# A Model for Astronomical Cross-Matching Disambiguation

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HOSCAR Meeting Bordeaux - France





# Agenda



#### • Introduction

- Fundamentals
  - Indexing structures
  - Algorithms
  - Experiments and evaluation
- Motivation and goal
- A disambiguation model proposal
- Conclusion

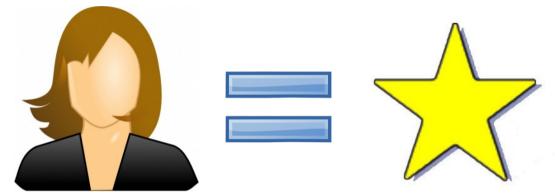




#### What is a catalog?









#### **Spectroscopic Survey** :









### Different Astronomical Surveys (Catalogs)









# Introduction



- Surveys produce catalogs with intersections in the covered area of the sky;
- Problem:
  - Getting an integrated view provided by different catalogs requires data cross-matching
  - How to identify celestial objects that appear in different catalogs with descriptive variations?





### Introduction



- Problem identified as "Entity Resolution"
  - Identify instances of objects from different databases that match the same real world entity
- Alternatives for entity resolution in the "crossmatching catalogs" problem:
  - use the position of the objects in the sky (coordinate system based on RA, DEC);
  - use other attributes to help treating the ambiguities.



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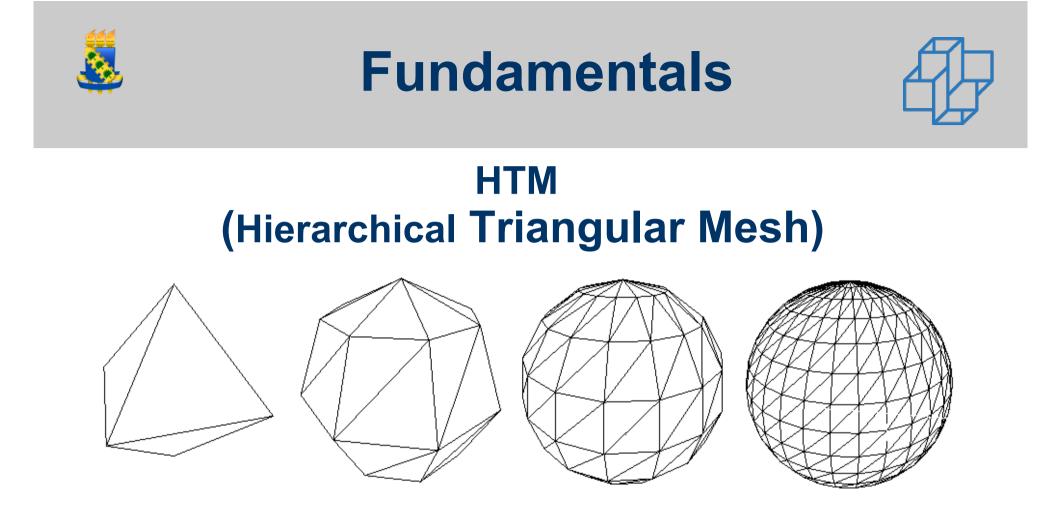
### **Current solutions**

- Support cross-matching by location
  - Main Strategy: represent the sky making the use of data structures in order to facilitate the localization of the stars and their neighbors in space
  - use the spatial indexing structures to support cross-matching:
    - HTM (Hierarchical Triangular Mesh)
    - Q3C (Quad Tree Cube)
    - Zones



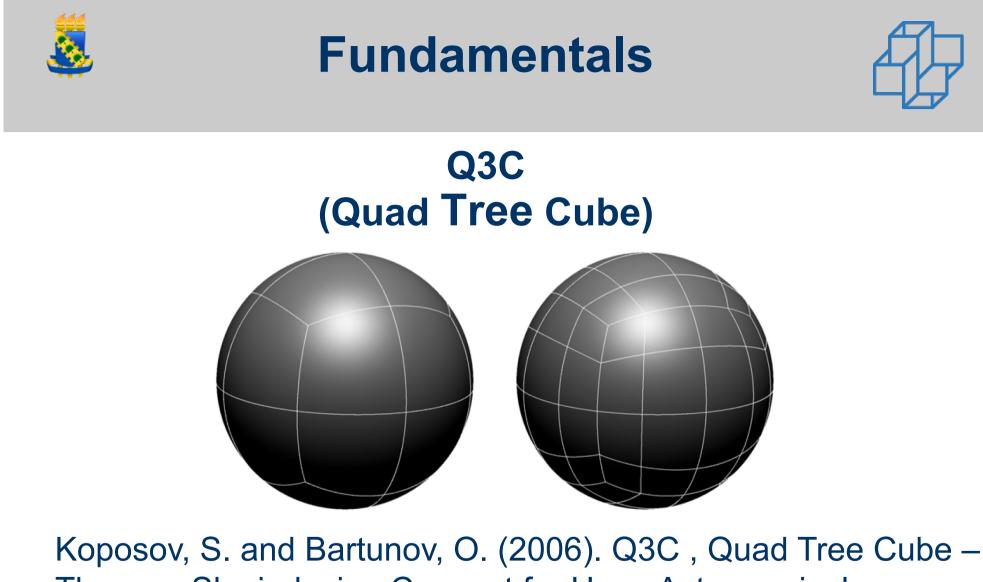
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Kunszt, P. Z., Szalay, A. S., and Thakar, A. R. (2001). The hierarchical triangular mesh.



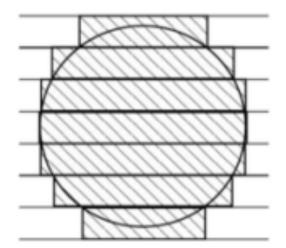


The new Sky-indexing Concept for Huge Astronomical Catalogues and its Realization for Main Astronomical









[Gray, J., Szalay, A. S., Thakar, A. R., and et al. (2004). There goes the neighborhood: Relational algebra for spatial data search.]

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**Binary cross-matching of catalogs** 

Support the cross-matching by location

Algorithms

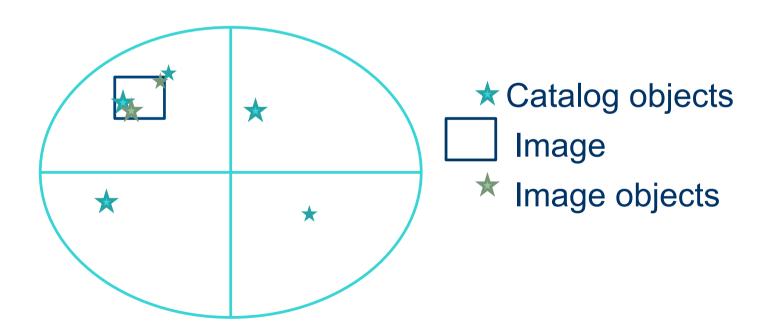
- Fast Approximate matching
- Q3c Join





#### **Algorithms**

[Fu et al. 2012] Fast approximate matching of astronomical objects.



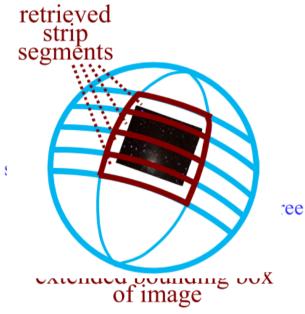






#### **Matching process**

- 1. Index by strips
- 2. Given an image, retrieve the catalog objects that are similar to the image objects;
- 3. A matching catalog object must be within 1 arcsec from the image object
  Extend the image bounding box of 1 arcsec in order to find all objects.



4. Retrieve in memory all the catalog objects contained in the area of the extended image

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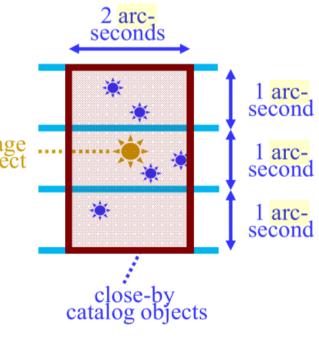






Matching process (cont.)

- 5. Recovered strips are divided into thinner sub-strips (1 arc-second)
  - For each image object:
    - Retrieve the catalog objects that are at 1 arcsec distant, considering only their own substrips, as well as the ones in their twobject adjacent substrips.
    - The condition to achieve the match is:
      For each image object p, it is assumed that its nearest catalog object is q, and the q nearest object is also p. Then q is the match for p.









Some considerations about the Fast approximate cross-matching

- Positive points
  - Indexing relatively fast
    - Index a catalog of 2 billion objects in less than two hours
  - Fast matching
    - Match a catalog of 2 billion objects and an image of 100,000 objects in 4 seconds.
- Negative point
  - Matching using more than 2 catalogs does not generate symmetric results







#### **Cross-Matching using Q3C**

Implementation in PostgreSQL (q3c\_join())

Select \*

from table1, table2

where q3c\_join(table1.ra, table1.dec, table2.ra, table2.dec, 0.001);

- Supposing there is a Q3C index over table2
  - function q3c\_join():
    - Defines 4 range queries to approximate the crossmatch circles:
      - If the object of table1 is within these ranges, then the matching is achieved.







#### **Experiments using Q3C**

- Goal: To evaluate the quality of binary crossmatching based on a spatial criterion
- Test environment:
  - Catalogs Involved:
    - 2MASS (470,992.70 objects)
    - BCC v.05 (1,376,582,713 objects)
  - Radius 0.001 degree
- Result
  - Matching of 17,701,306 objects
  - Processing: 142 seconds







#### **Evaluating the matching**

- Difficulty No prior knowledge
- Test environment :
  - Catalogs Involved:
    - 2MASS (470,992,970 objects)
    - 2MASS (470,992,970 objects)
  - Radius 0.001 degree
- Result:
  - The obvious would return 470,992,970 elements
  - Returned 483,197,616 objects
  - 12,204,646 (2.6% of the total) were near objects with different positions
  - 6,102,323 (1.3% of the total) were the real number of false positives two equal tables
- Why did it happen?







#### Ambiguity

- Although these mistakes represent just 1.3% of the total, the amount of ambiguous matchings was very high (millions of objects).
- Matching ambiguity is an open problem and needs to be explored.





# **Motivation**



#### • Ambiguity

- Binary matching does not generate symmetric results using more than 2 catalogs
- There are no solutions to n-way matching
- The best attribute which identifies the astronomical objects is its position, but it isn't precise
- All these characteristics produce ambiguities









• Measure this ambiguity and propose a better solution to ambiguous n-way matching





# Model proposal

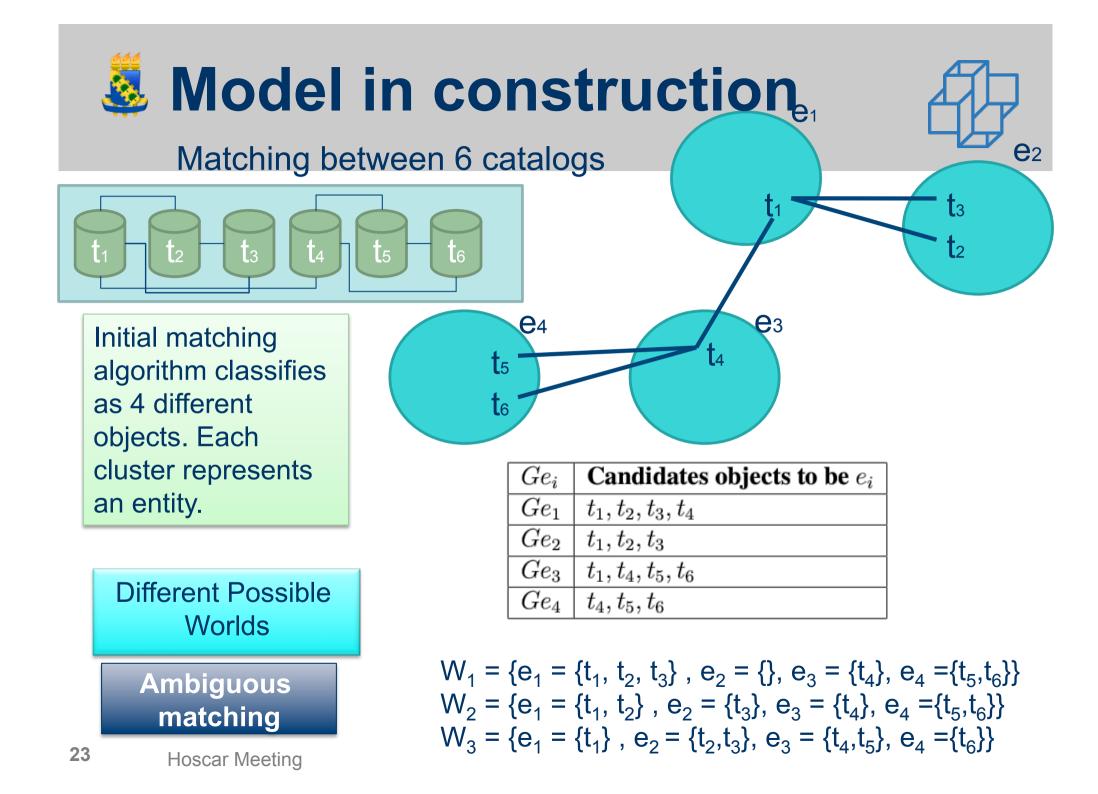


#### • Use a probabilistic model

# - Associate a probability distribution for possible matchings

#### Produce all possible worlds and calculate its probability based on a model proposed in [Ayat, N., Akbarinia, R., Afsarmanesh, H., Valduriez, P. Entity Resolution for Uncertain Data. (2012)]







# Main directives



#### • Problems:

- Probabilistic model to calculate the probability of each world;
- Efficient algorithm for choosing the best world
- Expectations:
  - Generate a n-way more precise matching algorithm
  - Solve ambiguities;
- How to evaluate the quality of the result?



# Model in construction

	Object	Entity	Probability	
	$t_1$	$e_1$	0.2	$\triangleright$
	$t_1$	$e_2$	0.3	
	$t_1$	$e_3$	0.5	
$\frown$	$t_2$	$e_1$	0.7	$\triangleright$
	$t_2$	$e_2$	0.3	]
<	$t_3$	$e_1$	0.4	$\triangleright$
	$t_3$	$e_2$	0.6	1
<	$t_4$	$e_1$	0.2	
	$t_4$	$e_3$	0.3	1
	$t_4$	$e_4$	0.5	
	$t_5$	$e_3$	0.3	
	$t_5$	$e_4$	0.7	
	$t_6$	$e_3$	0.6	
	$t_6$	$e_4$	0.4	

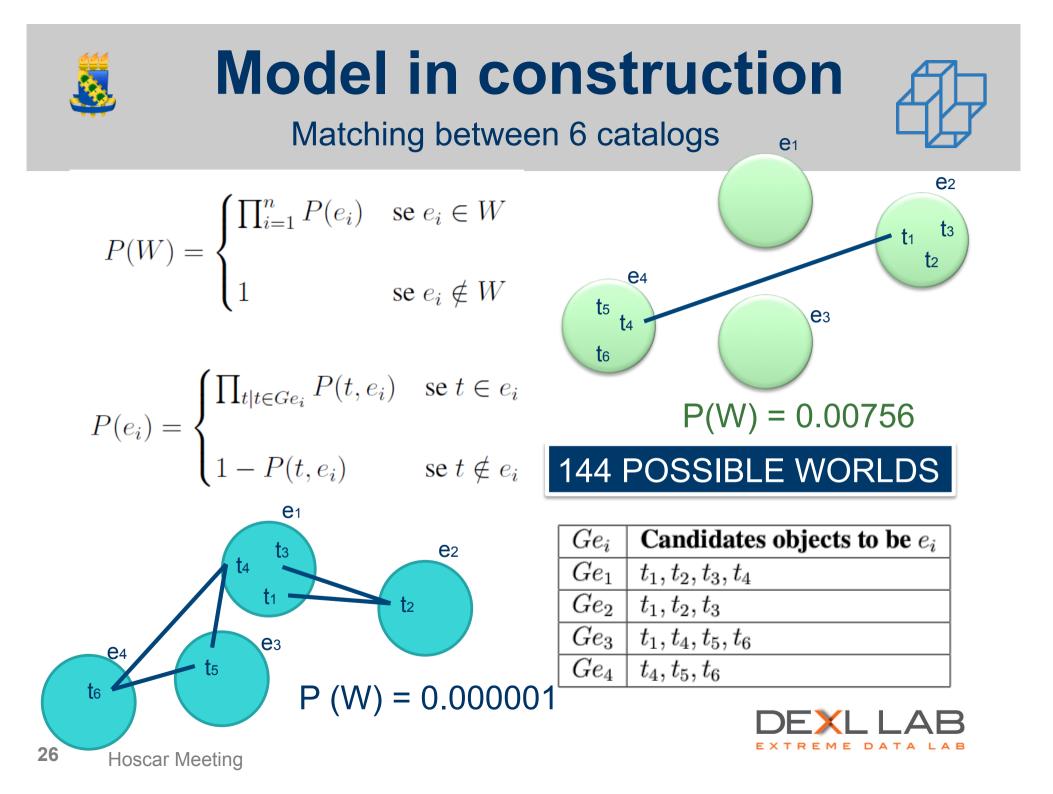
		ATT			
$W_1 = \{e_1 = \{t_1, t_2, t_3\}, e_2 = \{\}, e_3 = \{t_4\}, e_4 = \{t_5, t_6\}\}$					
$P(W) - \epsilon$	$ \left\{ \prod_{i=1}^{n} P(e_i) \right. $ 1	se $e_i \in W$			
$I(vv_i) =$	1	se $e_i \notin W$			

$$P(e_i) = \begin{cases} \prod_{t|t \in Ge_i} P(t, e_i) & \text{se } t \in e_i \\ \\ 1 - P(t, e_i) & \text{se } t \notin e_i \end{cases}$$

Probability of the object  $t_i$  represent the entity  $e_j$ 

 $P(W_1) = 0.0448 \times 1 \times 0.042 \times 0.14 = 0.0003$ 

	$e_i$	$P(e_i)$	$Ge_i$	Candidates objects to be $e_i$
	$e_1$	$0.2 \ge 0.7 \ge 0.4 \ge (1 - 0.2) = 0.0448$	$Ge_1$	$t_1, t_2, t_3, t_4$
	$e_2$	1	$Ge_2$	$t_1, t_2, t_3$
	$e_3$	$0.3 \ge (1 - 0.5) \ge (1 - 0.3) \ge (1 - 0.6) = 0.042$	$Ge_3$	$t_1, t_4, t_5, t_6$ DEXLLAB
25	$e_4$	$0.7 \ge 0.4 \ge (1 - 0.5) = 0.14$	$Ge_4$	$t_4, t_5, t_6$ extreme data LAB





# Conclusion



- To develop this work, it is necessary:
  - Decide which initial algorithm should be used
  - Find a way to calculate the probability of an object belonging to an entity
  - Define the best probabilistic model to calculate the probability of each world;
  - Develop an efficient algorithm for choosing the best world
  - Find a way to evaluate the quality of the result







# Merci

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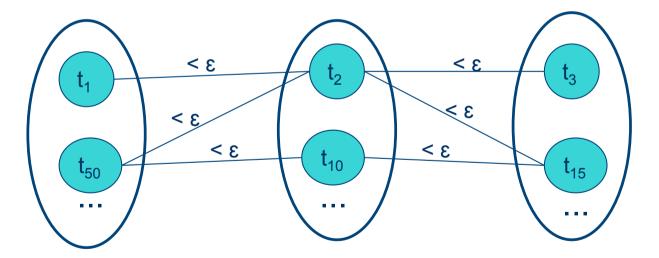


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Hypergraph Hypernode = Catalog Edge means possibility of matching





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# [Fu et al. 2012] Fast approximate matching of astronomical objects.

- For each object *p* in the image, we are looking for an object *q* in the catalog such that:
  - Among all objects in the image, *p* is the nearest to *q*.
  - Among all objects in the catalog, *q* is the nearest to *p*.
  - The distance between *p* and *q* in the twodimensional spherical coordinates is at most 1 arc second, which is 1/3600 of a degree.

