

A queuing analysis of packet dropping over a wireless link with retransmissions

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Abstract. One of the most important performance measure for real time applications is the packet loss probability. In a wireless environment losses are first due to noisy channels (fading, shadowing etc). In order to improve the radio link, one often retransmits packets that have not been well received (using the Automatic Retransmission reQuest - ARQ). This however may lead to queuing phenomena due to retransmissions and to losses of packets due to buffer overflow. We present a queuing analysis in order to compute the combined effect of the noisy channel, the retransmissions and the buffer overflow. We use advanced spectral methods of quasi-birth and death process in order to analyze the system for general multistate Markov channels and provide an approximation approach based on the theory of singular perturbation.

1 Introduction

The study of the performance of wireless channels and their effect on the quality of service is a major research issue as it provides tools for design of link layer protocols and dimension link layer buffers, see e.g. [12, 13] and references therein. We study in this paper wireless channels that use retransmission of lost packets: the Automatic Retransmission reQuest (ARQ) protocol. We note however that retransmissions cause extra delay which results in larger queues and therefore in possible losses due to buffer overflow. In order to study both the delay statistics as well as the global loss process, we provide a queuing analysis of the radio link.

Our starting point is a link model provided in [12]. The radio channel considered in [12] is modeled using a two state Markov chain: a perfectly "good" state in which transmission succeeds with probability one, and a perfectly "bad" state in which transmission fails with probability one. Our first goal is to allow for a much more rich class of models for the channel state. The need for richer models is motivated by the following considerations:

- It has been shown in [9, 11] that more than two states are needed for describing some radio channels; in fact it was shown that one needs typically at least four states.
- Even in the "best" channel state it might be too optimistic to expect losses to occur with probability zero. Conversely, it may be too pessimistic to assume that in a bad state all packets are systematically lost.
- Assume that for a given packet loss probability p we assign a single state of the channel. This means (if a discrete time model as in [12] is used) that the time during which a packet suffers a loss probability of p has a geometric distribution. If we wish to be able to model other distributions then more than one state should be assigned to a given loss probability!

Our analysis approach is based on modeling of the system as a two dimensional Markov chain, where the first corresponds to the channel state and the second to the number of packets queued at the buffer. In order to obtain the steady state distribution of this system we make use of recent advanced spectral methods of analysis of quasi

birth and death (QBD) processes [5, 6]. These methods allow us to represent the steady state solution as a simple function of the eigenvectors and eigenvalues of some matrices. For the case that the channel is described by a two dimensional chain, we obtain an *explicit solution* which has not been available in [12] even for the simpler special case of "perfectly" good and "perfectly" bad channel states.

The second objective of the paper is to study a singular perturbation solution approach for approximating the steady state solution. This approach allows to obtain simple approximations for the case that the channel state evolves much slower than the queue size. In fact, if the speed ratio is given by a parameter ϵ , then the singular perturbation approach allows to represent the steady state probabilities as a Taylor series of ϵ . The first term in the series already gives a very good approximation when ϵ is sufficiently small, and this term is much simpler to compute than by using the previous exact approach.

The structure of the paper is as follows. We present the model in section 2. Section 3 provides a detailed spectral analysis and give the steady state channel-queue length for the finite queue model. We concluded by sample numerical results. In section 4 We study a singular perturbation solution approach for approximating the steady state solution and we compare the two methods presented in section 3 and section 4 by numerical results.

2 Model

We consider the following queuing system. In each slot, a packet is generated at the transmitter according to a Bernoulli process with arrival probability $p > 0$. The channel state can be described by the variable $l \in E := \{1, 2, \dots, L\}$, where $l = i$ denote that the packet can be transmitted with probability q_i , with $q_i \in [0, 1]$. If a transmission fails we assume that it will be attempted again in the next time unit. As in the first model studied in [12], we assume that the number of retransmissions of a packet by the ARQ protocol is not bounded: retransmissions will be attempted until a success.

$$A = \begin{pmatrix} \lambda_{11} & \lambda_{12} & \cdots & \lambda_{1L} \\ \lambda_{21} & \lambda_{21} & \cdots & \lambda_{2L} \\ \vdots & \vdots & \vdots & \vdots \\ \lambda_{L1} & \lambda_{22} & \cdots & \lambda_{LL} \end{pmatrix} \quad (1)$$

be the transition matrix for the error process, where $\sum_{j=1}^L \lambda_{ij} = 1$. We shall denote $\nu = [\nu_1, \nu_2, \dots, \nu_L]$ the probability distribution of the transition matrix A . We assume throughout that A is irreducible and aperiodic.

The state of the system at time n is described by $X(n) = (l(n), i(n))$, where $l(n)$ is the state of channel in slot n and $i(n)$ is the number of packets in the queue at the beginning of slot t . The arrival process in slot m is independent of $X(n)$ for all $m \geq n$. Arrivals are assumed to occur at the end of the slot so that, in particular, they cannot depart in the same slot.

The system has a finite queue of size B . A loss occurs whenever an arriving packet finds the queue full. It can be shown that the Markov channel evolution and the assumption of independent arrivals imply that $X(n)$ is a Markov chain with value in a finite state space $\mathcal{B} = \{(i, j), i = 1, \dots, L, j = 0, \dots, B\}$

Since we consider discrete time, we have to make some convention on the order of events within each time unit. We assume that at each time unit, if the queue is nonempty then first a transmission is attempted, and only if it succeeds then there is a departure; then if there have been an arrival during the slot then it joins the queue. We call this a service-arrival order. This order means in particular that if the queue

is full at the beginning of the slot then an arriving packet during the slot need not be lost: it is only lost if the transmission has failed.¹

The transition probability $L(B+1) \times L(B+1)$ matrix of the Markov chain $X(n)$ are given by

$$P = \begin{pmatrix} B_0 & B_1 & 0 & \cdots & & & & 0 \\ A_2 & A_1 & A_0 & 0 & \cdots & & & 1 \\ 0 & A_2 & A_1 & A_0 & 0 & \cdots & & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & \ddots & & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & A_1 & A_0 & \vdots \\ 0 & \cdots & & 0 & A_2 & F_1 & & B \end{pmatrix}$$

P consists of $B+1$ block rows and block columns. The matrices $B_0, B_1, A_0, A_1, A_2,$ and F_1 are all $L \times L$ non-negative matrices, where $B_0 = \bar{p}\Lambda, B_1 = p\Lambda, A_0 = \text{diag}(p\bar{q}_i, i=1, \dots, L)\Lambda, A_1 = \text{diag}(pq_i + \bar{p}\bar{q}_i, i=1, \dots, L)\Lambda, A_2 = \text{diag}(\bar{p}q_i, i=1, \dots, L)\Lambda$ and $F_1 = \text{diag}(pq_i + \bar{q}_i, i=1, \dots, L)\Lambda$. Since Λ is irreducible and aperiodic, and since $p > 0$, the Markov chain $X(n)$ is aperiodic and it contains a single ergodic class. Hence it has a unique stationary (or steady state) probability distribution defined as

$$\pi(i, j) = \lim_{n \rightarrow \infty} Pr(l(n) = i, i(n) = j); \quad i = 1, 2, \dots, L; \quad j = 0, 1, \dots, B$$

The objective of this analysis is to determine the blocking probabilities, the expected delays, and in fact the whole probability distribution $\pi(i, j)$ in terms of the parameters of the system. $\pi(i, j)$ is the unique solution to the equation

$$\pi = \pi P, \quad \pi \bar{e}_L = 1 \tag{2}$$

where the notation $\bar{e}_k, k \geq 0$ is used to denote a column vector of ones of length $k(B+1)$. For convenience, we partition the stationary vector π as $[\pi_0, \pi_1, \dots, \pi_B]$, where π_k of size $1 \times L$, will called the solution vector of level $k, 0 \leq k \leq B$.

3 Analysis of the system

3.1 Spectral analysis

In the analysis a special role is played by the number of states of the server for which the transmission probability is zero or one. We denote these numbers by

$$m_0 := |\{i : q_i = 1\}|, \quad m_2 := |\{i : q_i = 0\}|$$

The structure of the transition probabilities would be similar to the one studied [6] if we did not have the first row and column. We handle this problem by reducing the state space, by removing the level zero and in the second phase, we construct the stationary probability of the reduced chain by using the spectral analysis. This will allow us to obtain the steady state probabilities of the original system.

Let \mathcal{B}_1 be subset of \mathcal{B} , where $\mathcal{B}_1 = \{(i, j) / i = 1, \dots, L; j = 1, \dots, B\}$. We define the restricted process $\{X_n^0, n \geq 0\}$ as follows: $X_n^0 = X(s_n)$ for $n \geq 0$, so that X_n^0 is the actual state visited when the Markov chain is not for the n th time in level 0.

¹ If in contrast, a real system operates in the arrival-service order i.e. first an arrival occurs (and is lost if the queue is full) and then a transmission is attempted, then the expressions for the steady state probability obtained for the service-arrival order are still valid but they apply to the state of the queue just *after* the arrival occurred in the arrival-service order system. From this it is then easy to derive the steady state distribution at the beginning of a time slot.

Remark 1. Recall m_0 and m_2 are the number of state $i \in \{1, 2, \dots, L\}$ for which $q_i = 1$ and $q_i = 0$. Then after some permutation of the rows in matrix A_i , $i = 0, 2$, the matrix A_i $i = 0, 2$ can be writing as the following block triangular matrix²

$$A_i = \begin{pmatrix} A_i^1 & A_i^2 \\ \mathbf{0} & \mathbf{0} \end{pmatrix} \quad (9)$$

The eigenvalues of A_i , $i = 0, 2$ are those of A_i^1 with those of matrix null $\mathbf{0} \in M_{m_i}$, counting multiplicities. Then if A_i^1 , $i = 0, 2$ has eigenvalue *zero* with multiplicity n_i , 0 is eigenvalue of A_i with algebraic multiplicity r_i where $r_i = m_i + n_i$, $i = 0, 2$.

In the sequel, we assume that the sum of all principal minus of $T(z)$ of order $L - n_0$ is nonzero. We deduce from these results the following

Lemma 3. (i)- The polynomial $\det[T(z)]$ has a root with geometric multiplicity r_0 located at $z = 0$, and (ii)- $2L - r_0 - r_2$ roots. For their location we distinguish two cases:

1. If $p < \sum_{i=1}^L \nu_i q_i$, $L - r_0$ of them lie in disc $|z| < 1$, one at $z = 1$, and $L - 1 - r_2$ in disc $|z| \geq 1$ with $z \neq 1$.
2. If $p > \sum_{i=1}^L \nu_i q_i$, $L - 1 - r_0$ of them lie in disc $|z| < 1$, one at $z = 1$, and $L - r_2$ in disc $|z| \geq 1$ and $z \neq 1$

Proof see the full version of the paper [4] ■

Now, we return to the steady-state for the restricted process X_n^0 . It was already showed that the matrices R and S have a full set of eigenvectors. Consequently the expression (5) can thus rewritten in spectral-expansion form :

$$\pi_j^0 = \sum_{k=1}^L \beta_k (\phi_k)^{j-1} w_k + \sum_{k=L+1}^{2L} \beta_k (1/\phi_k)^{B-j} w_k, \quad 1 \leq j \leq B \quad (10)$$

with ϕ_1, \dots, ϕ_L the eigenvalues of the matrix R and w_1, w_2, \dots, w_L the corresponding left eigenvectors, and $1/\phi_{L+1}, \dots, 1/\phi_{2L}$ are the eigenvalues of S with corresponding left eigenvectors w_{L+1}, \dots, w_{2L} . If $m_2 + n_2 > 0$, then we set the corresponding ϕ_k equal to ∞ and we write $1/\phi_k = 0$. Moreover by convention we set $0^0 = 1$.

Here the vectors π_j , $j = 1, 2, \dots, L$ are know if β_k 's are known. We derive an alternative set of equations which uniquely determines the coefficients β_k by combining the relation (10) with the equations (3), (4) and the normalization condition.

3.2 Examples and numerical results

In the full version of the paper [4], we illustrate how to apply this method to a channel that can be described by two states.

Numerical investigation In this paragraph, we present some numerical results obtained based on the analysis described in previous section 3. We investigate the relationship between the dropping probability and arrival rate p . In particular, we compare the values of the performance measures for two different schemes: The first schemes is perfectly good/absolutely bad state in which the transmission succeeds with probability 1 or 0, and second scheme is a perfectly good/fairly bad in which the transmission succeeds with probability 1 or 0.2.

² Suppose that there exists $i \in \{1, 2, \dots, L\}$ such that $q_i = 0$ (resp $q_i = 1$). Then all entries in row i of the matrix S (resp R) equal 0, so that 1_i with the i^{th} entry equal to 1 and all other entries equal to 0 is a left nullvector of S (resp R). Indeed, when the state of server is i in which $q_i = 0$, no packet can be transmitted. Hence starting from (j, i) , $\forall j$, no state state of level $j - 1$ can ever be visited before another state in the level j is visited.

In figures 1 and 2, the dropping probability is plotted versus the arrival rate for first scheme (resp. scheme 2).

As shown in figure 1-2, the performance measures (dropping probability) for two schemes are different. It is true that the modification on the transmission in bad state will influence the performance, but more surprisingly, the difference between the performance measures is of the order 10^{-1} . This actually confirms that using the simple model of scheme 1 (e.g. [12, 14, 13]) does not adequately capture the process of packet errors. In order to verify this, we calculate the performance measures with a small change in the transmission probability q_2 in the region near 0 for scheme 1 (see figure 3), we observe a large effects on the performance measures when the arrival rate is more than 0.3. From this observation, it is clear that the values of the transmission probabilities have an important role to model the process packet errors. We conclude that we have to be careful when using the simplify model to represent the wireless the analytic system.

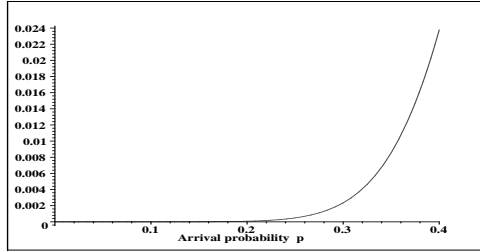


Fig. 1. The dropping probability of a packet as function of the arrival probability p for $q_1 = 1$, $q_2 = 0$, $\lambda_{11} = 0.3$, $\lambda_{22} = 0.5$ and buffer size $B = 6$

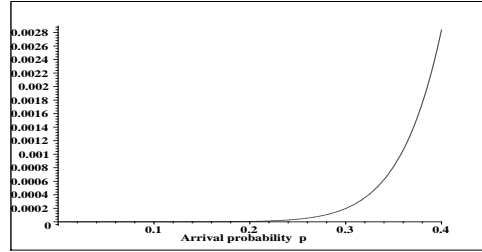


Fig. 2. The dropping probability of a packet as function of the arrival probability p for $q_1 = 1$, $q_2 = 0.2$, $\lambda_{11} = 0.3$, $\lambda_{22} = 0.5$ and buffer size $B = 6$

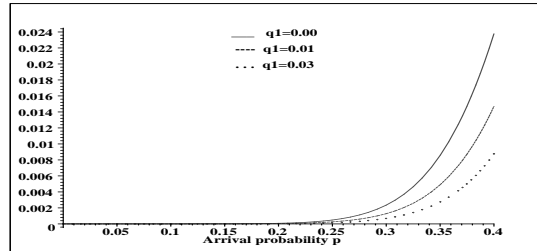


Fig. 3. The dropping probability of a packet as function of the arrival probability p for $q_1 = 1$, $\lambda_{11} = 0.3$, $\lambda_{22} = 0.5$ and buffer size $B = 10$ and $q_2 = 0, 0.01, 0.03$

4 Approximations using singular perturbation

4.1 Preliminaries

Singular perturbation is a powerful approach to obtain simple approximations of steady state probabilities of Markov chains in which we can identify two time scales: there are several groups of states such that there are rare transitions between groups and frequent transitions within each group (see e.g. [2, 3, 10]). We shall apply this method to the case when the transitions of the queue size occur much more frequently than those of the channel states. This will typically be the case in high speed

channels. The singular perturbation approach can also be useful for the opposite case in which the channel states evolve much faster than the queue size process.

We thus represent the transitions of the channel state as $\Lambda(\epsilon) = I + \epsilon A_1$ where A_1 is a generator (so that $\Lambda(\epsilon)$ represents indeed transition probabilities) and ϵ is a small parameter. Hence the transition matrix of the Markov chain $X(n)$ becomes

$$P(\epsilon) = Q_0 + \epsilon Q_1$$

where Q_0 is the probability transition matrix of the unperturbed Markov chain corresponding to strong interaction, and Q_1 is the generator corresponding to weak interaction, i.e.,

$$Q_0 = \begin{pmatrix} \bar{p}I & pI & 0 & \cdots \\ \bar{A}_2 & \bar{A}_1 & \bar{A}_0 & 0 & \cdots \\ 0 & \bar{A}_2 & \bar{A}_1 & \bar{A}_0 & 0 & \cdots \\ \vdots & \vdots & \vdots & \vdots & \vdots & \vdots \\ \vdots & \vdots & \vdots & \vdots & \vdots & \bar{A}_0 \\ 0 & \cdots & 0 & \bar{A}_2 & \bar{F}_1 \end{pmatrix} \quad \text{and} \quad Q_1 = \begin{pmatrix} \tilde{B}_0 & \tilde{B}_1 & 0 & \cdots \\ \tilde{A}_2 & \tilde{A}_1 & \tilde{A}_0 & 0 & \cdots \\ 0 & \tilde{A}_2 & \tilde{A}_1 & \tilde{A}_0 & 0 & \cdots \\ \vdots & \vdots & \vdots & \vdots & \vdots & \vdots \\ \vdots & \vdots & \vdots & \vdots & \vdots & \tilde{A}_0 \\ 0 & \cdots & 0 & \tilde{A}_2 & \tilde{F}_1 \end{pmatrix}$$

where $\bar{A}_2 = \text{diag}(\bar{p}q_i, i = 1.., L)$, $\bar{A}_1 = \text{diag}(pq_i + \bar{p}\bar{q}_i, i = 1.., L)$, $\bar{A}_0 = \text{diag}(p\bar{q}_i, i = 1.., L)$, $\bar{F}_1 = \text{diag}(pq_i + \bar{q}_i, i = 1.., L)$, $\tilde{B}_0 = \bar{p}(A_1)$, $\tilde{B}_1 = p(A_1)$, $\tilde{A}_2 = \text{diag}(\bar{p}q_i, i = 0, 1.., L)A_1$, $\tilde{A}_1 = \text{diag}(pq_i\bar{p}\bar{q}_i, i = 1.., L)A_1$, $\tilde{A}_0 = p\text{diag}(p\bar{q}_i, i = 1.., L)A_1$, and $\tilde{F}_1 = \text{diag}(pq_i + \bar{q}_i, i = 1.., L)A_1$.

Let $\pi^\epsilon(k, i)$, $i = 1..L$, $j = 0, 1, \dots, B$, be the probability distribution of the state of the channel and the number of packets in the system in steady-state. This is exactly the perturbation problem under the assumption that there exists a ergodic class (i.e., has exactly one closed communicating set of states), and Q_0 contains L subchains.

The partition of state set of Markov chain associated with the probability matrix Q_0 depends of the number of service for which the transmission probability is zero or one. Indeed, if $m_0 = m_2 = 0$, the L recurrent classes ξ_i , $i = 1, \dots, L$ are given by

$$\xi_i = \{(i, j), j = 0, 1, \dots, B\}$$

and the unperturbed chain admits no transient state (block diagonal structure). Now if $m_0 > 0$ or $m_2 > 0$, let $E_B = \{i / q_i = 0\}$ and $E_G = \{i / q_i = 1\}$. In this case, the L recurrent classes are given by : $\xi_i = \{(i, 0), (i, 1)\}$ if $i \in \{j / q_j = 1\}$, $\xi_i = \{(i, B)\}$ if $i \in \{j / q_j = 0\}$ and $\xi_i = \{(i, j) j = 0, 1, \dots, B\}$ otherwise. The set of transient states T is given by $T = \{(i, j) / j \neq 0, 1 \text{ and } q_i = 1\} \cup \{(i, j) / j \neq B \text{ and } q_i = 0\}$. Without loss of generality, the channel state in which the transmission probability is zero or one, are described by the variable $l \in \{1, \dots, r_0\}$ where $r_0 = m_0 + m_2$.

4.2 Taylor series expansion for steady state

The above assumptions still implies that the stationary probability π^ϵ has the following property

Lemma 4. *The stationary probability π^ϵ of the perturbed chain has a Taylor series expansion : $\pi^\epsilon = \sum_{n=0}^{\infty} \pi^{(n)} \epsilon^n$. Moreover $\pi^{(n)}$, $n \geq 0$ are recursively calculated as follows:*

$$\pi^{(0)}[I - Q_0] = 0, \quad \sum_{(i,j)} \pi^{(0)}(i, j) = 1 \quad (11)$$

$$\pi^{(n)}[I - Q_0] = \pi^{(n-1)}Q_1, \quad \sum_{(i,j)} \pi^{(n)}(i, j) = 0 \quad (12)$$

Proof see Schweitzer [10], ■

We consider the unperturbed Markov chain associated with the transition probability matrix Q_0 . There exists a partition of states set of Markov chain $X(n)$ into of family of L recurrent classes ξ_i , $i = 1, \dots, L$. To each class $\bar{\xi}$ is associated the invariant measure (row vector) $m_{\bar{\xi}}$ of the recurrent subchain defined on the class $\bar{\xi} \in \bar{S} = \{\xi_1, \dots, \xi_L\}$.

Let us consider the natural numbering the states of \mathcal{B} after the grouping defined by the partition \bar{S} . With this numbering, the transition probability matrix Q_0 has the following block structure:

$$Q_0 = \begin{pmatrix} P_1 & \cdots & 0 & 0 & \cdots & 0 & 0 & \cdots & 0 \\ \vdots & \ddots & \vdots & & \ddots & & & & \ddots \\ & & P_r & 0 & \cdots & 0 & 0 & \cdots & 0 \\ 0 & \cdots & & P_{r+1} & \cdots & 0 & 0 & \cdots & 0 \\ & & \vdots & \vdots & \ddots & \vdots & & & \ddots \\ 0 & \cdots & 0 & 0 & 0 & P_L & 0 & \cdots & 0 \\ F_1 & \cdots & 0 & 0 & \cdots & 0 & Y_1 & \cdots & 0 \\ \vdots & \ddots & \vdots & \vdots & \ddots & \vdots & \vdots & \ddots & \vdots \\ 0 & \cdots & F_r & 0 & \cdots & 0 & 0 & \cdots & Y_r \end{pmatrix}$$

where we distinguish three cases : For $i \notin \{1, \dots, r\}$, P_i is a $(B+1) \times (B+1)$ matrix :

$$P_i = \begin{pmatrix} \bar{p} & +p & \cdots & & & \\ \bar{p}q_i & pq_i + \bar{p}q_i & & p\bar{q}_i & & 0 \\ 0 & \cdots & \cdots & & \cdots & \\ \vdots & & & & & \\ 0 & \cdots & & pq_i + \bar{p}q_i & & p\bar{q}_i \\ & & & \bar{p}q_i & & pq_i + \bar{q}_i \end{pmatrix}$$

For $i \in E_G$, P_i is 2×2 matrix, F_i is $(B-1) \times 2$ and Y_i is a $(B-1) \times (B-1)$ matrix :

$$P_i = \begin{pmatrix} \bar{p} & p \\ \bar{p} & p \end{pmatrix}, \quad F_i = \begin{pmatrix} 0 & \bar{p} \\ \vdots & 0 \\ \vdots & \\ 0 & 0 \end{pmatrix}, \quad Y_i = \begin{pmatrix} p & 0 & \cdots & 0 \\ \bar{p} & p & \cdots & 0 \\ \vdots & \vdots & \ddots & \vdots \\ 0 & \cdots & \bar{p} & p \end{pmatrix}$$

For $i \in E_B$, the matrix $P_i = 1$, F_i is a B vector given by $F^T = (0, \dots, 0, p)$ and Y_i is a $B \times B$ matrix, where

$$Y_i = \begin{pmatrix} \bar{p} & p & \cdots & 0 \\ 0 & \cdots & \cdots & 0 \\ \vdots & \cdots & \bar{p} & p \\ 0 & \cdots & 0 & \bar{p} \end{pmatrix}.$$

Let us define the matrix $\bar{Q}_1 = (\bar{Q}_1)_{\xi\xi'}$ by: $(\bar{Q}_1)_{\xi_i\xi_j} := \sum_{k=1}^B m_{\xi_i}(k) \sum_{k'=1}^B (Q_1)_{(i,k),(j,k')}$

We note that \bar{Q}_1 is a generator of an aggregated Markov chain on a state space \bar{S} . Now let $\bar{\nu} = [\bar{\nu}_1, \dots, \bar{\nu}_L]$ the probability distribution of the aggregated Markov chain. Note in view of definition of \bar{Q}_1 , we observe that the generator of aggregated Markov chain is Λ_1 . Hence $\bar{\nu}$ is the unique stationary distribution of ergodic Markov chain $l(n)$, i.e.,

$$\bar{\nu} \Lambda_1 = 0, \quad \bar{\nu} \mathbf{1} = 1. \quad (13)$$

Following Delebecque[3], the unperturbed stationary probability is given by

$$\pi^{(0)}(i, j) = \bar{\nu}_i m_{\xi_i}(j) \quad (14)$$

where m_{ξ_i} is the probability distribution of the recurrent class which given by the following lemma.

Lemma 5. For each $i = 1, 2, \dots, L$, the stationary distribution m_{ξ_i} of recurrent class ξ_i is given by:

1- If $i \notin E_G \cup E_B$, $m_{\xi_i}(j) = \frac{p^j(1-q_i)^{j-1}}{q_i^j(1-p)^j} m_{\xi_i}(0)$, $j = 1, \dots, B$, where $m_{\xi_i}(0) = \frac{q_i - p}{q_i - p\alpha^B}$

with $\alpha = \frac{p(1-q_i)}{q_i(1-p)}$.

2- If $i \in E_G$, $m_{\xi_i}(0) = 1 - p$ and $m_{\xi_i}(1) = p$.

3- If $i \in E_B$, $m_{\xi_i}(B) = 1$.

Proof see the full version of the paper [4]. ■

We shall denote q_{ξ_i} the right eigenvector corresponding to the one eigenvalue of the matrix P_i . Recall that the right eigenvector provide the probability to end in the class ξ_i starting from i . In our model, the vector q_{ξ_i} is given as follows: For each $j = 0, 1, \dots, B$, $q_{\xi_i}(k, j) = 1$, if $k = i$ and $q_{\xi_i}(k, j) = 0$ otherwise.

Following the analysis in [1, Chap. 3], we can obtain the stationary distribution of the perturbed system by another equivalent system with the matrix coefficients of smaller dimension.

Theorem 1. A solution of the equations (11)-(12) with the normalization conditions is given by the following recursive formulae:

1. $\pi^{(0)} = \bar{\nu}M$, where $\bar{\nu}$ is the stationary distribution of aggregated Markov chain $A_1 \in \mathbb{R}^{L+1 \times L+1}$ and $M \in \mathbb{R}^{(L+1) \times (L+1)(B+1)}$ is a matrix whose rows are stationary distribution of the Markov chain Q_0 and $N \in \mathbb{R}^{(L+1)(B+1) \times (L+1)}$ is a matrix of eigenvectors corresponding to the zero eigenvalue of the unperturbed generator $Q_0 - I$.

2. $\pi^{(n)} = \pi^{(n-1)}Q_1H_0 + \pi_1^{(n)}M$, $\pi_1^{(0)} = \nu$, $\pi_1^{(n)} = \pi_1^{(n-1)}MQ_1H_0Q_1NH_1$, where H_0 is a deviation matrix of the unperturbed Markov chain $H_0 = [I - Q_0 + NM]^{-1} - NM$ and H_1 is a deviation of the aggregated Markov chain $H_1 = [-A_1 + \Lambda^*]^{-1} - \Lambda^*$, where Λ^* is the ergodic projection of transition matrix $\Lambda - I$, i.e., $\Lambda^* = \mu_L \nu$, with $\mu_L = [1, \dots, 1]^T \in \mathbb{R}^{L+1}$.

Observe that the stationary distribution is obtained here more efficiently than by using the spectral analysis given in section 3. In this reduced system, we need only to compute the stationary distribution of aggregated Markov chain A since the stationary distributions of each recurrent class of the unperturbed matrix transition Q_0 are given by Theorem 1.

4.3 Examples and numerical results

Consider the same example as the one presented in subsection 3.2. Also without loss of generality, we assume that $q_1 \neq q_2$. Since the channel state describes by two states, the Markov chain associated with probability matrix Q_0 has 2 recurrent classes and transient states if $m_0 > 0$ or $m_2 > 0$. The stationary probability distribution of the

aggregated Markov chain A_1 is given by $\bar{\nu}_1 = \frac{b}{a+b}$, $\bar{\nu}_2 = \frac{a}{a+b}$, where $A_1 = \begin{pmatrix} -a & +a \\ b & -b \end{pmatrix}$

Now, we are going to illustrate how to apply this method and obtain a simple approximation for steady state solution. This approximation can be presented by first term in the Taylor series of ϵ in which the first term is given by $\pi^{(0)} = [\bar{\nu}_1, \bar{\nu}_2] \begin{bmatrix} m_{\xi_1} \\ m_{\xi_2} \end{bmatrix}$

In figures 4-5 we compare the dropping probability for various values of the buffer size and for $a = 1/5$, $b = 1/2$, $q_1 = 1$ and $q_2 = 0.2$, as obtained via the two methods presented in section 3 and section 4.

From plots in figures 4-5, we observe that the first term in the Taylor series gives a good approximation when ϵ is small ($\epsilon = 10^{-3}$). Moreover, we observe that when

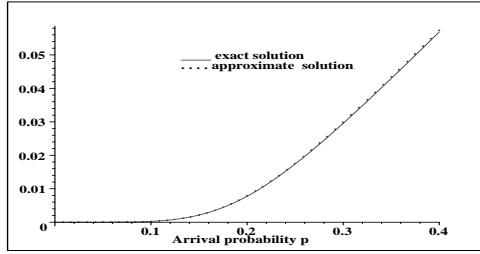


Fig. 4. Dropping probability versus arrival probability for $\epsilon = 10^{-3}$, $a = 1/5$, $b = 1/2$, $q_1 = 1$, $q_2 = 0$ and $B = 5$. Comparison of two methods: spectral analysis (exact solution) and approximate solution.

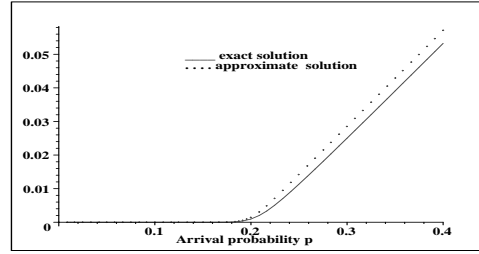


Fig. 5. Dropping probability versus arrival probability for $\epsilon = 10^{-3}$, $a = 1/5$, $b = 1/2$, $q_1 = 1$, $q_2 = 0.2$ and $B = 30$. Comparison of two methods: spectral analysis (exact solution) and approximate solution.

the buffer size increases the error between the exact solution and the approximation solution increases. An intuitive explanation is as follows: when we use the first term of the Taylor series, we get the limiting behavior as ϵ tends to zero, i.e. as the transitions of the queue sizes occur much faster than the transitions of the channel state. However, as the buffer size increases, the probability of reaching its boundaries become smaller (it converges to zero when the buffer size converges to infinity for the arrival probability that we chose) which means that the time between transitions to the state of full queue do not occur much faster than the transitions of the channel state.

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