

SSReflect - Logics & Basic tactics

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SSReflect – Reminder

(SSR = *Small Scale Reflection*)

SSREFLECT: extension of **Coq**

- developed while formalizing the Four Color Theorem (2004),
- now used for the Odd Order Theorem.

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Changes with standard **COQ**:

- Vernacular (Commands) and Gallina are mostly unchanged (e.g., `Definition`, `Lemma`, `forall`, `match with`);
- standard tactics are still available
- some tactics are superseded (e.g., `apply`, `rewrite`)
- new libraries are provided (e.g., `nat`, `seq`)

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- Add some ways to **structure** the scripts, so that breakages are easier to understand.
- Force the user to **explicitly name** things.
- Ease the use of **boolean reflection**.

Outline

- 1 Logics
- 2 Tactics, Tacticals
- 3 Proof Structure

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 - Booleans
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Minimal Propositional Logic

- Propositional variables: P Q $R \dots$
- Propositions: $(\text{even } 4)$ $(x < 10)$ $(7 \leq 2)$
- Implication: \rightarrow
- Formulas: $(P \rightarrow Q) \rightarrow (Q \rightarrow R) \rightarrow P \rightarrow R$

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- Propositional are of sort `Prop` : $(P : \text{Prop})$.
- Declaring variables: `Variables P Q R :Prop.`

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- Propositional are of sort `Prop` : $(P : \text{Prop})$.
- Declaring variables: `Variables` P Q $R : \text{Prop}$.
- Any term of type P ($p : P$) is a proof of P .

State and Proof a theorem

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Lemma imp_trans : (P -> Q) -> (Q -> R) -> P -> R.  
Proof. (* start the proof of a Lemma *)
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$$\left. \begin{array}{l} \vdots \\ P : Prop \\ Q : Prop \\ R : Prop \end{array} \right\} \text{named hypotheses (Context)}$$

$$\underbrace{(P \rightarrow Q) \rightarrow (Q \rightarrow R) \rightarrow P \rightarrow R}_{\text{Assumptions}} \underbrace{\}_{\text{Conclusion}} \text{ current goal}$$

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Tactic: any operation that allows the simplification, decomposition into subgoals, or resolution of a goal.

Proof

Theorem command:

Lemma `imp_trans` : $(P \rightarrow Q) \rightarrow (Q \rightarrow R) \rightarrow P \rightarrow R$.

Proof. (* start the proof of a Lemma *)

`move=>` `Hpq`.

$P : Prop$

$Q : Prop$

$R : Prop$

$Hpq : (P \rightarrow Q)$

$(Q \rightarrow R) \rightarrow P \rightarrow R$

Proof

Theorem command:

Lemma `imp_trans` : (P → Q) → (Q → R) → P → R.

Proof. (* start the proof of a Lemma *)

`move=>` `Hpq Hqr p`.

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Hpq : (P → Q)

Hqr : (Q → R)

p : P

R

Proof

Theorem command:

Lemma `imp_trans` : (P → Q) → (Q → R) → P → R.

Proof. (* start the proof of a Lemma *)

`move=>` `Hpq Hqr p`.

`apply:` `Hqr`.

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Q : Prop

R : Prop

Hpq : (*P* → *Q*)

p : *P*

Q

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`move=>` `Hpq Hqr p`.

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`apply:` `Hpq`.

`exact:` `p`.

Proof completed.

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Theorem command:

Lemma `imp_trans` : $(P \rightarrow Q) \rightarrow (Q \rightarrow R) \rightarrow P \rightarrow R$.

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Proof completed.

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Theorem command:

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move=> Hpq Hqr p.
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apply: Hqr.
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exact: (Hpq p).
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```
Qed.
```

Minimal Propositional Logic with universal quantifier

- `forall (P Q R :Prop), (P ->Q)-> (Q -> R) -> P -> R`

Minimal Propositional Logic with universal quantifier

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`apply:H. apply:(H A B). or ...`

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 - `move=>n.`

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Propositional Logic, Conjunction

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 (* by choosing the left part *)

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$$\frac{}{A \vee B} \rightarrow \frac{}{A}$$

- `right`. (* Prove a disjunction $:A \vee B$ *)
(* by choosing the right part *)

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 - `move=>B.` (* To prove the goal $(\sim B)$ *)

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- Then `apply:H.` (*for a $(H : \sim C)$ in the context*)

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Existential: `exists n:nat, P n`

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- `case: Hex`.

(* To break the (Hex:exists n, P n)hypothesis *)

(* combined with (move=>n Hn.)*)

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$$\frac{\frac{\text{Hex: exists } n, P n}{G}}{\text{exists } n, P n} \rightarrow \frac{\frac{n : \text{nat} \quad \text{Hn} : P n}{G}}{\text{exists } n, P n}$$

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Booleans

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b_1	b_2	$b_1 \ \&\& \ b_2$	$b_1 \ \ b_2$	$b_1 \ ==> \ b_2$	$b_1 \ (+) \ b_2$
T	T	T	T	T	F
T	F	F	T	F	T
F	T	F	T	T	T
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- Operators: `"&&"`, `"||"`, `"~~"`, `"==>"`, `"(+)"`.

- Some notations
 - `"[&& b1 , b2 , .. , bn & c]" := (b1 && (b2 && .. (bn && c)..))`
 - `"[|| b1 , b2 , .. , bn | c]" := (b1 || (b2 || .. (bn || c)..))`

Booleans

- **Inductive** `bool := true | false.`
- Operators: `"&&"`, `"||"`, `"~~"`, `"==>"`, `"(+)"`.

- `is_true : bool -> Prop.`
 - `fun b : bool => b = true.`
 - Notation: `"x 'is_true'" := (is_true x)`

Booleans in proofs

- Reason by case on a boolean:

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b : bool

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b : bool

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`false (+) b = (false && ~~b) || (~~false && b)`

Booleans in proofs(2)

- Compute, simplify:

```
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`b : bool`

`~ b = ~ b || false`

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Tactics / Tacticals

- **Tactic**: any operation that allows the simplification, decomposition into subgoals, or resolution of a goal.
- **Tactical**: any function of tactics (eg. `;` the composition of two tactics).

Tactics and Tacticals

- `move=>`
- `by`
- `apply:`
- `exact:`
- `case:`
- `elim:`
- `rewrite`

tactical

Introduction Tactic

- `move=> a b c.`
`pops` the top 3 elements of the goal, and it puts them into the context with `names` *a*, *b*, and *c*.
- `move=> _.`
`pops` the first top element of the goal, without putting it in the context.
- `move=> a _ c.`

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$$\begin{array}{ccc}
 \vdots & & \vdots \\
 \frac{H: P \rightarrow Q}{\underline{\underline{Q}}} & \rightarrow & \frac{\vdots}{\underline{\underline{P}}}
 \end{array}$$

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- "apply:H" applies H to the goal.

$$\frac{\vdots}{\frac{H: P \rightarrow Q}{Q}} \rightarrow \frac{\vdots}{P}$$

- "exact:H" performs "by apply: H"

Tactics case: / elim:

Performs a **case analysis** / **inductive elimination** on the element given as an argument.

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Lemma P_of_n forall n : nat, P n.  
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case:n.

1. $\frac{}{P\ 0}$
2. $\frac{}{\text{forall } n : \text{nat}, P\ (S\ n)}$

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Basic Rewriting tactic

Tactic "`rewrite items...`" modifies subterms of the goal:

- "`/name`" unfolds a definition
- "`-/name`" folds a definition
- "`term`" rewrites (left to right) with a lemma or an hypothesis which conclusion is an equality

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$$\frac{\text{Eqab: } a = b}{P \ a} \quad \rightarrow \quad \frac{\text{Eqab: } a = b}{P \ b}$$

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$$\frac{\text{Eqab: } a = b}{P \ a} \quad \rightarrow \quad \frac{\text{Eqab: } a = b}{P \ b}$$

- "`-term`" rewrites right to left

Multiple Rewriting and Occurrence selection

rewrite -*multiplicity* term

- " ? " : as many times as possible, possibly none,
- " ! " : as many times as possible, at least once,
- " n ? " : at most n times,
- " n ! " : exactly n times.

Multiple Rewriting and Occurrence selection

rewrite -*multiplicity* term

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Lemma dbl a b : 2 * (a + b) = (b + a) + (a + b).

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Proof.

a : nat

b : nat

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rewrite -!addnA.

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b : nat

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b : nat

2 * (a + b) = (b + (a + (b + a)))

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Forward Reasoning

- have
- suffices (suff)

Forward Reasoning: `have` / `suffices`

`have` H : *intermediate_goal*
performs a **logical cut**.

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Variable f : nat -> nat.
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```
Variable P : nat -> Prop.
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Lemma P_of_3: P 3.
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have H: exists x, f x = 3.
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Lemma P_of_3: P 3.
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`Proof`.

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have H: exists x, f x = 3.
```

$$1. \frac{}{\text{exists } x, f x = 3} \quad 2. \frac{H: \text{exists } x, f x = 3}{P 3}$$

Forward Reasoning: `have` / `suffices`

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have H : intermediate_goal  
  performs a logical cut.
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Variable f : nat -> nat.  
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Lemma P_of_3: P 3.
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Proof.  
have H: exists x, f x = 3.
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Tactic "suff" also performs a logical cut, but it produces the two subgoals in the opposite order.

Proof control flow

- **Tabulation** (depending on the number of subgoals number)

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Proof control flow

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- **Bullets** -, +, *
- **Proof terminators** : by , **exact**:

A proof example

```

Emacs@zeniba-2
case E1: (abezoutn _ _) => [[| k1] [| k2]].
- rewrite !muln0 !gexpn0 mulg1 => H1.
  move/eqP: (sym_equal F0); rewrite -H1 orderg1 eqn_mull.
  by case/andP; move/eqP.
- rewrite muln0 gexpn0 mulg1 => H1.
  have F1: t %| t * S k2.+1 - 1.
    apply: (@dvdn_trans (orderg x)); first by rewrite F0; exact: dvdn_mull.
    rewrite orderg_dvd; apply/eqP; apply: (mulgI x).
    rewrite -{1}(gexpn1 x) mulg1 gexpn_add leq_add_sub //.
    by move: P1; case t.
    rewrite dvdn_subr in F1; last by exact: dvdn_mulr.
  + rewrite H1 F0 -{2}(muln1 (p ^ 1)); congr (_ * _).
    by apply/eqP; rewrite -dvdn1.
  + by move: P1; case: (t) => [| [| s1]].
- rewrite muln0 gexpn0 mulg1 => H1.□

-1:--- f.v          All L15      (coq Holes)-----

```

Subgoal Selectors

- Solving one subgoal with a single tactic:
 - *tactic* ; first **by** *tactic*
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- Changing the order of subgoals:
 - *tactic* ; first last (or last first)