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# On the development of a reliable integrated computational environment

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# Reliable Integrated Computational Environment



- compiler
- operating system
- hardware



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# Floating-point standardization

## IEEE 754-854

$$\mathbb{F}_0(\beta, t, L, U) = \{\pm d_0.d_1 \dots d_{t-1} \times \beta^e \mid 0 \leq d_i \leq \beta - 1, d_0 \neq 0, L \leq e \leq U\}$$

(base  $\beta$ , precision  $t$  and normalized exponent range  $[L, U]$ )

**Principle 1:** exact rounding

Rounding  $\bigcirc : \mathbb{R}_0 \rightarrow X$  satisfies

$$\begin{aligned} \bigcirc(x) &= x & \forall x \in X \\ x \leq y \Rightarrow \bigcirc(x) &\leq \bigcirc(y) & \forall x, y \in \mathbb{R}_0 \end{aligned}$$

Roundings: round to nearest, round up, round down, truncate

**Principle 2:** exactly rounded operations

$$x \circledast y = \bigcirc(x * y) \quad \forall x, y \in X$$

Operations:  $+, -, \times, /, \sqrt{\phantom{x}}$ , remainder, I/O and conversions between number sets  $X_i$ .

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## IEEE 754-854

**Principle 3:** closed number system

$$\begin{aligned}\overline{\mathbb{R}} &= \mathbb{R}_0 \cup S\mathbb{R}_0 \\ &= \mathbb{R}_0 \cup \{0, (\pm)\infty, \text{INV}\}\end{aligned}$$

$$\overline{X} \stackrel{?}{=} X \cup SX$$

- denormal numbers  
 $\pm d_0.d_1 \dots d_{t-1} \times \beta^L, d_0 = 0, \exists d_i \neq 0$   
→ no underflow exception in addition
- signed infinities:  $\pm\infty$   
→ affine arithmetic
- signed zeroes:  $\pm 0$

$$+0 \approx +\epsilon, -0 \approx -\delta, (+0) \oplus (-0) = ?$$

– IEEE-based arithmetic:

$$(+0) \oplus (-0) = +0$$

– TI-based arithmetic:  $(+0) \oplus (-0) = 0$

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## IEEE 754-854

- NaN: result of invalid operation

$$+\infty + (-\infty) = \text{NaN}$$

$$+\infty - (+\infty) = \text{NaN}$$

$$-\infty - (-\infty) = \text{NaN}$$

$$\frac{\pm\infty}{\pm\infty} = \text{NaN}$$

$$\pm 0 \times \pm\infty = \text{NaN}$$

$$\frac{\pm 0}{\pm 0} = \text{NaN}$$

$$x \text{REM } \pm 0 = \text{NaN}$$

$$\pm\infty \text{REM } y = \text{NaN}$$

$$\sqrt{x} = \text{NaN} \text{ when } x < 0$$

Conclusion:

$$\overline{\mathbb{F}} = \mathbb{F}_0 \cup \{\pm 0, (\pm)\infty, \text{denormals}, \text{NaN}\} = \mathbb{F}_0 \cup S\mathbb{F}_0$$

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## Strange output?

[Rump]:

$$a = 77617, b = 33096$$

$$\begin{aligned} z &= 333.75b^6 + a^2(11a^2b^2 - b^6 - 121b^4 - 2) \\ x &= 5.5b^8 \\ y &= z + x + \frac{a}{2b} \\ &\stackrel{?}{=} 5.76461\dots \times 10^{17} \\ &\stackrel{?}{=} 6.33825\dots \times 10^{29} \\ &\stackrel{?}{=} 1.1726\dots \\ &\stackrel{?}{=} -0.827396\dots \end{aligned}$$

---

## Strange output?

INTEL Pentium (Borland C++ compiler):

- (1) literals and variables: t=24

intermediate results: t=64

$$\tilde{y} = 5.76461\dots \times 10^{17}$$

- (2) literals and variables: t=24

intermediate results: t=24

$$\tilde{y} = 6.33825\dots \times 10^{29}$$

- (3) literals and variables: t=53

intermediate results: t=64

$$\tilde{y} = 5.76461\dots \times 10^{17}$$

- (4) literals and variables: t=53

intermediate results: t=53

$$\tilde{y} = 1.1726\dots$$

- (5) literals and variables: t=64

intermediate results: t=64

$$\tilde{y} = 5.76461\dots \times 10^{17}$$

---

## Strange output?

SUN Sparc (Sun f77, Sun cc):

(1) literals and variables: t=24

intermediate results: t=64

$$\tilde{y} = \text{unavailable}$$

(2) literals and variables: t=24

intermediate results: t=24

$$\tilde{y} = 6.33825\dots \times 10^{29}$$

(3) literals and variables: t=53

intermediate results: t=64

$$\tilde{y} = \text{unavailable}$$

(4) literals and variables: t=53

intermediate results: t=53

$$\tilde{y} = 1.1726\dots$$

(5) literals and variables: t=113

intermediate results: t=113

$$\tilde{y} = 1.1726\dots$$

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# Machine epsilon

```
#include <iostream.h>
int main()
{ double epsilonZero, epsilonOne, epsTmp;
  int i;

  epsilonZero = 1.0; i = 0;
  while (epsilonZero > 0.0)
  {
    i = i+1;
    epsTmp = epsilonZero;
    epsilonZero = epsilonZero/2.0;
  }
  cout << i << " " << epsTmp << "\n";

  epsilonOne = 1.0; i = 0;
  while (1.0+epsilonOne > 1.0)
  {
    i = i+1;
    epsTmp = epsilonOne;
    epsilonOne = epsilonOne/2.0;
  }
  cout << i << " " << epsTmp << "\n"; }

=1= CC -o macheps.x -fast macheps.cpp
=2= macheps.x
    1023      2.22507e-308
    53        2.22045e-16
=3= CC -o macheps.x macheps.cpp
=4= macheps.x
    1075      4.94066e-324
    53        2.22045e-16
```

---

## Counting to six

[Higham]:

$$2 - 1$$

$$\left( \frac{1}{\cos(100\pi + \pi/4)} \right)^2$$

$$3 \times \frac{\tan(\arctan(10000))}{10000}$$

$$\left( \left( \dots \left( \sqrt{ \sqrt{ \dots \sqrt{4}} }^2 \dots \right)^2 \right)^2 \right)^2$$

$$5 \times \frac{(1 + e^{-100}) - 1}{(1 + e^{-100}) - 1}$$

$$\frac{\ln(e^{6000})}{1000}$$

---

## Counting to six

one = 1.000000000000000

two = 2.0000000000001110

three = 2.9999999999971618

four = 2.7182818081824731

five = NaN

six = Infinity

---

# Exception handling

[Kahan]:

$$\forall x \in \mathbb{R}_0 \cup \{0, \pm\infty\} : x^0 = 1$$



$$\text{NaN}^0 = 1$$

[IEEE 754-854]: 

$$x_1 = +\infty$$

$$x_2 = -\infty$$

$$x = x_1 + x_2 \Rightarrow x = \text{NaN}$$

$$y = \sqrt{-1} \Rightarrow y = \text{NaN}$$

$$x^0 = \text{NaN}$$

$$y^0 = \text{NaN}$$

# Reliable Integrated Computational Environment



- libraries: numerics, graphics, statistics, ...
- Matlab, Octave, ...
- Maple, Mathematica, ...



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✓ • Floating-point standardization	5
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# The toolkit era

## Graphics

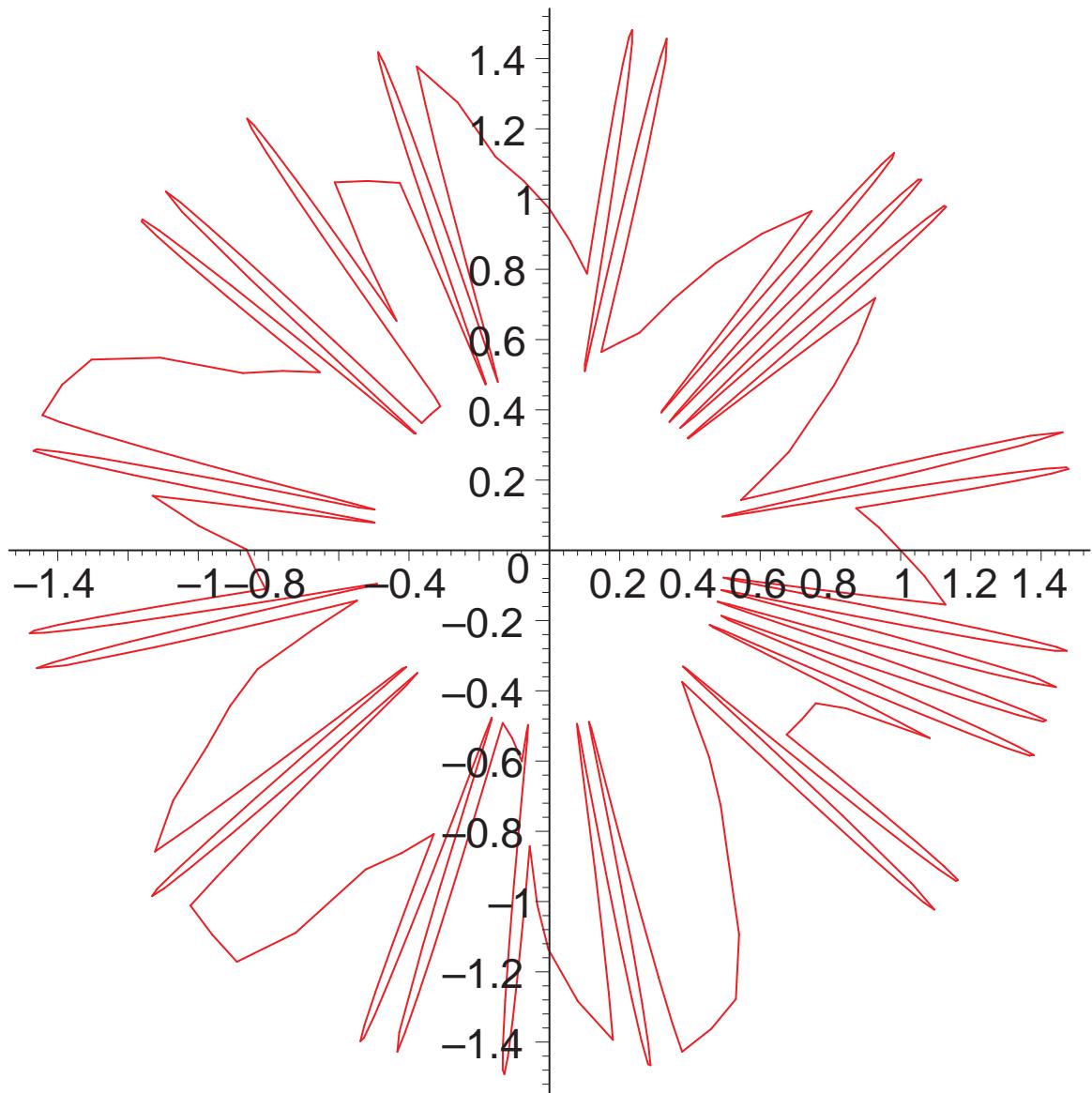
“daisy with 90 leaves” [Geulig & Krämer]:

$$\begin{aligned} r &= \frac{1}{2} \sin(90t) \\ x &= (1 + r) \cos(t) \\ y &= (1 + r) \sin(t) \quad t = 0 \dots 2\pi \end{aligned}$$

---

# Graphics

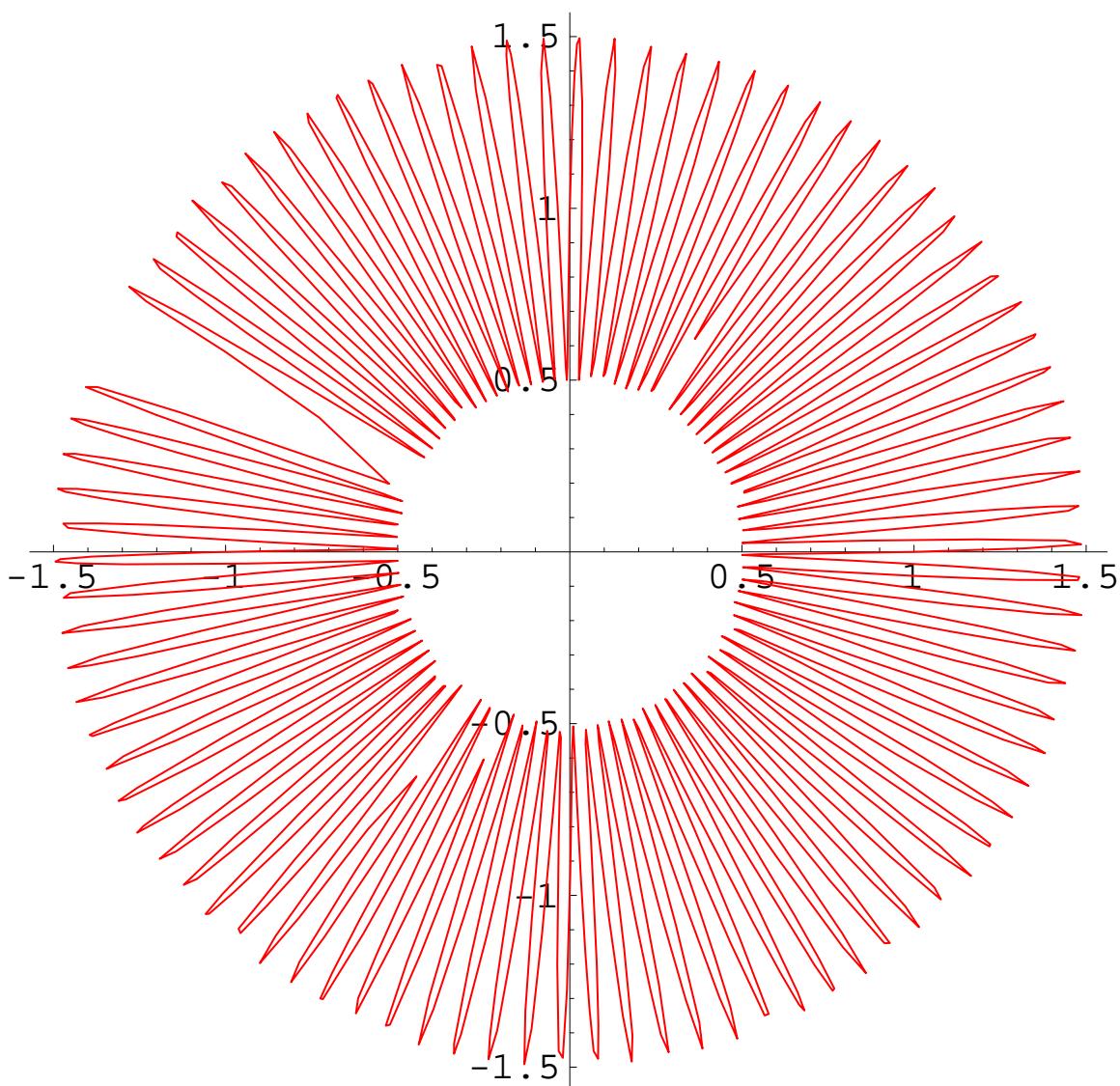
Maple 5.0:



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# Graphics

Mathematica 4.02:

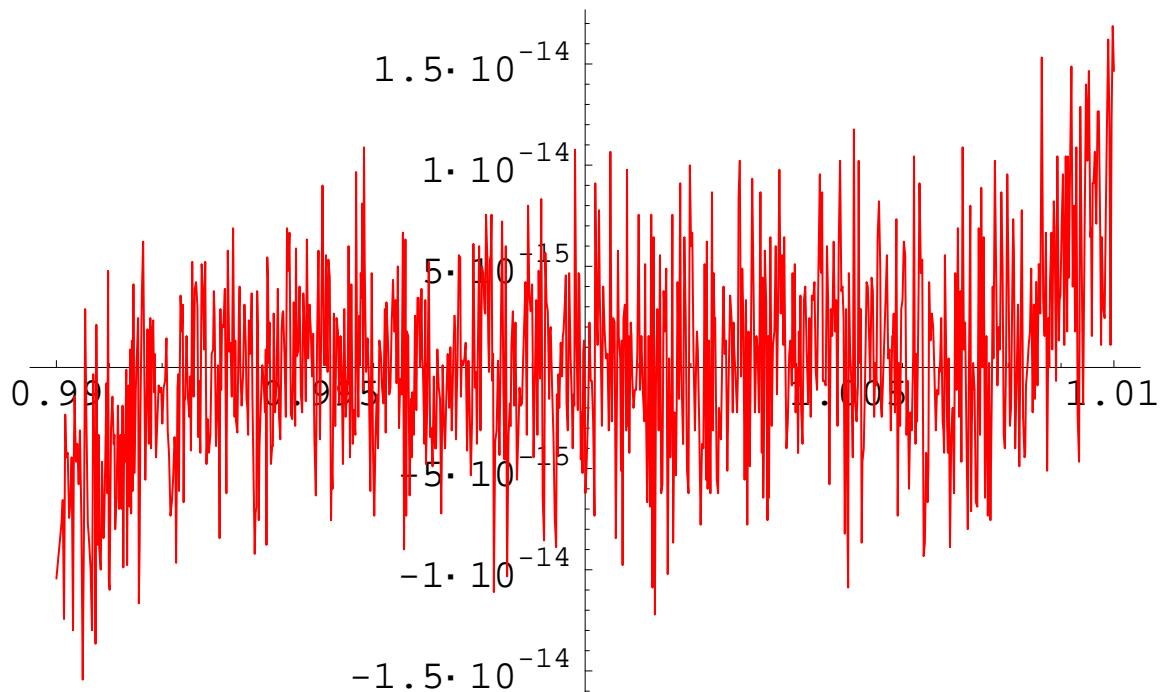


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# Graphics

$$f(x) = \sum_{i=0}^7 \binom{7}{i} (-1)^{7-i} x^i = (x-1)^7$$

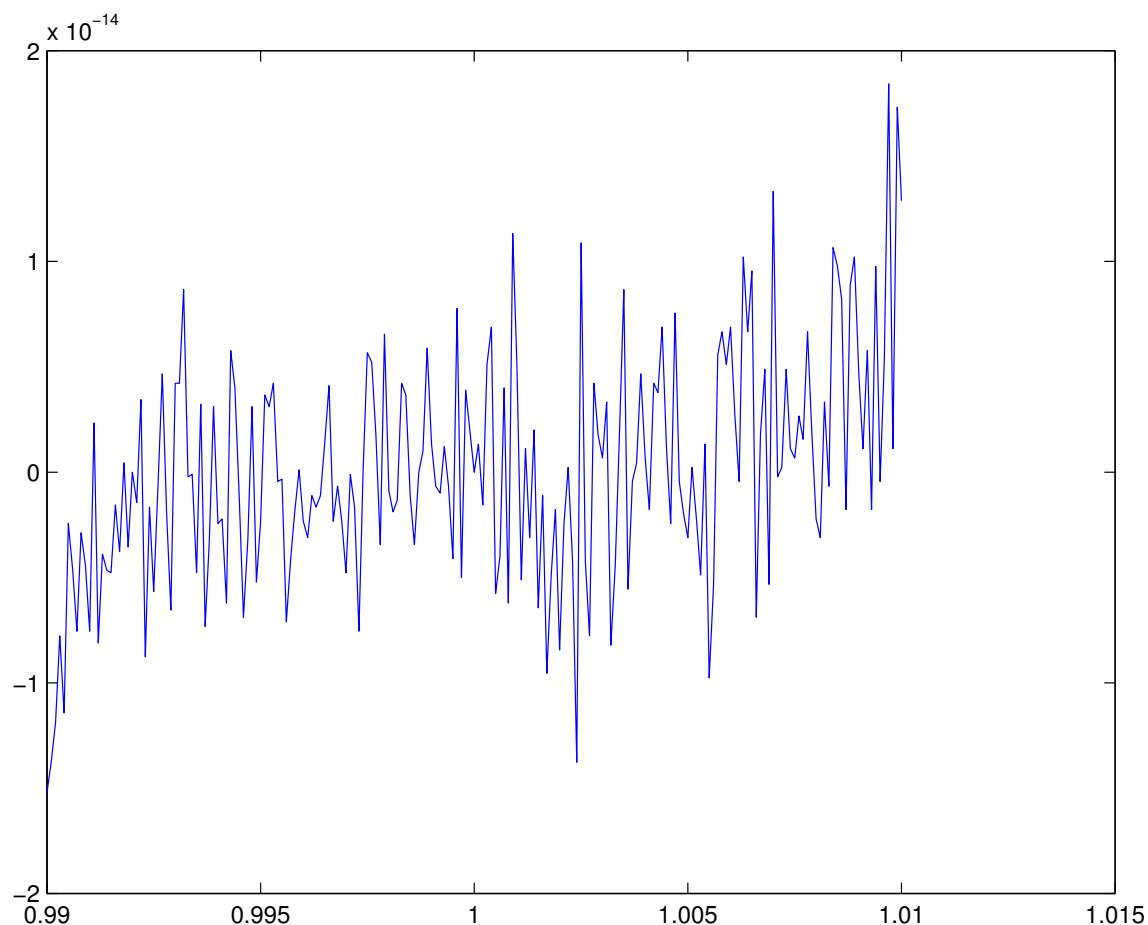
Mathematica 4.02:



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# Graphics

Matlab 5.03:



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# Tough sign problem

## Wilkinson polynomial

$$\begin{aligned} p_0(x) &= (x - 1) \dots (x - 20) \\ &= 2432902008176640000 - 8752948036761600000x + \\ &\quad 13803759753640704000x^2 - 12870931245150988800x^3 + \\ &\quad 8037811822645051776x^4 - 3599979517947607200x^5 + \\ &\quad 1206647803780373360x^6 - 311333643161390640x^7 + \\ &\quad 63030812099294896x^8 - 10142299865511450x^9 + \\ &\quad 1307535010540395x^{10} - 135585182899530x^{11} + \\ &\quad 11310276995381x^{12} - 756111184500x^{13} + \\ &\quad 40171771630x^{14} - 1672280820x^{15} + 53327946x^{16} \\ &\quad - 1256850x^{17} + 20615x^{18} - 210x^{19} + x^{20} \end{aligned}$$

Real roots:

$$z_i = i \quad i = 1, \dots, 20$$

## Modified Wilkinson polynomial

$$p_1(x) = p_0(x) - 2^{-23}x^{19}$$

Real roots:

$$\begin{aligned} z_1 &= 1.0000000 \dots & z_2 &= 2.0000000 \dots \\ z_3 &= 3.0000000 \dots & z_4 &= 4.0000000 \dots \\ z_5 &= 4.9999993 \dots & z_6 &= 6.0000069 \dots \\ z_7 &= 6.9996972 \dots & z_8 &= 8.0072676 \dots \\ z_9 &= 8.9172502 \dots & z_{10} &= 20.846908 \dots \end{aligned}$$

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## Tough sign problem

Aim: Obtain reliable bounds for real roots of

$$p_1(x) = 0$$

False Position Method

Given:

$$(x_\ell, f_\ell), (x_h, f_h) \quad f_\ell < 0, f_h > 0$$

Compute:

$$x_n = x_h + f_h \frac{x_h - x_\ell}{f_\ell - f_h}$$

$$f_n = f(x_n)$$

if  $f_n < 0$  then  $x_l = x_n$

else  $x_h = x_n$

Mathematically:

$$z_i \in [x_l, x_h]$$

---

## Tough sign problem

Implementation: IEEE arithmetic

- data error in representation of coefficients of  $p_1(x)$
- rounding error in evaluation of  $p_1(x)$

Implementation yields floating-point bounds

$$[\tilde{x}_l, \tilde{x}_h]_{t=53}$$

$$z_i \in [\tilde{x}_l, \tilde{x}_h]?$$

No guarantee!

$$z_9 = 8.917250249\dots$$

$$z_9 \notin [\tilde{x}_i, \tilde{x}_j]_{53} = [8.9172245567595, 8.9172249342899]$$

$$\tilde{p}_1(\tilde{x}_i)_{t=53} = +5.369856e7$$

$$\tilde{p}_1(\tilde{x}_j)_{t=53} = -2.626468e8$$

$$p_1(\tilde{x}_j) = +4.87759e7$$

$$\tilde{p}_1(\tilde{x}_j)_{t=53} \in [-1.177340e9, +1.324868e9]$$

# Reliable Integrated Computational Environment

- IEEE compliant multiprecision floats
- rational arithmetic
- complex arithmetic
- sharp interval arithmetic
- reliable graphics
- calculator, parser, precompiler, GUI



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Strange output unraveled	
Reliable graphics	
Advanced exception handling	
Reliable false position	

# Reliability in view

## Strange output unraveled

[Cuyt&Verdonk]:

- multiprecision float ( $t = 122$ ) :

```

z      =      -7.917111340668961361101134701524942850e36
x      =      +7.917111340668961361101134701524942848e36
z + x =      -2.00000000000000000000000000000000000000000000000000000
y      =      -8.27396059946821368141165095479816292e-1

```

- single precision interval ( $t = 24$ ) :

$$\begin{aligned} z &= [-7.917116e36; -7.917105e36] \\ x &= [7.917108e36; 7.917112e36] \\ z + x &= [-7.605904e30; 6.338254e30] \\ y &= [-7.605904e30; 6.338254e30] \end{aligned}$$

- double precision interval ( $t = 53$ ) :

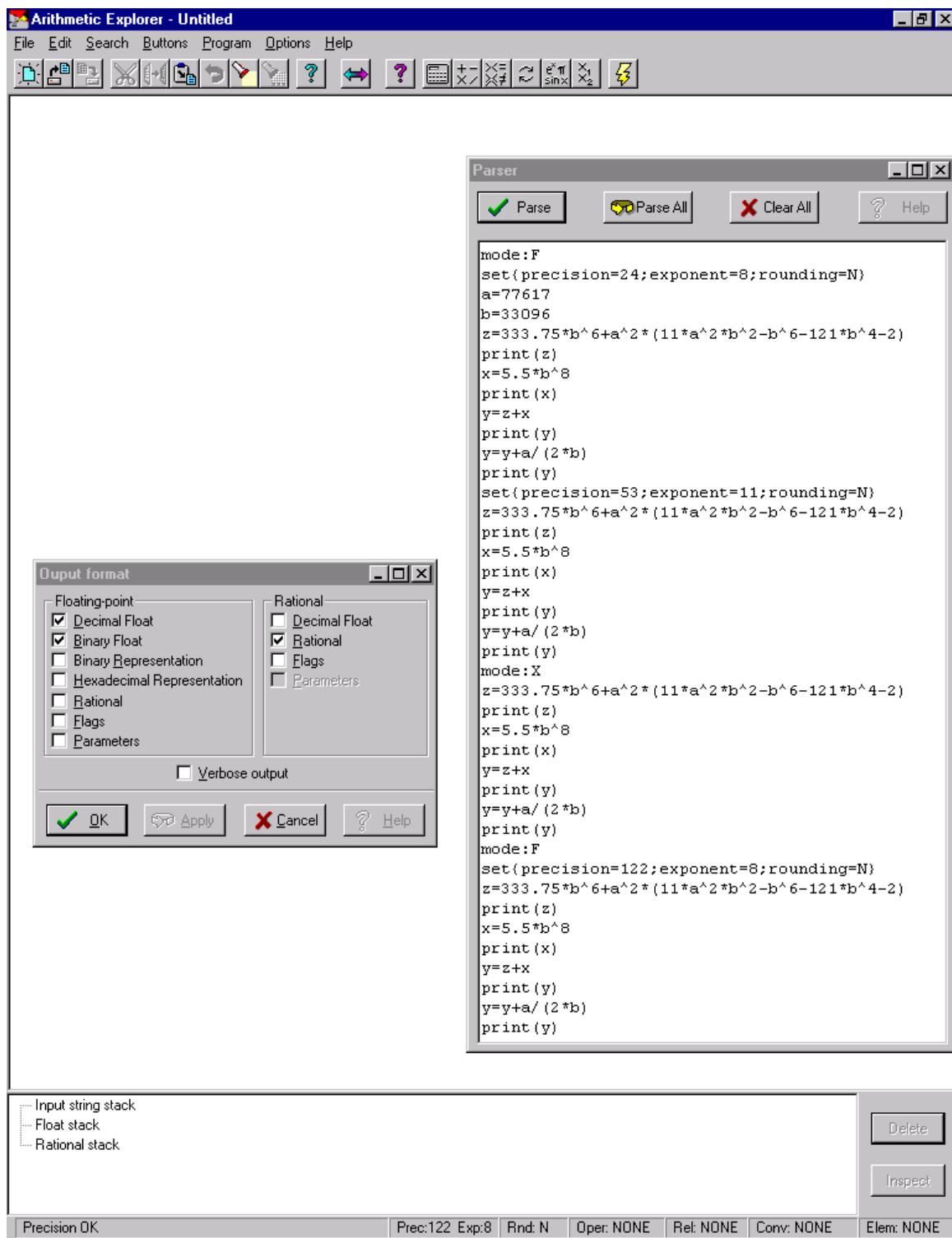
$$\begin{aligned} z &= [-7.91711134066897e36; -7.91711134066895e36] \\ x &= [7.91711134066895e36; 7.91711134066896e36] \\ z + x &= [-8.26414134502188e21; 7.08354972430447e21] \\ y &= [-8.26414134502188e21; 7.08354972430447e21] \end{aligned}$$

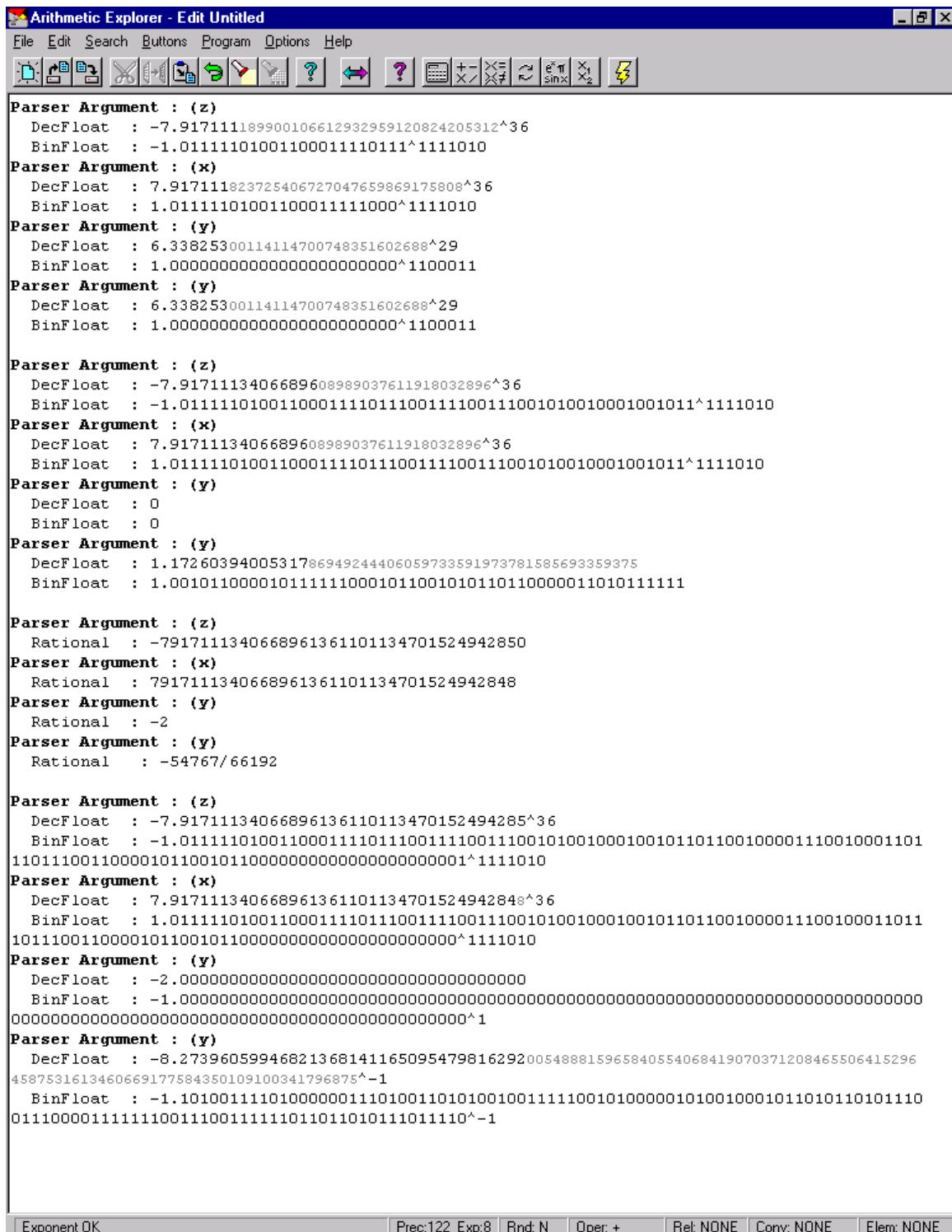
- multiprecision interval ( $t = 122$ ) :

$$z = [-7.9171113406689613611011347015249428 \frac{4}{5} e36]$$

$$x = [7.9171113406689613611011347015249428 \frac{5}{4} e36]$$

$$y = [-8.2739605994682136814116509547981629 \begin{matrix} 0 \\ 1 \end{matrix} e - 1]$$

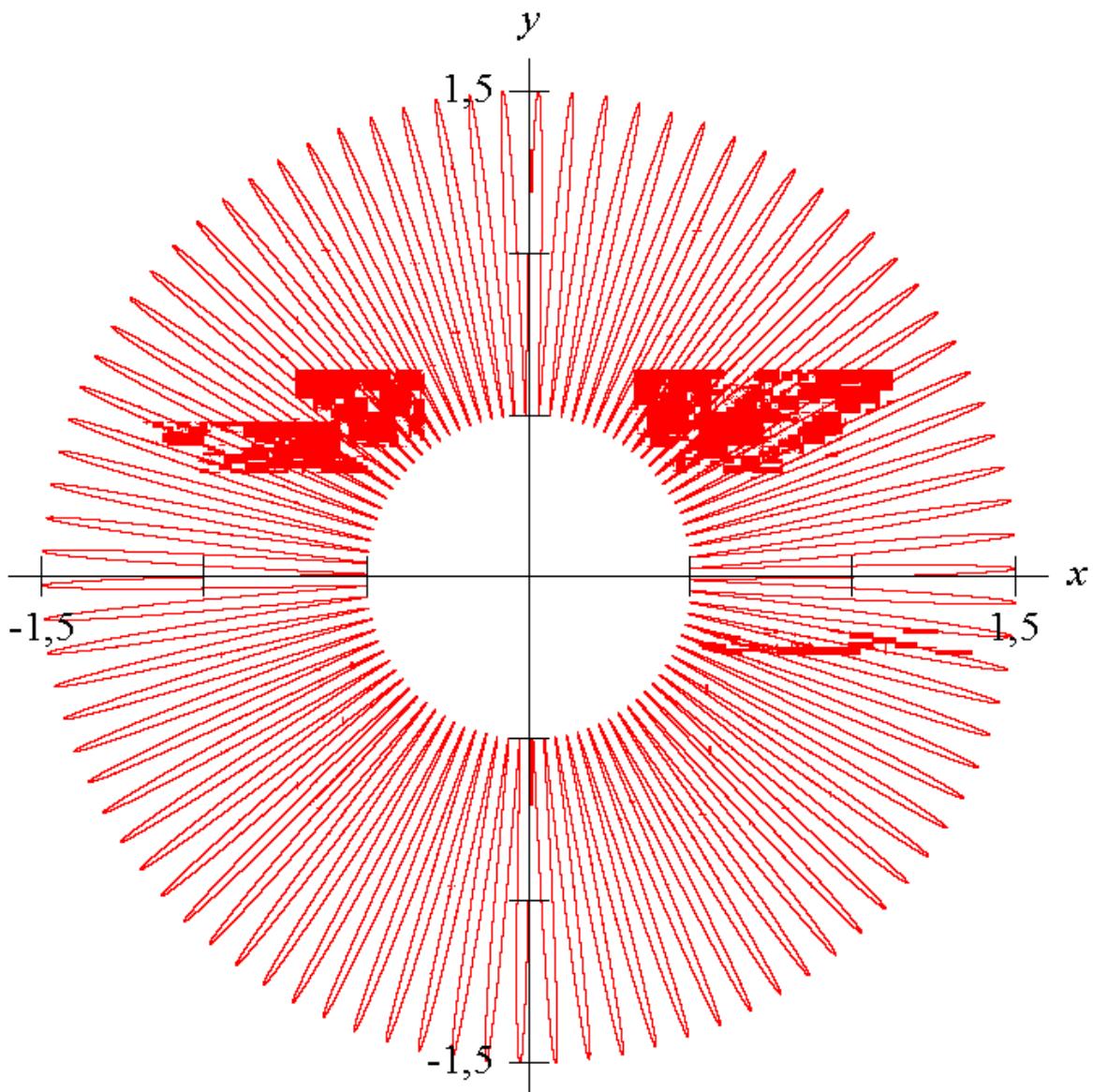




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# Reliable graphics

[Tupper]: generalized interval arithmetic



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## Advanced exception handling

[ARITHMOS]:

$$x_1 = +\infty$$

$$x_2 = -\infty$$

$$x = x_1 + x_2 \Rightarrow x = \text{NaN}$$

$$y = \sqrt{-1} \Rightarrow y = \text{IaC}_0$$

IaY: result can only be represented in  $Y \supset X$

$\rightarrow \text{IaC}, \text{IaC}_0, \text{IaC} \sim \text{Inf}, \text{IaC}_0 \sim \text{Inf}$

$$\forall x \in Y \cup \{0, \pm\infty\} \quad x^0 = 1$$

$$\Downarrow$$

$$(\text{IaY})^0 = 1$$

$$x^0 = \text{NaN}$$

$$y^0 = 1$$

---

# Reliable false position

```
#include "MpIeee.h"
#include "Rational.h"
....
Rational coef[NR_OF_COEF];

MpIeee f(MpIeee xx)
{ Rational result, x(xx);

  result=coef[DEGREE];
  for(int i=DEGREE-1;i>=0;i--) {
    result=x*result+coef[i];
  }
  return result.toMpIeee(53,11);
}

void regfalsi(MpIeee x1,MpIeee x2,
              MpIeee (*f)(MpIeee x),MpIeee eps)
{ MpIeee xh, xl, fl, fh, xnew, fxnew;

  ...
  for (int count=0;count<MAXIT;count++) {
    xnew=xh+fh*(xh-xl)/(fl-fh);
    fxnew=f(xnew);
    if (fxnew<0) {
      xl=xnew;
      fl=fxnew;
    }
    else {
      xh=xnew;
      fh=fxnew;
    }
    ...
  }
}
```

---

# Reliable false position

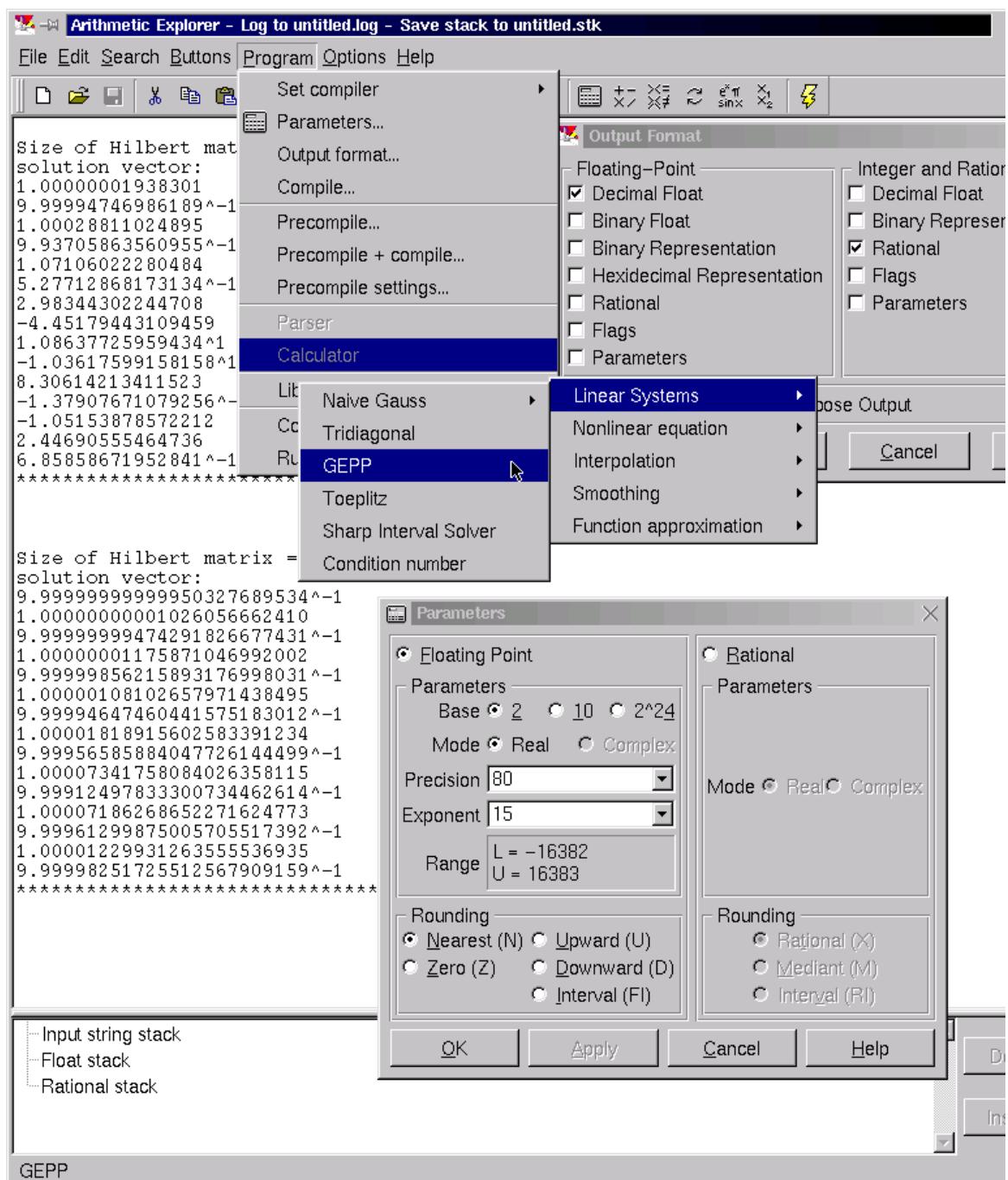
[ARITHMOS]:

```
main()
{ MpIeee (*fptr)(MpIeee);
MpIeee startl, starth, eps;

coef[0]=Rational("2432902008176640000.0");
coef[1]=Rational("-8752948036761600000.0");
coef[2]=Rational("13803759753640704000.0");
coef[3]=Rational("-12870931245150988800.0");
coef[4]=Rational("8037811822645051776.0");

...
coef[19]=Rational("-210.0")-pow(2,-23);
coef[20]=Rational("1.0");
...

regfalsi(startl,starth,fptr,eps);
}
```



**Arithmetic Explorer - Edit F:\ArExp\rnc2000\linsys.cc**

The screenshot shows a window titled "Arithmetic Explorer - Edit F:\ArExp\rnc2000\linsys.cc". The menu bar includes File, Edit, Search, Buttons, Program, Options, and Help. A toolbar with various icons is visible above the code area. A context menu is open over the code, showing options like Set Compiler, Parameters..., Output format..., Compile..., Calculator..., Parser..., Numerical Library..., Command line arguments..., and Run... The main code area contains three functions: rational\_solveTri, rational\_forward\_elim, and rational\_back\_subst.

```

for (i = matrix_size - 1; i >= 0; i--) {
    sum = B[i];
    for (j = i+1; j < matrix_size; j++)
        sum -= A[i][j] * X[j];
    X[i] = sum / A[i][i];
}

void rational_solveTri(int n, rational *a, rational *d, rational *c,
                      rational *b, rational *x) {
    int i;

    rational xmult;

    for (i = 1; i < n; i++) {
        xmult = a[i-1] / d[i-1];
        d[i] = d[i] - xmult * c[i-1];
        b[i] -= xmult * b[i-1];
    }
    x[n-1] = b[n-1] / d[n-1];
    for (i = n-2; i >= 0; i--)
        x[i] = (b[i] - c[i] * x[i+1]) / d[i];
}

void rational_forward_elim (int matrix_size, rational **A, rational *B,
                           rational *X) {
    int i, j, k;
    rational xmult;

    for (k = 1; k < matrix_size; k++) {
        for (i = k + 1; i <= matrix_size; i++) {
            xmult = A[i][k] / A[k][k];
            A[i][k] = xmult;
            for (j = k + 1; j <= matrix_size; j++)
                A[i][j] -= xmult * A[k][j];
            B[i] -= xmult * B[k];
        }
    }
    X[matrix_size] = B[matrix_size] / A[matrix_size][matrix_size];
}

void rational_back_subst (int matrix_size, rational **A, rational *B,
                         rational *X) {
    int i, j;
    rational sum;

    for (i = matrix_size-1; i >= 1; i--) {
        sum = B[i];
        for (j = i+1; j <= matrix_size; j++)
            sum -= A[i][j] * X[j];
        X[i] = sum / A[i][i];
    }
}

```

Specify the 'cout' options for compiled C++ programs

---

## References

- [1] A. Cuyt and B. Verdonk. A remarkable example of catastrophic cancellation unraveled. Submitted for publication, <ftp://wins.uia.ac.be/pub/CANT/Arithmos/catacanc.pdf>, 1999.
- [2] I. Geulig and W. Krämer. Intervallrechnung in Maple - die erweiterung `intpakx` zum paket `intpak` der Share-library. Preprint nr. 99/2, Universität Karlsruhe, 1999.
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- [6] J.A. Tupper. Graphing equations with generalized interval arithmetic. thesis, Toronto, 1996.