \mathbb{R}, \underline{x} \leq \bar{x}$$

Notation  $\mathbb{IR}$  – set of intervals

## Classification

- thin interval  $\underline{x} = \bar{x}$
- thick interval  $\underline{x} < \bar{x}$

Associated quantities

$$\text{midpoint} \quad x_c := \frac{\underline{x} + \bar{x}}{2} \qquad \text{radius} \quad \Delta_x := \frac{\bar{x} - \underline{x}}{2}$$

$$x = [x_c - \Delta_x, x_c + \Delta_x]$$

# Basic interval operations

$$x + z := \square\{\tilde{x} + \tilde{z} \mid \tilde{x} \in x, \tilde{z} \in z\}$$

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$$-x := \square\{-\tilde{x} \mid \tilde{x} \in x\} = \{-\tilde{x} \mid \tilde{x} \in x\} = [-\bar{x}, -\underline{x}]$$

$$\begin{aligned}xz &:= \square\{\tilde{x}\tilde{z} \mid \tilde{x} \in x, \tilde{z} \in z\} = \{\tilde{x}\tilde{z} \mid \tilde{x} \in x, \tilde{z} \in z\} = \\ &= [\min(\underline{x}\underline{z}, \underline{x}\bar{z}, \bar{x}\underline{z}, \bar{x}\bar{z}), \max(\underline{x}\underline{z}, \underline{x}\bar{z}, \bar{x}\underline{z}, \bar{x}\bar{z})]\end{aligned}$$

Principle:

Correctness is more important than accuracy.

$$\pi - \pi = 0$$

$$[3.14, 3.15] - [3.14, 3.15] = [-0.01, 0.01]$$

Techniques to increase accuracy (avoid decorrelation)

- e.g., bisection

# Rounded interval arithmetic

## Usage

- in theory:  $[\underline{x}, \bar{x}]$  with  $\underline{x}, \bar{x} \in \mathbb{R}$
- in practice:  $[\underline{x}, \bar{x}]$  with  $\underline{x}, \bar{x} \in M$ ,  
where  $M$  is a machine representable subset of  $\mathbb{R}$

## Outward rounding

$$\diamond x := [\nabla \underline{x}, \Delta \bar{x}]$$

$$x \subseteq \diamond x$$

$$x + \diamond z = \diamond(x + z)$$

# Rounded interval arithmetic

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## Example

$$[-3.15, -3.14] * [1.41, 1.42] = [-4.473, -4.4274]$$

$M$  : decimal numbers with 2 digits

$$[-3.15, -3.14] * \diamond [1.41, 1.42] = [-4.48, -4.42]$$

Ideal arithmetic

$$x + z = \{\tilde{x} + \tilde{z} \mid \tilde{x} \in x, \tilde{z} \in z\} = [\underline{x} + \underline{z}, \bar{x} + \bar{z}]$$

Rounded arithmetic

$$x +^\diamond z = \diamond[\underline{x} + \underline{z}, \bar{x} + \bar{z}]$$

$$\{\tilde{x} + \tilde{z} \mid \tilde{x} \in x, \tilde{z} \in z\} \subseteq x +^\diamond z$$

# Interval arithmetic in proof assistants

## Nature of interval methods

- interval arithmetic was born to safely deal with errors

## Usage

- interval arithmetic appears in critical software
- certified computation

## Formalizations

- CoQ, PVS, Isabelle
- focus on computation efficiency and automation of techniques

# Computation driven formalizations

- basic operations
- elementary functions
- techniques to increase accuracy
- rounded interval arithmetic
- automated procedures to compute and prove bounds for expressions
- computations by external tools

# Formalizing more “theoretical” results

# Formalizing more “theoretical” results

- solving systems of linear equations with interval coefficients

## Exercise

Consider the following system:

$$\begin{cases} [1, 2]x_1 + [2, 4]x_2 = [-1, 1] \\ [2, 4]x_1 + [1, 2]x_2 = [1, 2] \end{cases}$$

Find a box that contains all pairs  $(x_1, x_2) \in \mathbb{R}^2$  that satisfy the equations for some choice of coefficients in their respective intervals.

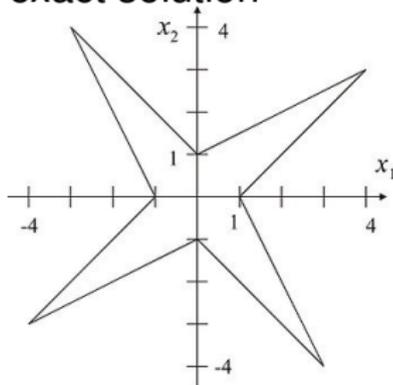
- correctness of methods for solving these systems is based on more involved theoretical results
- application: robot movement

# Solving systems of linear interval equations

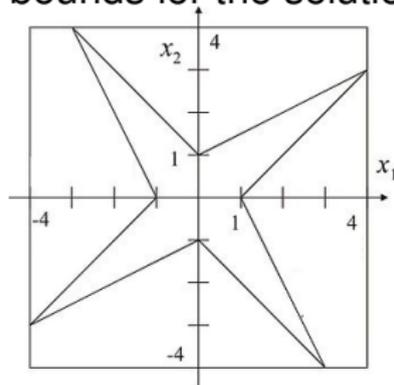
Two steps:

- 1 checking regularity of the associated interval matrix
- 2 computing bounds of the solution set

exact solution



bounds for the solution set

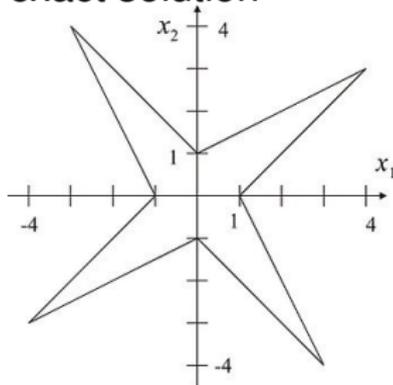


# Solving systems of linear interval equations

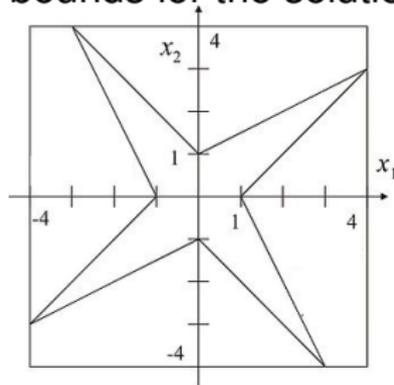
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exact solution



bounds for the solution set



## Definition

$$A = [A_{ij}]_{m \times n}, \quad A_{ij} \in \mathbb{IR}.$$

## Characterization

$$A = \{\tilde{A} \in M(\mathbb{R})_{m \times n} \mid \tilde{A}_{ij} \in A_{ij}, i = 1, \dots, m, j = 1, \dots, n\}.$$

## Associated real matrices

$$\underline{A} := [\underline{A}_{ij}] \quad \bar{A} := [\bar{A}_{ij}]$$

$$A_c := [(A_{ij})_c] \quad \Delta_A := [\Delta_{A_{ij}}]$$

## Addition

$$A + B := \square\{\tilde{A} + \tilde{B} \mid \tilde{A} \in A, \tilde{B} \in B\}$$

# Operations on interval matrices

## Addition

$$A + B := \square\{\tilde{A} + \tilde{B} \mid \tilde{A} \in A, \tilde{B} \in B\} = \{\tilde{A} + \tilde{B} \mid \tilde{A} \in A, \tilde{B} \in B\}$$

$$(A + B)_{ij} = A_{ij} + B_{ij}$$

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## Multiplication

$$AB = \square\{\tilde{A}\tilde{B} \mid \tilde{A} \in A, \tilde{B} \in B\} \neq \{\tilde{A}\tilde{B} \mid \tilde{A} \in A, \tilde{B} \in B\}$$

$$(AB)_{ij} = \sum_k A_{ik} B_{kj}$$

Special case: multiplication by a scalar vector

$$A\tilde{x} = \{\tilde{A}\tilde{x} \mid \tilde{A} \in A\}$$

# Regularity of interval matrices

An interval matrix  $A$  is called **regular** iff  $\forall \tilde{A} \in A, \det \tilde{A} \neq 0$

and it is called **singular** otherwise ( $\exists \tilde{A}, \tilde{A} \in A \wedge \det \tilde{A} = 0$ ).

! Notice the classical nature of the concepts we manipulate.

# Systems of linear interval equations

A system of linear interval equations with coefficient matrix  $A \in M(\mathbb{IR})_{m \times n}$  and right-hand side  $b \in \mathbb{IR}^m$  is defined as the family of linear systems of equations

$$\tilde{A}\tilde{x} = \tilde{b} \text{ with } \tilde{A} \in A, \tilde{b} \in b$$

The *solutions set* of such a system is given by:

$$\Sigma(A, b) := \{\tilde{x} \in \mathbb{R}^n \mid \exists \tilde{A} \in A, \exists \tilde{b} \in b \text{ such that } \tilde{A}\tilde{x} = \tilde{b}\}$$

# Proof example

## Theorem

$$\Sigma(A, b) = \{\tilde{x} \in \mathbb{R}^n \mid A\tilde{x} \cap b \neq \emptyset\}$$

## Proof excerpt.

We show:  $\{\tilde{x} \in \mathbb{R}^n \mid A\tilde{x} \cap b \neq \emptyset\} \subseteq \Sigma(A, b)$ .

Consider  $\tilde{x}$  such that  $A\tilde{x} \cap b \neq \emptyset$ .

Then  $A\tilde{x} \cap b$  contains some  $\tilde{b} \in \mathbb{R}^m$ .

Clearly  $\tilde{b} \in b$ .

Also,  $\tilde{b} \in A\tilde{x}$  and by relation (1),  $\tilde{b} = \tilde{A}\tilde{x}$  for some  $\tilde{A} \in A$ .

Therefore  $\tilde{x} \in \Sigma(A, b)$ . □

$$A\tilde{x} = \{\tilde{A}\tilde{x} \mid \tilde{A} \in A\} \tag{1}$$

# Setting up the formalization

We need to talk about

- real numbers
- matrices

We use

- COQ standard library **Reals**
- SSREFLECT library **matrix**

**Mix SSREFLECT and standard COQ !**

## in SSREFLECT

- types with decidable equality and a choice operator
- hierarchy of algebraic structures
- abstract matrices, but operations when elements are from a ring

## in COQ's Reals library

- axiom of trichotomy  $\Rightarrow$  decidable equality

**Axiom** total\_order\_T:  $\forall r1\ r2: R, \{r1 < r2\} + \{r1 = r2\} + \{r1 > r2\}$ .

- choice operator by choice and extensionality axioms (for now)
- ring structure

# Yet another formalization of intervals

## Definition

$x := [\underline{x}, \bar{x}] = \{\tilde{x} \in \mathbb{R} \mid \underline{x} \leq \tilde{x} \leq \bar{x}\}$ , where  $\underline{x}, \bar{x} \in \mathbb{R}, \underline{x} \leq \bar{x}$

**Structure** IR: **Type** := ClosedInt  
{ inf: R ; sup: R ; leq\_proof: inf  $\leq_b$  sup }.

## Intervals as sets

- coerce IR to  $\mathbb{R} \rightarrow \text{bool}$

## Equality of intervals

**Lemma** eq\_intervalP :  
 $\forall x z : \text{IR}, x = z \leftrightarrow \text{inf } x = \text{inf } z \wedge \text{sup } x = \text{sup } z.$

# leq\_proof: $\text{inf} \leq_b \text{sup}$

**Lemma** Rle\_dec:  $\forall r1\ r2, \{r1 \leq r2\} + \{\sim r1 \leq r2\}$ .

**Definition** Rleb r1 r2 :=  
 match (Rle\_dec r1 r2) with  
 | left \_  $\Rightarrow$  true  
 | right \_  $\Rightarrow$  false  
end.

$\text{inf} \leq_b \text{sup} \rightsquigarrow \text{Rleb inf sup} \rightsquigarrow \text{is\_true (Rleb inf sup)} \rightsquigarrow$   
 $\rightsquigarrow \text{Rleb inf sup} = \text{true}$

Boolean equality is decidable and therefore proof irrelevant.

$$\{\tilde{x} + \tilde{z} \mid \tilde{x} \in x, \tilde{z} \in z\} = [\underline{x} + \underline{z}, \bar{x} + \bar{z}]$$

Interval addition

- associative
- commutative
- has  $[0, 0]$  as neutral element

⇒ intervals with addition form a monoid

good news for work with big operators!

# Interval matrices

- use SSREFLECT library
- vectors are column matrices
- redefine operations on matrices as intervals do not have a ring structure

**Definition**  $\text{mmul}_i (A: 'M[IR]_(m, n)) (x: 'M[IR]_(n, 1)) :=$   
 $\backslash \text{col}_i \backslash \text{big}[\text{add}_i / 0 ]_j \text{mul}_i (A \ i \ j) (x \ j).$

- prove specific properties

$$A\tilde{x} = \{\tilde{A}\tilde{x} \mid \tilde{A} \in A\}$$

- associated real matrices

**Definition**  $\text{minf} (A: 'M[R]_(m, n)) := \backslash \text{matrix}_-(i, j) \text{inf} (A \ i \ j).$

# Results on real matrices

- norm for real matrices
- properties for symmetric and positive definite matrices
- eigenvalues for real matrices
  - Rayleigh quotients
- spectral radius
  - Perron Frobenius theorem

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# The issues

eigenvalues for real matrices:

- roots of the characteristic polynomial
- they can be complex
- Rayleigh quotient:  $\frac{x^T Ax}{x^T x}$ ,  $x \neq 0$ ,  $A$  – symmetric

$$\forall x \in \mathbb{R}^n, x \neq 0, \lambda_{\min}(A) \leq \frac{x^T Ax}{x^T x} \leq \lambda_{\max}(A)$$

spectral radius:  $\rho(A) = \max\{|\lambda(A)|\}$

## Theorem (Perron Frobenius)

If  $A \in \mathbb{R}^{n \times n}$  is nonnegative then the spectral radius  $\rho(A)$  is an eigenvalue of  $A$ , and there is a real, nonnegative vector  $x \neq 0$  with  $Ax = \rho(A)x$ .

# Formalized criteria of regularity

## Criterion

A is regular if and only if  $\forall \tilde{x} \in \mathbb{R}^n, 0 \in A\tilde{x} \Rightarrow \tilde{x} = 0$ .

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A is regular if and only if  $\forall \tilde{x} \in \mathbb{R}^n, |A_c \tilde{x}| \leq \Delta_A |\tilde{x}| \Rightarrow \tilde{x} = 0$ .

## Criterion (using positive definiteness)

If the matrix  $(A_c^T A_c - \|\Delta_A^T \Delta_A\| I)$  is positive definite for some consistent matrix norm  $\|\cdot\|$ , then A is regular.

## Criterion (using the midpoint inverse)

If the following inequality holds  $\rho(|I - RA_c| + |R|\Delta_A) < 1$  for an arbitrary matrix R, then A is regular.

## Criterion (using eigenvalues)

If the inequality  $\lambda_{\max}(\Delta_A^T \Delta_A) < \lambda_{\min}(A_c^T A_c)$  holds, then A is regular.

# How far from the real world

- adapt results for rounded arithmetic
- treat methods for bounding the solution set
- finish proving the admitted results

# Interesting References

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