

Sémantique de JUNIOR

$$\text{exec}(t, E) = (t', E', b)$$

t : instruction de départ
E : environnement de départ
t' : instruction d'arrivée
E' : environnement d'arrivée
b : status de terminaison

Environnement :

ensemble d'événements présents + 2 booléens : **eo**i et **move**

Status de terminaison :

TERM, STOP ou **SUSP**

Séquence :

$\text{exec}(\text{Seq}(t, u), E) =$
soit $\text{exec}(t, E) = (t', E', b)$;
si $b \neq \text{TERM}$ alors $(\text{Seq}(t', u), E', b)$ sinon $\text{exec}(u, E')$

Sémantique - 2

Stop et nothing :

$\text{exec}(\text{Nothing}, E) = (\text{Nothing}, E, \text{TERM})$

$\text{exec}(\text{Stop}, E) = (\text{Nothing}, E, \text{STOP})$

Loop :

$\text{exec}(\text{Loop}(t), E) = \text{exec}(\text{Seq}(t, \text{Loop}(t)), E)$

Generate :

$\text{exec}(\text{Generate}(S), E) = (\text{Nothing}, E + \{S\} + [\text{move} = \text{vrai}], \text{TERM})$

Présence :

$\text{exec}(\text{When}(S, t, u), E) =$
si S dans E alors $\text{exec}(t, E)$,
sinon si eoi alors (u, E, STOP) ,
sinon $(\text{When}(S, t, u), E, \text{SUSP})$

Sémantique - 3

Await :

$\text{exec}(\text{Await}(S), E) =$
si S dans E alors $(\text{Nothing}, E, \text{TERM})$,
sinon si eoi alors $(\text{Await}(S), E, \text{STOP})$,
sinon $(\text{Await}(S), E, \text{SUSP})$

Préemption :

$\text{exec}(\text{Until}(S, t), E) =$
si $\text{exec}(t, E) = (t', E', b)$ et $b = \text{STOP}$ alors $\text{exec}(\text{Waiting}(S, t'), E')$
sinon $(\text{Until}(S, t'), E', b)$

$\text{exec}(\text{Waiting}(S, t), E) =$
si S dans E alors $(\text{Nothing}, E, \text{TERM})$,
sinon si eoi alors $(\text{Until}(S, t), E, \text{STOP})$,
sinon $(\text{Waiting}(S, t), E, \text{SUSP})$

Sémantique - 4

Parallélisme :

$$\text{exec}(\text{Par}(t, u), E) = \text{exec}(\text{Par}(t, u, \text{SUSP}, \text{SUSP}), E)$$

$$\begin{aligned} \text{exec}(\text{Par}(t, u, b, \text{SUSP}), E) = \\ (\text{aux}(t, u', b, c), E', \text{stat}(b, c)) \\ \text{où } \text{exec}(u, E) = (u', E', c) \end{aligned}$$

stat	TERM	STOP	SUSP
TERM	TERM	STOP	SUSP
STOP	STOP	STOP	SUSP
SUSP	SUSP	SUSP	SUSP

$$\begin{aligned} \text{exec}(\text{Par}(t, u, \text{SUSP}, c), E) = \\ (\text{aux}(t', u, b, c), E', \text{stat}(b, c)) \\ \text{où } \text{exec}(t, E) = (t', E', b) \end{aligned}$$

$$\begin{aligned} \text{exec}(\text{Par}(t, u, \text{SUSP}, \text{SUSP}), E) = (\text{aux}(t', u', b, c), E'', \text{stat}(b, c)) \\ \text{où } \text{exec}(t, E) = (t', E', b) \text{ et } \text{exec}(u, E') = (u', E'', c) \end{aligned}$$

$$\begin{aligned} \text{aux}(t, u, b, c) = \text{Par}(t, u) \text{ si } b \text{ et } c \neq \text{SUSP}, \\ \text{Par}(t, u, b, c) \quad \text{sinon} \end{aligned}$$

Sémantique - 5

Événements :

$\text{exec}(\text{Event}(S, t), E) = \text{exec}(\text{AbsentEvent}(S, t), E)$

E'' est égal à E' , mais $E''(S) = E(S)$

$\text{exec}(\text{AbsentEvent}(S, t), E) =$
soit $\text{exec}(t, E - \{S\}) = (t', E', b)$;
si $b \neq \text{SUSP}$ alors $(\text{Event}(S, t'), E'', b)$
sinon si S dans E' alors $(\text{PresentEvent}(S, t'), E'', \text{SUSP})$
sinon $(\text{AbsentEvent}(S, t'), E'', \text{SUSP})$

$\text{exec}(\text{PresentEvent}(S, t), E) =$
soit $\text{exec}(t, E + \{S\}) = (t', E', b)$;
si $b \neq \text{SUSP}$ alors $(\text{Event}(S, t'), E'', b)$
sinon $(\text{PresentEvent}(S, t'), E'', \text{SUSP})$

Sémantique - 6

Instants :

`exec(instant(t),E) =`
 soit `exec(t,E[move = faux]) = (instant(t'),E',b) ;`
 si `b ≠ SUSP` alors `(instant(t'),E',b)`
 sinon si `move = vrai` alors `exec(instant(t'),E')`
 sinon `exec(instant(t'),E'[eoi = vrai])`

Implémentations

- **REWRITE = exec**
- **REPLACE = REWRITE + optimisation des créations de structures intermédiaires**
- **SIMPLE = implémentation adaptée à un grand nombre d'événements**
- ...

REWRITE

$\text{exec}(t,E) = (t',E',b)$

Instruction
MicroState s = t.rewrite(E);
b = s.flag
t' == s.term
E' est la nouvelle valeur de E

TERM, STOP, SUSP

```
abstract public class Instruction implements Flags, Program
{
    abstract public MicroState rewrite(EnvironmentImpl env);
    public MicroState unknown(){ return new MicroState(SUSP,this); }
}
```

```
public class Stop extends Instruction
{
    public MicroState rewrite(EnvironmentImpl env){
        return new MicroState(STOP,new Nothing());
    }
}
```

$\text{exec}(\text{Stop},E) =$
 $(\text{Nothing} ,E,\text{STOP})$

Séquence

```
public abstract class BinaryInstruction extends Instruction
{
    final public Instruction left,right;

    public BinaryInstruction(Program left,Program right){
        this.left = (Instruction)left; this.right = (Instruction)right;
    }
}
```

```
public class Seq extends BinaryInstruction
{
    public Seq(Program left,Program right){ super(left, right); }

    public MicroState rewrite(EnvironmentImpl env){
        MicroState s = left.rewrite(env);
        if (TERM == s.flag) return right.rewrite(env);
        return new MicroState(s.flag,new Seq(s.term,right));
    }
}
```

exec(Seq(t,u) ,E) =
soit exec(t,E) = (t' ,E',b) ;
si b ≠ TERM alors (Seq(t' ,u) ,E',b)
sinon exec(u,E')

création d'une nouvelle instruction



Événements

```
abstract public class Config implements Configuration
{
    abstract public boolean fixed(EnvironmentImpl env);
    abstract public boolean eval(EnvironmentImpl env);
    abstract public void reset();
    public boolean sat(EnvironmentImpl env){ return fixed(env) && eval(env); }
    public boolean unsat(EnvironmentImpl env){ return fixed(env) && !eval(env); }
}

public class Presence extends Config
{
    final public IdentifierWrapper wrapper;
    public boolean evaluated = false;
    public Identifier event;

    public Presence(IdentifierWrapper wrapper){ this.wrapper = wrapper; }
    public void reset(){ evaluated = false; }

    public boolean fixed(EnvironmentImpl env){
        if(evaluated == false){ event = wrapper.evaluate(env); evaluated = true; }
        return env.isGenerated(event) || env.eoi;
    }

    public boolean eval(EnvironmentImpl env){ return env.isGenerated(event); }
}

public class Await extends Instruction
{
    final public Config config;

    public Await(Configuration config){ this.config = (Config)config; }

    public MicroState rewrite(EnvironmentImpl env){
        if (config.sat(env)) return new MicroState((env.eoi ? STOP : TERM),new Nothing());
        if (config.unsat(env)) return new MicroState(STOP,this);
        return unknown();
    }
}
```

Parallélisme

Par implémenté par la classe Merge

```
public class Merge extends BinaryInstruction
{
    final public byte leftFlag, rightFlag;

    public Merge(Program left, Program right, byte leftFlag, byte rightFlag){
        super(left,right); this.leftFlag = leftFlag; this.rightFlag = rightFlag;
    }
    public Merge(Program left,Program right){ this(left,right,SUSP,SUSP); }

    public Instruction newTerm(Instruction l, Instruction r, byte lf, byte rf){
        return new Merge(l,r,lf,rf);
    }

    public MicroState result(Instruction l, Instruction r, byte lf, byte rf){
        byte b = SUSP, nlf = lf, nrf = rf;
        if(lf != SUSP && rf != SUSP){
            b = (lf==TERM && rf==TERM) ? TERM : STOP;
            if (lf==STOP) nlf = SUSP;
            if (rf==STOP) nrf = SUSP;
        }
        return new MicroState(b,newTerm(l,r,nlf,nrf));
    }
    public MicroState rewrite(EnvironmentImpl env){
        if (leftFlag == SUSP && rightFlag != SUSP){
            MicroState s = left.rewrite(env);
            return result(s.term,right,s.flag,rightFlag);
        }
        if (leftFlag != SUSP && rightFlag == SUSP){
            MicroState s = right.rewrite(env);
            return result(left,s.term,leftFlag,s.flag);
        }
        MicroState ls = left.rewrite(env), rs = right.rewrite(env);
        return result(ls.term,rs.term,ls.flag,rs.flag);
    }
}
```

stat	TERM	STOP	SUSP
TERM	TERM	STOP	SUSP
STOP	STOP	STOP	SUSP
SUSP	SUSP	SUSP	SUSP

Instants

```
public class Instant extends UnaryInstruction
{
    public Instant(Program body) { super(body); }

    public MicroState rewrite(EnvironmentImpl env) {
        MicroState s = body.rewrite(env);
        if (s.flag != SUSP) return new MicroState(s.flag, new Instant(s.term));
        if (env.move) env.move = false; else env.eoi = true;
        return new Instant(s.term).rewrite(env);
    }
}
```

exec(instant(t),E) =
soit **exec(t,E[move = faux]) = (instant(t'),E',b) ;**
si **b ≠ SUSP** alors **(instant(t') ,E',b)**
sinon si **move = vrai** alors **exec(instant(t'),E')**
sinon **exec(instant(t') ,E'[eoi = vrai])**

Machines non synchronisées

```
public class BasicContext implements Flags, UnSyncMachine
{
    public Instant instant;
    public EnvironmentImpl env;
    public Program toAdd = Jr.Nothing();
    public boolean somethingToAdd = false;

    public BasicContext(Program program) {
        this.instant = new Instant(program);
        buildEnvironment();
    }

    public void buildEnvironment() { env = new EnvironmentImpl(); }

    public void add(Program inst) {
        toAdd = Jr.Par(toAdd,inst);
        somethingToAdd = true;
    }

    protected void performAddings() {
        if (somethingToAdd == false) return;
        instant = new Instant(Jr.Par(toAdd,instant.body));
        toAdd = Jr.Nothing();
        somethingToAdd = false;
    }

    public void generate(Identifier event, Object obj) { env.generate(event,obj); }

    public Program getFrozen (Identifier event) { return env.getFrozen(event); }

    public boolean react() {
        performAddings();
        MicroState s = instant.rewrite(env);
        instant = (Instant)s.term;
        env.newInstant();
        return (TERM == s.flag);
    }
}
```

Machines synchronisées

```
public class ExecContext extends BasicContext implements Machine, SyncMachine
{
    public Vector toGenerate = new Vector();

    public ExecContext(Program program) { super(program); }

    public void add(Program inst) {
        synchronized(toAdd) { super.add(inst); }
    }
    protected void performAddings() {
        synchronized(toAdd) { super.performAddings(); }
    }

    public void generate(Identifier event, Object obj) {
        synchronized(toGenerate) {
            toGenerate.addElement(new GenerateOrder(event,obj));
        }
    }

    protected void performGenerations() {
        synchronized(toGenerate) {
            if (toGenerate.size() > 0) {
                Enumeration list = toGenerate.elements();
                while (list.hasMoreElements()) {
                    GenerateOrder order = (GenerateOrder)list.nextElement();
                    super.generate(order.identifier,order.value);
                }
                toGenerate.removeAllElements();
            }
        }
    }

    public synchronized boolean react() {
        performGenerations();
        return super.react();
    }
}
```