

Pôle 3 / RUNTIME

Action d'envergure C2S@Exa

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EPI RUNTIME
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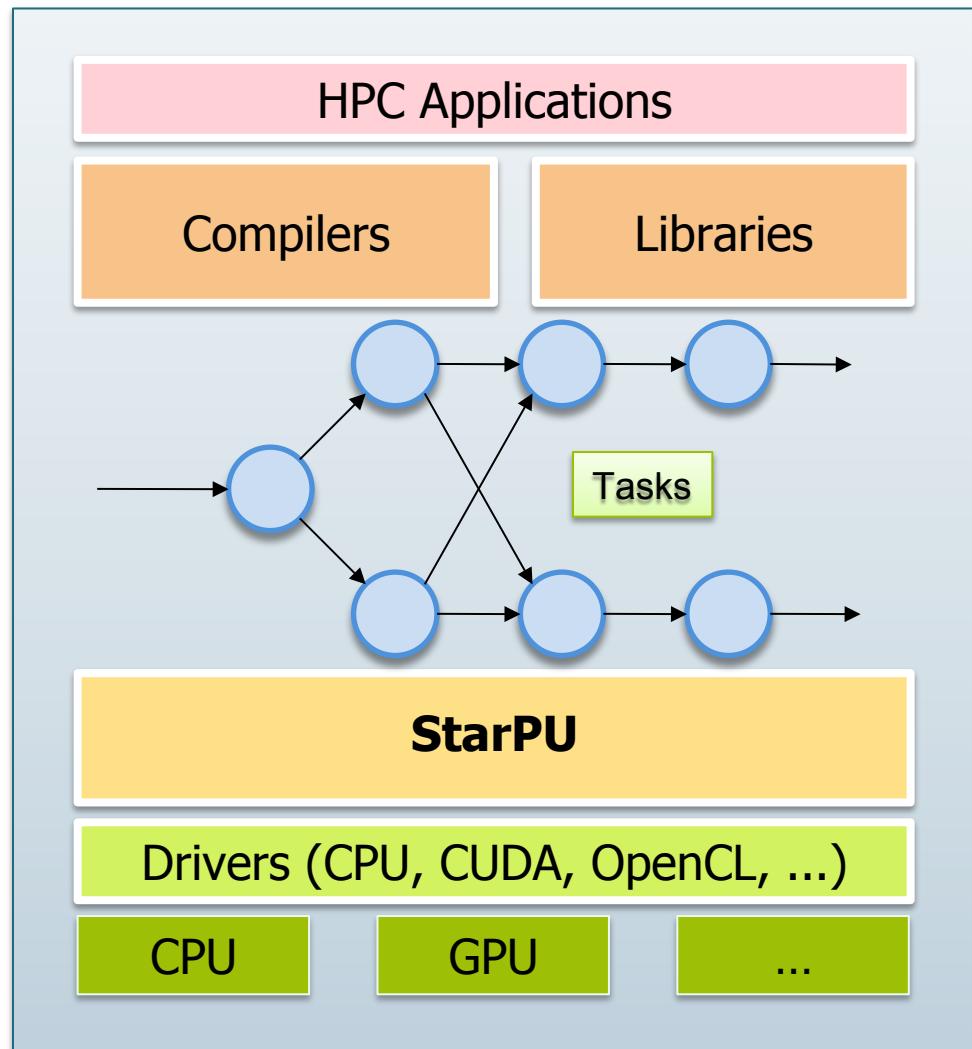
Sophia
12/2013



Overview of StarPU

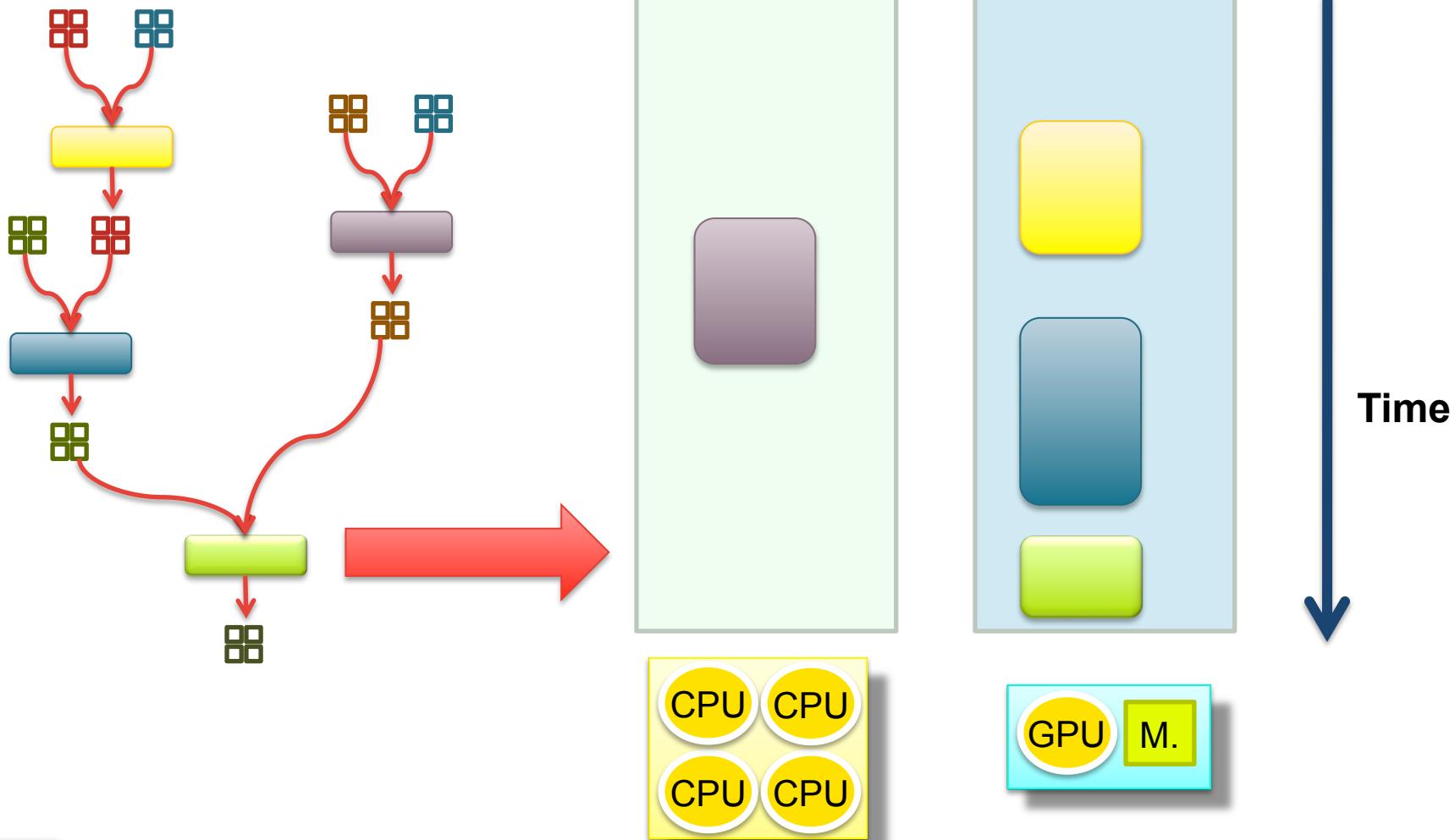
Information needed from applications

- Tasks
 - Implementation(s)
 - CPU
 - Regular
 - MMX, SSE, AVX, ...
 - Cuda, OpenCL
- Data
 - Type, layout
 - Vector, matrix, ...
 - Partitioning
- Task/Data Relationships
 - Dependencies
 - R, W, R/W, reduction, ...
- (Optional) Custom Scheduler
 - Open Scheduling Platform



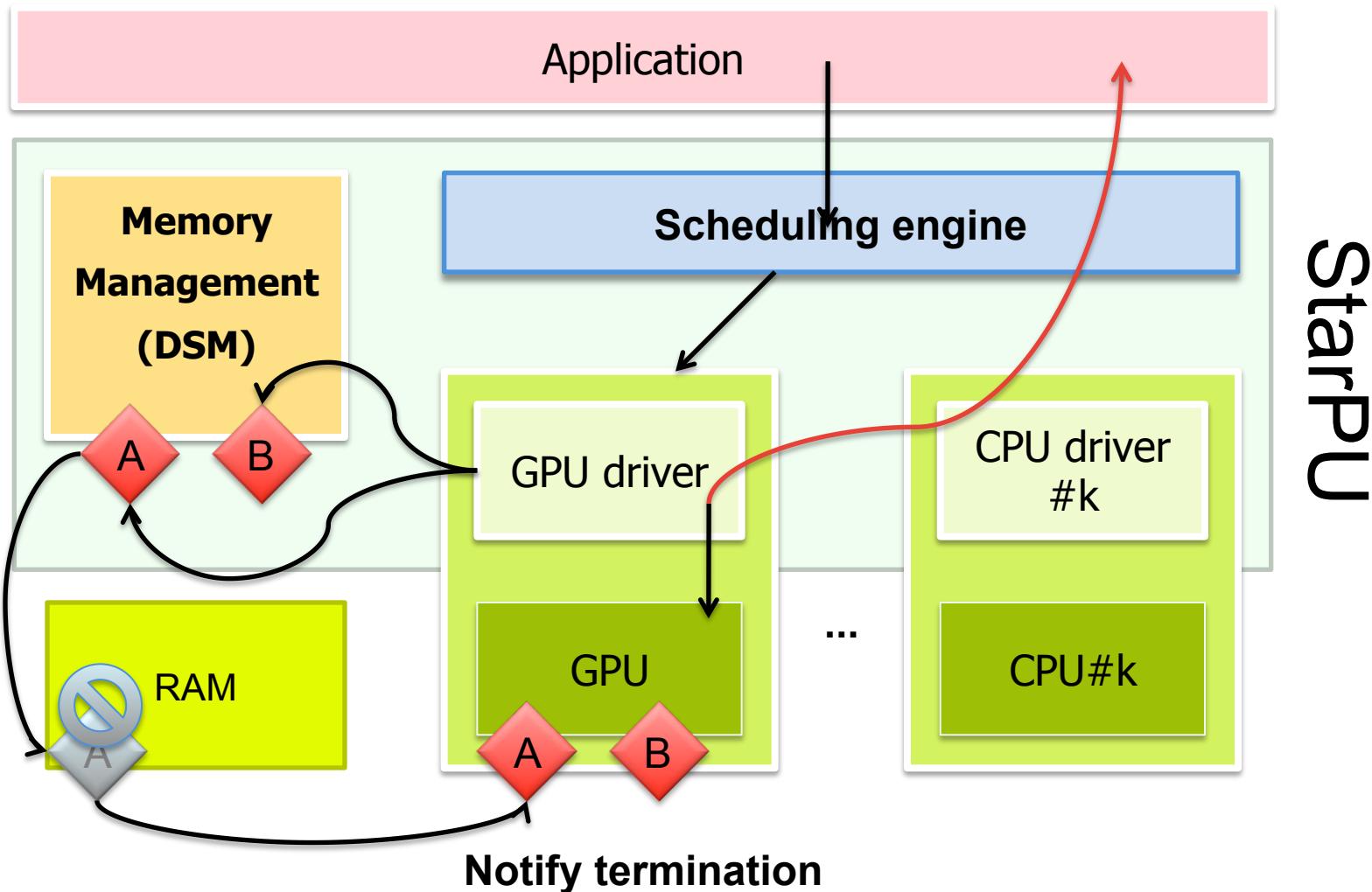
Heterogeneous Task Scheduling

Mapping a graph of tasks (DAG) on a heterogeneous platform



The StarPU runtime system

Execution model



Recent and On-going Developments

StarPU – Task Scheduling on Heterogeneous platforms

- StarPU / MPI cooperation
 - Parallel vs distributed relationship support
- Out of core
 - Disks as memory nodes
- Support for Manycores
 - Intel Xeon PHI, Intel SCC
- StarPU as an OpenCL backend
 - Use StarPU from OpenCL API
- Composition
 - Scheduling contexts

StarPU interfacing with MPI

- Automatically handles task/communications dependencies
 - Automatically handles all needed CPU/GPU transfers
- Automatically overlaps MPI communications, CPU/GPU communications, and CPU/GPU computations
 - Thanks to the data transfer requests mechanism
- StarPU/MPI support could be useful to XcalableMP/StarPU port in simplifying the interface between node-local management and distributed management
- People
 - Samuel Thibault, Nathalie Furmento, Olivier Aumage (permanents)
 - Marc Sergent (Master M2 completed, now PhD 1st year)

Out of core

Using disks as StarPU memory nodes

- Large workloads
 - Enable StarPU to evict temporarily unused data to disk
- Integration with general StarPU's memory management layer
 - StarPU data handles
- Multiple disk drivers supported
 - Legacy stdio/unistd methods
 - Google's LevelDB (key/value database library)
- Readily available in StarPU's SVN repository for the FP3C project
- People
 - Samuel Thibault (permanent)
 - M1 internship completed, proposal for a M2 internship

Support for Manycores

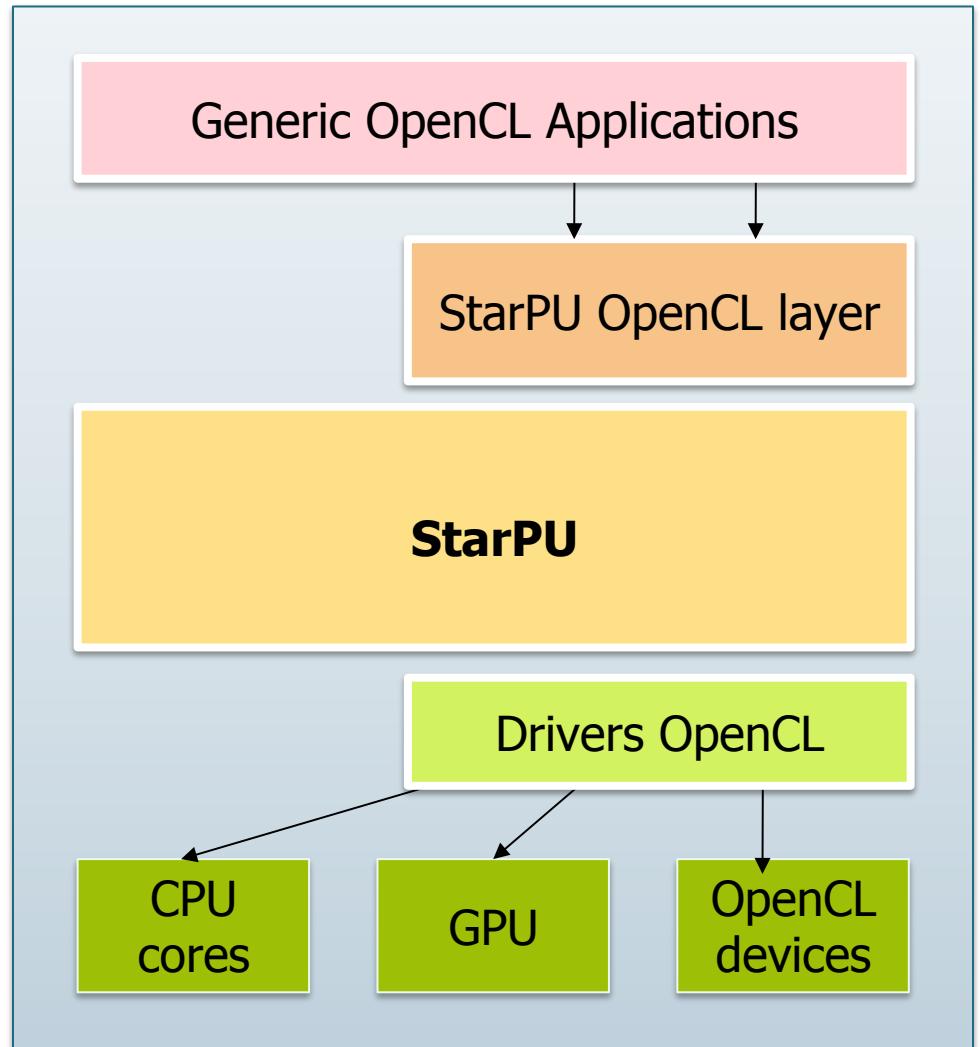
Intel Xeon Phi & SCC, Kalray MPPA

- Common support for shared aspects of Xeon Phi and SCC
 - Source = the main CPU node
 - Sink = the manycore device node
- Scheduling performed by the main StarPU instance
- Tested successfully on University of Tokyo's KNCC cluster of Intel MICs, within the context of FP3C project
- People
 - Samuel Thibault
 - Completed M1 internships (Intel Xeon Phi, SCC, and also Kalray MPPA)

StarPU as backend for OpenCL

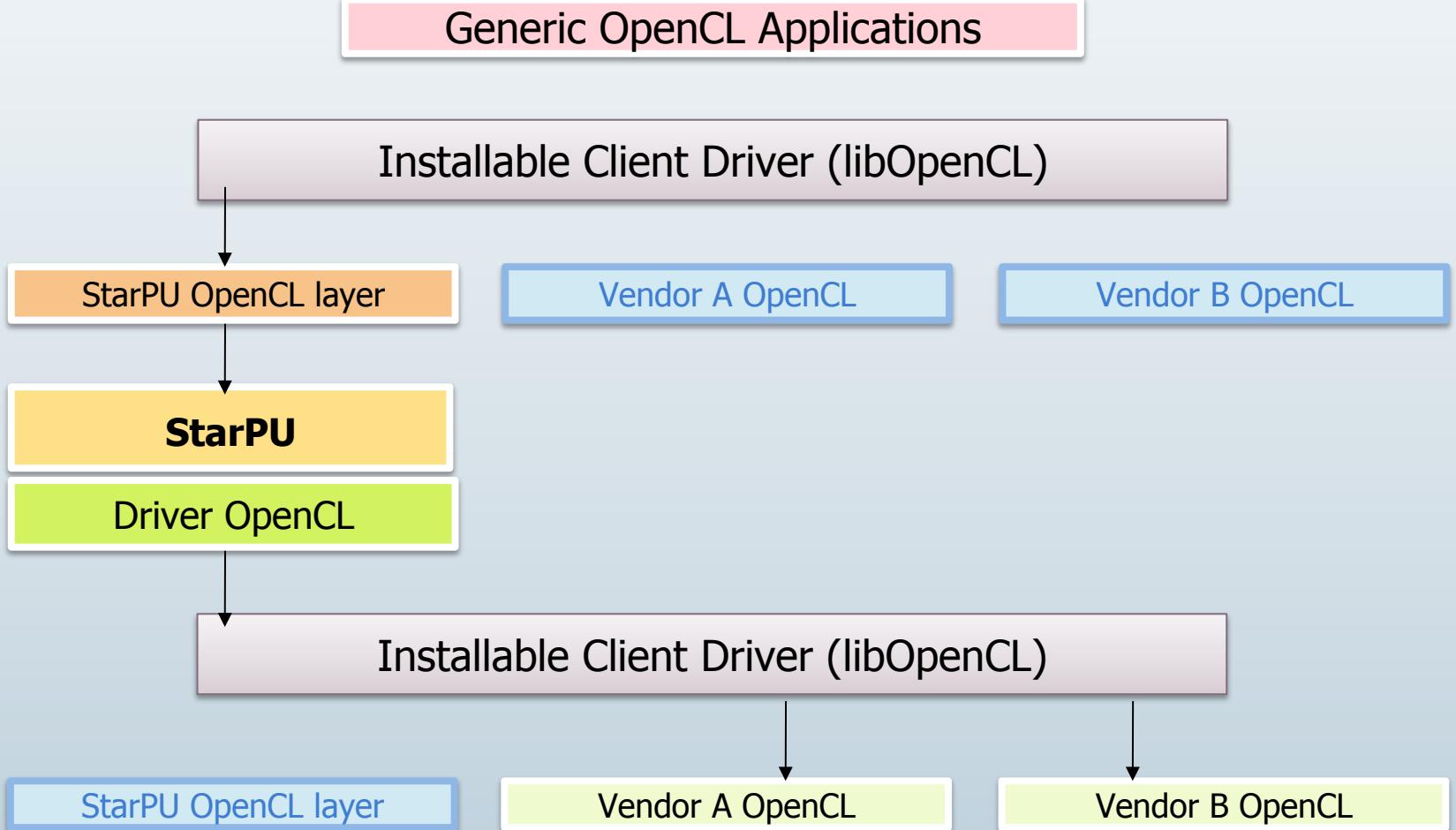
Portably programming StarPU through OpenCL

- Run generic OpenCL codes on top of StarPU



Built on OpenCL's ICD extension (`cl_khr_icd`)

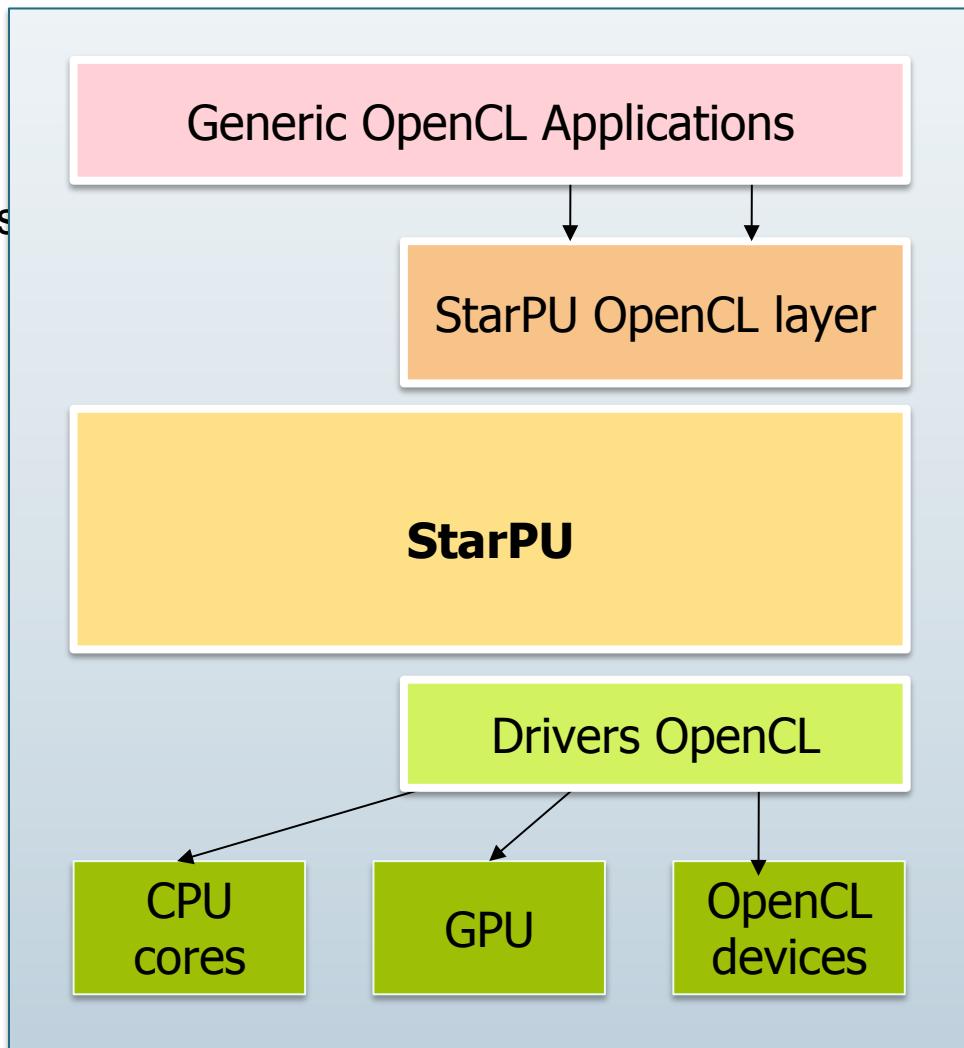
ICD = Installable Client Driver



StarPU as backend for OpenCL

Portably programming StarPU through OpenCL

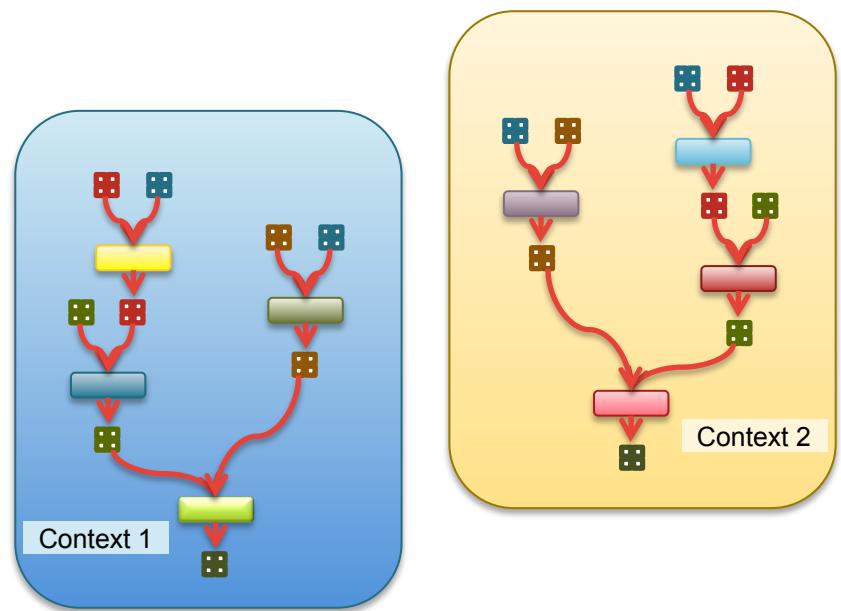
- Run generic OpenCL codes on top of StarPU
- Events can be used between all devices
 - Synchronizations
- Buffers can be shared by all devices
 - Data transfers handled by StarPU
- Compatibility with StarPU's contexts
 - Dedicated command queues
 - Dedicated schedulers
- People
 - Sylvain Henry, PhD completed Nov, 2013



Composition

Scheduling Contexts

- Managing computing resources for multiple kernels simultaneously
 - Map kernels on subsets of computing units
 - Isolate competing kernels or library calls
 - OpenMP kernel, Intel MKL, etc.
 - Select scheduling policy per context
- Two flavors
 - Explicit worker set mapping
 - Initial worker set
 - Worker add/remove ops
 - Managed worker set
 - Hypervisor
- Compatible with StarPU as OpenCL backend
- Can be used to interoperate XcalableMP codes with existing OpenCL codes
- People
 - Raymond Namyst, Pierre-André Wacrenier, Abdou Guermouche
 - CDD: Andra Hugo, Phd. Thesis (starting 3rd year)



Actions

ADT conjointe MOAIS (Pole 4) – RUNTIME (Pole 3)

- ADT « K'Star »
 - Portée par Thierry Gautier, EPI MOAIS
 - Inria Grenoble (Thierry Gautier) + Inria Bordeaux (Olivier Aumage, Samuel Thibault, Nathalie Furmento)
- Compilation *source-to-source* d'un langage standard...
 - OpenMP 4.0
- ... vers des routines XKaapi ou StarPU

- Objectif pour l'AE HPC
 - Faciliter l'adaptation portable des codes au-dessus des supports exécutifs
- Objectif pour XKaapi & StarPU
 - Simplifier l'utilisation
 - Augmenter la visibilité
- CDD
 - Pierrick Brunet, ingénieur (2 ans) à Montbonnot
 - + 1 ingénieur IJD (1an) à recruter à Bordeaux lors de la campagne 2014

Actions

ANR SOLHAR (Pole 1 + Pole 3)

- ANR « SOLHAR »
 - Portée par Abdou Guermouche, EPI HIEPACS
 - Inria Bordeaux (HIEPACS, RUNTIME), Inria Rhone-Alpes (ROMA)
 - CNRS, EADS, CEA
- Solveurs directs en algèbre linéaire creuse
- ... au dessus du support exécutif **StarPU**
- Intérêt pour l'AE HPC
 - Mise en œuvre directe du modèle de l'AE
 - Runtime + Solveur + Méthode numérique + Application
- CDD, 1 nouveau doctorant sur la thématique de l'adaptation de granularité des tâches
- Projet Hi-BOX (DGA), présenté par Guillaume Sylvand
 - EADS, IMACS, INRIA HIEPACS & RUNTIME
- Projet connexe ANR ANEMOS
 - INRIA HIEPACS (Pierre Ramet, Xavier Lacoste), CEA (Guillaume Latu)
- Nous sommes intéressés pour distribuer les tâches de la nouvelle version « light » de Gysela sur CPU+MIC avec StarPU !
- Voir également le logiciel Hwloc pour les questions d'affinité threads/mémoire

En résumé, au niveau de StarPU

2 axes de travail

- Faciliter l'adaptation de codes au-dessus de StarPU
 - Programmation par OpenCL et OpenMP 4.0 : avec pôle 4
 - Coopérations avec les solveurs (algèbre linéaire) : avec pôle 1
- Apporter de nouvelles fonctionnalités
 - Ordonnancement de tâches sur machines distribuées et many-cores
 - exploiter des clusters hétérogènes
 - Out-of-core
 - travailler sur de grands volumes de données
 - Composition
 - associer des codes ou des kernels parallèles
 - Granularité
 - adapter la charge de calcul demandée aux ressources de calcul disponibles
- Interlocuteurs au niveau de RUNTIME
 - Raymond Namyst, Olivier Aumage

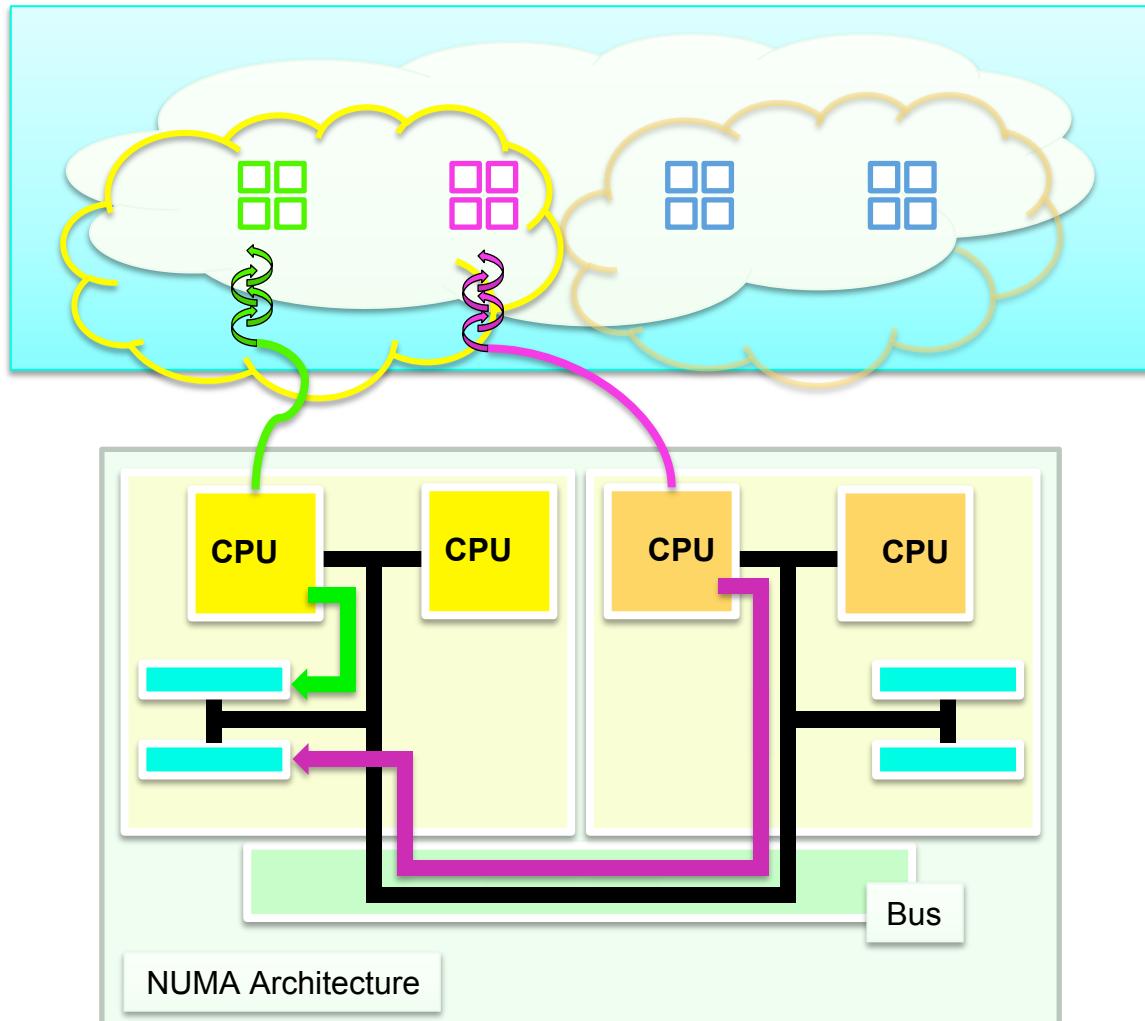
Hwloc – Hardware Locality

**Discovering Hardware Topology
Binding Computing Resources**



Olivier Aumage
RUNTIME group

Why is thread binding needed?



Why thread binding may be difficult?

- Lack of generic, uniform interface
 - OS specific
 - Distribution specific
 - Low-level APIs
- Undocumented, non deterministic behaviour
 - Resource identification scheme
 - Identification stability over time
 - BIOS dependent settings
 - Reboot dependent settings
- Evolving technology
 - E.g.: AMD Bulldozer “half-cores”, Intel SCC “tiles” (2 cores, Hz) & groups (8 cores, V)
- Need generic tools & abstract API
 - Logical resource identification

- Brice Goglin, Samuel Thibault, EPI RUNTIME
- Joint development
 - Runtime Group + Open MPI/Cisco
 - libtopology
 - PLPA
- Two parts
 - Set of command line tools
 - lstopo
 - hwloc-bind, hwloc-calc, etc.
 - API + library
- <http://www.open-mpi.org/projects/hwloc/>

Hwloc – Lstopo

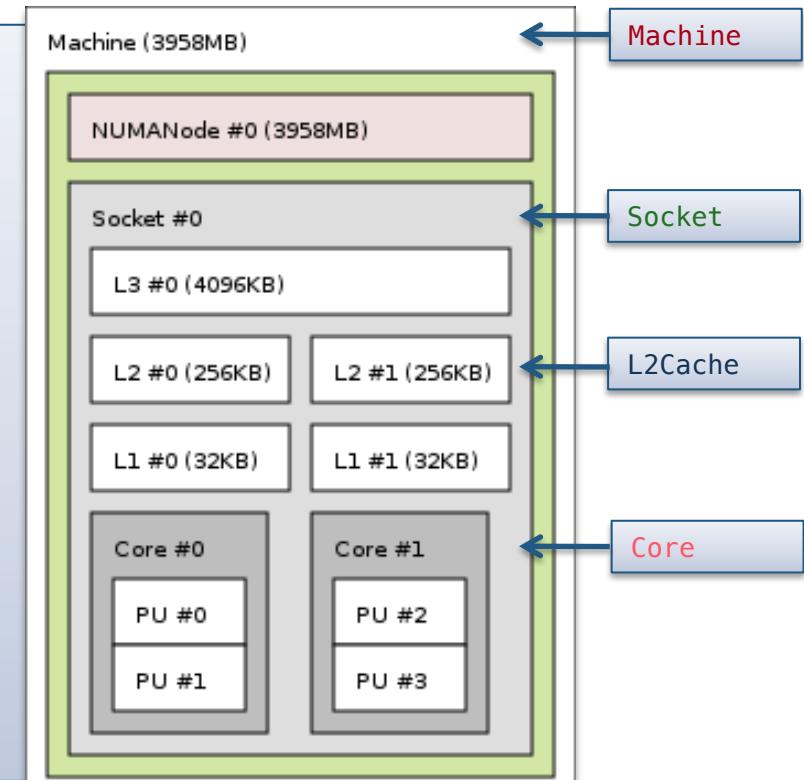
Displays the hierarchical topology map of the current system

```
Machine          (total=4052948KB)
NUMANode #0     (phys=0 local=4052948KB)

Socket #0       (phys=0)
L3Cache #0      (4096KB)

L2Cache #0      (256KB)
| L1Cache #0    (32KB)
| Core #0       (phys=0)
|   PU #0        (phys=0)
|   PU #1        (phys=2)

L2Cache #1      (256KB)
| L1Cache #1    (32KB)
| Core #1       (phys=2)
|   PU #2        (phys=1)
|   PU #3        (phys=3)
```

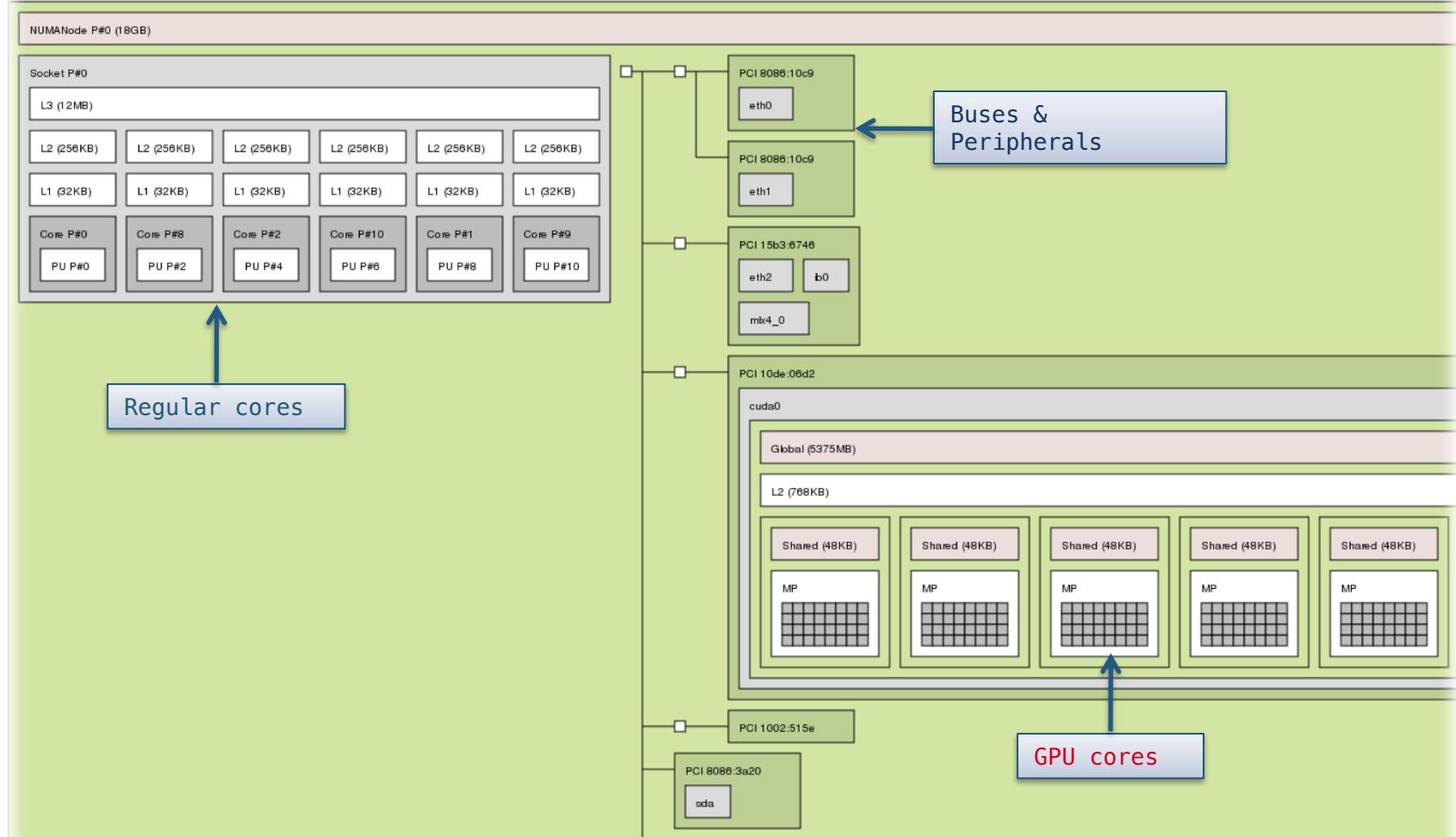


Textual representation

Graphical representation

Hwloc – Lstopo

Machine (36GB)



NUMANode P#1 (18GB)

Socket P#1

Hwloc – hwloc-*

Other command line tools

- hwloc-bind
 - Binds processes to specific hardware objects
- hwloc-calc
 - Creates bitmap strings to pass to hwloc-bind
- hwloc-distrib
 - Generates a set of uniformly distributed bitmap strings
- hwloc-ps
 - Displays the bindings of currently running processes
- hwloc-gather-topology
 - Saves the relevant topology files of the current machine
 - Allows offline simulation or debugging
 - Linux only

Hwloc – API

Programming interface, run-time library

- Topology information
 - CPU, caches, memory, NUMA nodes, I/O devices
 - Query, traversal
- CPU binding
 - Processes binding
 - Threads binding
- Memory binding
 - Get/set memory binding policy
 - Allocation of memory bound to specific nodes
- Bitmap API
 - Sets of resources
 - Common set operations