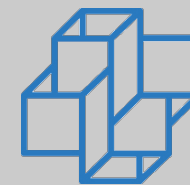


HaQoop: scientific workflows over BigData

3rd HOSCAR Meeting
Bordeaux, Sept. 2013

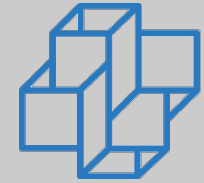
Fabio Porto, Douglas Ericson Oliveira
Matheus Bandini, Henrique Kloh
Reza Akbarinia, Patrick Valduriez



Outline

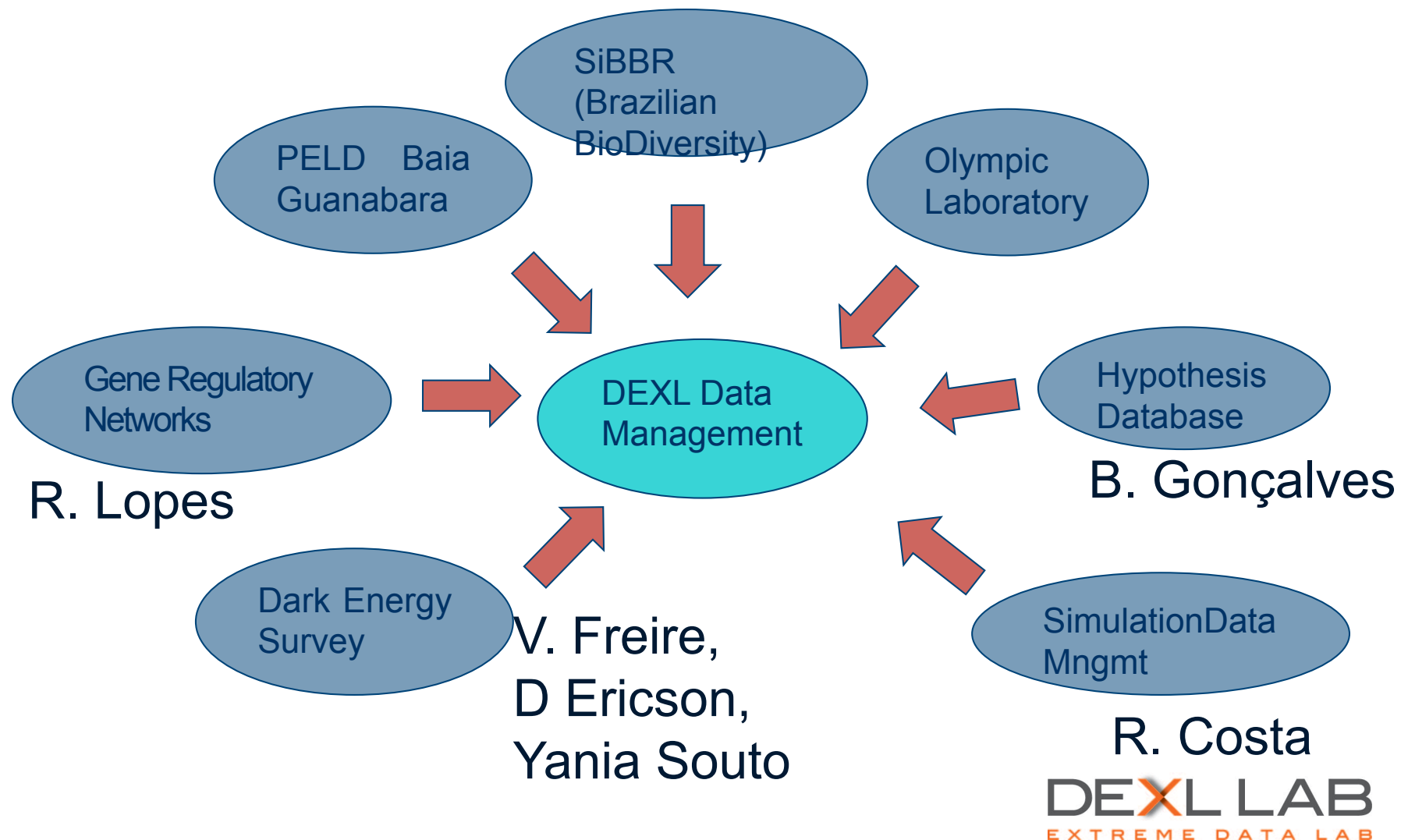
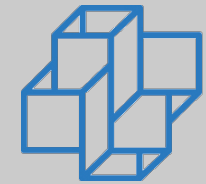
- Introduction
- Previous work in the collaboration
- HaQoop
- Initial experiments
- Final Comments

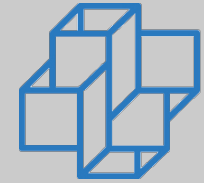
The Data eXtreme Lab (DEXL) Mission



- To support *in-silico* science with data management techniques;
 - To develop interdisciplinary research with contributions on data modelling, design and management;
 - To develop tools and systems in support to in-silico science;
- Currently
 - 3 researchers
 - 8 PhD students
 - 10 engineers
- Projects
 - Astronomy
 - Medicine
 - Sports Science
 - Biology, Ecology
 - Biodiversity

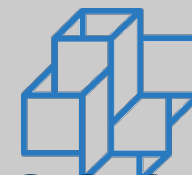
Current projects



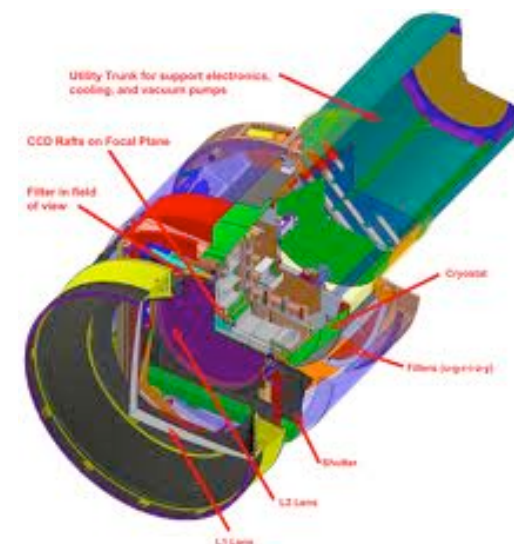


BigData for the *10s*

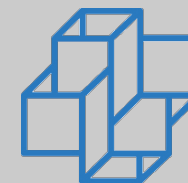
- BigData Processing and Analyses
 - Concerns with Obtaining
 - Volume, Variety, Velocity
 - Concerns with Usage
 - Sparse, infrequent
 - Exploratory, hypotheses driven
- Interested in processing scientific BigData




LSST – Large Synoptic Survey Telescope



- 800 images p/ night during 10 years !!
- 3D Map of the Universe
- 30 TeraBytes per night
- **100 PetaBytes in 10 years**
 - 10^5 disks of 1 TB




Skyserver – Sloan Project

 Sloan Digital Sky Survey / SkyServer

DR8

HomeDataSchemaEducationAstronomySDSSContact UsDownloadSite SearchHelp

DR8 Tools



Getting Started

Famous places

Get Images

Scrolling sky

Visual Tools

Search

- Radial

- Rectangular

- SQL

- Imaging Query

- Spectro Query

Object Crossid

CasJobs

Radial Search

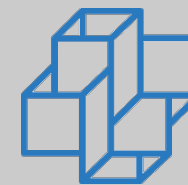
NOTE: To be fair to other users, queries run from SkyServer search tools are restricted in how long they can run and how much output they return, by **timeouts** and **row limits**. Please see the [Query Limits help page](#). To run a query that is not restricted by a timeout or number of rows returned, please use the [CasJobs batch query service](#).

ra		114.5	
dec		10.89	
radius [arcmins]		3	

	Min		Max
<input type="checkbox"/>	0	u	20
<input type="checkbox"/>	0	g	20
<input type="checkbox"/>	0	r	20
<input type="checkbox"/>	0	i	20
<input type="checkbox"/>	0	z	20

SubmitReturn: ☐ all rows ☒ max 10Format: ☒ HTML ☐ CSVReset

Enter the **ra** and **dec** either in degrees or in h:m:s, d:m:s notation. The search **radius** is measured in arcminutes. Check the magnitudes you would like to constrain in your query. If you prefer not to use specific attributes, leave those rows unchecked. (If you do not insert constraints and select all entries, you will receive many records!)



Dark Energy survey - pipelines

The Dark Energy Survey

Patricia Bittencourt Egeland ▾
Welcome to the Ana portal.

Science Portal

Home My Workspace Pipelines Tools Data Server Documentation Help Credits Other Portals Logout

Code Viewer
List of available software

Pipeline Pipeline Components Source Codes Executables Libraries

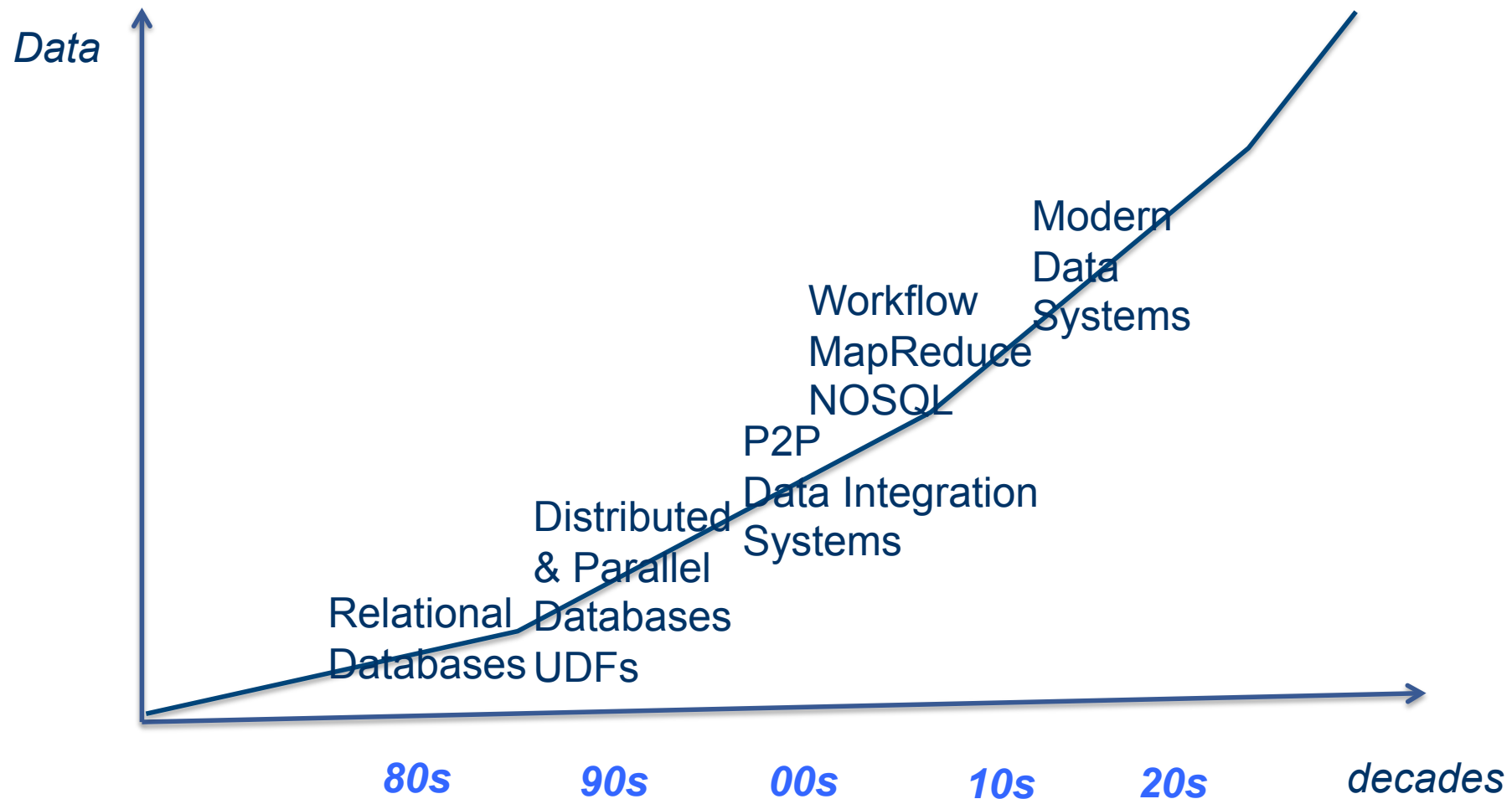
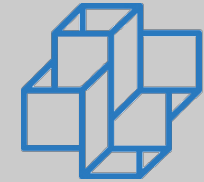
Show 10 entries

Name	Owner	Version	Date	Status	Description	Components	Release Notes	XML
AddClus	Eduardo Balbinot	---	---	Available		+		
Angular 2-point Correlation	Fernando Simoni	---	---	Available		+		
Angular Correlation	Fernando Simoni	---	---	Available		+		
BAOs	Fernando Simoni	---	---	Available		+		
Cluster Matching	Ricardo Ogando	---	---	Available		+		
Color Catalog	Bruno Rossetto	---	---	Not Available		+		
Compute Mass	Beatriz Ramos	---	---	Available		+		
Correlation Function	Fernando Simoni	---	---	Available		+		
Fit CF	Fernando Simoni	---	---	Not Available		+		
Fit LF	Fernando Simoni	---	---	Not Available		+		

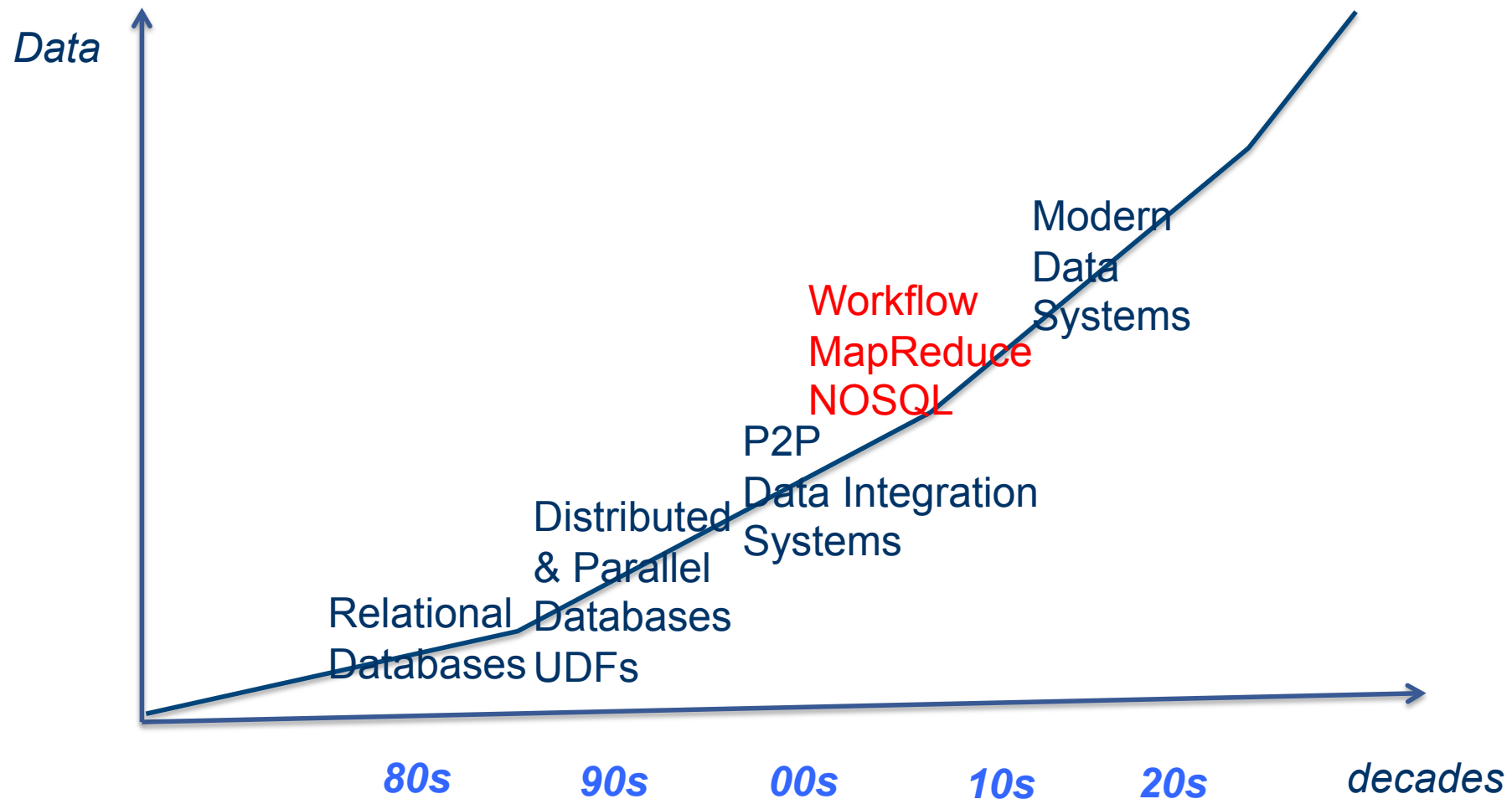
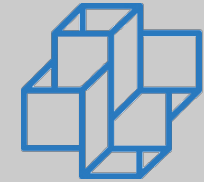
Showing 1 to 10 of 30 entries

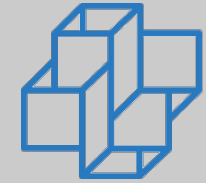
First Previous 1 2 3 Next Last

Data Processing Systems: an Evolution



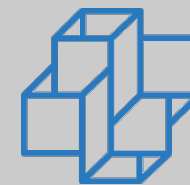
Data Processing Systems: an Evolution





Data Processing Pillars

- Reduce the number of data retrieval operations
- Efficient iterative processing over elements of sets;
- Parallelism obtained by partitioning data;
 - Or pipelining data through parallel execution of operators
- Explore the semantics of data operations;
- Automatic decisions based on data statistics;
- Data consumed by humans
- Data of simple structure/semantics



General Model

$$\mathcal{R} \text{-----} f(x) \rightarrow \mathcal{R}'$$

$$\mathcal{R}_1 \text{-----} f(x) \rightarrow \mathcal{R}_1'$$

$$\mathcal{R}_2 \text{-----} f(x) \rightarrow \mathcal{R}_2'$$

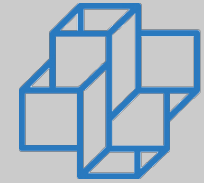
...

$$\mathcal{R}_n \text{-----} f(x) \rightarrow \mathcal{R}_n'$$

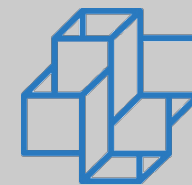
WHAT CHANGES?

\mathcal{R}''

Processing BigData



- Reduced data is still Big: millions of elements;
 - Access patterns less predictable
- Data may be:
 - Incomplete
 - Uncertain
 - Ambiguous
- Operation semantics are unknown (black box modules)
 - User code implementation
 - Arbitrary f (a workflow)
- Some operations are blocking, with respect to the consumption and production of data
 - Parallel MPI based programs
 - Prevent data-driven parallelism
- Consumption
 - Data analysis



Big Data Model

$$X \text{-----} \langle T_1, M_2, \dots, V_k \rangle \text{----} \rightarrow Z'$$

$$X_1 \text{-----} \langle T_1, M_2, \dots, V_k \rangle \text{----} \rightarrow X_1'$$

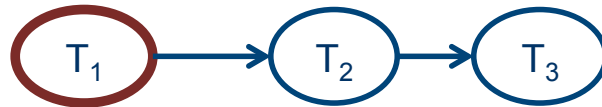
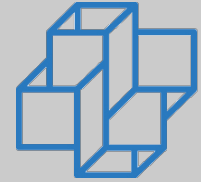
$$X_2 \text{-----} \langle T_1, M_2, \dots, V_k \rangle \text{----} \rightarrow X_2'$$

...

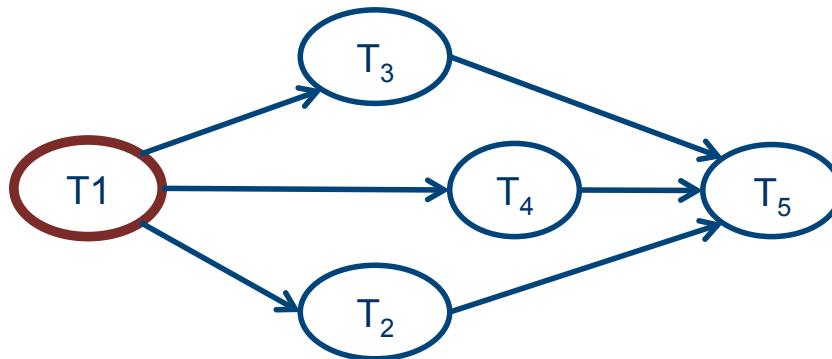
$$X_n \text{-----} \langle T_1, M_2, \dots, V_k \rangle \text{----} \rightarrow X_n'$$

$$\bigcup_{i=1,n} X_i' \text{-----} g(y) \text{----} \rightarrow X''$$

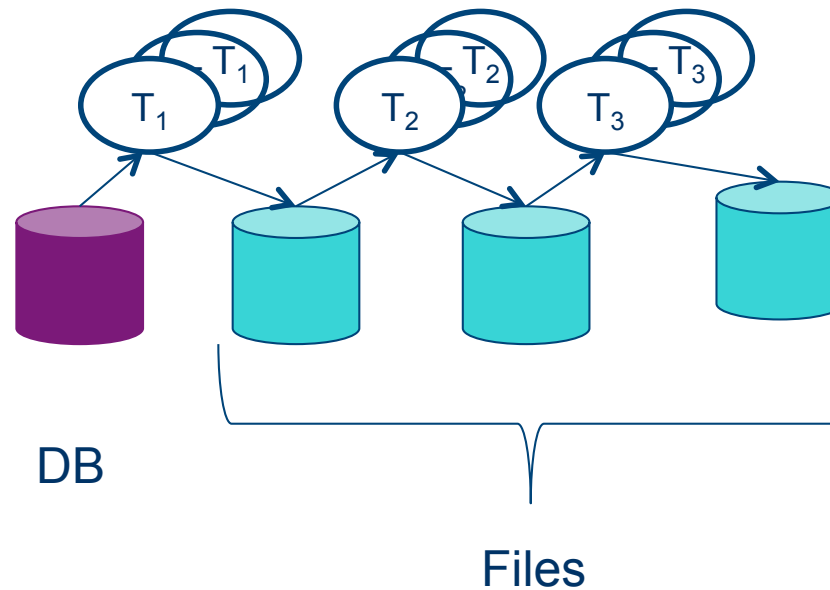
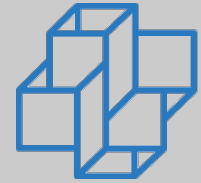
Workflow - Partial Ordering of \mathcal{T}

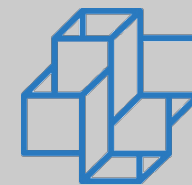


Where each  is an activity



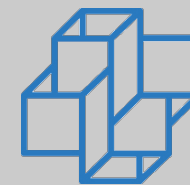
Workflow DB – complete picture





General Problem

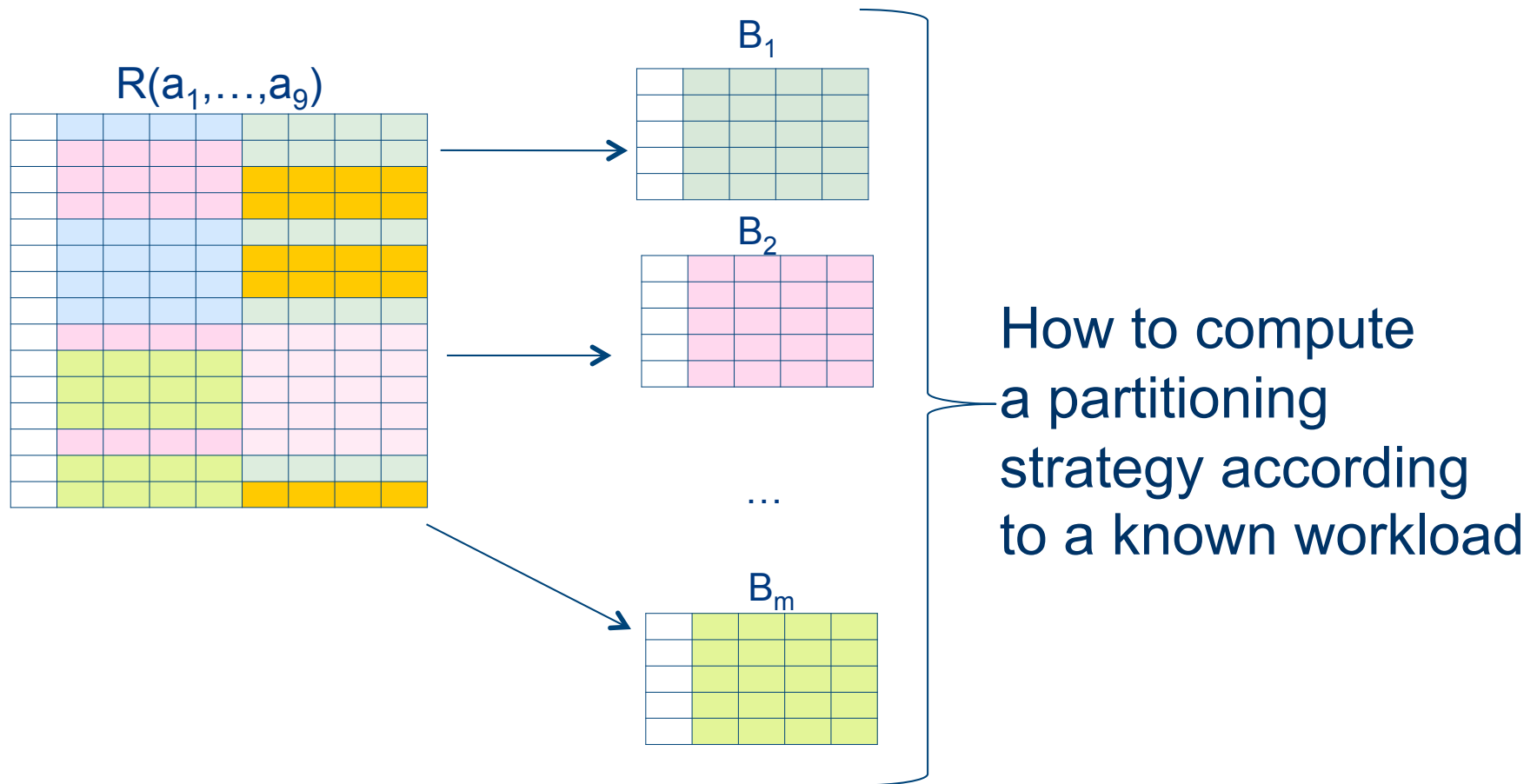
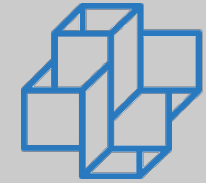
- To Conceive an efficient and robust workflow execution strategy that considers data retrieved from databases and files produced in intermediate steps



PREVIOUS WORK IN THE COLLABORATION: LNCC, COPPE-UFRJ, INRIA - ZENITH

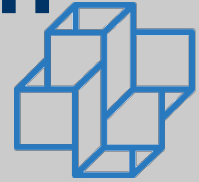
Partitioning the DB into Blocks

Work with: Miguel Liroz-Gistau,
Esther, Patrick, Reza



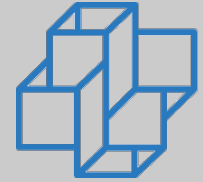
Workflow algebra and optimization

Eduardo Ogasawara,
Marta Mattoso, Patrick Valduriez

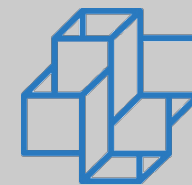


- Scientific workflow definition mapped to a known data model
 - Input/output modelled as relations
 - workflow activities mapped to operators in a generic algebra;
 - Algebra operators describe input/output ratio
 - Enables automatic analysis of workflow definition according to type of applied data transformation
 - Enables automatic workflow transformation

Objective

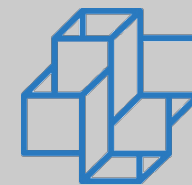


- Processing big data by scientific workflows shall benefit from known data processing techniques
 - Activities semantics ✨ ✨
 - Process to data locality ✨
 - Optimize data and files distribution
 - Use generic MapReduce parallelism paradigm



Approach

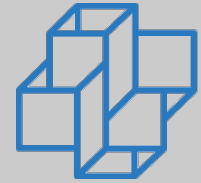
- Use MapReduce paradigm to run scientific workflow
- Define a allocation strategy that considers:
 - The number of database partitions
 - The number of map tasks
 - The input/output semantics of workflow activities
 - The number of reduce tasks



Three scenarios evaluated

- Exploring experimentally variations on $|P|$, $|T|$, $|F|$ as the basis for the model:
 - a) $|P| = 1$, $|T| \gg 1$
 - b) $|P| = |T| \gg 1$, D is a distributed database
 - c) $|P| \leq |T|$, $|P|, |T| \gg 1$
- Which data processing parallel strategy leads to best results in workflow execution?

Parallel workflow evaluation on BigData



HaDooPDB

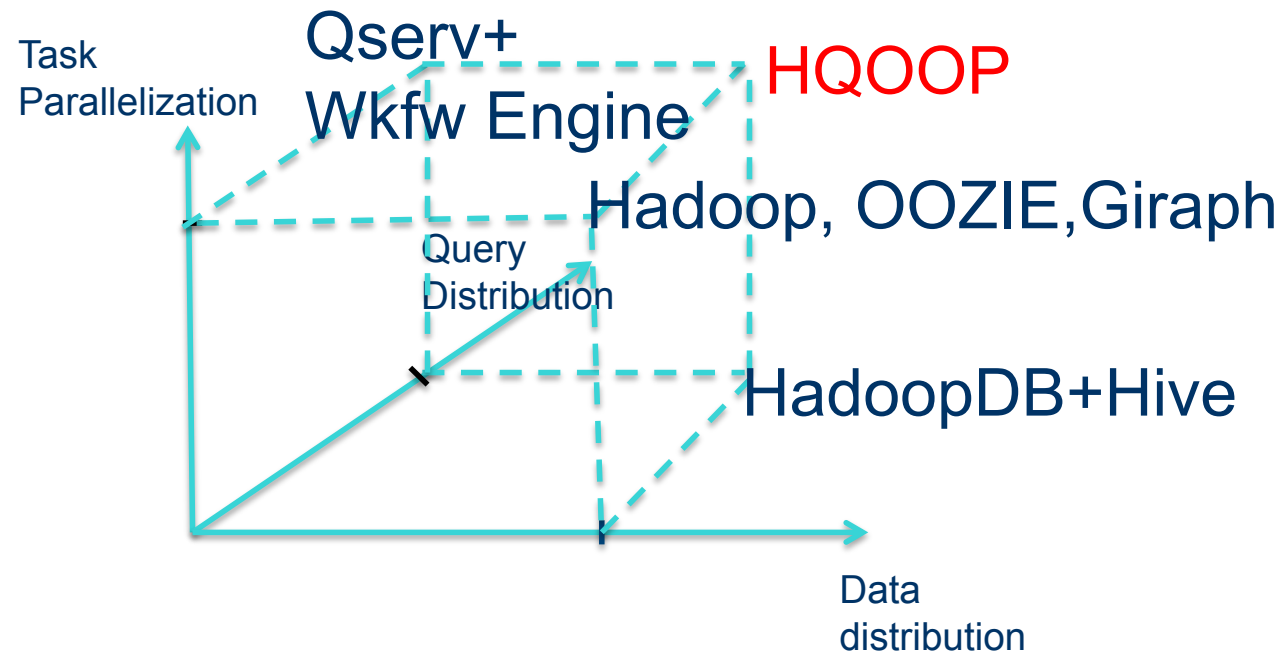
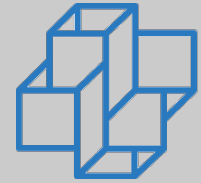


Dryad
LINQ
MS Research

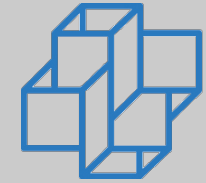


DEXL LAB
EXTREME DATA LAB

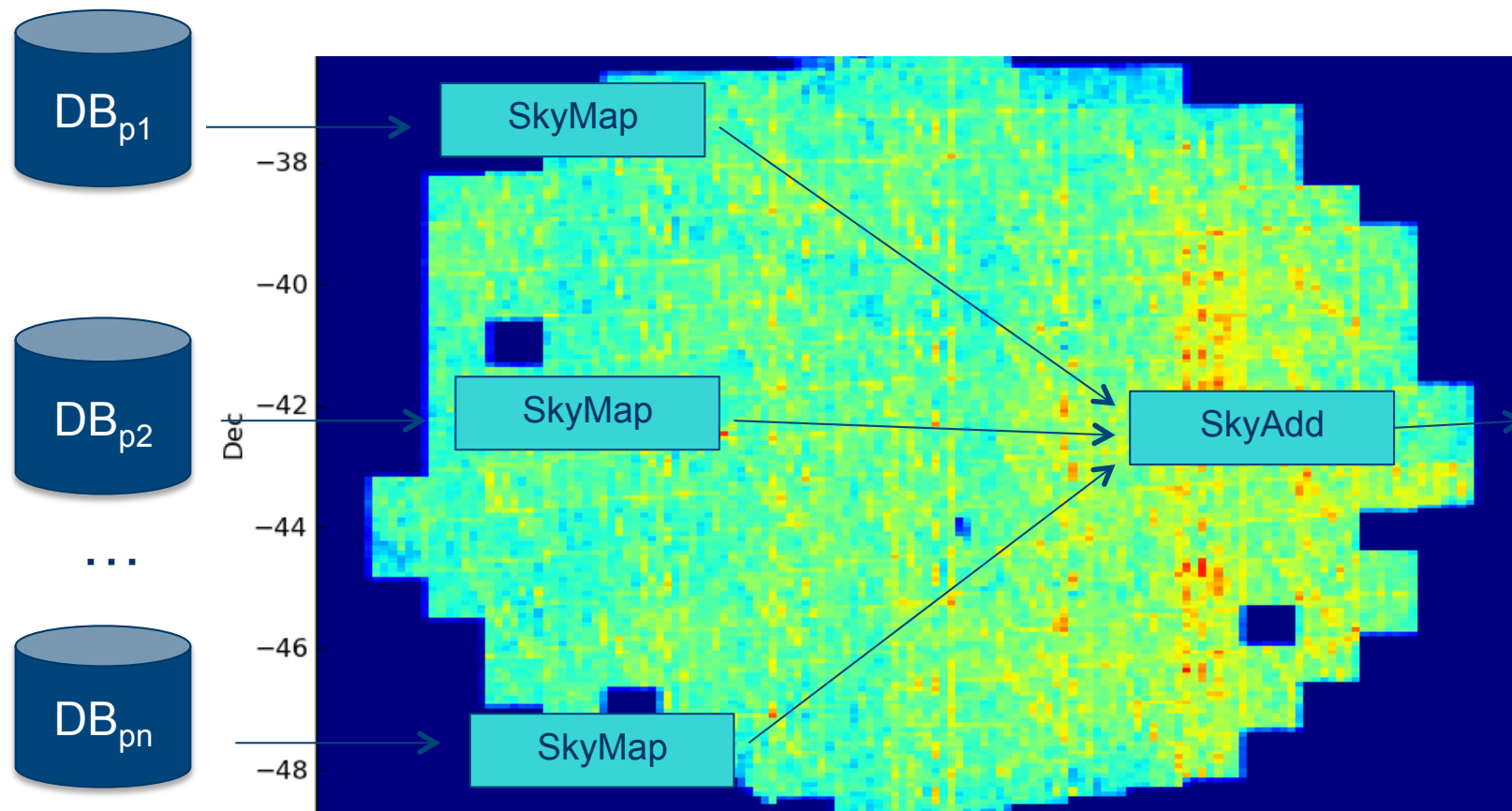
Architectural Viewpoint

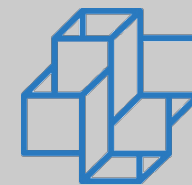


Parallel workflow execution over Dark Energy Survey Catalog



Partitioned catalogue stored on PostgreSQL

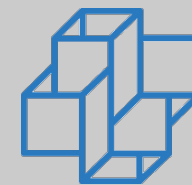




HaQoop

- Hadoop – Open Source apache project
 - A state of the art task parallelization framework for Big Data processing
 - Split computation into two steps
 - Map (remember f ?)
 - Reduce (remember g ?)
- To reuse Hadoop scalability, fault tolerance
- To extend Hadoop with workflow expressions
 - Make f a general workflow engine (QEF)
- Restricted workflow expressions



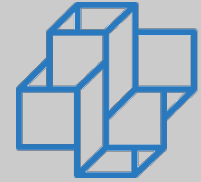


QEF – Data Processing System



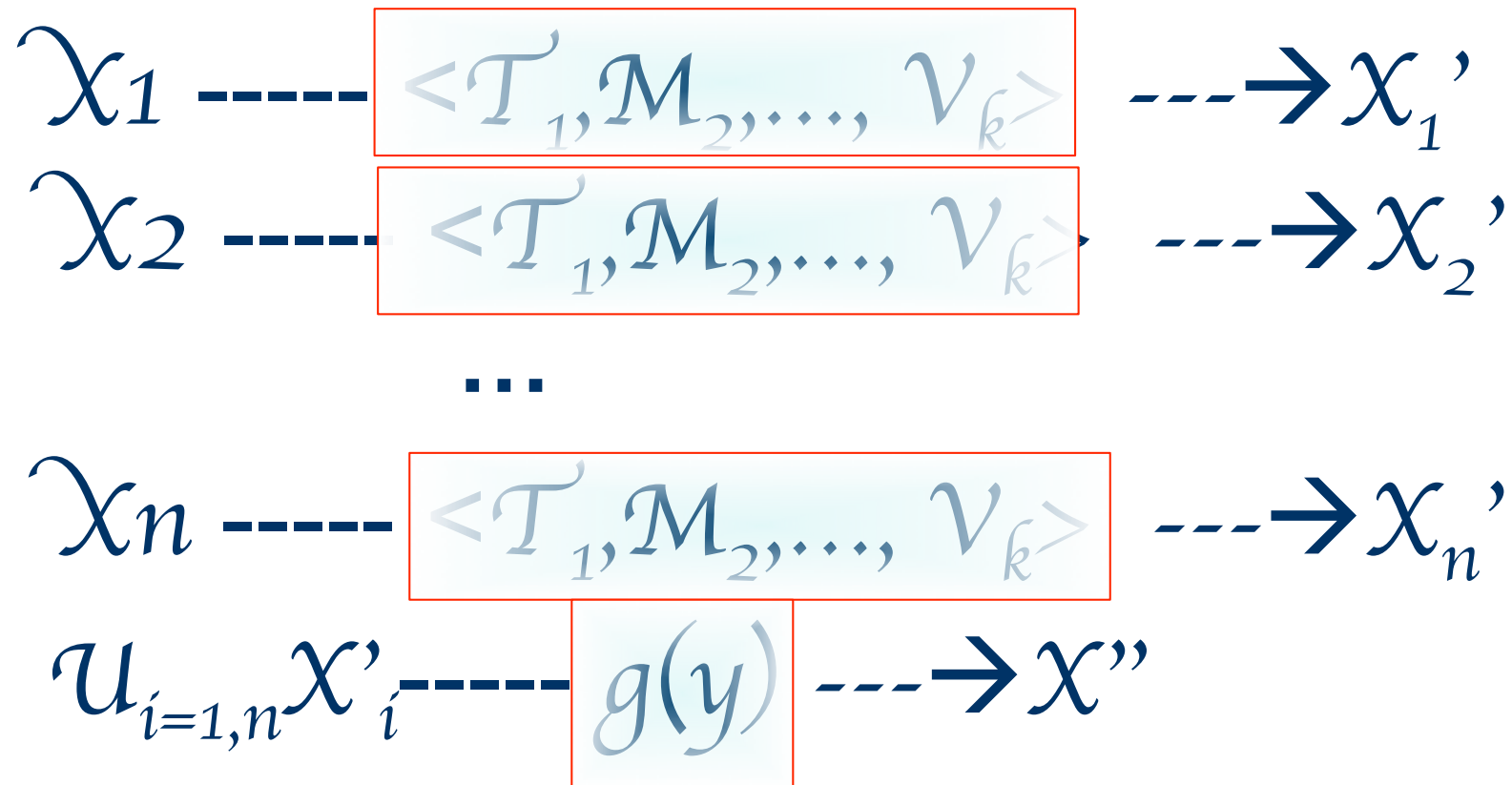
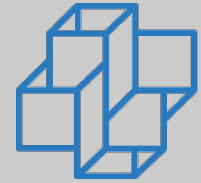
- Designed based on principles of modern database query engines;
- Extendable for any user code
- Extendable for any data structure
- Can be downloaded: <http://dexl.incc.br/qef>

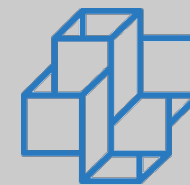
Main technical characteristics



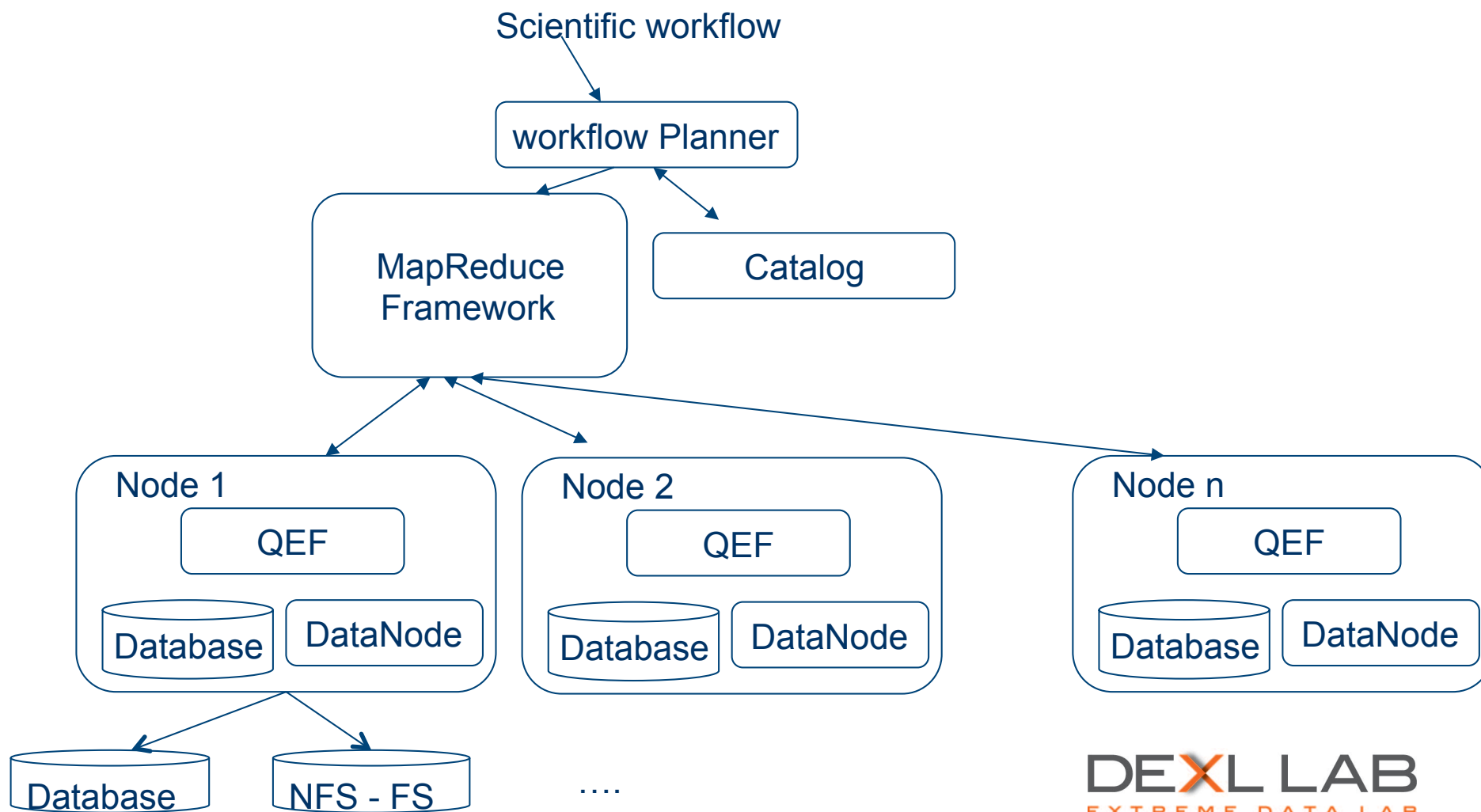
- Pipeline (iterator execution model)
- Iterations
- Algebraic/control operations
 - Allows both in-memory data exchange as file-based i/o
 - Run in both CPUs and GPUs
 - Push and pull data execution (using control operations)
- Dynamic optimization
 - Block-size computation
- Global and local state
 - Control tuples
- Catalog
 - Environment
 - Statistics
 - Metadata
- Synchronous and asynchronous execution

QEF as a Mappers & Reduce Job on Hadoop

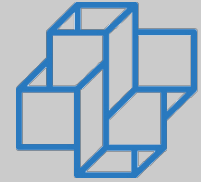




HaQoop architecture



Example: SkyMap Workflow



Select ra, dec
From Catalog
Where ra between 330 and 333 and
dec between -42 and -43



SkyMapAdd

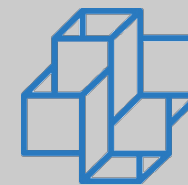


.pkl files

SkyMapAdd

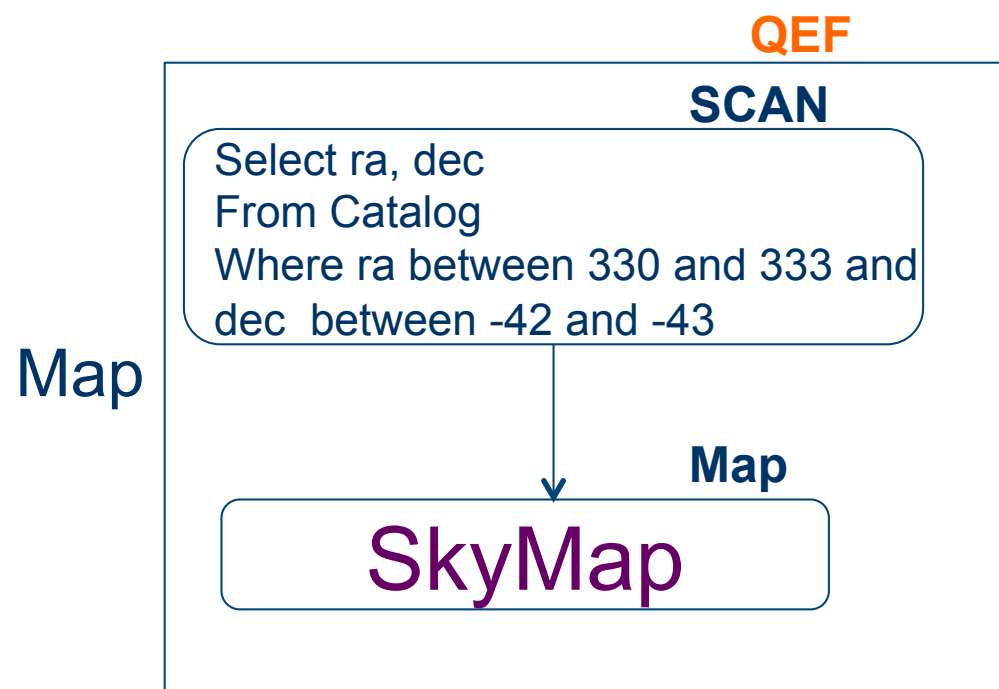
Catalog table

- query returns 200 million sky objects
- uniformly distributed through nodes
- centralized mode
 - each tuple is logically partitioned



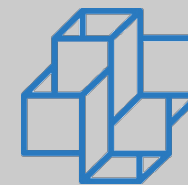
Example

a) Catalog Table uniformly partitioned



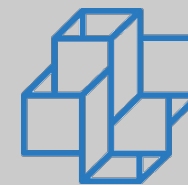
Reduce

SkyMapAdd

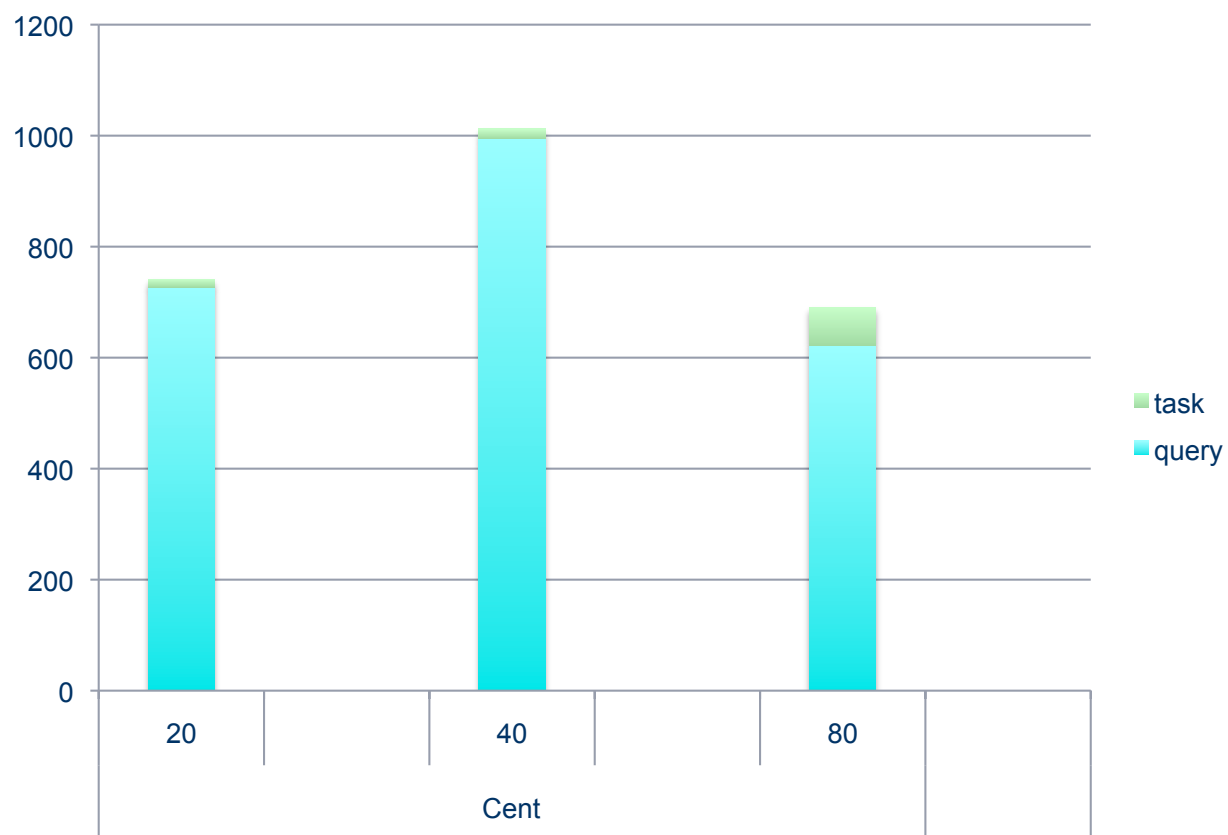


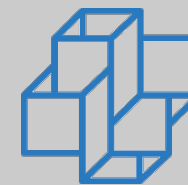
Initial Experiments

- Initial experiments
 - Skymap scenario;
- Cluster SGI
 - Configurations: 20, 40 and 80 nodes;
 - Each node:
 - 2 proc. Intel Zeon – X5650, 6 cores, 2.67 GHz
 - 24 GB RAM
 - 500 GB HD
- Data
 - DES Catalog DC6B
- Tasks
 - Python
- HAQOOP
- Centralized version
 - PostgreSQL 9.1
- Distributed
 - Pg_pool
- Partirioned
 - Multiple postgresQL

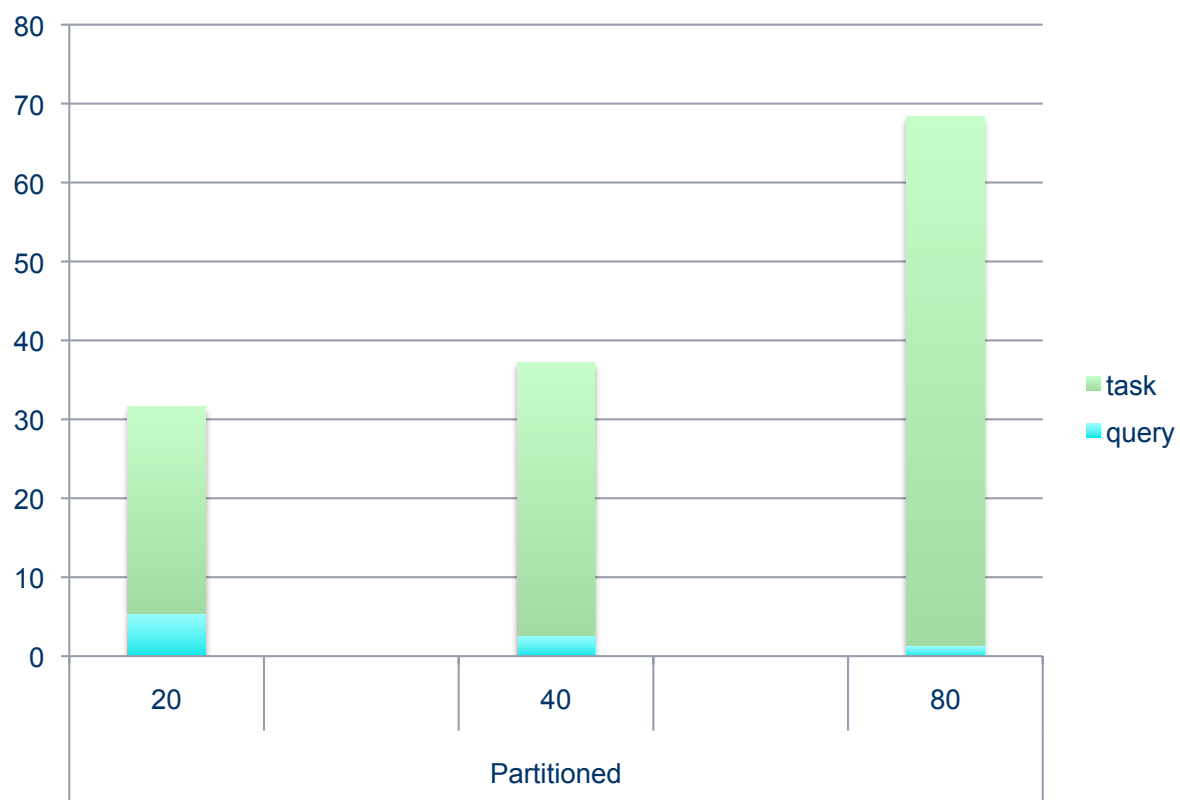


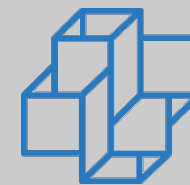
Centralized – Elapsed-time (s)





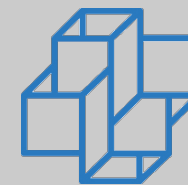
Partitioned DB – Elapsed-time (s)





Final comments

- Collaboration with Zenith-Inria team
- Probable PhD student exchange in 2014



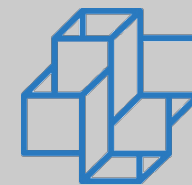
MERCI – OBRIGADO

fporto@lncc.br



Laboratório Interinstitucional de e-Astronomia

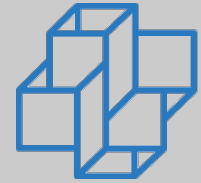




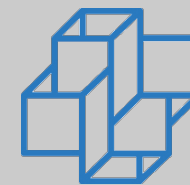
Processing Scientific Workflows

- Analytical Workflows process a large part of Catalog data
 - Catalogs are supported by few indexes, thus most queries scan tens-to-hundreds of millions of tuples
- Parallelization comes as a rescue to reduce analyses elapsed-time, but
 - Compromise between:
 - Data partitioning and degree of parallelization;
 - Current solutions consider:
 - Centralized files to be distributed through nodes (MapReduce)
 - [Alagianins, SIGMOD, 2012] NoDB – reading raw files without data ingestion;
 - Distributed databases (Qserv) to serve Workflow engines
 - [Wang.D.L,2011], Qserv: A Distributed Shared-Nothing Database for the LSST catalog;
 - Centralized databases to serve Workflow Engine (Orchestration LineA)
 - Partitioned database to serve distributed queries (HadoopDB)

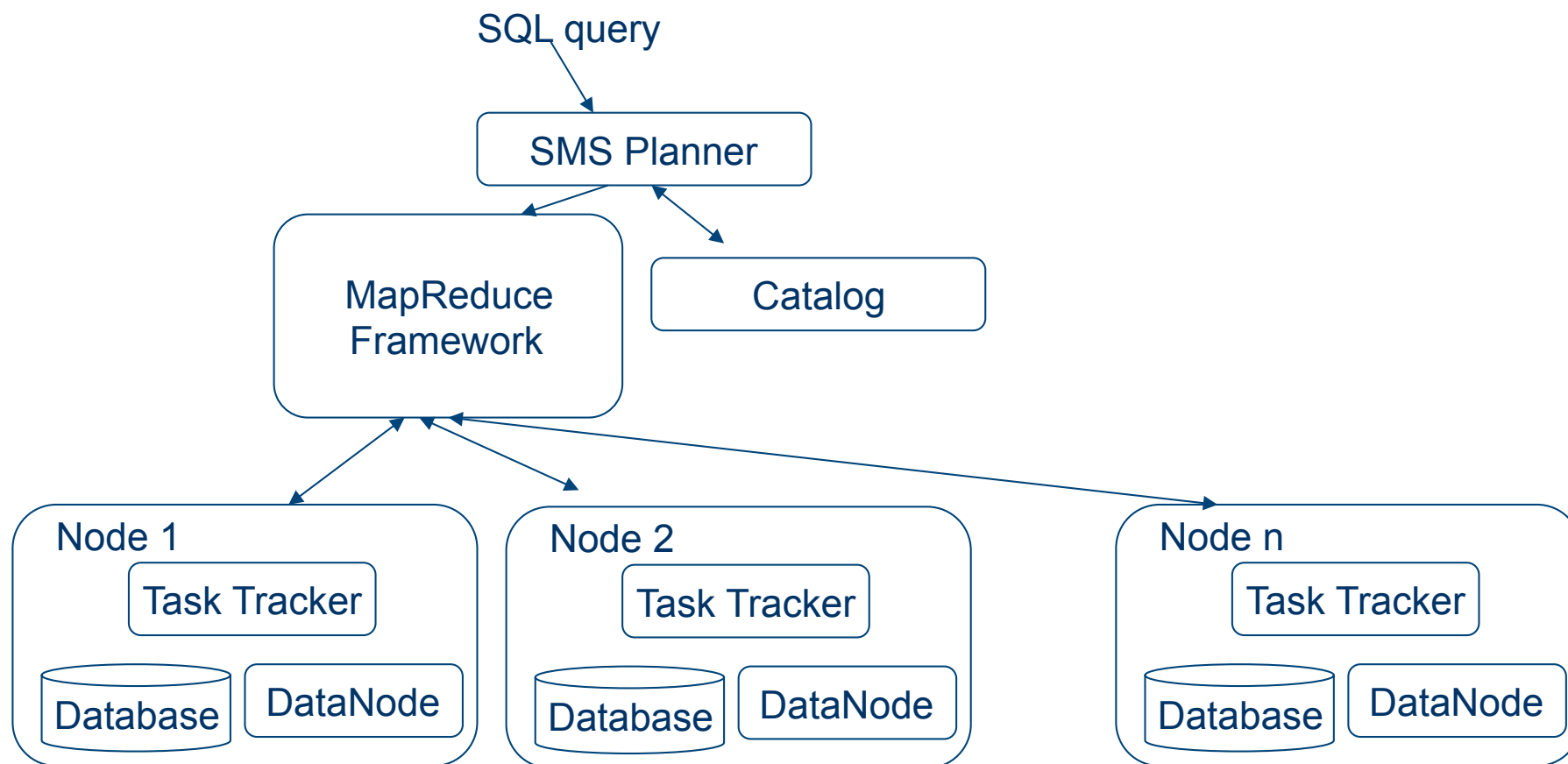
HadoopDB - a step in between [Abouzeid09]

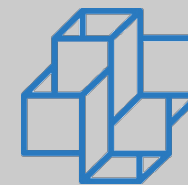


- Offers parallelism and fault tolerance as Hadoop, with SQL queries pushed-down to PostgreSQL DBMS;
- Pushed-down queries are implemented as Map-reduce functions;
- Data are partitioned through nodes.
 - Partitioning information stored in the catalog
 - Distributed through the N nodes



HadoopDB architecture





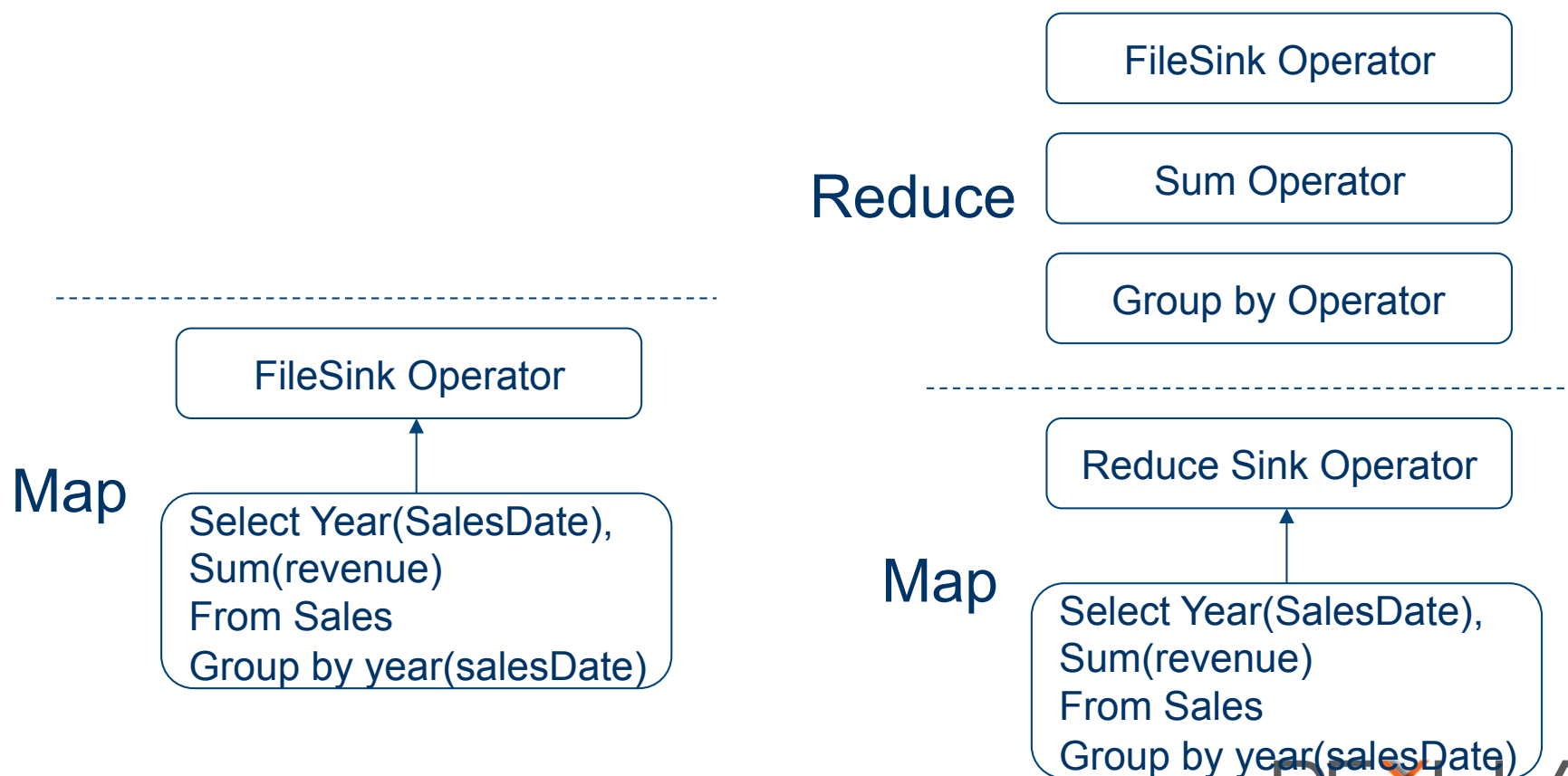
Example

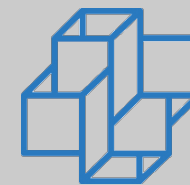
Select year(SalesDate),sum(revenue)

From Sales

Group by year(salesDate)

a) Table partitioned by year(SalesDate) b) no partitioning by year(SalesDate)





Processing Astronomy data

User access

- Ad-hoc queries
- downloads



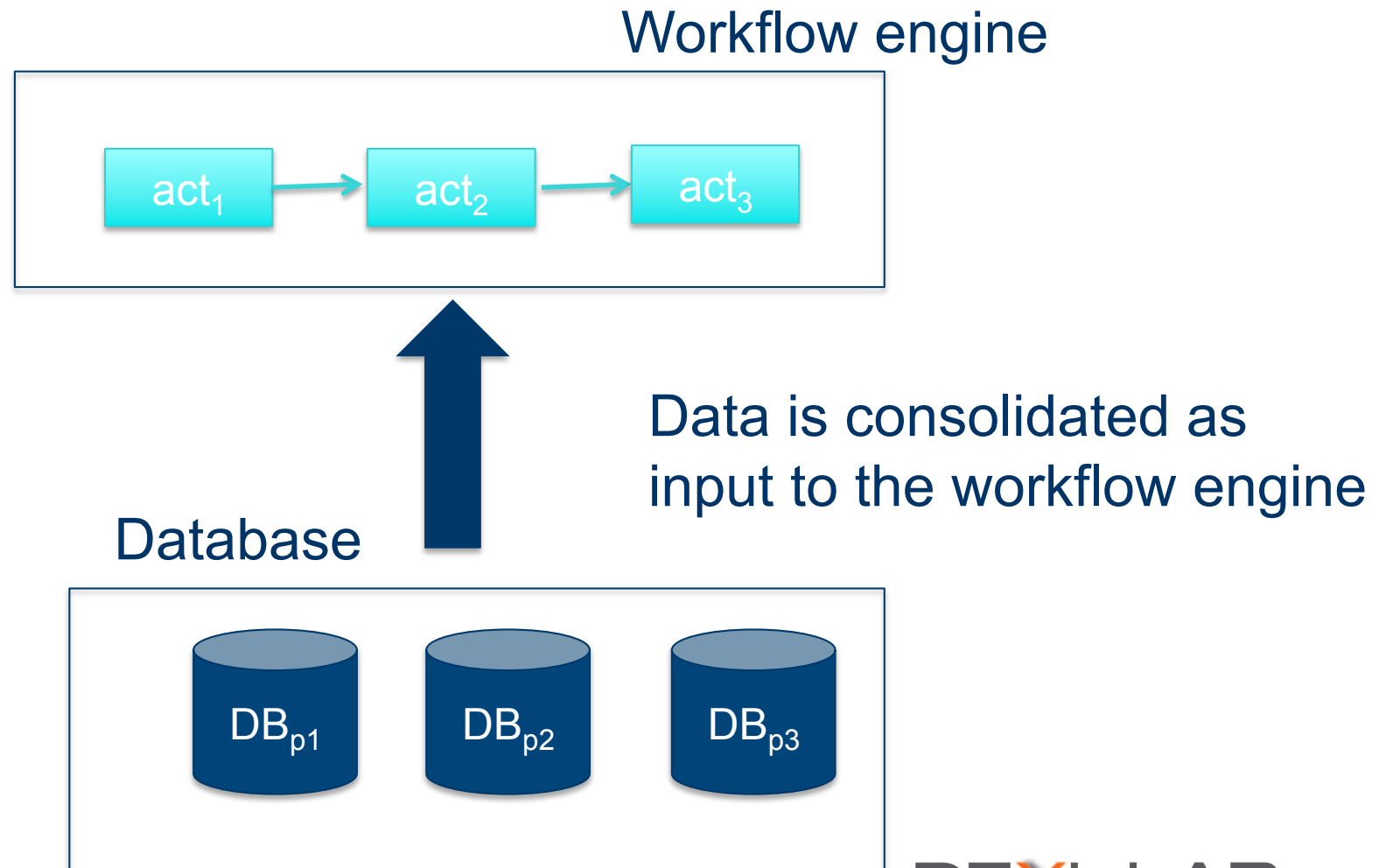
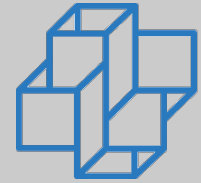
Scientific workflows

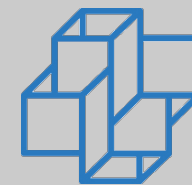
- Analysis



Astronomy
catalogs

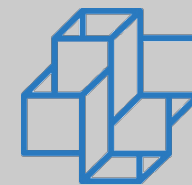
Traditional WF–Database decoupled architecture





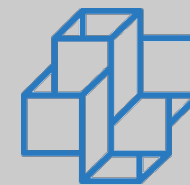
Problems

- Data locality
 - Workflow activities run in remote nodes wrt the partitioned data;
- Load Balance
 - Local processes facing different processing time



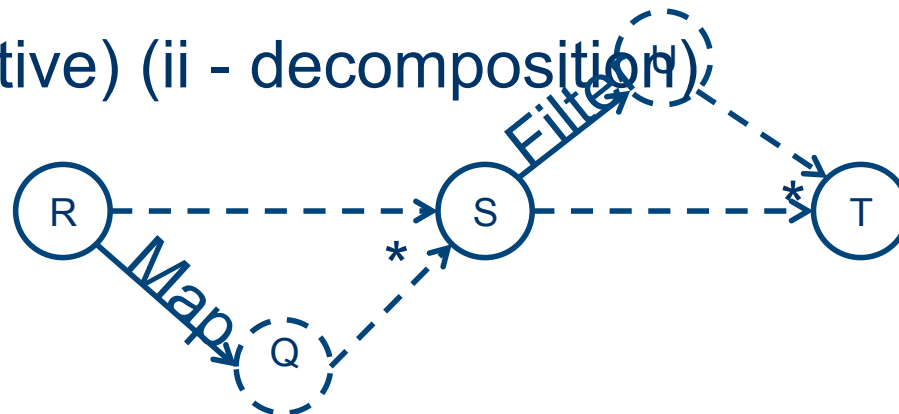
Data locality

- Traditional distributed query processing pushes operations through joins and unions so that can be done close to the data partitions;
- Can we “localize” workflow activities?
 - Moving activities in workflows require operation semantics to be exposed
 - Mapping of workflow activities to a known algebra
 - Equivalence of algebra expressions enabling pushing down operations

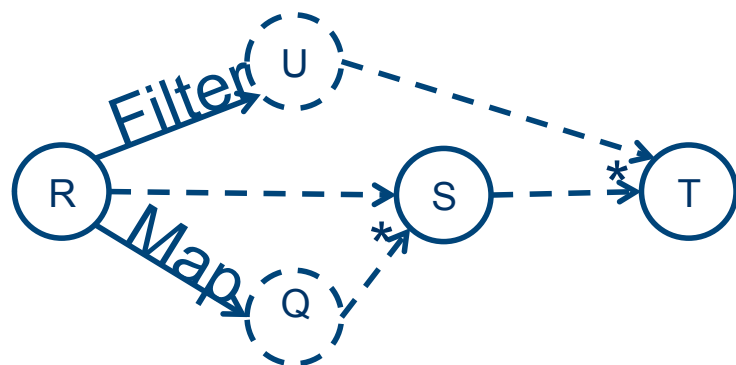


Algebraic transformation

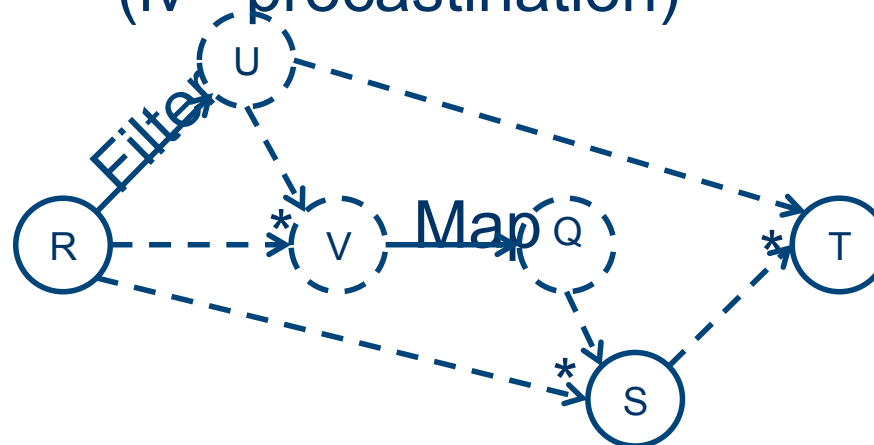
(i - *workflow* – relation perspective) (ii - *decomposition*)

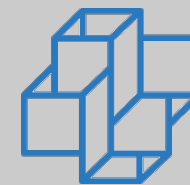


(iii - *anticipation*)

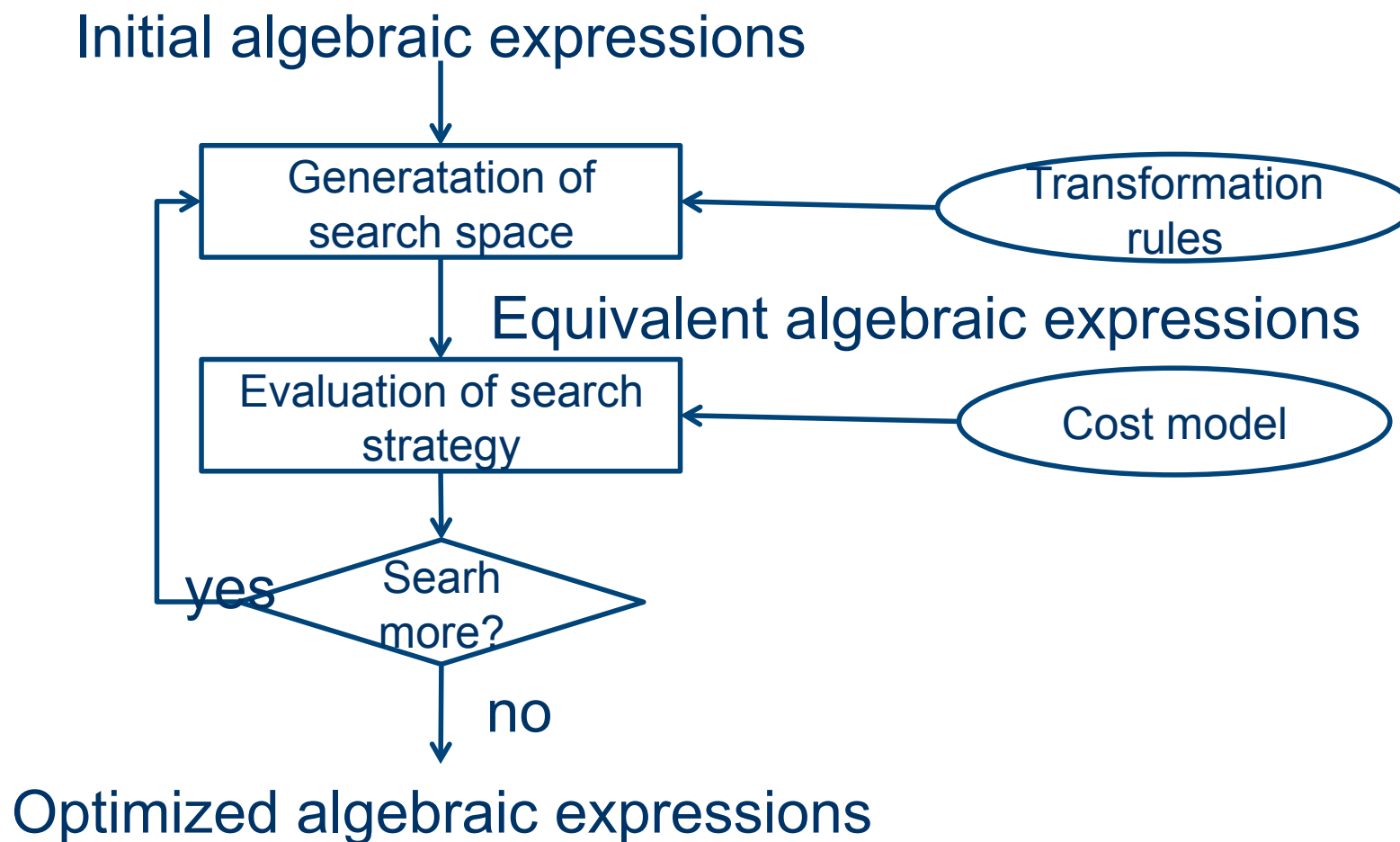


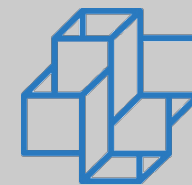
(iv - *procastination*)





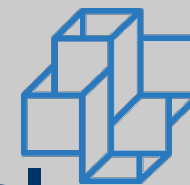
Workflow optimization process



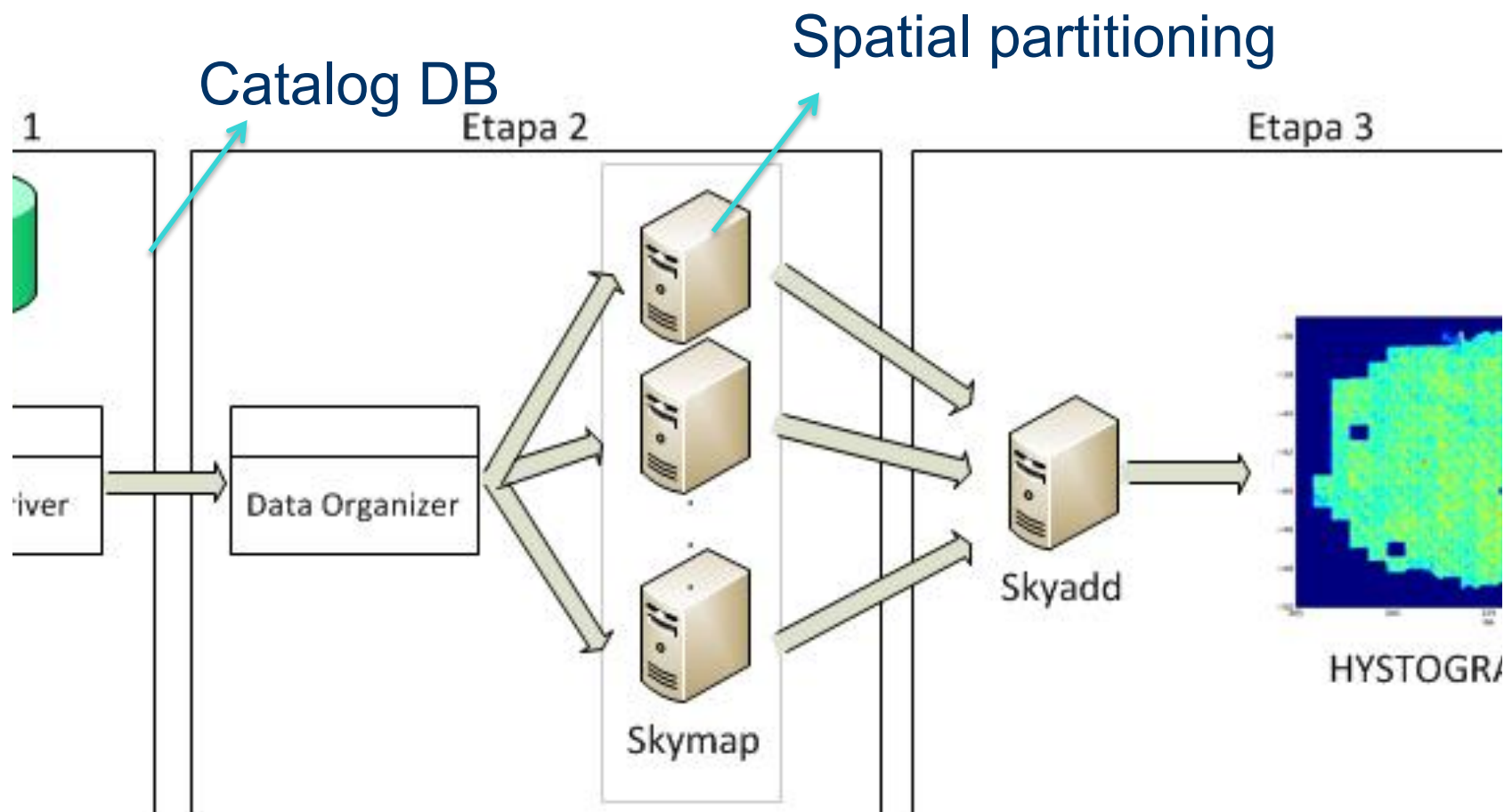


Pushing down workflow activities

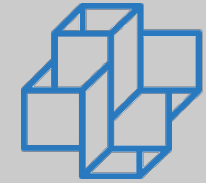
- A first naïve attempt
 - Push down all operations before a Reduce;
- Use a MapReduce implementation where
 - Mappers execute the “pushed-down” operations close to the data



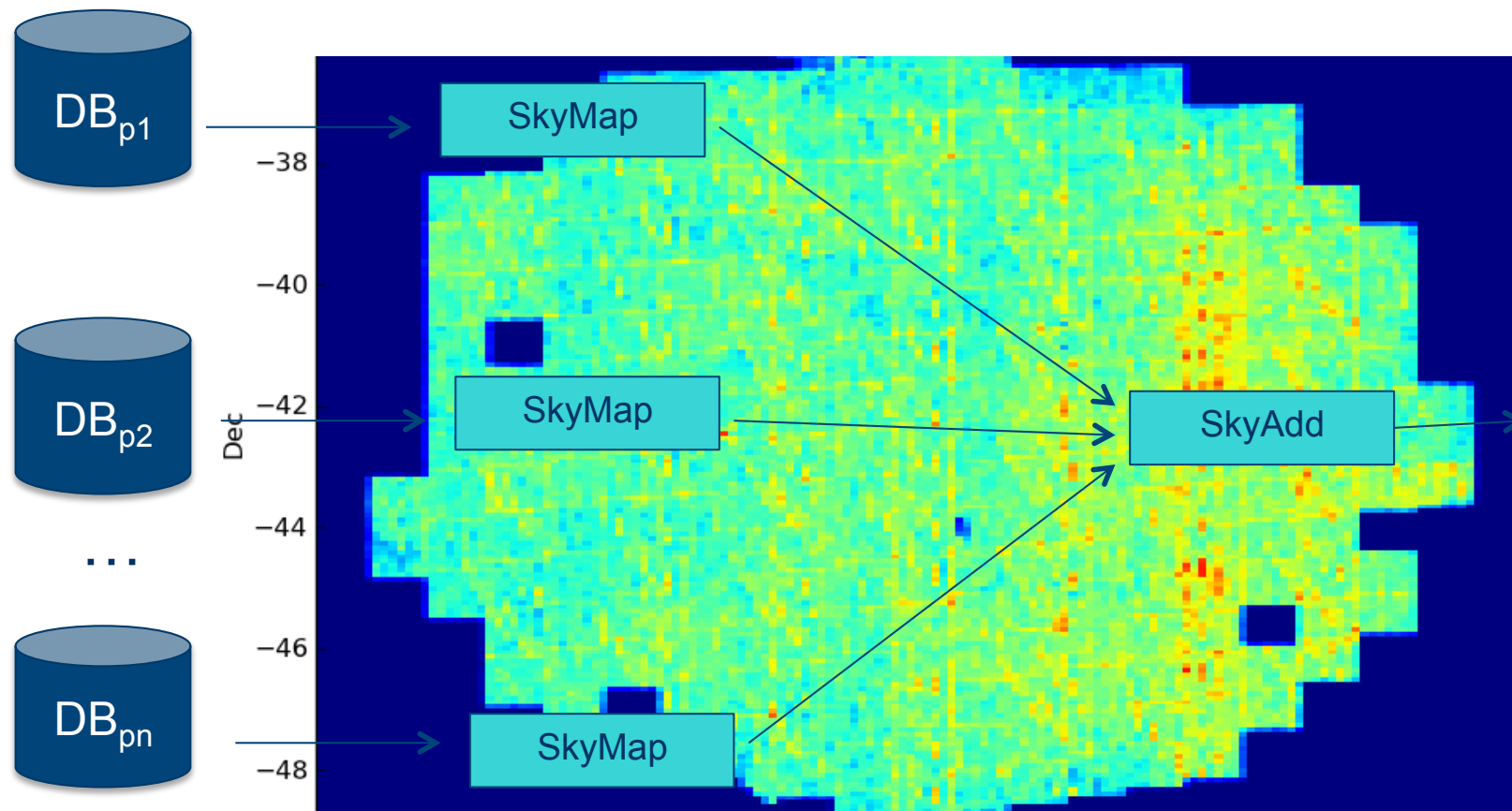
Typical Implementation at LineA Portal



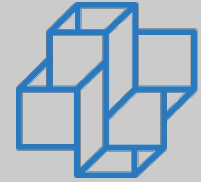
Parallel workflow over partitioned data



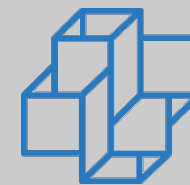
Partitioned catalogue stored on PostgreSQL



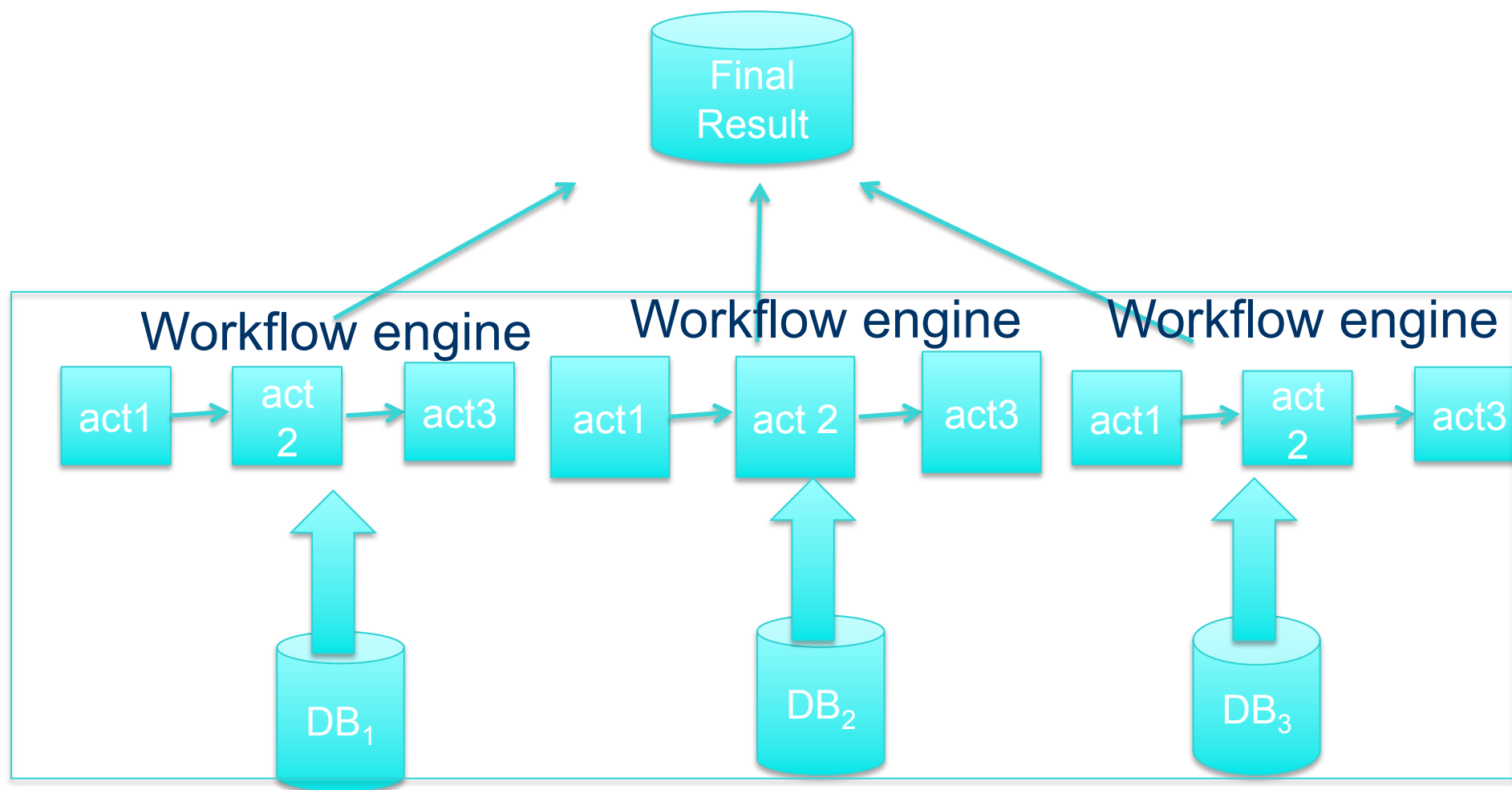
HQOOP - Parallelizing Pushed-down Scientific Workflows

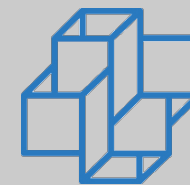


- Partition of data across cluster nodes
 - Partitioning criteria
 - Spatial (currently used and necessary for some applications)
 - Random (possible in SkyMap)
 - Based on query workload (Miguel Liroz-Gestau's Work)
- Process the workflow close to data location
 - Reduce data transfer
- Use Apache/Hadoop Implementation to manage parallel execution
 - Widely used in Big Data processing;
 - Implements Map-Reduce programming paradigm;
 - Fault Tolerance of failed Map processes;
- Use QEF as workflow Engine
 - Implements Mapper interface
 - Run workflows in Hadoop seamlessly;



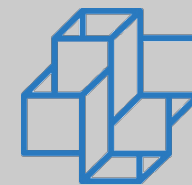
Integrated architecture





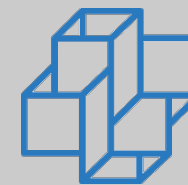
Experiment Set-up

- Cluster SGI
 - Configurations: 1, 47 and 95 nodes;
 - Each node:
 - 2 proc. Intel Zeon – X5650, 6 cores, 2.67 GHz
 - 24 GB RAM
 - 500 GB HD
- Data
 - Catalog DC6B
- Hadoop
 - QEF workflow engine

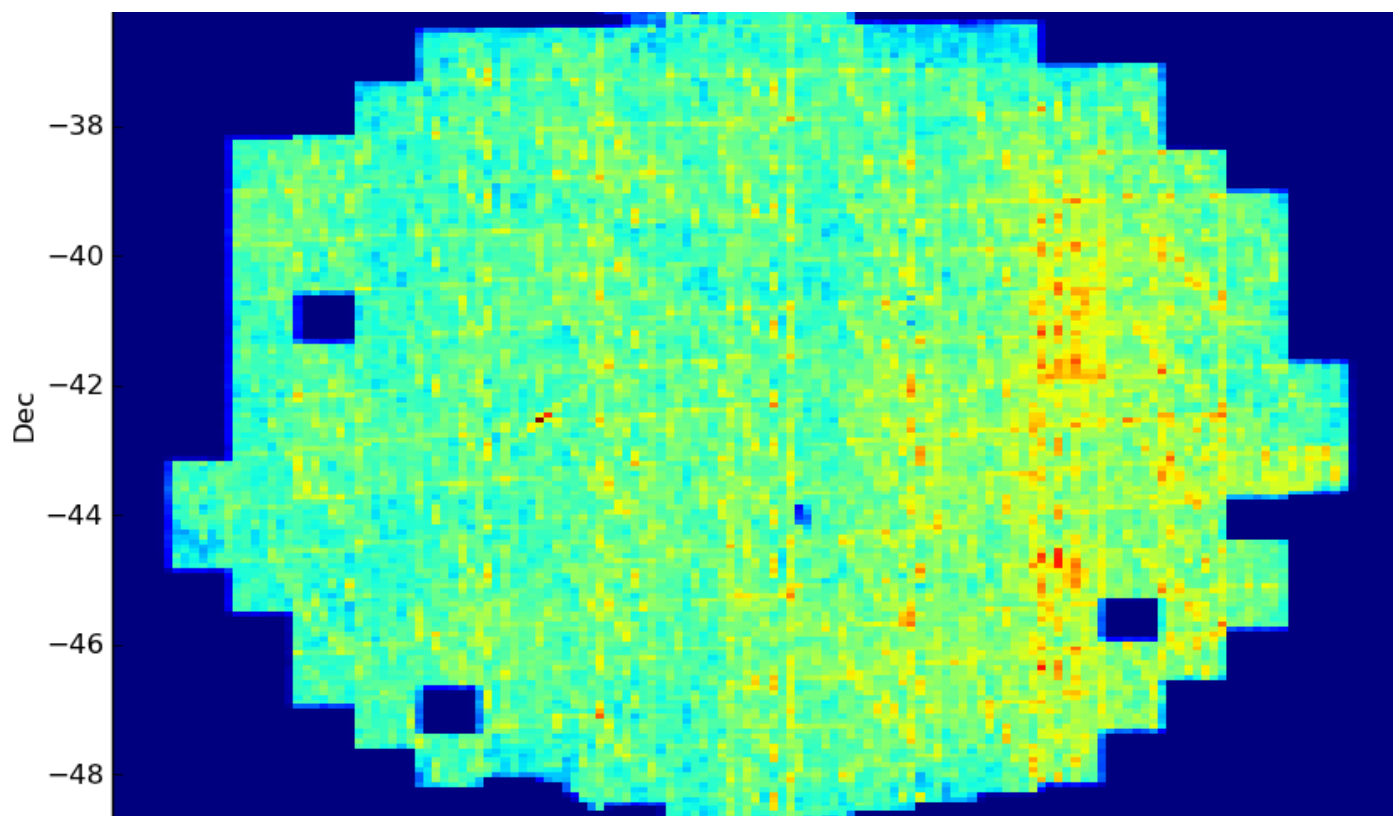


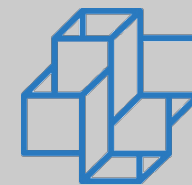
Preliminary Results

- Preliminary results are encouraging:
 - Baseline Orchestration layer (234 nodes) – approx. 46 min
 - 1 node HQOOP – approx. 35 min
 - 4 nodes HQOOP – approx. 12.3 min
 - 95 nodes (94 workers) HQOOP – approx. 2.10 min
 - 95 nodes (94 workers) Hadoop+Python – approx. 2.4 min



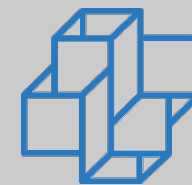
Resulting Image





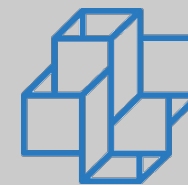
Conclusions

- Big data users (scientists) are in Big Trouble;
 - Too much data, too fast, too complex;
- Different expertise required to cooperate towards Big Data Management;
- Adapted software development methods based on workflows;
- Complete support to scientific exploration life-cycle
- Efficient workflow execution on Big Data



Collaborators

- LNCC Researchers
 - Ana Maria de C. Moura
 - Bruno R. Schulze
 - Antonio Tadeu Gomes
- PhD Students
 - Bernardo N. Gonçalves
 - Rocio Millagros
 - Douglas Ericson de Oliveira
 - Miguel Liroz-Gistau (INRIA)
 - Vinicius Pires (UFC)



Collaborators

- ON
 - Angelo Fausti
 - Luiz Nicolaci da Costa
 - Ricardo Ogando
- COPPE-UFRJ
 - Marta Mattoso
 - Jonas Dias (Phd Student)
 - Eduardo Ogasawara (CEFET-RJ)
- UFC
 - Vania Vidal
 - José Antonio F. de Macedo
- PUC-Rio
 - Marco Antonio Casanova
- INRIA-Montpellier
 - Patrick Valduriez group
- EPFL
 - Stefano Spaccapietra

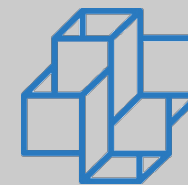
EMC Summer School on BIG DATA – NCE/UFRJ

Big Data in Astronomy

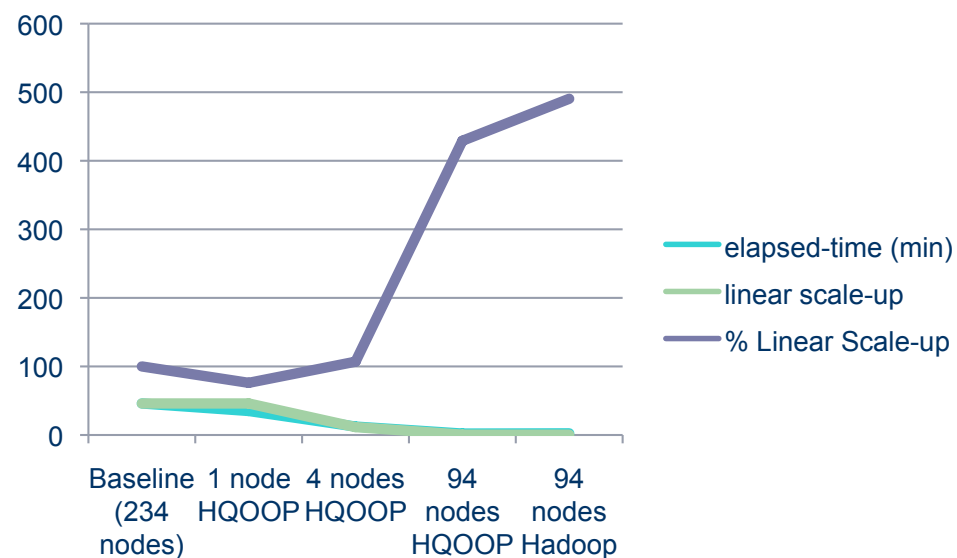
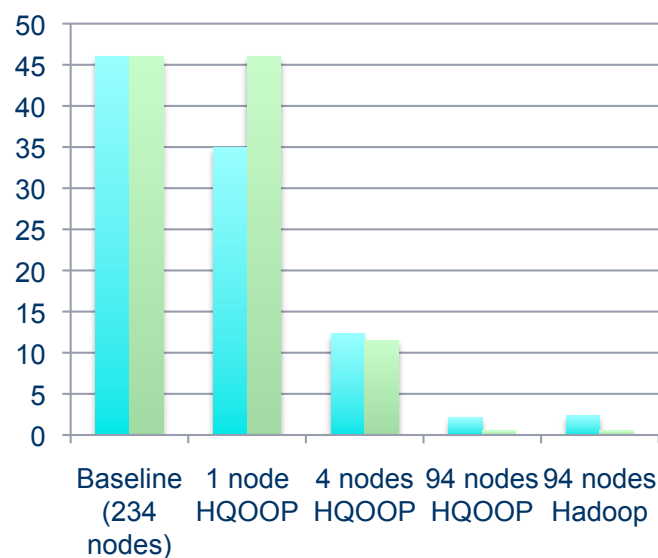
Fabio Porto (fporto@lncc.br)

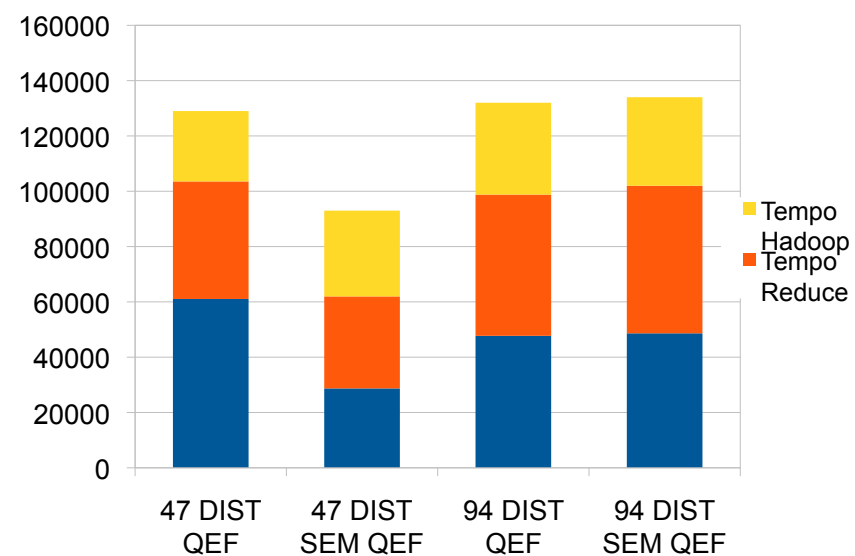
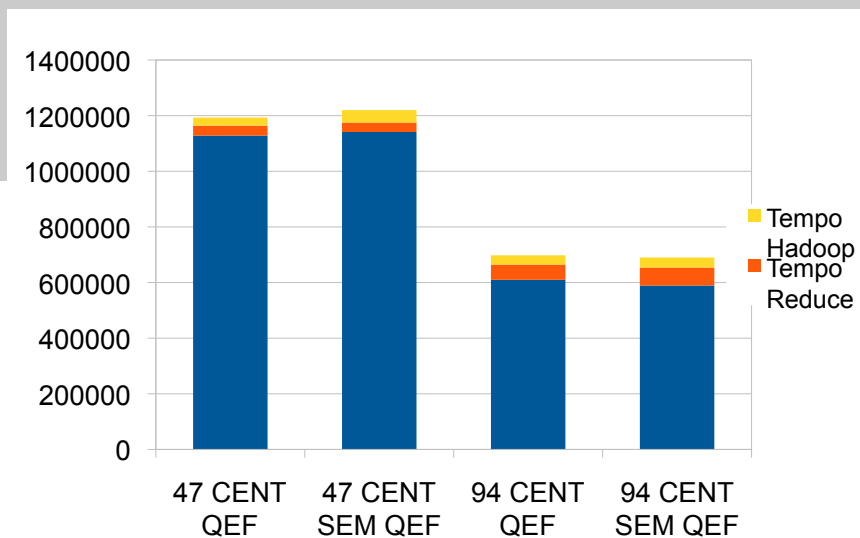
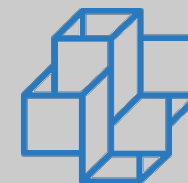
LNCC – MCTI

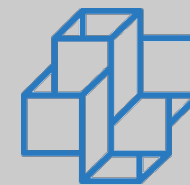
DEXL Lab (dexl.lncc.br)



Overall performance

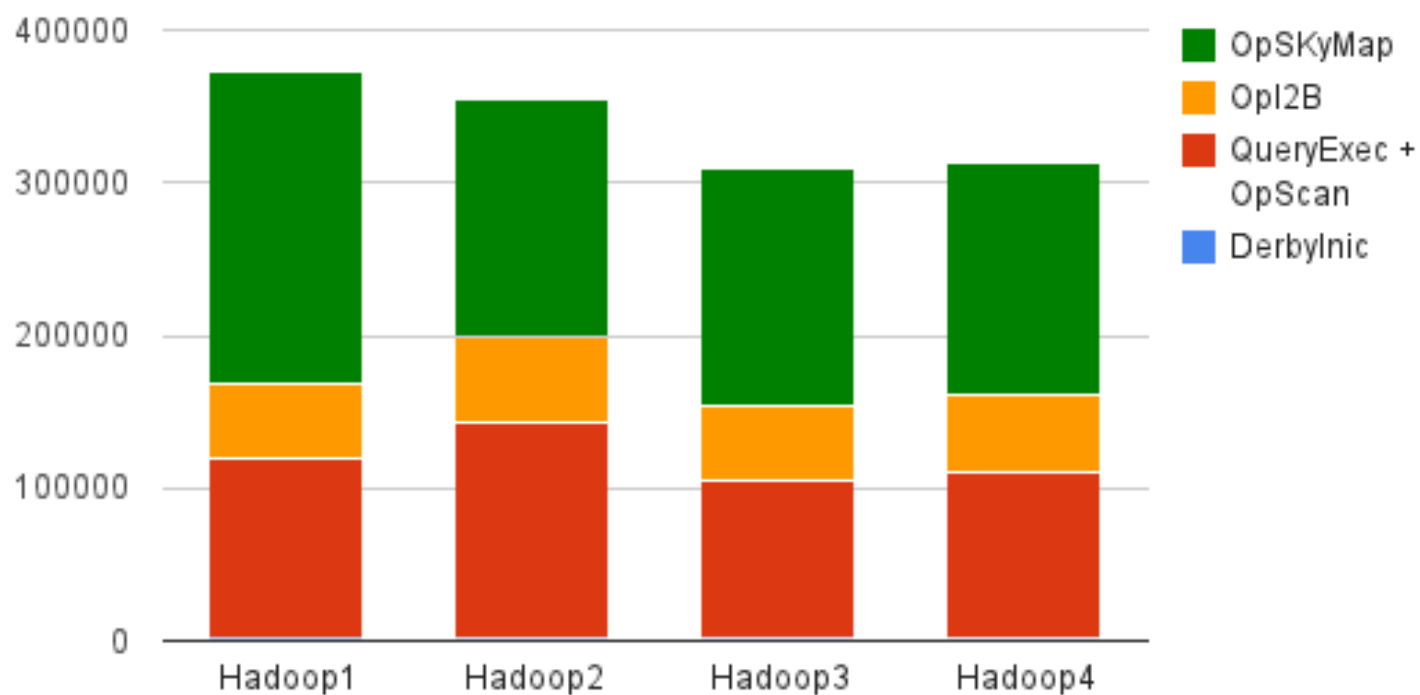




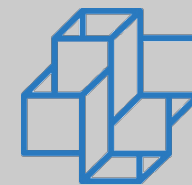


Execution with 4 nodes

Elapsed-time total: 11.27 min

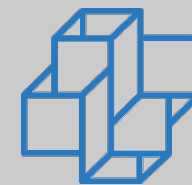






Adaptive and Extensible Query Engine

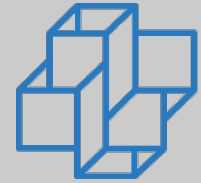
- Extensible to **data types**
- Extensible to **application algebra**
- Extensible to **execution model**
- Extensible to **heterogeneous data sources**



Objective

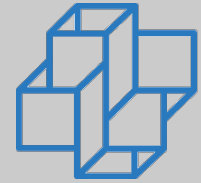
- Offer a query processing framework that can be extended to adapt to data centric application needs;
- Offer transparency in using resources to answer queries;
 - Query optimization transparently introduced
 - Standardize remote communication using web services even when dealing with large amount of unstructured data
 - Run-time performance monitoring and decision

Control Operators

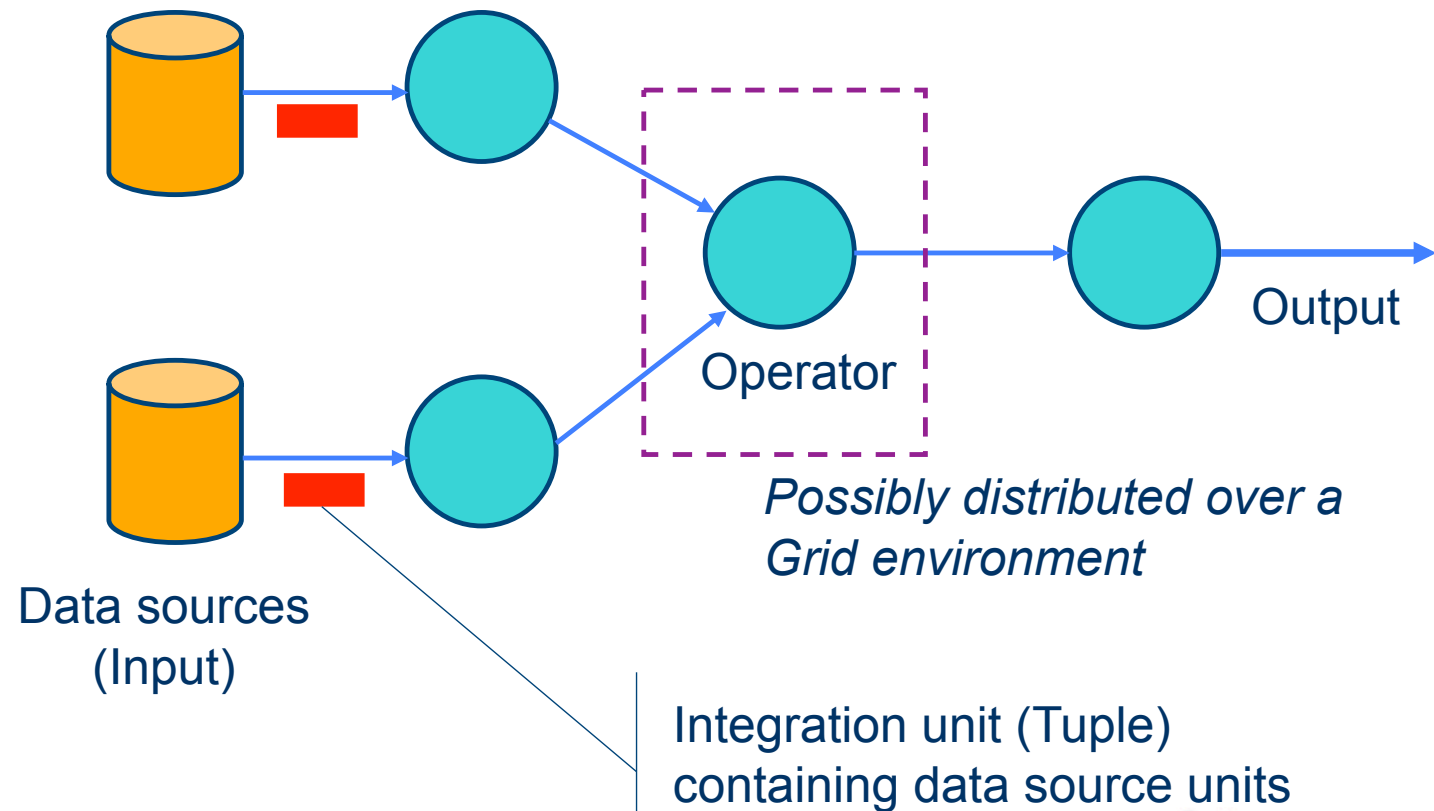


- Add data-flow and transformation operators
- Isolate application oriented operators from execution model data-flow concerns
- parallel grid based execution model:
 - **Split/Merge** - controls the routing of tuples to parallel nodes and the corresponding unification of multiple routes to a single flow
 - **Send/Receive** - marshalling/ unmarshalling of tuples and interface with communication mechanisms
 - **B2I/I2B** - blocks and unblocks tuples
 - **Orbit** - implements loop in a data-flow
 - **Fold/Unfold** - logical serialization of complex structures (e.g. PointList to Points)

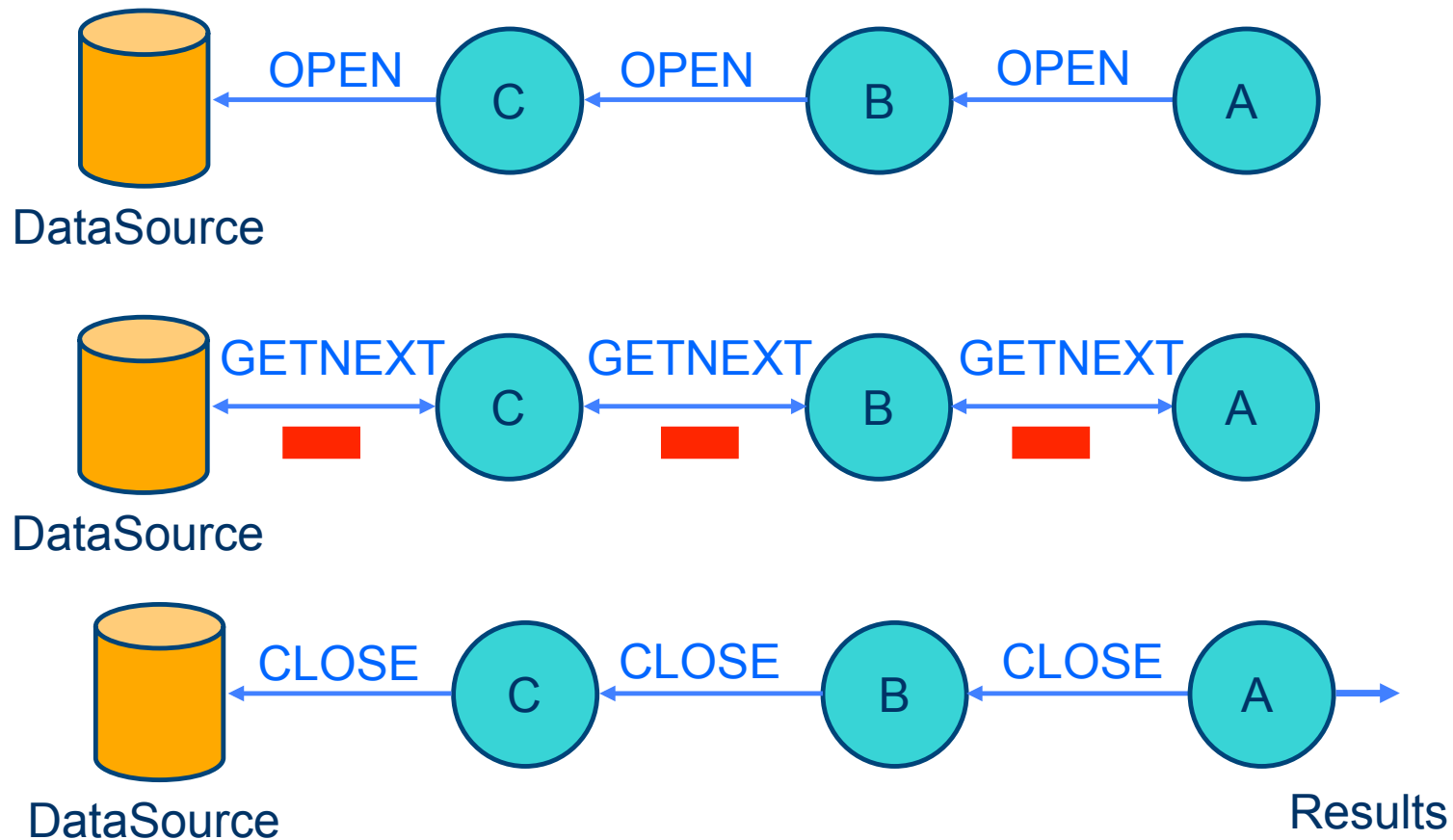
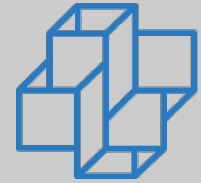
The Execution Model



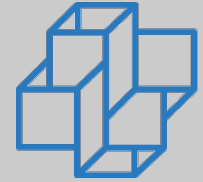
Example of simple QEF Workflow



Iteration Model

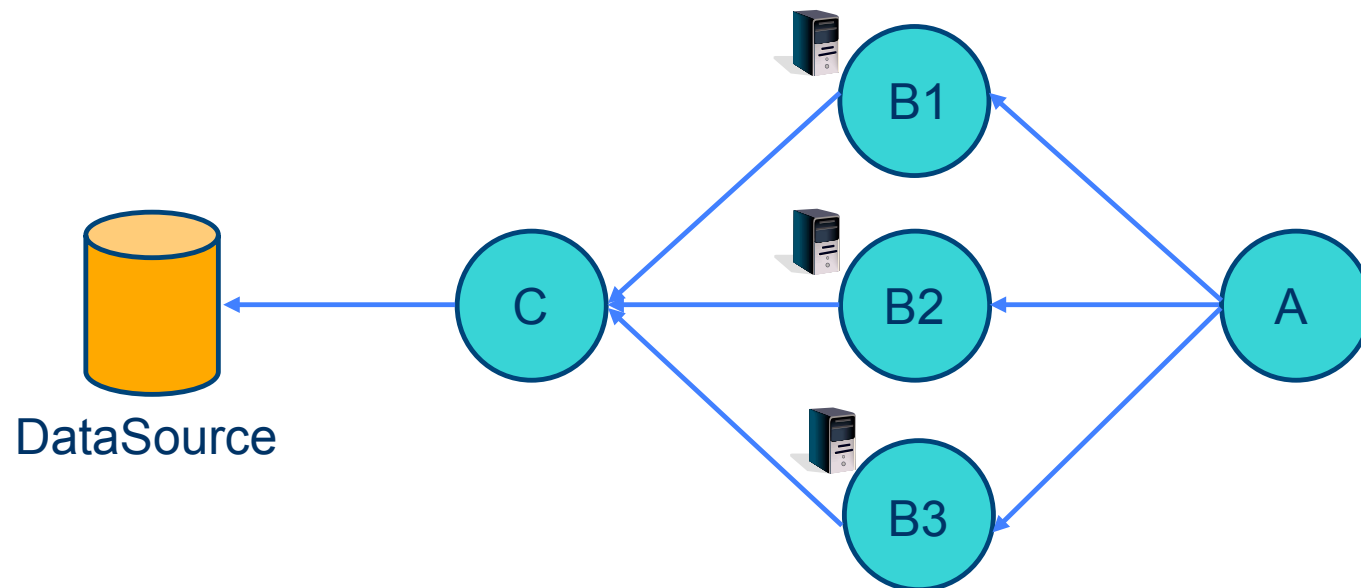


Distribution and Parallelization

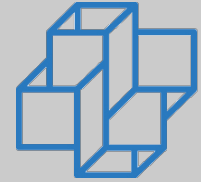


Operator distribution

A Query Optimizer selects a set of operators in the QEP to execute over a Grid environment.

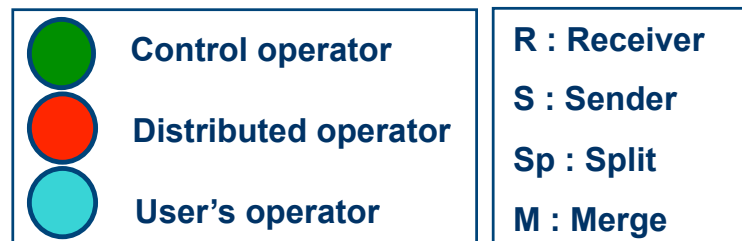
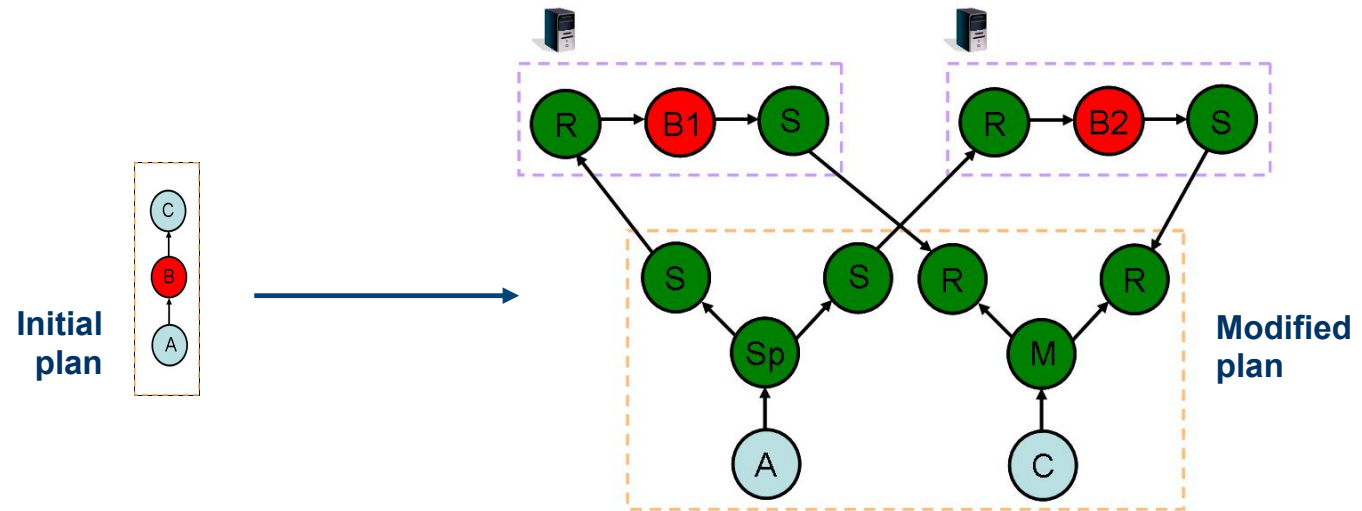


General Parallel Execution Model

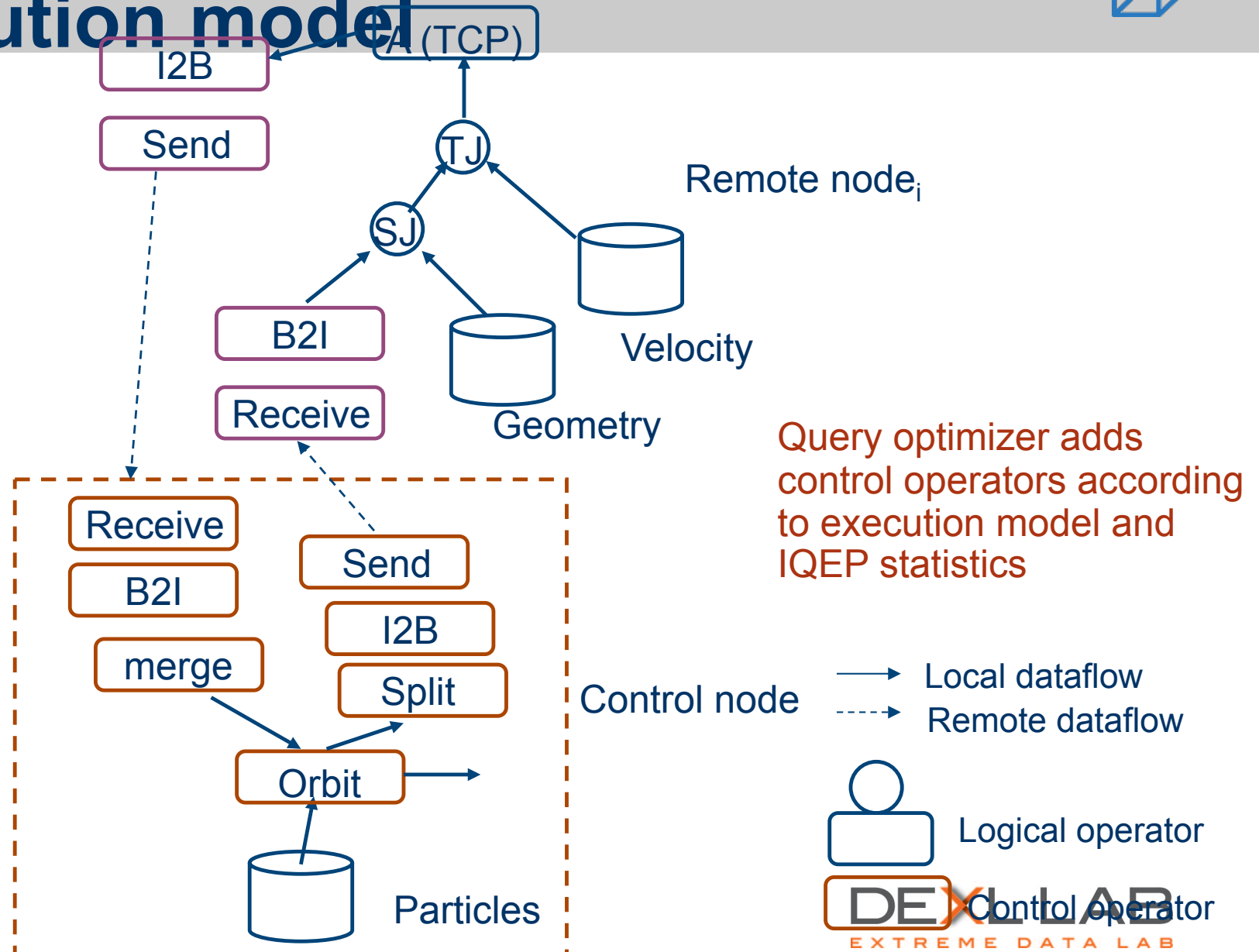
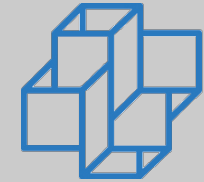


Remote QEP

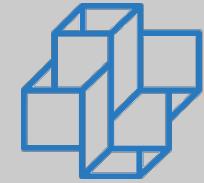
In order to parallelize an execution, the initial QEP is modified and sent to remote nodes to handle the distributed execution.



Modifying IQEP to adapt to execution model



Grid node allocation algorithm (G2N)



Introduction

Grid Greedy Node scheduling algorithm (G2N)

Principles

- Offers maximum usage of scheduled resources during query evaluation.

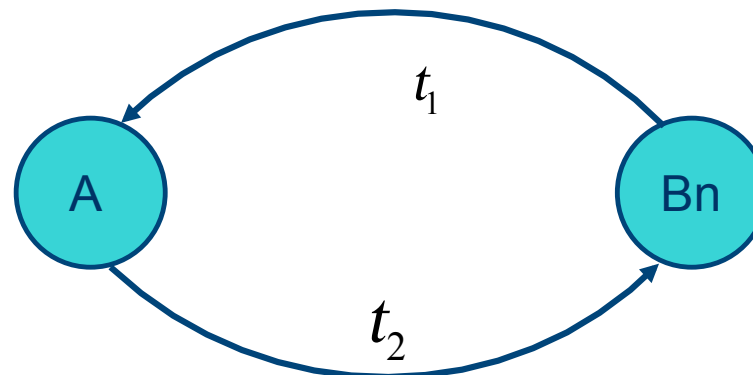
Application

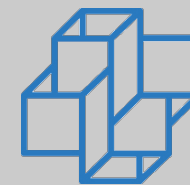
- Basic idea : “an optimal parallel allocation strategy for an independent query operator ... is the one in which the computed elapsed-time of its execution is as close as possible to the maximum sequential time in each node evaluating an instance of the operator”.

Architecture

Implem.

Conclusion

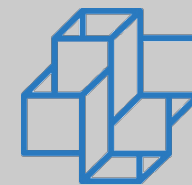




Implementation

- Core development in Java 1.5.
- Globus toolkit 4.
- Derby DBMS (catalog).
- Tomcat, AJAX and Google Web Toolkit for user interface.
- Runs on Windows, Unix and Linux.
- source code, demo, user guide available at:

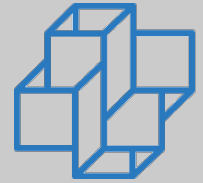
<http://dexl.incc.br>



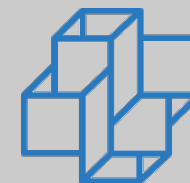
Summing-up

- HadoopDB extends Hadoop with expressive query language, supported by DBMSs
- Keeps Hadoop MapReduce framework
- Queries are mapped to MapReduce tasks
- For scientific applications is a question to be answered whether or not scientists will enjoy writing SQL queries
- Algebraic like languages may seem more natural (eg. Pig Latin)

Pig Latin - an high-level language alternative to SQL

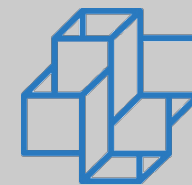


- The use of high-level languages such as SQL may not please scientific community;
- Pig Latin tries to give an answer by providing a procedural language where primitives are Relational algebra operations;
- Pig Latin: A not-so-foreign language for data processing, Christopher Olson, Benjamin Reed et al., SIGMOD08;



Example

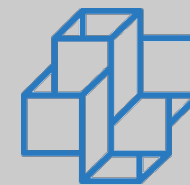
- Urls (url, category, pagerank)
- In SQL
 - Select category, avg (pagerank)
from urls where pagerank > 0.2
group by category
having count(*) > 10⁶
- In PIG
 - Groupurls = FILTER urls by Pagerank > 0.2;
 - Groups= Group good-urls by category;
 - Big-group=FILTER groups BY count(good_urls) > 10⁶
 - Output = FOREACH big-groups GENERATE
category, avg(good_urls_pagerank);



Pig Latin

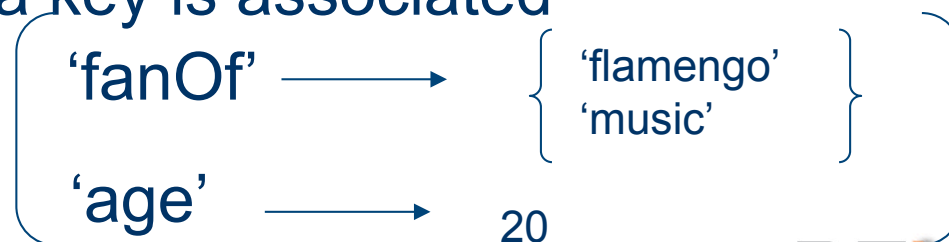
- Program is a sequence of steps
 - Each step executes one data transformation
- Optimizations among steps can be dynamically generated, example:
 - 1) spam-urls= FILTER urls BY isSpam(url);
 - 2) Highrankurl = FILTER spam-url BY pagerank > 0.8;



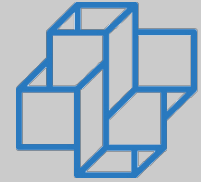


Data Model

- Types:
 - Atom - a single atomic value;
 - Tuple - a sequence of fields, eg.('DB','Science',7)
 - Bag - a collection of tuples with possible duplicates;
 - Map - a collection of data items where for each data item a key is associated



Operations



- Per tuple processing: Foreach
 - Allows the specification of iterations over bags
 - Ex:
 - Expanded-queries=FOREACH queries generate `userId`, `expandedQuery (queryString)`;
 - Each tuple in a bag should be independent of all others, so parallelization is possible;
 - Flatten
 - Permits flattening of nested-tuples

$$\left(\text{alice}, \begin{Bmatrix} \{ \text{Ipod, nano} \} \\ \{ \text{Ipod, shuffle} \} \end{Bmatrix} \right) \xrightarrow{\text{flatten}} \begin{matrix} [\text{alice}, \text{ipod}, \text{nano}] \\ [\text{alice}, \text{ipod}, \text{shuffle}] \end{matrix}$$

Olympic Laboratory

